Adding an assignment statement to the syntax

Simple Expression	SE	::=	identifier number
Adding Expression	AE	::=	SE ((<u>+ -</u>) SE)*
Statement	STMT	::=	<pre>print AE ; identifier = AE ;</pre>
Program	PROG	::=	STMT end-of-file

Nothing changes except that an extra case is added to read_statement and execute.

read statement(): Use next() to get first symbol if it is "print": SAME AS BEFORE if it is an identifier: create a new node to represent this identifier, save it as ID (a good way to do this is to call back(), then let read simple expr() do the real work for you) use next() to check for "=" symbol if there is no "=": error and return NULL use read adding expression() to read the AE, save it as VAL use next() to check for semi-colon if semi-colon not found: error message, return NULL; make a assignment-statement node containing ID and VAL, return that node (pointer) as result otherwise SAME AS BEFORE execute(node * t) if t is NULL: SAME AS BEFORE else if t->kind is "print" SAME AS BEFORE else if t->kind is "assignment" follow the pointers to get the variable name: string varname = t->ptr1->detail use value of to evaluate the expression int val = value of($t \rightarrow ptr2$) mem.set(varname, val) else SAME AS BEFORE

Allowing parentheses in expressions

Just realise that "(" followed by any expression, followed by ")" behaves like a very simple basic expression, so add one clause to SE:

Simple Expression	SE	::=	identifier
			number
			<u>(</u> number <u>)</u>

This is implemented by adding one new case to read_simple_expr(). After checking for an identifier or a number, check for an opening parenthesis:

use next() to get one symbol if it is a number or identifier create appropriate node and return pointer otherwise if it is "(" call read_adding_expression to do its job, save result as E call next() to check for ")" if ")" not present, error message and return NULL otherwise return E otherwise error message and return NULL.

No other additions are needed. Parentheses in expressions just change the way the parser builds the tree.

Defining a **block**, or sequence of statements, which now becomes the main thing in a program:

Block	BLOCK ::=	<u>{</u> STMT * <u>}</u>
Program	PROG ::=	BLOCK end-of-file

This requires a new parsing method, perhaps called read_block():

```
use next() to check for "{"
if "{" not present, error message, return NULL
L = NULL
enter loop:
       use next() to check for "}"
       if "}" is seen:
               break from loop.
       use read statement() to read just one statement, save result as S
       if L is still NULL
               set L = S
       otherwise
               create new node labelled "sequence" with pointers L and S
               set L = that new node
after end of loop:
       if L is still NULL
               replace L with new node labelled "empty statement", no content
       return L as result.
```

Also add a case to execute() to handle these two new kinds of node:

```
if t->kind is "empty statement":
    don't do anything, the program was just "{ }".
    if t->kind is "sequence":
        do the first step - execute(t->ptr1)
        do the second step - execute(t->ptr2)
        that's it.
```

To allow a block to appear as a kind of statement:

Statement

Fortunately, a block always begins with "{", which is distinct from the existing cases, so just add a new case to read_statement:

```
Use next() to get first symbol

if it is "print":

SAME AS BEFORE

if it is an identifier:

SAME AS BEFORE

if it is "{":

use back(). The "{" is block's responsibility.

call read_block(), return whatever it gives you.

otherwise

SAME AS BEFORE
```

You may like to exercise your minds by thinking about how new operators may be added, such as *, /, <, >, etc., and then about how an if statement could be invented.