

Disjunctive Normal Form - DNF

$\langle \text{literal} \rangle = \langle \text{variable} \rangle \mid \neg \langle \text{variable} \rangle$

$\langle \text{term} \rangle = \langle \text{literal} \rangle (\wedge \langle \text{literal} \rangle)^*$

i.e. $\langle \text{literal} \rangle \wedge \langle \text{literal} \rangle \wedge \langle \text{literal} \rangle \dots$

$\langle \text{formula} \rangle = \langle \text{term} \rangle (\vee \langle \text{term} \rangle)^*$

"OR of ANDs"

e.g.
 $F = (A \wedge \neg B \wedge \neg C) \vee (A \wedge B \wedge C) \vee (\neg A \wedge \neg B \wedge C)$

DNF-SAT is trivial $O(N)$
N is number of variables

Conjunctive Normal Form - CNF

just switch \wedge and \vee in previous definition

"AND of ORs"

e.g.

$$F = (A \vee \neg B \vee \neg C) \wedge (A \vee B \vee C) \wedge (\neg A \vee \neg B \vee C)$$

CNF-SAT is difficult $O(2^N)$

But De Morgan's theorems allow you to
convert each into the other.

So does plain old associativity.

$(A \vee \bar{B} \vee \bar{C}) \wedge (A \vee B \vee C) \wedge (\bar{A} \vee \bar{B} \vee C)$ in CNF

$$(A \vee B \vee C) \wedge (\bar{A} \vee \bar{B} \vee C)$$

$$= A \wedge (\bar{A} \vee \bar{B} \vee C) \vee B \wedge (\bar{A} \vee \bar{B} \vee C) \vee C \wedge (\bar{A} \vee \bar{B} \vee C)$$

$$= \frac{A \wedge \bar{A}}{F} \vee A \wedge \bar{B} \vee A \wedge C \vee B \wedge \bar{A} \vee \frac{B \wedge \bar{B}}{F} \vee B \wedge C \vee C \wedge \bar{A} \vee C \wedge \bar{B} \vee \frac{C \wedge C}{C}$$

$$= \bar{A} \bar{B} \vee AC \vee \bar{B} \bar{A} \vee BC \vee \bar{C} \bar{A} \vee \bar{C} \bar{B}$$

$\therefore (A \vee \bar{B} \vee \bar{C}) \wedge$ that

$$= A \wedge \text{that} \vee \bar{B} \wedge \text{that} \vee \bar{C} \wedge \text{that}$$

$$= \underline{A} \underline{\bar{A}} \vee \underline{A} \underline{C} \vee \underline{\bar{A}} \underline{\bar{A}} \vee ABC \vee \underline{AC} \bar{A} \vee AC \bar{B}$$

$$\vee \underline{\bar{B}} \underline{A} \bar{B} \vee \underline{\bar{B}} \underline{A} C \vee \underline{\bar{B}} \underline{\bar{B}} \underline{A} \vee \underline{\bar{B}} \underline{B} \underline{C} \vee \bar{B} \bar{C} \bar{A} \vee \bar{B} \bar{C} \bar{B}$$

$$\vee \bar{C} \underline{A} \bar{B} \vee \underline{\bar{C}} \underline{A} \underline{C} \vee \bar{C} \underline{B} \underline{A} \vee \underline{\bar{C}} \underline{B} \underline{C} \vee \bar{C} \underline{C} \underline{A} \vee \underline{\bar{C}} \underline{C} \underline{B}$$

all = $A \bar{B} \vee A C \vee A \wedge B \wedge C \vee A C \bar{B} \vee \bar{B} A \vee \bar{B} \wedge A \wedge C \vee \bar{B} \wedge C \wedge A \vee \bar{B} C \vee \bar{C} A \bar{B} \vee \bar{C} B A$ in DNF

If the conversion between domains is non-P, the transformation is not helpful.