

*2·14. $\vdash \sim(\sim p) \supset p$

Dem.

$$\begin{aligned} & \left[\text{Perm } \frac{\sim\{\sim(\sim p)\}}{q} \right] \vdash : p \vee \sim\{\sim(\sim p)\} \supset \sim\{\sim(\sim p)\} \vee p \quad (1) \\ & [(1), *2·13, *1·11] \quad \vdash \sim\{\sim(\sim p)\} \vee p \\ & [(2).(*1·01)] \quad \vdash \sim(\sim p) \supset p \quad (2) \end{aligned}$$

*2·15. $\vdash : \sim p \supset q \supset \sim q \supset p$

Dem.

$$\begin{aligned} & \left[*2·05 \frac{\sim p, \sim(\sim q)}{p, r} \right] \vdash : q \supset \sim(\sim q) \supset : \sim p \supset q \supset \sim p \supset \sim(\sim q) \quad (1) \\ & \left[*2·12 \frac{q}{p} \right] \vdash : q \supset \sim(\sim q) \quad (2) \\ & [(1).(2).*1·11] \quad \vdash : \sim p \supset q \supset \sim p \supset \sim(\sim q) \quad (3) \\ & \left[*2·03 \frac{\sim p, \sim q}{p, q} \right] \vdash : \sim p \supset \sim(\sim q) \supset \sim q \supset \sim(p) \quad (4) \\ & \left[*2·05 \frac{\sim q, \sim(\sim p), p}{p, q, r} \right] \vdash : \sim(\sim p) \supset p \supset : \sim q \supset \sim(\sim p) \supset \sim q \supset p \quad (5) \\ & [(5).*2·14, *1·11] \quad \vdash : \sim q \supset \sim(\sim p) \supset \sim q \supset p \quad (6) \\ & \left[*2·05 \frac{\sim p \supset q, \sim p \supset \sim(\sim q), \sim q \supset \sim(\sim p)}{p, q, r} \right] \vdash : \sim p \supset \sim(\sim q) \supset \sim q \supset \sim(\sim p) : \vdots \quad (7) \\ & \quad \sim p \supset \sim(\sim q) \supset \sim q \supset \sim(\sim p) : \vdots \\ & \quad \sim p \supset q \supset \sim p \supset \sim(\sim q) \supset : \sim p \supset q \supset \sim q \supset \sim(p) \quad (8) \\ & [(4).(7).*1·11] \vdash : \sim p \supset q \supset \sim p \supset \sim(\sim q) : \vdots \\ & \quad \sim p \supset q \supset \sim q \supset \sim(\sim p) \quad (9) \\ & [(3).(8).*1·11] \vdash : \sim p \supset q \supset \sim q \supset \sim(\sim p) \quad (10) \\ & \left[*2·05 \frac{\sim p \supset q, \sim q \supset \sim(\sim p), \sim q \supset p}{p, q, r} \right] \vdash : \sim q \supset \sim(\sim p) \supset \sim q \supset p \quad (11) \\ & \quad \sim q \supset \sim(\sim p) \supset \sim p \supset \sim q \supset p \quad (12) \\ & [(9).(11).*1·11] \vdash : \sim p \supset q \supset \sim q \supset p \quad (13) \end{aligned}$$

*Note on the proof of *2·15.* In the above proof, it will be seen that (3), (4), (6) are respectively of the forms $p_1 \supset p_2$, $p_2 \supset p_3$, $p_3 \supset p_4$, where $p_i \supset p_j$ is the proposition to be proved. From $p_1 \supset p_2$, $p_2 \supset p_3$, $p_3 \supset p_4$ the proposition $p_1 \supset p_4$ results by repeated applications of *2·05 or *2·06 (both of which are called "Syll"). It is tedious and unnecessary to repeat this process every time it is used; it will therefore be abbreviated into

"[Syll] $\vdash : (a) \cdot (b) \cdot (c) \supset \vdash : (d)$,"

where (a) is of the form $p_1 \supset p_2$, (b) of the form $p_2 \supset p_3$, (c) of the form $p_3 \supset p_4$, and (d) of the form $p_1 \supset p_4$. The same abbreviation will be applied to a series of any length.

Also where we have " $\vdash : p_1$ " and " $\vdash : p_1 \supset p_2$," and p_2 is the proposition to be proved, it is convenient to write simply

$$\begin{aligned} & \vdash : p_1 \supset \\ & \vdash : p_2, \end{aligned}$$

[etc.] where "etc." will be a reference to the previous propositions in virtue of which the implication " $p_1 \supset p_2$ " holds. This form embodies the use of *1·11 or *1·1, and makes many proofs at once shorter and easier to follow. It is used in the first two lines of the following proof.

*2·16. $\vdash : p \supset q \supset \sim q \supset \sim p$

Dem.

$$\begin{aligned} & [*2·12] \quad \vdash : q \supset \sim(\sim q) \supset \\ & [*2·05] \quad \vdash : p \supset q \supset p \supset \sim(\sim q) \quad (1) \\ & [*2·03 \frac{\sim q}{q}] \quad \vdash : p \supset \sim(\sim q) \supset \sim q \supset \sim p \quad (2) \\ & [\text{Syll}] \quad \vdash : (1) \cdot (2) \supset : p \supset q \supset \sim q \supset \sim p \end{aligned}$$

Note. The proposition to be proved will be called "Prop," and when a proof ends, like that of *2·16, by an implication between asserted propositions, of which the consequent is the proposition to be proved, we shall write " $\vdash . \text{etc.} \supset \vdash . \text{Prop}$ ". Thus " $\supset \vdash . \text{Prop}$ " ends a proof, and more or less corresponds to "Q.E.D."

*2·17. $\vdash : \sim q \supset \sim p \supset p \supset q$

Dem.

$$\begin{aligned} & [*2·03 \frac{\sim q, p}{p, q}] \quad \vdash : \sim q \supset \sim p \supset p \supset \sim(\sim q) \quad (1) \\ & [*2·14] \quad \vdash : \sim(\sim q) \supset q : \supset \\ & [*2·05] \quad \vdash : p \supset \sim(\sim q) \supset p \supset q \quad (2) \\ & [\text{Syll}] \quad \vdash : (1) \cdot (2) \supset \vdash . \text{Prop} \end{aligned}$$

*2·15, *2·16 and *2·17 are forms of the principle of transposition, and will be all referred to as "Transp."

*2·18. $\vdash : \sim p \supset p \supset p$

Dem.

$$\begin{aligned} & [*2·12] \quad \vdash : p \supset \sim(\sim p) \supset \\ & [*2·05] \quad \vdash : \sim p \supset p \supset \sim p \supset \sim(\sim p) \quad (1) \\ & [*2·01 \frac{\sim p}{p}] \quad \vdash : \sim p \supset \sim(\sim p) \supset \sim(\sim p) \quad (2) \\ & [\text{Syll}] \quad \vdash : (1) \cdot (2) \supset : \sim p \supset p \supset \sim(\sim p) \quad (3) \\ & [*2·14] \quad \vdash : \sim(\sim p) \supset p \quad (4) \\ & [\text{Syll}] \quad \vdash : (3) \cdot (4) \supset \vdash . \text{Prop} \end{aligned}$$

This is the complement of the principle of the *reductio ad absurdum*. It

- *23·45. $\vdash : R \subseteq S . R \subseteq T . \supset . R \subseteq S \wedge T$
- *23·46. $\vdash : xRy . R \subseteq S . \supset . xSy$
- *23·47. $\vdash : R \subseteq T . \supset . R \wedge S \subseteq T$
- *23·48. $\vdash : R \subseteq S . \supset . R \wedge T \subseteq S \wedge T$
- *23·49. $\vdash : P \subseteq Q . R \subseteq S . \supset . P \wedge R \subseteq Q \wedge S$
- *23·5. $\vdash . R \wedge R = R$
- *23·51. $\vdash . R \wedge S = S \wedge R$
- *23·52. $\vdash . (R \wedge S) \wedge T = R \wedge (S \wedge T)$
- *23·53. $R \wedge S \wedge T = (R \wedge S) \wedge T$ Df
- *23·54. $\vdash : . R = S . \supset : R \subseteq T . \equiv . S \subseteq T$
- *23·55. $\vdash : . R = S . \supset : T \subseteq R . \equiv . T \subseteq S$
- *23·551. $\vdash : R = S . \supset . R \cup T = S \cup T$
- *23·56. $\vdash . R \cup R = R$
- *23·57. $\vdash . R \cup S = S \cup R$
- *23·58. $\vdash . R \subseteq R \cup S . S \subseteq R \cup S$
- *23·59. $\vdash : R \subseteq T . S \subseteq T . \equiv . R \cup S \subseteq T$
- *23·6. $\vdash : . x(R \cup S)y . \equiv : R \subseteq T . S \subseteq T . \supset_T . xTy$
- *23·61. $\vdash : R \subseteq S . \supset . R \subseteq S \cup T$
- *23·62. $\vdash : R \subseteq S . \equiv . R \cup S = S$
- *23·621. $\vdash : R \subseteq S . \equiv . R \wedge S = R$
- *23·63. $\vdash . R \cup (R \wedge S) = R$
- *23·631. $\vdash . R \wedge (R \cup S) = R$
- *23·632. $\vdash : R = S . \supset . R = R \wedge S$
- *23·633. $\vdash : R \subseteq S . \supset . R \cup T = (R \wedge S) \cup T$
- *23·64. $\vdash : . R \subseteq T . v . S \subseteq T : \supset . R \wedge S \subseteq T$
- *23·65. $\vdash : . R \subseteq S . v . R \subseteq T : \supset . R \subseteq S \cup T$
- *23·66. $\vdash : R \subseteq S . \supset . R \cup T \subseteq S \cup T$
- *23·68. $\vdash . (R \wedge S) \cup (R \wedge T) = R \wedge (S \cup T)$
- *23·69. $\vdash . (R \cup S) \wedge (R \cup T) = R \cup (S \wedge T)$
- *23·7. $\vdash . (R \cup S) \cup T = R \cup (S \cup T)$
- *23·71. $R \cup S \cup T = (R \cup S) \cup T$ Df
- *23·72. $\vdash : P \subseteq R . Q \subseteq S . \supset . P \cup Q \subseteq R \cup S$
- *23·73. $\vdash : P = R . Q = S . \supset . P \cup Q = R \cup S$
- *23·74. $\vdash : P \wedge Q \subseteq R . P \wedge R \subseteq Q . \equiv . P \wedge Q = P \wedge R$
- *23·8. $\vdash . \neg(\neg R) = R$
- *23·81. $\vdash : R \subseteq S . \equiv . \neg S \subseteq \neg R$
- *23·811. $\vdash : R \subseteq S . \equiv . S \subseteq \neg R$
- *23·82. $\vdash : R \wedge S \subseteq T . \equiv . R \wedge T \subseteq S$
- *23·83. $\vdash : R = S . \equiv . \neg R = \neg S$

- *23·831. $\vdash : R = \neg S . \equiv . S = \neg R$
- *23·84. $\vdash . \neg(R \wedge S) = \neg R \vee \neg S$
- *23·85. $\vdash . R \wedge S = \neg(\neg R \vee \neg S)$
- *23·86. $\vdash . \neg(\neg R \wedge \neg S) = R \vee S$
- *23·87. $\vdash . \neg R \wedge \neg S = \neg(R \vee S)$
- *23·88. $\vdash . (x, y) . x(R \vee \neg R)y$
- *23·89. $\vdash . (x, y) . \sim\{x(R \vee \neg R)y\}$
- *23·9. $\vdash . (R \vee S) \neg S = R \neg S$
- *23·91. $\vdash . R \vee S = R \vee (S \neg R)$
- *23·92. $\vdash : R \subseteq S . \supset . S = R \vee (S \neg R)$
- *23·93. $\vdash . R \neg S = R \neg(R \wedge S)$
- *23·94. $\vdash : (R) . fR . \equiv . (R) . f(\neg R)$
- *23·95. $\vdash : (\exists R) . fR . \equiv . (\exists R) . f(\neg R)$