The inference rule

 $A \lor B \lor C \lor ... \lor Z$ $\mbox{\small$\tt D$}\ \vee\ \mbox{\small$\tt E$}\ \vee\ \mbox{\small$\tt F$}\ \vee\ \ldots\ \vee\ \neg \mbox{\small$\tt Z$}$ --------------------------- $A \lor B \lor C \lor D \lor E \lor F \lor \ldots$

Part of the original Prolog example:

```
eats(X, Y) : cat(X), mouse(Y).
chases(X, Y) :- eats(X, Y).
chases(X, Y) : - dog(X), cat(Y).
 cat(tom). 
 mouse(jerry).
```
Those are all reversed implications, e.g.

if cat(X) and mouse(Y) then eats(X, Y)

So what we've got is five things that are all true

 $(cat(X) \wedge mouse(Y) \Rightarrow eats(X, Y)) \wedge$ $\text{(eats(X, Y)} \Rightarrow \text{chases(X, Y)}) \land$ $(dog(X) \wedge cat(Y) \Rightarrow chases(X, Y)) \wedge$ $(cat(tom)) \wedge$ (mouse(jerry))

Propositional logic can't handle those sorts of formulae, so simplify:

```
(cx \wedge my \Rightarrow exy) \wedge(exy \Rightarrow chxy) \land(dx \wedge cy \Rightarrow chxy) \wedge(cx) \wedge(my)
```
And recall these equivalences:

 $A \implies Z = Z \lor \neg A$

so

 $A \wedge B \wedge C \Rightarrow Z = Z \vee \neg (A \wedge B \wedge C) = Z \vee \neg A \vee \neg B \vee \neg C$

Our rule base is therefore

```
(exy \lor \neg cx \lor \neg my) \land(\text{chxy} \lor \neg \text{exy}) \land(\text{chxy} \lor \neg \text{dx} \lor \neg \text{cy}) \land(cx) \wedge(my)
```
All nicely in Conjunctive Normal Form (CNF)

```
(exy \lor \neg cx \lor \neg my) \land(\text{chxy} \lor \neg \text{exy}) \land(\text{chxy} \lor \neg \text{dx} \lor \neg \text{cy}) \land(cx) \wedge(my)
```
Suppose we want to know whether chxy or not.

Resolve cx with exy \vee \neg cx \vee \neg my giving exy \vee \neg my Resolve that with chxy \vee \neg exy giving \neg my \vee chxy Resolve that with my giving chxy Q.E.D.

The absolute technical truth.

 Logic programming actually works through proof by contradiction, attempt to prove the negation of what you want, and hope to fail.

We wanted to prove chxy, so we and \neg chxy with our rule base:

```
(exy \lor \neg cx \lor \neg my) \land(\text{chxy} \lor \neg \text{exy}) \land(\text{chxy} \lor \neg \text{dx} \lor \neg \text{cy}) \land(cx) \wedge(my) \wedge(\neg \text{chxy})
```
 and let resolution proceed exactly as above. At the end, when we were left with chxy, there is now one more step.

Resolve that with $\neg \text{chxy}$, resulting in nothing at all, which is failure.

That means that (our knowledge base) \land ($\neg \text{chxy}$) is false.

Our knowledge base is the definition of truth in this program, so it must be \neg chxy that is false, so according to our knowledge base, chxy is true.

The only trick is knowing which two facts to apply resolution to at each step.