Intel® 64 and IA-32 Architectures Software Developer's Manual

Volume 2A: Instruction Set Reference, A-M

NOTE: The Intel 64 and IA-32 Architectures Software Developer's Manual consists of five volumes: *Basic Architecture*, Order Number 253665; *Instruction Set Reference A-M*, Order Number 253666; *Instruction Set Reference N-Z*, Order Number 253667; *System Programming Guide, Part 1*, Order Number 253668; *System Programming Guide, Part 2*, Order Number 253669. Refer to all five volumes when evaluating your design needs.

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CONTENTS

PAGE

CHAPTER 3
INSTRUCTION SET REFERENCE, A-M

PAGE

PAGE

PAGE

CHAPTER 4
INSTRUCTION SET REFERENCE, N-Z

PAGE

PAGE

PAGE

CHAPTER 5
VMX INSTRUCTION REFERENCE

APPENDIX B
INSTRUCTION FORMATS AND ENCODINGS

FIGURES

PAGE

TABLES

PAGE

CONTENTS

PAGE

CHAPTER 1
ABOUT THIS MANUAL

The *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volumes 2A & 2B*: *Instruction Set Reference* (order numbers 253666 and 253667) are part of a set that describes the architecture and programming environment of all Intel 64 and IA-32 architecture processors. Other volumes in this set are:

- The *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1: Basic Architecture* (Order Number 253665).
- The *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volumes 3A & 3B: System Programming Guide* (order numbers 253668 and 253669).

The *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1,* describes the basic architecture and programming environment of Intel 64 and IA-32 processors. The *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volumes 2A & 2B*, describe the instruction set of the processor and the opcode structure. These volumes apply to application programmers and to programmers who write operating systems or executives. The *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volumes 3A & 3B,* describe the operating-system support environment of Intel 64 and IA-32 processors. These volumes target operating-system and BIOS designers. In addition, the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3B*, addresses the programming environment for classes of software that host operating systems.

1.1 IA-32 PROCESSORS COVERED IN THIS MANUAL

This manual set includes information pertaining primarily to the most recent Intel 64 and IA-32 processors, which include:

- Pentium® processors
- P6 family processors
- Pentium® 4 processors
- Pentium® M processors
- Intel® Xeon® processors
- Pentium® D processors
- Pentium® processor Extreme Editions
- 64-bit Intel® Xeon® processors
- Intel® Core™ Duo processor
- Intel® Core™ Solo processor
- Dual-Core Intel® Xeon® processor LV

- Intel® Core™2 Duo processor
- Intel® Core™2 Quad processor Q6000 series
- Intel® Xeon® processor 3000, 3200 series
- Intel® Xeon® processor 5000 series
- Intel® Xeon® processor 5100, 5300 series
- Intel® Core™2 Extreme processor X7000 and X6800 series
- Intel® Core™2 Extreme QX6000 series
- Intel® Xeon® processor 7100 series
- Intel® Pentium® Dual-Core processor
- Intel® Xeon® processor 7200, 7300 series
- Intel® Xeon® processor 5200, 5400, 7400 series
- Intel® Core™2 Extreme processor QX9000 and X9000 series
- Intel® Core™2 Quad processor Q9000 series
- Intel® Core™2 Duo processor E8000, T9000 series
- Intel® Atom™ processor family
- Intel® Core™ i7 processor
- Intel® Core™ i5 processor

P6 family processors are IA-32 processors based on the P6 family microarchitecture. This includes the Pentium® Pro, Pentium® II, Pentium® III, and Pentium® III Xeon® processors.

The Pentium® 4, Pentium® D, and Pentium® processor Extreme Editions are based on the Intel NetBurst® microarchitecture. Most early Intel® Xeon® processors are based on the Intel NetBurst® microarchitecture. Intel Xeon processor 5000, 7100 series are based on the Intel NetBurst® microarchitecture.

The Intel® Core™ Duo, Intel® Core™ Solo and dual-core Intel® Xeon® processor LV are based on an improved Pentium® M processor microarchitecture.

The Intel® Xeon® processor 3000, 3200, 5100, 5300, 7200, and 7300 series, Intel® Pentium® dual-core, Intel® Core™2 Duo, Intel® Core™2 Quad, and Intel® Core™2 Extreme processors are based on Intel® Core™ microarchitecture.

The Intel® Xeon® processor 5200, 5400, 7400 series, Intel® Core™2 Quad processor Q9000 series, and Intel® Core™2 Extreme processors QX9000, X9000 series, Intel® Core™2 processor E8000 series are based on Enhanced Intel® Core™ microarchitecture.

The Intel® Atom™ processor family is based on the Intel® Atom™ microarchitecture and supports Intel 64 architecture.

The Intel® Core™ i7 processor and the Intel® Core™ i5 processor are based on the Intel® microarchitecture code name Nehalem and support Intel 64 architecture.

Processors based on Intel® microarchitecture code name Westmere support Intel 64 architecture.

P6 family, Pentium® M, Intel® Core™ Solo, Intel® Core™ Duo processors, dual-core Intel® Xeon® processor LV, and early generations of Pentium 4 and Intel Xeon processors support IA-32 architecture. The Intel® Atom™ processor Z5xx series support IA-32 architecture.

The Intel® Xeon® processor 3000, 3200, 5000, 5100, 5200, 5300, 5400, 7100, 7200, 7300, 7400 series, Intel® Core™2 Duo, Intel® Core™2 Extreme, Intel® Core™2 Quad processors, Pentium® D processors, Pentium® Dual-Core processor, newer generations of Pentium 4 and Intel Xeon processor family support Intel® 64 architecture.

IA-32 architecture is the instruction set architecture and programming environment for Intel's 32-bit microprocessors.

Intel® 64 architecture is the instruction set architecture and programming environment which is the superset of Intel’s 32-bit and 64-bit architectures. It is compatible with the IA-32 architecture.

1.2 OVERVIEW OF VOLUME 2A AND 2B: INSTRUCTION SET REFERENCE

A description of *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volumes 2A & 2B*, content follows:

Chapter 1 — About This Manual. Gives an overview of all five volumes of the *Intel® 64 and IA-32 Architectures Software Developer’s Manual*. It also describes the notational conventions in these manuals and lists related Intel® manuals and documentation of interest to programmers and hardware designers.

Chapter 2 — Instruction Format. Describes the machine-level instruction format used for all IA-32 instructions and gives the allowable encodings of prefixes, the operand-identifier byte (ModR/M byte), the addressing-mode specifier byte (SIB byte), and the displacement and immediate bytes.

Chapter 3 — Instruction Set Reference, A-M. Describes Intel 64 and IA-32 instructions in detail, including an algorithmic description of operations, the effect on flags, the effect of operand- and address-size attributes, and the exceptions that may be generated. The instructions are arranged in alphabetical order. General-purpose, x87 FPU, Intel MMX™ technology, SSE/SSE2/SSE3/SSSE3/SSE4 extensions, and system instructions are included.

Chapter 4 — Instruction Set Reference, N-Z. Continues the description of Intel 64 and IA-32 instructions started in Chapter 3. It provides the balance of the alphabetized list of instructions and starts *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 2B*.

Chapter 5 — VMX Instruction Reference. Describes the virtual-machine extensions (VMX). VMX is intended for a system executive to support virtualization of processor hardware and a system software layer acting as a host to multiple guest software environments.

Chapter 6— Safer Mode Extensions Reference. Describes the safer mode extensions (SMX). SMX is intended for a system executive to support launching a measured environment in a platform where the identity of the software controlling the platform hardware can be measured for the purpose of making trust decisions.

Appendix A — Opcode Map. Gives an opcode map for the IA-32 instruction set.

Appendix B — Instruction Formats and Encodings. Gives the binary encoding of each form of each IA-32 instruction.

Appendix C — Intel® C/C++ Compiler Intrinsics and Functional Equivalents. Lists the Intel® C/C++ compiler intrinsics and their assembly code equivalents for each of the IA-32 MMX and SSE/SSE2/SSE3 instructions.

1.3 NOTATIONAL CONVENTIONS

This manual uses specific notation for data-structure formats, for symbolic representation of instructions, and for hexadecimal and binary numbers. A review of this notation makes the manual easier to read.

1.3.1 Bit and Byte Order

In illustrations of data structures in memory, smaller addresses appear toward the bottom of the figure; addresses increase toward the top. Bit positions are numbered from right to left. The numerical value of a set bit is equal to two raised to the power of the bit position. IA-32 processors are “little endian” machines; this means the bytes of a word are numbered starting from the least significant byte. Figure 1-1 illustrates these conventions.

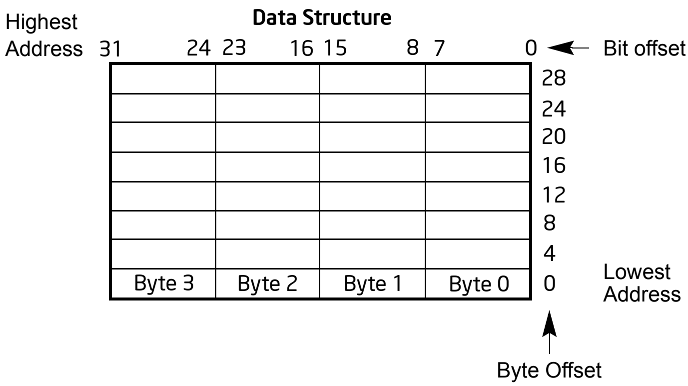


Figure 1-1. Bit and Byte Order

1.3.2 Reserved Bits and Software Compatibility

In many register and memory layout descriptions, certain bits are marked as **reserved**. When bits are marked as reserved, it is essential for compatibility with future processors that software treat these bits as having a future, though unknown, effect. The behavior of reserved bits should be regarded as not only undefined, but unpredictable. Software should follow these guidelines in dealing with reserved bits:

- Do not depend on the states of any reserved bits when testing the values of registers which contain such bits. Mask out the reserved bits before testing.
- Do not depend on the states of any reserved bits when storing to memory or to a register.
- Do not depend on the ability to retain information written into any reserved bits.
- When loading a register, always load the reserved bits with the values indicated in the documentation, if any, or reload them with values previously read from the same register.

NOTE

Avoid any software dependence upon the state of reserved bits in IA-32 registers. Depending upon the values of reserved register bits will make software dependent upon the unspecified manner in which the processor handles these bits. Programs that depend upon reserved values risk incompatibility with future processors.

1.3.3 Instruction Operands

When instructions are represented symbolically, a subset of the IA-32 assembly language is used. In this subset, an instruction has the following format:

```
label: mnemonic argument1, argument2, argument3
```

where:

- A **label** is an identifier which is followed by a colon.
- A **mnemonic** is a reserved name for a class of instruction opcodes which have the same function.
- The operands *argument1*, *argument2*, and *argument3* are optional. There may be from zero to three operands, depending on the opcode. When present, they take the form of either literals or identifiers for data items. Operand identifiers are either reserved names of registers or are assumed to be assigned to data items declared in another part of the program (which may not be shown in the example).

When two operands are present in an arithmetic or logical instruction, the right operand is the source and the left operand is the destination.

For example:

```
LOADREG: MOV EAX, SUBTOTAL
```

In this example, LOADREG is a label, MOV is the mnemonic identifier of an opcode, EAX is the destination operand, and SUBTOTAL is the source operand. Some assembly languages put the source and destination in reverse order.

1.3.4 Hexadecimal and Binary Numbers

Base 16 (hexadecimal) numbers are represented by a string of hexadecimal digits followed by the character H (for example, F82EH). A hexadecimal digit is a character from the following set: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, and F.

Base 2 (binary) numbers are represented by a string of 1s and 0s, sometimes followed by the character B (for example, 1010B). The “B” designation is only used in situations where confusion as to the type of number might arise.

1.3.5 Segmented Addressing

The processor uses byte addressing. This means memory is organized and accessed as a sequence of bytes. Whether one or more bytes are being accessed, a byte address is used to locate the byte or bytes in memory. The range of memory that can be addressed is called an **address space**.

The processor also supports segmented addressing. This is a form of addressing where a program may have many independent address spaces, called **segments**.

For example, a program can keep its code (instructions) and stack in separate segments. Code addresses would always refer to the code space, and stack addresses would always refer to the stack space. The following notation is used to specify a byte address within a segment:

Segment-register:Byte-address

For example, the following segment address identifies the byte at address FF79H in the segment pointed by the DS register:

DS:FF79H

The following segment address identifies an instruction address in the code segment. The CS register points to the code segment and the EIP register contains the address of the instruction.

CS:EIP

1.3.6 Exceptions

An exception is an event that typically occurs when an instruction causes an error. For example, an attempt to divide by zero generates an exception. However, some exceptions, such as breakpoints, occur under other conditions. Some types of exceptions may provide error codes. An error code reports additional information about the error. An example of the notation used to show an exception and error code is shown below:

#PF(fault code)

This example refers to a page-fault exception under conditions where an error code naming a type of fault is reported. Under some conditions, exceptions which produce error codes may not be able to report an accurate code. In this case, the error code is zero, as shown below for a general-protection exception:

#GP(0)

1.3.7 A New Syntax for CPUID, CR, and MSR Values

Obtain feature flags, status, and system information by using the CPUID instruction, by checking control register bits, and by reading model-specific registers. We are moving toward a new syntax to represent this information. See Figure 1-2.

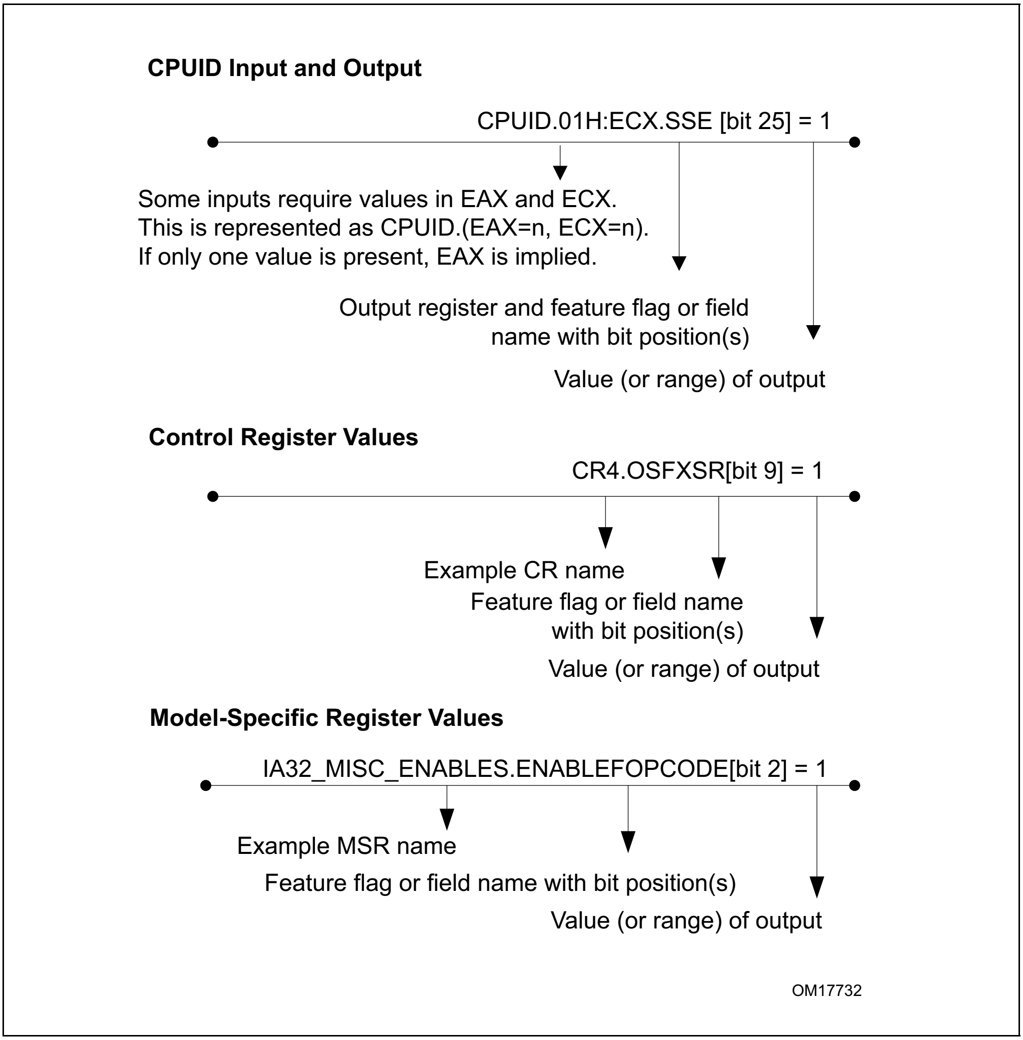


Figure 1-2. Syntax for CPUID, CR, and MSR Data Presentation

1.4 RELATED LITERATURE

Literature related to Intel 64 and IA-32 processors is listed on-line at:

http://developer.intel.com/products/processor/manuals/index.htm

Some of the documents listed at this web site can be viewed on-line; others can be ordered. The literature available is listed by Intel processor and then by the following literature types: applications notes, data sheets, manuals, papers, and specification updates.

See also:

- The data sheet for a particular Intel 64 or IA-32 processor
- The specification update for a particular Intel 64 or IA-32 processor
- Intel® C++ Compiler documentation and online help
http://www.intel.com/cd/software/products/asmo-na/eng/index.htm
- Intel® Fortran Compiler documentation and online help
http://www.intel.com/cd/software/products/asmo-na/eng/index.htm
- Intel® VTune™ Performance Analyzer documentation and online help
http://www.intel.com/cd/software/products/asmo-na/eng/index.htm
- Intel® 64 and IA-32 Architectures Software Developer's Manual (in five volumes)
http://developer.intel.com/products/processor/manuals/index.htm
- Intel® 64 and IA-32 Architectures Optimization Reference Manual
http://developer.intel.com/products/processor/manuals/index.htm
- Intel® Processor Identification with the CPUID Instruction, AP-485
http://www.intel.com/support/processors/sb/cs-009861.htm
- Intel 64 Architecture x2APIC Specification:

  http://developer.intel.com/products/processor/manuals/index.htm
- Intel 64 Architecture Processor Topology Enumeration:

  http://softwarecommunity.intel.com/articles/eng/3887.htm
- Intel® Trusted Execution Technology Measured Launched Environment Programming Guide, http://www.intel.com/technology/security/index.htm
- Intel® SSE4 Programming Reference,
http://developer.intel.com/products/processor/manuals/index.htm
- Developing Multi-threaded Applications: A Platform Consistent Approach
http://cache-www.intel.com/cd/00/00/05/15/51534_developing_multithreaded_applications.pdf
- Using Spin-Loops on Intel Pentium 4 Processor and Intel Xeon Processor MP
http://www3.intel.com/cd/ids/developer/asmo-na/eng/dc/threading/knowledgebase/19083.htm

More relevant links are:

- Software network link:

  http://softwarecommunity.intel.com/isn/home/
- Developer centers:

  http://www.intel.com/cd/ids/developer/asmo-na/eng/dc/index.htm
- Processor support general link:

  http://www.intel.com/support/processors/

- Software products and packages:
  http://www.intel.com/cd/software/products/asmo-na/eng/index.htm
- Intel 64 and IA-32 processor manuals (printed or PDF downloads):
  http://developer.intel.com/products/processor/manuals/index.htm
- Intel® Multi-Core Technology:
  http://developer.intel.com/multi-core/index.htm
- Intel® Hyper-Threading Technology (Intel® HT Technology):
  http://developer.intel.com/technology/hyperthread/

CHAPTER 2
INSTRUCTION FORMAT

This chapter describes the instruction format for all Intel 64 and IA-32 processors. The instruction format for protected mode, real-address mode and virtual-8086 mode is described in Section 2.1. Increments provided for IA-32e mode and its sub-modes are described in Section 2.2.

2.1 INSTRUCTION FORMAT FOR PROTECTED MODE, REAL-ADDRESS MODE, AND VIRTUAL-8086 MODE

The Intel 64 and IA-32 architectures instruction encodings are subsets of the format shown in Figure 2-1. Instructions consist of optional instruction prefixes (in any order), primary opcode bytes (up to three bytes), an addressing-form specifier (if required) consisting of the ModR/M byte and sometimes the SIB (Scale-Index-Base) byte, a displacement (if required), and an immediate data field (if required).

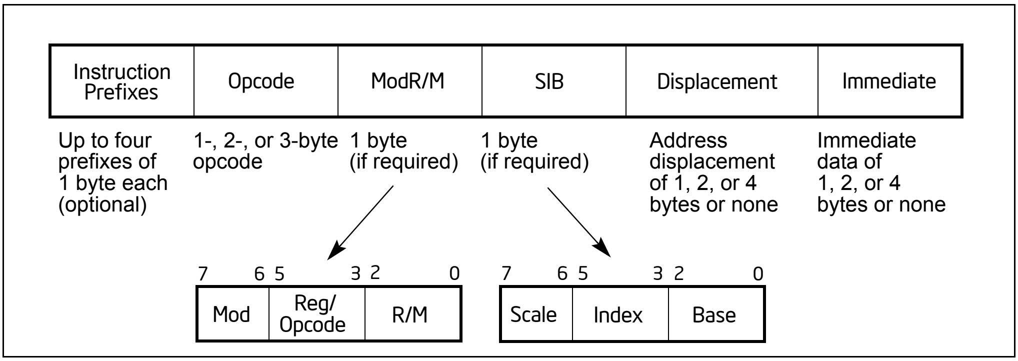


Figure 2-1. Intel 64 and IA-32 Architectures Instruction Format

2.1.1 Instruction Prefixes

Instruction prefixes are divided into four groups, each with a set of allowable prefix codes. For each instruction, it is only useful to include up to one prefix code from each of the four groups (Groups 1, 2, 3, 4). Groups 1 through 4 may be placed in any order relative to each other.

- Group 1
  - Lock and repeat prefixes:

    - LOCK prefix is encoded using F0H
    - REPNE/REPNZ prefix is encoded using F2H. Repeat-Not-Zero prefix applies only to string and input/output instructions. (F2H is also used as a mandatory prefix for some instructions)
    - REP or REPE/REPZ is encoded using F3H. Repeat prefix applies only to string and input/output instructions.(F3H is also used as a mandatory prefix for some instructions)
- Group 2
  - — Segment override prefixes:
    - 2EH—CS segment override (use with any branch instruction is reserved)
    - 36H—SS segment override prefix (use with any branch instruction is reserved)
    - 3EH—DS segment override prefix (use with any branch instruction is reserved)
    - 26H—ES segment override prefix (use with any branch instruction is reserved)
    - 64H—FS segment override prefix (use with any branch instruction is reserved)
    - 65H—GS segment override prefix (use with any branch instruction is reserved)
  - — Branch hints:
    - 2EH—Branch not taken (used only with J*cc* instructions)
    - 3EH—Branch taken (used only with J*cc* instructions)
- Group 3
    - Operand-size override prefix is encoded using 66H (66H is also used as a mandatory prefix for some instructions).
- Group 4
    - 67H—Address-size override prefix

The LOCK prefix (F0H) forces an operation that ensures exclusive use of shared memory in a multiprocessor environment. See “LOCK—Assert LOCK# Signal Prefix” in Chapter 3, “Instruction Set Reference, A-M,” for a description of this prefix.

Repeat prefixes (F2H, F3H) cause an instruction to be repeated for each element of a string. Use these prefixes only with string and I/O instructions (MOVS, CMPS, SCAS, LODS, STOS, INS, and OUTS). Use of repeat prefixes and/or undefined opcodes with other Intel 64 or IA-32 instructions is reserved; such use may cause unpredictable behavior.

Some instructions may use F2H,F3H as a mandatory prefix to express distinct functionality. A mandatory prefix generally should be placed after other optional prefixes (exception to this is discussed in Section 2.2.1, “REX Prefixes”)

Branch hint prefixes (2EH, 3EH) allow a program to give a hint to the processor about the most likely code path for a branch. Use these prefixes only with conditional branch instructions (J*cc*). Other use of branch hint prefixes and/or other undefined opcodes with Intel 64 or IA-32 instructions is reserved; such use may cause unpredictable behavior.

The operand-size override prefix allows a program to switch between 16- and 32-bit operand sizes. Either size can be the default; use of the prefix selects the non-default size.

Some SSE2/SSE3/SSSE3/SSE4 instructions and instructions using a three-byte sequence of primary opcode bytes may use 66H as a mandatory prefix to express distinct functionality. A mandatory prefix generally should be placed after other optional prefixes (exception to this is discussed in Section 2.2.1, “REX Prefixes”)

Other use of the 66H prefix is reserved; such use may cause unpredictable behavior.

The address-size override prefix (67H) allows programs to switch between 16- and 32-bit addressing. Either size can be the default; the prefix selects the non-default size. Using this prefix and/or other undefined opcodes when operands for the instruction do not reside in memory is reserved; such use may cause unpredictable behavior.

2.1.2 Opcodes

A primary opcode can be 1, 2, or 3 bytes in length. An additional 3-bit opcode field is sometimes encoded in the ModR/M byte. Smaller fields can be defined within the primary opcode. Such fields define the direction of operation, size of displacements, register encoding, condition codes, or sign extension. Encoding fields used by an opcode vary depending on the class of operation.

Two-byte opcode formats for general-purpose and SIMD instructions consist of:

- An escape opcode byte 0FH as the primary opcode and a second opcode byte, or
- A mandatory prefix (66H, F2H, or F3H), an escape opcode byte, and a second opcode byte (same as previous bullet)

For example, CVTDQ2PD consists of the following sequence: F3 0F E6. The first byte is a mandatory prefix (it is not considered as a repeat prefix).

Three-byte opcode formats for general-purpose and SIMD instructions consist of:

- An escape opcode byte 0FH as the primary opcode, plus two additional opcode bytes, or
- A mandatory prefix (66H, F2H, or F3H), an escape opcode byte, plus two additional opcode bytes (same as previous bullet)

For example, PHADDW for XMM registers consists of the following sequence: 66 0F 38 01. The first byte is the mandatory prefix.

Valid opcode expressions are defined in Appendix A and Appendix B.

2.1.3 ModR/M and SIB Bytes

Many instructions that refer to an operand in memory have an addressing-form specifier byte (called the ModR/M byte) following the primary opcode. The ModR/M byte contains three fields of information:

- The *mod* field combines with the r/m field to form 32 possible values: eight registers and 24 addressing modes.
- The *reg/opcode* field specifies either a register number or three more bits of opcode information. The purpose of the reg/opcode field is specified in the primary opcode.
- The *r/m* field can specify a register as an operand or it can be combined with the mod field to encode an addressing mode. Sometimes, certain combinations of the *mod* field and the *r/m* field is used to express opcode information for some instructions.

Certain encodings of the ModR/M byte require a second addressing byte (the SIB byte). The base-plus-index and scale-plus-index forms of 32-bit addressing require the SIB byte. The SIB byte includes the following fields:

- The *scale* field specifies the scale factor.
- The *index* field specifies the register number of the index register.
- The *base* field specifies the register number of the base register.

See Section 2.1.5 for the encodings of the ModR/M and SIB bytes.

2.1.4 Displacement and Immediate Bytes

Some addressing forms include a displacement immediately following the ModR/M byte (or the SIB byte if one is present). If a displacement is required; it be 1, 2, or 4 bytes.

If an instruction specifies an immediate operand, the operand always follows any displacement bytes. An immediate operand can be 1, 2 or 4 bytes.

2.1.5 Addressing-Mode Encoding of ModR/M and SIB Bytes

The values and corresponding addressing forms of the ModR/M and SIB bytes are shown in Table 2-1 through Table 2-3: 16-bit addressing forms specified by the ModR/M byte are in Table 2-1 and 32-bit addressing forms are in Table 2-2. Table 2-3 shows 32-bit addressing forms specified by the SIB byte. In cases where the reg/opcode field in the ModR/M byte represents an extended opcode, valid encodings are shown in Appendix B.

In Table 2-1 and Table 2-2, the Effective Address column lists 32 effective addresses that can be assigned to the first operand of an instruction by using the Mod and R/M fields of the ModR/M byte. The first 24 options provide ways of specifying a memory location; the last eight (Mod = 11B) provide ways of specifying general-purpose, MMX technology and XMM registers.

The Mod and R/M columns in Table 2-1 and Table 2-2 give the binary encodings of the Mod and R/M fields required to obtain the effective address listed in the first column. For example: see the row indicated by Mod = 11B, R/M = 000B. The row identifies the general-purpose registers EAX, AX or AL; MMX technology register MM0; or XMM register XMM0. The register used is determined by the opcode byte and the operand-size attribute.

Now look at the seventh row in either table (labeled "REG ="). This row specifies the use of the 3-bit Reg/Opcode field when the field is used to give the location of a second operand. The second operand must be a general-purpose, MMX technology, or XMM register. Rows one through five list the registers that may correspond to the value in the table. Again, the register used is determined by the opcode byte along with the operand-size attribute.

If the instruction does not require a second operand, then the Reg/Opcode field may be used as an opcode extension. This use is represented by the sixth row in the tables (labeled "/digit (Opcode)"). Note that values in row six are represented in decimal form.

The body of Table 2-1 and Table 2-2 (under the label "Value of ModR/M Byte (in Hexadecimal)") contains a 32 by 8 array that presents all of 256 values of the ModR/M byte (in hexadecimal). Bits 3, 4 and 5 are specified by the column of the table in which a byte resides. The row specifies bits 0, 1 and 2; and bits 6 and 7. The figure below demonstrates interpretation of one table value.

	Mod	11
	RM	000
/digit (Opcode);	REG =	001
	C8H	11001000

Figure 2-2. Table Interpretation of ModR/M Byte (C8H)

Table 2-1. 16-Bit Addressing Forms with the ModR/M Byte

r8(/r) r16(/r) r32(/r) mm(/r) xmm(/r) (In decimal) /digit (Opcode) (In binary) REG =			AL AX EAX MM0 XMM0 0 000	CL CX ECX MM1 XMM1 1 001	DL DX EDX MM2 XMM2 2 010	BL BX EBX MM3 XMM3 3 011	AH SP ESP MM4 XMM4 4 100	CH BP1 EBP MM5 XMM5 5 101	DH SI ESI MM6 XMM6 6 110	BH DI EDI MM7 XMM7 7 111
Effective Address	**Mod**	**R/M**	**Value of ModR/M Byte (in Hexadecimal)**							
[BX+SI]	00	000	00	08	10	18	20	28	30	38
[BX+DI]		001	01	09	11	19	21	29	31	39
[BP+SI]		010	02	0A	12	1A	22	2A	32	3A
[BP+DI]		011	03	0B	13	1B	23	2B	33	3B
[SI]		100	04	0C	14	1C	24	2C	34	3C
[DI]		101	05	0D	15	1D	25	2D	35	3D
disp16[2]		110	06	0E	16	1E	26	2E	36	3E
[BX]		111	07	0F	17	1F	27	2F	37	3F
[BX+SI]+disp8[3]	01	000	40	48	50	58	60	68	70	78
[BX+DI]+disp8		001	41	49	51	59	61	69	71	79
[BP+SI]+disp8		010	42	4A	52	5A	62	6A	72	7A
[BP+DI]+disp8		011	43	4B	53	5B	63	6B	73	7B
[SI]+disp8		100	44	4C	54	5C	64	6C	74	7C
[DI]+disp8		101	45	4D	55	5D	65	6D	75	7D
[BP]+disp8		110	46	4E	56	5E	66	6E	76	7E
[BX]+disp8		111	47	4F	57	5F	67	6F	77	7F
[BX+SI]+disp16	10	000	80	88	90	98	A0	A8	B0	B8
[BX+DI]+disp16		001	81	89	91	99	A1	A9	B1	B9
[BP+SI]+disp16		010	82	8A	92	9A	A2	AA	B2	BA
[BP+DI]+disp16		011	83	8B	93	9B	A3	AB	B3	BB
[SI]+disp16		100	84	8C	94	9C	A4	AC	B4	BC
[DI]+disp16		101	85	8D	95	9D	A5	AD	B5	BD
[BP]+disp16		110	86	8E	96	9E	A6	AE	B6	BE
[BX]+disp16		111	87	8F	97	9F	A7	AF	B7	BF
EAX/AX/AL/MM0/XMM0	11	000	C0	C8	D0	D8	E0	E8	F0	F8
ECX/CX/CL/MM1/XMM1		001	C1	C9	D1	D9	EQ	E9	F1	F9
EDX/DX/DL/MM2/XMM2		010	C2	CA	D2	DA	E2	EA	F2	FA
EBX/BX/BL/MM3/XMM3		011	C3	CB	D3	DB	E3	EB	F3	FB
ESP/SP/AHMM4/XMM4		100	C4	CC	D4	DC	E4	EC	F4	FC
EBP/BP/CH/MM5/XMM5		101	C5	CD	D5	DD	E5	ED	F5	FD
ESI/SI/DH/MM6/XMM6		110	C6	CE	D6	DE	E6	EE	F6	FE
EDI/DI/BH/MM7/XMM7		111	C7	CF	D7	DF	E7	EF	F7	FF

NOTES:

1. The default segment register is SS for the effective addresses containing a BP index, DS for other effective addresses.
2. The disp16 nomenclature denotes a 16-bit displacement that follows the ModR/M byte and that is added to the index.
3. The disp8 nomenclature denotes an 8-bit displacement that follows the ModR/M byte and that is sign-extended and added to the index.

Table 2-2. 32-Bit Addressing Forms with the ModR/M Byte

r8(/r) r16(/r) r32(/r) mm(/r) xmm(/r) (In decimal) /digit (Opcode) (In binary) REG =			AL AX EAX MM0 XMM0 0 000	CL CX ECX MM1 XMM1 1 001	DL DX EDX MM2 XMM2 2 010	BL BX EBX MM3 XMM3 3 011	AH SP ESP MM4 XMM4 4 100	CH BP EBP MM5 XMM5 5 101	DH SI ESI MM6 XMM6 6 110	BH DI EDI MM7 XMM7 7 111
Effective Address	**Mod**	**R/M**	**Value of ModR/M Byte (in Hexadecimal)**							
[EAX]	00	000	00	08	10	18	20	28	30	38
[ECX]		001	01	09	11	19	21	29	31	39
[EDX]		010	02	0A	12	1A	22	2A	32	3A
[EBX]		011	03	0B	13	1B	23	2B	33	3B
[--][--][1]		100	04	0C	14	1C	24	2C	34	3C
disp32[2]		101	05	0D	15	1D	25	2D	35	3D
[ESI]		110	06	0E	16	1E	26	2E	36	3E
[EDI]		111	07	0F	17	1F	27	2F	37	3F
[EAX]+disp8[3]	01	000	40	48	50	58	60	68	70	78
[ECX]+disp8		001	41	49	51	59	61	69	71	79
[EDX]+disp8		010	42	4A	52	5A	62	6A	72	7A
[EBX]+disp8		011	43	4B	53	5B	63	6B	73	7B
[--][--]+disp8		100	44	4C	54	5C	64	6C	74	7C
[EBP]+disp8		101	45	4D	55	5D	65	6D	75	7D
[ESI]+disp8		110	46	4E	56	5E	66	6E	76	7E
[EDI]+disp8		111	47	4F	57	5F	67	6F	77	7F
[EAX]+disp32	10	000	80	88	90	98	A0	A8	B0	B8
[ECX]+disp32		001	81	89	91	99	A1	A9	B1	B9
[EDX]+disp32		010	82	8A	92	9A	A2	AA	B2	BA
[EBX]+disp32		011	83	8B	93	9B	A3	AB	B3	BB
[--][--]+disp32		100	84	8C	94	9C	A4	AC	B4	BC
[EBP]+disp32		101	85	8D	95	9D	A5	AD	B5	BD
[ESI]+disp32		110	86	8E	96	9E	A6	AE	B6	BE
[EDI]+disp32		111	87	8F	97	9F	A7	AF	B7	BF
EAX/AX/AL/MM0/XMM0	11	000	C0	C8	D0	D8	E0	E8	F0	F8
ECX/CX/CL/MM/XMM1		001	C1	C9	D1	D9	E1	E9	F1	F9
EDX/DX/DL/MM2/XMM2		010	C2	CA	D2	DA	E2	EA	F2	FA
EBX/BX/BL/MM3/XMM3		011	C3	CB	D3	DB	E3	EB	F3	FB
ESP/SP/AH/MM4/XMM4		100	C4	CC	D4	DC	E4	EC	F4	FC
EBP/BP/CH/MM5/XMM5		101	C5	CD	D5	DD	E5	ED	F5	FD
ESI/SI/DH/MM6/XMM6		110	C6	CE	D6	DE	E6	EE	F6	FE
EDI/DI/BH/MM7/XMM7		111	C7	CF	D7	DF	E7	EF	F7	FF

NOTES:

1. The [--][--] nomenclature means a SIB follows the ModR/M byte.
2. The disp32 nomenclature denotes a 32-bit displacement that follows the ModR/M byte (or the SIB byte if one is present) and that is added to the index.
3. The disp8 nomenclature denotes an 8-bit displacement that follows the ModR/M byte (or the SIB byte if one is present) and that is sign-extended and added to the index.

Table 2-3 is organized to give 256 possible values of the SIB byte (in hexadecimal). General purpose registers used as a base are indicated across the top of the table, along with corresponding values for the SIB byte's base field. Table rows in the body of the table indicate the register used as the index (SIB byte bits 3, 4 and 5) and the scaling factor (determined by SIB byte bits 6 and 7).

Table 2-3. 32-Bit Addressing Forms with the SIB Byte

r32 (In decimal) Base = (In binary) Base =			EAX 0 000	ECX 1 001	EDX 2 010	EBX 3 011	ESP 4 100	[*] 5 101	ESI 6 110	EDI 7 111
Scaled Index	**SS**	**Index**	**Value of SIB Byte (in Hexadecimal)**							
[EAX]	00	000	00	01	02	03	04	05	06	07
[ECX]		001	08	09	0A	0B	0C	0D	0E	0F
[EDX]		010	10	11	12	13	14	15	16	17
[EBX]		011	18	19	1A	1B	1C	1D	1E	1F
none		100	20	21	22	23	24	25	26	27
[EBP]		101	28	29	2A	2B	2C	2D	2E	2F
[ESI]		110	30	31	32	33	34	35	36	37
[EDI]		111	38	39	3A	3B	3C	3D	3E	3F
[EAX*2]	01	000	40	41	42	43	44	45	46	47
[ECX*2]		001	48	49	4A	4B	4C	4D	4E	4F
[EDX*2]		010	50	51	52	53	54	55	56	57
[EBX*2]		011	58	59	5A	5B	5C	5D	5E	5F
none		100	60	61	62	63	64	65	66	67
[EBP*2]		101	68	69	6A	6B	6C	6D	6E	6F
[ESI*2]		110	70	71	72	73	74	75	76	77
[EDI*2]		111	78	79	7A	7B	7C	7D	7E	7F
[EAX*4]	10	000	80	81	82	83	84	85	86	87
[ECX*4]		001	88	89	8A	8B	8C	8D	8E	8F
[EDX*4]		010	90	91	92	93	94	95	96	97
[EBX*4]		011	98	89	9A	9B	9C	9D	9E	9F
none		100	A0	A1	A2	A3	A4	A5	A6	A7
[EBP*4]		101	A8	A9	AA	AB	AC	AD	AE	AF
[ESI*4]		110	B0	B1	B2	B3	B4	B5	B6	B7
[EDI*4]		111	B8	B9	BA	BB	BC	BD	BE	BF
[EAX*8]	11	000	C0	C1	C2	C3	C4	C5	C6	C7
[ECX*8]		001	C8	C9	CA	CB	CC	CD	CE	CF
[EDX*8]		010	D0	D1	D2	D3	D4	D5	D6	D7
[EBX*8]		011	D8	D9	DA	DB	DC	DD	DE	DF
none		100	E0	E1	E2	E3	E4	E5	E6	E7
[EBP*8]		101	E8	E9	EA	EB	EC	ED	EE	EF
[ESI*8]		110	F0	F1	F2	F3	F4	F5	F6	F7
[EDI*8]		111	F8	F9	FA	FB	FC	FD	FE	FF

NOTES:

1. The [*] nomenclature means a disp32 with no base if the MOD is 00B. Otherwise, [*] means disp8 or disp32 + [EBP]. This provides the following address modes:

MOD bits	Effective Address
00	[scaled index] + disp32
01	[scaled index] + disp8 + [EBP]
10	[scaled index] + disp32 + [EBP]

2.2 IA-32E MODE

IA-32e mode has two sub-modes. These are:

- **Compatibility Mode.** Enables a 64-bit operating system to run most legacy protected mode software unmodified.
- **64-Bit Mode.** Enables a 64-bit operating system to run applications written to access 64-bit address space.

2.2.1 REX Prefixes

REX prefixes are instruction-prefix bytes used in 64-bit mode. They do the following:

- Specify GPRs and SSE registers.
- Specify 64-bit operand size.
- Specify extended control registers.

Not all instructions require a REX prefix in 64-bit mode. A prefix is necessary only if an instruction references one of the extended registers or uses a 64-bit operand. If a REX prefix is used when it has no meaning, it is ignored.

Only one REX prefix is allowed per instruction. If used, the REX prefix byte must immediately precede the opcode byte or the escape opcode byte (0FH). When a REX prefix is used in conjunction with an instruction containing a mandatory prefix, the mandatory prefix must come before the REX so the REX prefix can be immediately preceding the opcode or the escape byte. For example, CVTDQ2PD with a REX prefix should have REX placed between F3 and 0F E6. Other placements are ignored. The instruction-size limit of 15 bytes still applies to instructions with a REX prefix. See Figure 2-3.

Legacy Prefixes	REX Prefix	Opcode	ModR/M	SIB	Displacement	Immediate
Grp 1, Grp 2, Grp 3, Grp 4 (optional)	(optional)	1-, 2-, or 3-byte opcode	1 byte (if required)	1 byte (if required)	Address displacement of 1, 2, or 4 bytes	Immediate data of 1, 2, or 4 bytes or none

Figure 2-3. Prefix Ordering in 64-bit Mode

2.2.1.1 Encoding

Intel 64 and IA-32 instruction formats specify up to three registers by using 3-bit fields in the encoding, depending on the format:

- ModR/M: the reg and r/m fields of the ModR/M byte
- ModR/M with SIB: the reg field of the ModR/M byte, the base and index fields of the SIB (scale, index, base) byte
- Instructions without ModR/M: the reg field of the opcode

In 64-bit mode, these formats do not change. Bits needed to define fields in the 64-bit context are provided by the addition of REX prefixes.

2.2.1.2 More on REX Prefix Fields

REX prefixes are a set of 16 opcodes that span one row of the opcode map and occupy entries 40H to 4FH. These opcodes represent valid instructions (INC or DEC) in IA-32 operating modes and in compatibility mode. In 64-bit mode, the same opcodes represent the instruction prefix REX and are not treated as individual instructions.

The single-byte-opcode form of INC/DEC instruction not available in 64-bit mode. INC/DEC functionality is still available using ModR/M forms of the same instructions (opcodes FF/0 and FF/1).

See Table 2-4 for a summary of the REX prefix format. Figure 2-4 though Figure 2-7 show examples of REX prefix fields in use. Some combinations of REX prefix fields are invalid. In such cases, the prefix is ignored. Some additional information follows:

- Setting REX.W can be used to determine the operand size but does not solely determine operand width. Like the 66H size prefix, 64-bit operand size override has no effect on byte-specific operations.
- For non-byte operations: if a 66H prefix is used with prefix (REX.W = 1), 66H is ignored.
- If a 66H override is used with REX and REX.W = 0, the operand size is 16 bits.
- REX.R modifies the ModR/M reg field when that field encodes a GPR, SSE, control or debug register. REX.R is ignored when ModR/M specifies other registers or defines an extended opcode.
- REX.X bit modifies the SIB index field.
- REX.B either modifies the base in the ModR/M r/m field or SIB base field; or it modifies the opcode reg field used for accessing GPRs.

Table 2-4. REX Prefix Fields [BITS: 0100WRXB]

Field Name	Bit Position	Definition
-	7:4	0100
W	3	0 = Operand size determined by CS.D
		1 = 64 Bit Operand Size
R	2	Extension of the ModR/M reg field
X	1	Extension of the SIB index field
B	0	Extension of the ModR/M r/m field, SIB base field, or Opcode reg field

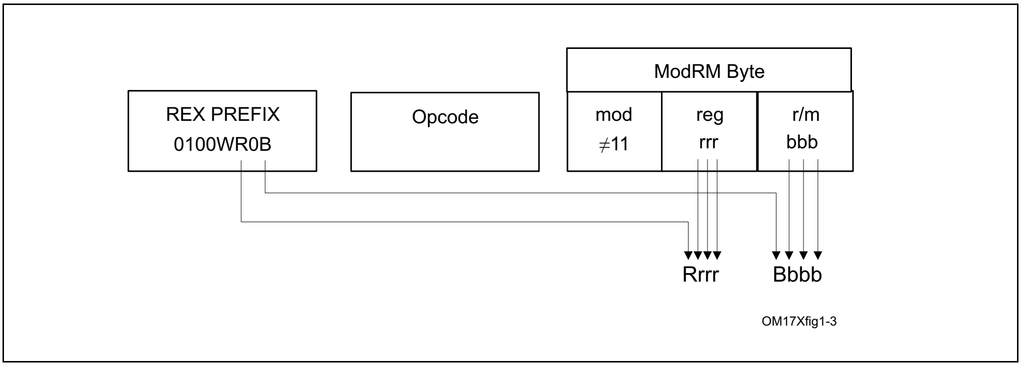


Figure 2-4. Memory Addressing Without an SIB Byte; REX.X Not Used

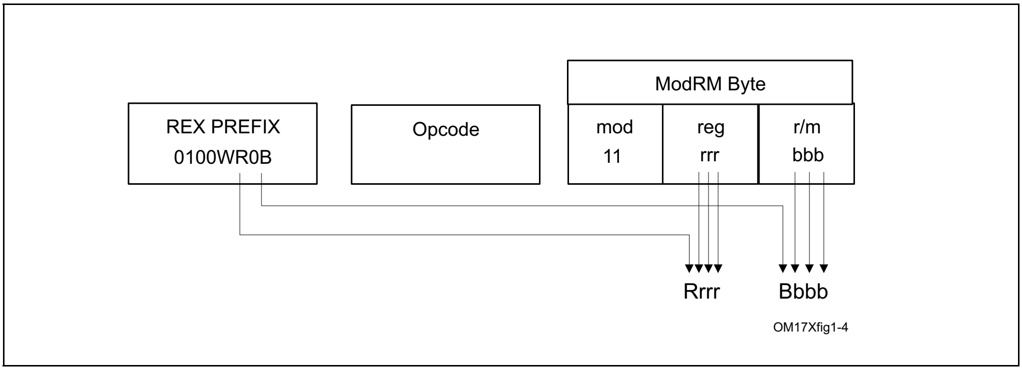


Figure 2-5. Register-Register Addressing (No Memory Operand); REX.X Not Used

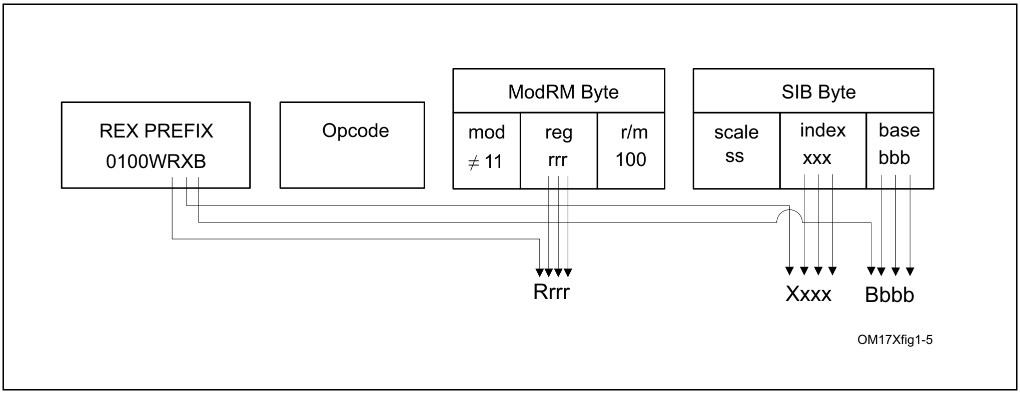


Figure 2-6. Memory Addressing With a SIB Byte

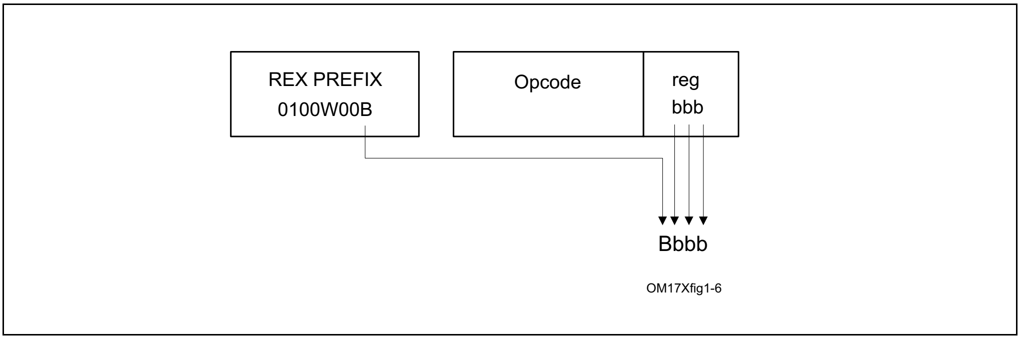


Figure 2-7. Register Operand Coded in Opcode Byte; REX.X & REX.R Not Used

In the IA-32 architecture, byte registers (AH, AL, BH, BL, CH, CL, DH, and DL) are encoded in the ModR/M byte's reg field, the r/m field or the opcode reg field as registers 0 through 7. REX prefixes provide an additional addressing capability for byte-registers that makes the least-significant byte of GPRs available for byte operations.

Certain combinations of the fields of the ModR/M byte and the SIB byte have special meaning for register encodings. For some combinations, fields expanded by the REX prefix are not decoded. Table 2-5 describes how each case behaves.

Table 2-5. Special Cases of REX Encodings

ModR/M or SIB	Sub-field Encodings	Compatibility Mode Operation	Compatibility Mode Implications	Additional Implications
ModR/M Byte	mod != 11 r/m = b*100(ESP)	SIB byte present.	SIB byte required for ESP-based addressing.	REX prefix adds a fourth bit (b) which is not decoded (don't care). SIB byte also required for R12-based addressing.
ModR/M Byte	mod = 0 r/m = b*101(EBP)	Base register not used.	EBP without a displacement must be done using mod = 01 with displacement of 0.	REX prefix adds a fourth bit (b) which is not decoded (don't care). Using RBP or R13 without displacement must be done using mod = 01 with a displacement of 0.
SIB Byte	index = 0100(ESP)	Index register not used.	ESP cannot be used as an index register.	REX prefix adds a fourth bit (b) which is decoded. There are no additional implications. The expanded index field allows distinguishing RSP from R12, therefore R12 can be used as an index.
SIB Byte	base = 0101(EBP)	Base register is unused if mod = 0.	Base register depends on mod encoding.	REX prefix adds a fourth bit (b) which is not decoded. This requires explicit displacement to be used with EBP/RBP or R13.

NOTES:

* Don't care about value of REX.B

2.2.1.3 Displacement

Addressing in 64-bit mode uses existing 32-bit ModR/M and SIB encodings. The ModR/M and SIB displacement sizes do not change. They remain 8 bits or 32 bits and are sign-extended to 64 bits.

2.2.1.4 Direct Memory-Offset MOVs

In 64-bit mode, direct memory-offset forms of the MOV instruction are extended to specify a 64-bit immediate absolute address. This address is called a moffset. No prefix is needed to specify this 64-bit memory offset. For these MOV instructions, the size of the memory offset follows the address-size default (64 bits in 64-bit mode). See Table 2-6.

Table 2-6. Direct Memory Offset Form of MOV

Opcode	Instruction
A0	MOV AL, moffset
A1	MOV EAX, moffset
A2	MOV moffset, AL
A3	MOV moffset, EAX

2.2.1.5 Immediates

In 64-bit mode, the typical size of immediate operands remains 32 bits. When the operand size is 64 bits, the processor sign-extends all immediates to 64 bits prior to their use.

Support for 64-bit immediate operands is accomplished by expanding the semantics of the existing move (MOV reg, imm16/32) instructions. These instructions (opcodes B8H – BFH) move 16-bits or 32-bits of immediate data (depending on the effective operand size) into a GPR. When the effective operand size is 64 bits, these instructions can be used to load an immediate into a GPR. A REX prefix is needed to override the 32-bit default operand size to a 64-bit operand size.

For example:

```
48 B8   8877665544332211   MOV RAX,1122334455667788H
```

2.2.1.6 RIP-Relative Addressing

A new addressing form, RIP-relative (relative instruction-pointer) addressing, is implemented in 64-bit mode. An effective address is formed by adding displacement to the 64-bit RIP of the next instruction.

In IA-32 architecture and compatibility mode, addressing relative to the instruction pointer is available only with control-transfer instructions. In 64-bit mode, instructions that use ModR/M addressing can use RIP-relative addressing. Without RIP-relative addressing, all ModR/M instruction modes address memory relative to zero.

RIP-relative addressing allows specific ModR/M modes to address memory relative to the 64-bit RIP using a signed 32-bit displacement. This provides an offset range of ±2GB from the RIP. Table 2-7 shows the ModR/M and SIB encodings for RIP-relative addressing. Redundant forms of 32-bit displacement-addressing exist in the current ModR/M and SIB encodings. There is one ModR/M encoding and there are several SIB encodings. RIP-relative addressing is encoded using a redundant form.

In 64-bit mode, the ModR/M Disp32 (32-bit displacement) encoding is re-defined to be RIP+Disp32 rather than displacement-only. See Table 2-7.

Table 2-7. RIP-Relative Addressing

ModR/M and SIB Sub-field Encodings		Compatibility Mode Operation	64-bit Mode Operation	Additional Implications in 64-bit mode
ModR/M Byte	mod = 00	Disp32	RIP + Disp32	Must use SIB form with normal (zero-based) displacement addressing
	r/m = 101 (none)			
SIB Byte	base = 101 (none)	if mod = 00, Disp32	Same as legacy	None
	index = 100 (none)			
	scale = 0, 1, 2, 4			

The ModR/M encoding for RIP-relative addressing does not depend on using prefix. Specifically, the r/m bit field encoding of 101B (used to select RIP-relative addressing) is not affected by the REX prefix. For example, selecting R13 (REX.B = 1, r/m = 101B) with mod = 00B still results in RIP-relative addressing. The 4-bit r/m field of REX.B combined with ModR/M is not fully decoded. In order to address R13 with no displacement, software must encode R13 + 0 using a 1-byte displacement of zero.

RIP-relative addressing is enabled by 64-bit mode, not by a 64-bit address-size. The use of the address-size prefix does not disable RIP-relative addressing. The effect of the address-size prefix is to truncate and zero-extend the computed effective address to 32 bits.

2.2.1.7 Default 64-Bit Operand Size

In 64-bit mode, two groups of instructions have a default operand size of 64 bits (do not need a REX prefix for this operand size). These are:

- Near branches
- All instructions, except far branches, that implicitly reference the RSP

2.2.2 Additional Encodings for Control and Debug Registers

In 64-bit mode, more encodings for control and debug registers are available. The REX.R bit is used to modify the ModR/M reg field when that field encodes a control or debug register (see Table 2-4). These encodings enable the processor to address CR8-CR15 and DR8- DR15. An additional control register (CR8) is defined in 64-bit mode. CR8 becomes the Task Priority Register (TPR).

In the first implementation of IA-32e mode, CR9-CR15 and DR8-DR15 are not implemented. Any attempt to access unimplemented registers results in an invalid-opcode exception (#UD).

2.3 INTEL® ADVANCED VECTOR EXTENSIONS (INTEL® AVX)

Intel AVX instructions are encoded using an encoding scheme that combines prefix bytes, opcode extension field, operand encoding fields, and vector length encoding capability into a new prefix, referred to as VEX. In the VEX encoding scheme, the VEX prefix may be two or three bytes long, depending on the instruction semantics. Despite the two-byte or three-byte length of the VEX prefix, the VEX encoding format provides a more compact representation/packing of the components of encoding an instruction in Intel 64 architecture. The VEX encoding scheme also allows more head-room for future growth of Intel 64 architecture.

2.3.1 Instruction Format

Instruction encoding using VEX prefix provides several advantages:

- Instruction syntax support for three operands and up-to four operands when necessary. For example, the third source register used by VBLENDVPD is encoded using bits 7:4 of the immediate byte.
- Encoding support for vector length of 128 bits (using XMM registers) and 256 bits (using YMM registers)
- Encoding support for instruction syntax of non-destructive source operands.
- Elimination of escape opcode byte (0FH), SIMD prefix byte (66H, F2H, F3H) via a compact bit field representation within the VEX prefix.
- Elimination of the need to use REX prefix to encode the extended half of general-purpose register sets (R8-R15) for direct register access, memory addressing, or accessing XMM8-XMM15 (including YMM8-YMM15).
- Flexible and more compact bit fields are provided in the VEX prefix to retain the full functionality provided by REX prefix. REX.W, REX.X, REX.B functionalities are provided in the three-byte VEX prefix only because only a subset of SIMD instructions need them.
- Extensibility for future instruction extensions without significant instruction length increase.

Figure 2-8 shows the Intel 64 instruction encoding format with VEX prefix support. Legacy instruction without a VEX prefix is fully supported and unchanged. The use of VEX prefix in an Intel 64 instruction is optional, but a VEX prefix is required for Intel 64 instructions that operate on YMM registers or support three and four operand syntax. VEX prefix is not a constant-valued, “single-purpose” byte like 0FH, 66H, F2H, F3H in legacy SSE instructions. VEX prefix provides substantially richer capability than the REX prefix.

# Bytes	2,3	1	1	0,1	0,1,2,4	0,1
[Prefixes]	[VEX]	OPCODE	ModR/M	[SIB]	[DISP]	[IMM]

Figure 2-8. Instruction Encoding Format with VEX Prefix

2.3.2 VEX and the LOCK prefix

Any VEX-encoded instruction with a LOCK prefix preceding VEX will #UD.

2.3.3 VEX and the 66H, F2H, and F3H prefixes

Any VEX-encoded instruction with a 66H, F2H, or F3H prefix preceding VEX will #UD.

2.3.4 VEX and the REX prefix

Any VEX-encoded instruction with a REX prefix proceeding VEX will #UD.

2.3.5 The VEX Prefix

The VEX prefix is encoded in either the two-byte form (the first byte must be C5H) or in the three-byte form (the first byte must be C4H). The two-byte VEX is used mainly for 128-bit, scalar, and the most common 256-bit AVX instructions; while the three-byte VEX provides a compact replacement of REX and 3-byte opcode instructions (including AVX and FMA instructions). Beyond the first byte of the VEX prefix, it consists of a number of bit fields providing specific capability, they are shown in Figure 2-9.

The bit fields of the VEX prefix can be summarized by its functional purposes:

- Non-destructive source register encoding (applicable to three and four operand syntax): This is the first source operand in the instruction syntax. It is represented by the notation, VEX.vvvv. This field is encoded using 1's complement form (inverted form), i.e. XMM0/YMM0/R0 is encoded as 1111B, XMM15/YMM15/R15 is encoded as 0000B.
- Vector length encoding: This 1-bit field represented by the notation VEX.L. L= 0 means vector length is 128 bits wide, L=1 means 256 bit vector. The value of this field is written as VEX.128 or VEX.256 in this document to distinguish encoded values of other VEX bit fields.

- REX prefix functionality: Full REX prefix functionality is provided in the three-byte form of VEX prefix. However the VEX bit fields providing REX functionality are encoded using 1's complement form, i.e. XMM0/YMM0/R0 is encoded as 1111B, XMM15/YMM15/R15 is encoded as 0000B.
  - Two-byte form of the VEX prefix only provides the equivalent functionality of REX.R, using 1's complement encoding. This is represented as VEX.R.
  - Three-byte form of the VEX prefix provides REX.R, REX.X, REX.B functionality using 1's complement encoding and three dedicated bit fields represented as VEX.R, VEX.X, VEX.B.
  - Three-byte form of the VEX prefix provides the functionality of REX.W only to specific instructions that need to override default 32-bit operand size for a general purpose register to 64-bit size in 64-bit mode. For those applicable instructions, VEX.W field provides the same functionality as REX.W. VEX.W field can provide completely different functionality for other instructions.

  Consequently, the use of REX prefix with VEX encoded instructions is not allowed. However, the intent of the REX prefix for expanding register set is reserved for future instruction set extensions using VEX prefix encoding format.
- Compaction of SIMD prefix: Legacy SSE instructions effectively use SIMD prefixes (66H, F2H, F3H) as an opcode extension field. VEX prefix encoding allows the functional capability of such legacy SSE instructions (operating on XMM registers, bits 255:128 of corresponding YMM unmodified) to be encoded using the VEX.pp field without the presence of any SIMD prefix. The VEX-encoded 128-bit instruction will zero-out bits 255:128 of the destination register. VEX-encoded instruction may have 128 bit vector length or 256 bits length.
- Compaction of two-byte and three-byte opcode: More recently introduced legacy SSE instructions employ two and three-byte opcode. The one or two leading bytes are: 0FH, and 0FH 3AH/0FH 38H. The one-byte escape (0FH) and two-byte escape (0FH 3AH, 0FH 38H) can also be interpreted as an opcode extension field. The VEX.mmmmm field provides compaction to allow many legacy instruction to be encoded without the constant byte sequence, 0FH, 0FH 3AH, 0FH 38H. These VEX-encoded instruction may have 128 bit vector length or 256 bits length.

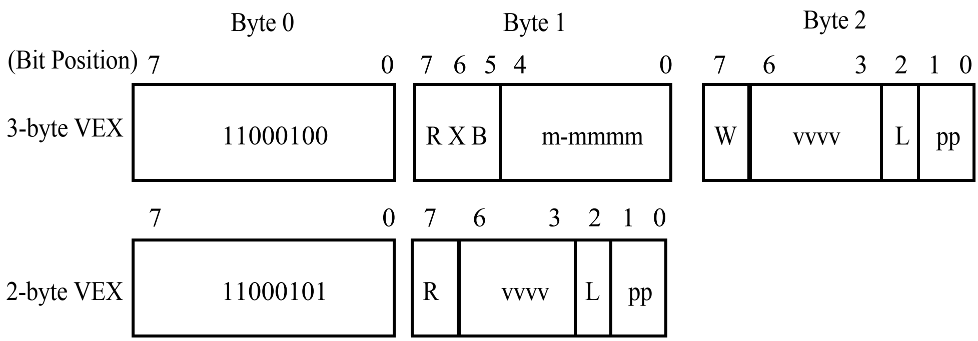
The VEX prefix is required to be the last prefix and immediately precedes the opcode bytes. It must follow any other prefixes. If VEX prefix is present a REX prefix is not supported.

The 3-byte VEX leaves room for future expansion with 3 reserved bits. REX and the 66h/F2h/F3h prefixes are reclaimed for future use.

VEX prefix has a two-byte form and a three byte form. If an instruction syntax can be encoded using the two-byte form, it can also be encoded using the three byte form of VEX. The latter increases the length of the instruction by one byte. This may be helpful in some situations for code alignment.

The VEX prefix supports 256-bit versions of floating-point SSE, SSE2, SSE3, and SSE4 instructions. Note, certain new instruction functionality can only be encoded with the VEX prefix.

The VEX prefix will #UD on any instruction containing MMX register sources or destinations.



R: REX.R in 1's complement (inverted) form
- 1: Same as REX.R=0 (must be 1 in 32-bit mode)
- 0: Same as REX.R=1 (64-bit mode only)

X: REX.X in 1's complement (inverted) form
- 1: Same as REX.X=0 (must be 1 in 32-bit mode)
- 0: Same as REX.X=1 (64-bit mode only)

B: REX.B in 1's complement (inverted) form
- 1: Same as REX.B=0 (Ignored in 32-bit mode).
- 0: Same as REX.B=1 (64-bit mode only)

W: opcode specific (use like REX.W, or used for opcode extension, or ignored, depending on the opcode byte)

m-mmmm:
- 00000: Reserved for future use (will #UD)
- 00001: implied 0F leading opcode byte
- 00010: implied 0F 38 leading opcode bytes
- 00011: implied 0F 3A leading opcode bytes
- 00100-11111: Reserved for future use (will #UD)

vvvv: a register specifier (in 1's complement form) or 1111 if unused.

L: Vector Length
- 0: scalar or 128-bit vector
- 1: 256-bit vector

pp: opcode extension providing equivalent functionality of a SIMD prefix
- 00: None
- 01: 66
- 10: F3
- 11: F2

Figure 2-9. VEX bitfields

The following subsections describe the various fields in two or three-byte VEX prefix:

2.3.5.1 VEX Byte 0, bits[7:0]

VEX Byte 0, bits [7:0] must contain the value 11000101b (C5h) or 11000100b (C4h). The 3-byte VEX uses the C4h first byte, while the 2-byte VEX uses the C5h first byte.

2.3.5.2 VEX Byte 1, bit [7] - 'R'

VEX Byte 1, bit [7] contains a bit analogous to a bit inverted REX.R. In protected and compatibility modes the bit must be set to '1' otherwise the instruction is LES or LDS.

This bit is present in both 2- and 3-byte VEX prefixes.

The usage of WRXB bits for legacy instructions is explained in detail section 2.2.1.2 of Intel 64 and IA-32 Architectures Software developer's manual, Volume 2A.

This bit is stored in bit inverted format.

2.3.5.3 3-byte VEX byte 1, bit[6] - 'X'

Bit[6] of the 3-byte VEX byte 1 encodes a bit analogous to a bit inverted REX.X. It is an extension of the SIB Index field in 64-bit modes. In 32-bit modes, this bit must be set to '1' otherwise the instruction is LES or LDS.

This bit is available only in the 3-byte VEX prefix.

This bit is stored in bit inverted format.

2.3.5.4 3-byte VEX byte 1, bit[5] - 'B'

Bit[5] of the 3-byte VEX byte 1 encodes a bit analogous to a bit inverted REX.B. In 64-bit modes, it is an extension of the ModR/M r/m field, or the SIB base field. In 32-bit modes, this bit is ignored.

This bit is available only in the 3-byte VEX prefix.

This bit is stored in bit inverted format.

2.3.5.5 3-byte VEX byte 2, bit[7] - 'W'

Bit[7] of the 3-byte VEX byte 2 is represented by the notation VEX.W. It can provide following functions, depending on the specific opcode.

- For AVX instructions that have equivalent legacy SSE instructions (typically these SSE instructions have a general-purpose register operand with its operand size attribute promotable by REX.W), if REX.W promotes the operand size attribute of the general-purpose register operand in legacy SSE instruction, VEX.W has same meaning in the corresponding AVX equivalent form. In 32-bit modes, VEX.W is silently ignored.

- For AVX instructions that have equivalent legacy SSE instructions (typically these SSE instructions have operands with their operand size attribute fixed and not promotable by REX.W), if REX.W is don't care in legacy SSE instruction, VEX.W is ignored in the corresponding AVX equivalent form irrespective of mode.
- For new AVX instructions where VEX.W has no defined function (typically these meant the combination of the opcode byte and VEX.mmmmm did not have any equivalent SSE functions), VEX.W is reserved as zero and setting to other than zero will cause instruction to #UD.

2.3.5.6 2-byte VEX Byte 1, bits[6:3] and 3-byte VEX Byte 2, bits [6:3]- 'vvvv' the Source or dest Register Specifier

In 32-bit mode the VEX first byte C4 and C5 alias onto the LES and LDS instructions. To maintain compatibility with existing programs the VEX 2nd byte, bits [7:6] must be 11b. To achieve this, the VEX payload bits are selected to place only inverted, 64-bit valid fields (extended register selectors) in these upper bits.

The 2-byte VEX Byte 1, bits [6:3] and the 3-byte VEX, Byte 2, bits [6:3] encode a field (shorthand VEX.vvvv) that for instructions with 2 or more source registers and an XMM or YMM or memory destination encodes the first source register specifier stored in inverted (1's complement) form.

VEX.vvvv is not used by the instructions with one source (except certain shifts, see below) or on instructions with no XMM or YMM or memory destination. If an instruction does not use VEX.vvvv then it should be set to 1111b otherwise instruction will #UD.

In 64-bit mode all 4 bits may be used. See Table 2-8 for the encoding of the XMM or YMM registers. In 32-bit and 16-bit modes bit 6 must be 1 (if bit 6 is not 1, the 2-byte VEX version will generate LDS instruction and the 3-byte VEX version will ignore this bit).

Table 2-8. VEX.vvvv to register name mapping

VEX.vvvv	Dest Register	Valid in Legacy/Compatibility 32-bit modes?
1111B	XMM0/YMM0	Valid
1110B	XMM1/YMM1	Valid
1101B	XMM2/YMM2	Valid
1100B	XMM3/YMM3	Valid
1011B	XMM4/YMM4	Valid
1010B	XMM5/YMM5	Valid
1001B	XMM6/YMM6	Valid
1000B	XMM7/YMM7	Valid
0111B	XMM8/YMM8	Invalid
0110B	XMM9/YMM9	Invalid
0101B	XMM10/YMM10	Invalid
0100B	XMM11/YMM11	Invalid
0011B	XMM12/YMM12	Invalid
0010B	XMM13/YMM13	Invalid
0001B	XMM14/YMM14	Invalid
0000B	XMM15/YMM15	Invalid

The VEX.vvvv field is encoded in bit inverted format for accessing a register operand.

2.3.6 Instruction Operand Encoding and VEX.vvvv, ModR/M

VEX-encoded instructions support three-operand and four-operand instruction syntax. Some VEX-encoded instructions have syntax with less than three operands, e.g. VEX-encoded pack shift instructions support one source operand and one destination operand).

The roles of VEX.vvvv, reg field of ModR/M byte (ModR/M.reg), r/m field of ModR/M byte (ModR/M.r/m) with respect to encoding destination and source operands vary with different type of instruction syntax.

The role of VEX.vvvv can be summarized to three situations:

- VEX.vvvv encodes the first source register operand, specified in inverted (1's complement) form and is valid for instructions with 2 or more source operands.
- VEX.vvvv encodes the destination register operand, specified in 1's complement form for certain vector shifts. The instructions where VEX.vvvv is used as a destination are listed in Table 2-9. The notation in the "Opcode" column in Table 2-9 is described in detail in section 3.1.1.

- VEX.vvvv does not encode any operand, the field is reserved and should contain 1111b.

Table 2-9. Instructions with a VEX.vvvv destination

Opcode	Instruction mnemonic
VEX.NDD.128.66.0F 73 /7 ib	VPSLLDQ xmm1, xmm2, imm8
VEX.NDD.128.66.0F 73 /3 ib	VPSRLDQ xmm1, xmm2, imm8
VEX.NDD.128.66.0F 71 /2 ib	VPSRLW xmm1, xmm2, imm8
VEX.NDD.128.66.0F 72 /2 ib	VPSRLD xmm1, xmm2, imm8
VEX.NDD.128.66.0F 73 /2 ib	VPSRLQ xmm1, xmm2, imm8
VEX.NDD.128.66.0F 71 /4 ib	VPSRAW xmm1, xmm2, imm8
VEX.NDD.128.66.0F 72 /4 ib	VPSRAD xmm1, xmm2, imm8
VEX.NDD.128.66.0F 71 /6 ib	VPSLLW xmm1, xmm2, imm8
VEX.NDD.128.66.0F 72 /6 ib	VPSLLD xmm1, xmm2, imm8
VEX.NDD.128.66.0F 73 /6 ib	VPSLLQ xmm1, xmm2, imm8

The role of ModR/M.r/m field can be summarized to two situations:

- ModR/M.r/m encodes the instruction operand that references a memory address.
- For some instructions that do not support memory addressing semantics, ModR/M.r/m encodes either the destination register operand or a source register operand.

The role of ModR/M.reg field can be summarized to two situations:

- ModR/M.reg encodes either the destination register operand or a source register operand.
- For some instructions, ModR/M.reg is treated as an opcode extension and not used to encode any instruction operand.

For instruction syntax that support four operands, VEX.vvvv, ModR/M.r/m, ModR/M.reg encodes three of the four operands. The role of bits 7:4 of the immediate byte serves the following situation:

- Imm8[7:4] encodes the third source register operand.

2.3.6.1 3-byte VEX byte 1, bits[4:0] - “m-mmmm”

Bits[4:0] of the 3-byte VEX byte 1 encode an implied leading opcode byte (0F, 0F 38, or 0F 3A). Several bits are reserved for future use and will #UD unless 0.

Table 2-10. VEX.m-mmmm interpretation

VEX.m-mmmm	Implied Leading Opcode Bytes
00000B	Reserved
00001B	0F
00010B	0F 38
00011B	0F 3A
00100-11111B	Reserved
(2-byte VEX)	0F

VEX.m-mmmm is only available on the 3-byte VEX. The 2-byte VEX implies a leading 0Fh opcode byte.

2.3.6.2 2-byte VEX byte 1, bit[2], and 3-byte VEX byte 2, bit [2]- “L”

The vector length field, VEX.L, is encoded in bit[2] of either the second byte of 2-byte VEX, or the third byte of 3-byte VEX. If “VEX.L = 1”, it indicates 256-bit vector operation. “VEX.L = 0” indicates scalar and 128-bit vector operations.

The instruction VZEROUPPER is a special case that is encoded with VEX.L = 0, although its operation zero’s bits 255:128 of all YMM registers accessible in the current operating mode.

See the following table.

Table 2-11. VEX.L interpretation

VEX.L	Vector Length
0	128-bit (or 32/64-bit scalar)
1	256-bit

2.3.6.3 2-byte VEX byte 1, bits[1:0], and 3-byte VEX byte 2, bits [1:0]- “pp”

Up to one implied prefix is encoded by bits[1:0] of either the 2-byte VEX byte 1 or the 3-byte VEX byte 2. The prefix behaves as if it was encoded prior to VEX, but after all other encoded prefixes.

See the following table.

Table 2-12. VEX.pp interpretation

pp	Implies this prefix after other prefixes but before VEX
00B	None
01B	66
10B	F3
11B	F2

2.3.7 The Opcode Byte

One (and only one) opcode byte follows the 2 or 3 byte VEX. Legal opcodes are specified in Appendix B, in color. Any instruction that uses illegal opcode will #UD.

2.3.8 The MODRM, SIB, and Displacement Bytes

The encodings are unchanged but the interpretation of reg_field or rm_field differs (see above).

2.3.9 The Third Source Operand (Immediate Byte)

VEX-encoded instructions can support instruction with a four operand syntax. VBLENDVPD, VBLENDVPS, and PBLENDVB use imm8[7:4] to encode one of the source registers.

2.3.10 AVX Instructions and the Upper 128-bits of YMM registers

If an instruction with a destination XMM register is encoded with a VEX prefix, the processor zeroes the upper bits (above bit 128) of the equivalent YMM register . Legacy SSE instructions without VEX preserve the upper bits.

2.3.11 AVX Instruction Length

The AVX instructions described in this document (including VEX and ignoring other prefixes) do not exceed 11 bytes in length, but may increase in the future. The maximum length of an Intel 64 and IA-32 instruction remains 15 bytes.

2.4 INSTRUCTION EXCEPTION SPECIFICATION

To look up the exceptions of legacy 128-bit SIMD instruction, 128-bit VEX-encoded instructions, and 256-bit VEX-encoded instruction, Table 2-13 summarizes the exception behavior into separate classes, with detailed exception conditions defined in sub-sections 2.4.1 through 2.4.9. For example, ADDPS contains the entry:

"See Exceptions Type 2"

In this entry, *"Type2"* can be looked up in Table 2-13.

The instruction's corresponding CPUID feature flag can be identified in the fourth column of the Instruction summary table.

Note: #UD on CPUID feature flags=0 is not guaranteed in a virtualized environment if the hardware supports the feature flag.

NOTE

Instructions that operate only with MMX, X87, or general-purpose registers are not covered by the exception classes defined in this section.

Table 2-13. Exception class description

Exception Class	Instruction set	Mem arg	Floating-Point Exceptions (#XM)
Type 1	AVX, Legacy SSE	16/32 byte explicitly aligned	none
Type 2	AVX, Legacy SSE	16/32 byte not explicitly aligned	yes
Type 3	AVX, Legacy SSE	< 16 byte	yes
Type 4	AVX, Legacy SSE	16/32 byte not explicitly aligned	no
Type 5	AVX, Legacy SSE	< 16 byte	no
Type 6	AVX (no Legacy SSE)	Varies	(At present, none do)
Type 7	AVX, Legacy SSE	none	none
Type 8	AVX	none	none
Type 9	AVX	4 byte	none

See Table 2-14 for lists of instructions in each exception class.

Table 2-14. Instructions in each Exception Class

Exception Class	Instruction
Type 1	(V)MOVAPD, (V)MOVAPS, (V)MOVDQA, (V)MOVNTDQ, (V)MOVNTDQA, (V)MOVNTPD, (V)MOVNTPS
Type 2	(V)ADDPD, (V)ADDPS, (V)ADDSUBPD, (V)ADDSUBPS, (V)CMPPD, (V)CMPPS, (V)CVTDQ2PS, (V)CVTPD2DQ, (V)CVTPD2PS, (V)CVTPS2DQ, (V)CVTTPD2DQ, (V)CVTTPS2DQ, (V)DIVPD, (V)DIVPS, (V)DPPD*, (V)DPPS*, (V)HADDPD, (V)HADDPS, (V)HSUBPD, (V)HSUBPS, (V)MAXPD, (V)MAXPS, (V)MINPD, (V)MINPS, (V)MULPD, (V)MULPS, (V)ROUNDPD, (V)ROUNDPS, (V)SQRTPD, (V)SQRTPS, (V)SUBPD, (V)SUBPS
Type 3	(V)ADDSD, (V)ADDSS, (V)CMPSD, (V)CMPSS, (V)COMISD, (V)COMISS, (V)CVTPS2PD, (V)CVTSD2SI, (V)CVTSD2SS, (V)CVTSI2SD, (V)CVTSI2SS, (V)CVTSS2SD, (V)CVTSS2SI, (V)CVTTSD2SI, (V)CVTTSS2SI, (V)DIVSD, (V)DIVSS, (V)MAXSD, (V)MAXSS, (V)MINSD, (V)MINSS, (V)MULSD, (V)MULSS, (V)ROUNDSD, (V)ROUNDSS, (V)SQRTSD, (V)SQRTSS, (V)SUBSD, (V)SUBSS, (V)UCOMISD, (V)UCOMISS
Type 4	(V)AESDEC, (V)AESDECLAST, (V)AESENC, (V)AESENCLAST, (V)AESIMC, (V)AESKEYGENASSIST, (V)ANDPD, (V)ANDPS, (V)ANDNPD, (V)ANDNPS, (V)BLENDPD, (V)BLENDPS, VBLENDVPD, VBLENDVPS, (V)LDDQU, (V)MASKMOVDQU, (V)PTEST, VTESTPS, VTESTPD, (V)MOVDQU*, (V)MOVSHDUP, (V)MOVSLDUP, (V)MOVUPD*, (V)MOVUPS*, (V)MPSADBW, (V)ORPD, (V)ORPS, (V)PABSB, (V)PABSW, (V)PABSD, (V)PACKSSWB, (V)PACKSSDW, (V)PACKUSWB, (V)PACKUSDW, (V)PADDB, (V)PADDW, (V)PADDD, (V)PADDQ, (V)PADDSB, (V)PADDSW, (V)PADDUSB, (V)PADDUSW, (V)PALIGNR, (V)PAND, (V)PANDN, (V)PAVGB, (V)PAVGW, (V)PBLENDVB, (V)PBLENDW, (V)PCMP(E/I)STRI/M, (V)PCMPEQB, (V)PCMPEQW, (V)PCMPEQD, (V)PCMPEQQ, (V)PCMPGTB, (V)PCMPGTW, (V)PCMPGTD, (V)PCMPGTQ, (V)PCLMULQDQ, (V)PHADDW, (V)PHADDD, (V)PHADDSW, (V)PHMINPOSUW, (V)PHSUBD, (V)PHSUBW, (V)PHSUBSW,
	(V)PMADDWD, (V)PMADDUBSW, (V)PMAXSB, (V)PMAXSW, (V)PMAXSD, (V)PMAXUB, (V)PMAXUW, (V)PMAXUD, (V)PMINSB, (V)PMINSW, (V)PMINSD, (V)PMINUB, (V)PMINUW, (V)PMINUD, (V)PMULHUW, (V)PMULHRSW, (V)PMULHW, (V)PMULLW, (V)PMULLD, (V)PMULUDQ, (V)PMULDQ, (V)POR, (V)PSADBW, (V)PSHUFB, (V)PSHUFD, (V)PSHUFHW, (V)PSHUFLW, (V)PSIGNB, (V)PSIGNW, (V)PSIGND, (V)PSLLW, (V)PSLLD, (V)PSLLQ, (V)PSRAW, (V)PSRAD, (V)PSRLW, (V)PSRLD, (V)PSRLQ, (V)PSUBB, (V)PSUBW, (V)PSUBD, (V)PSUBQ, (V)PSUBSB, (V)PSUBSW, (V)PUNPCKHBW, (V)PUNPCKHWD, (V)PUNPCKHDQ, (V)PUNPCKHQDQ, (V)PUNPCKLBW, (V)PUNPCKLWD, (V)PUNPCKLDQ, (V)PUNPCKLQDQ, (V)PXOR, (V)RCPPS, (V)RSQRTPS, (V)SHUFPD, (V)SHUFPS, (V)UNPCKHPD, (V)UNPCKHPS, (V)UNPCKLPD, (V)UNPCKLPS, (V)XORPD, (V)XORPS
Type 5	(V)CVTDQ2PD, (V)EXTRACTPS, (V)INSERTPS, (V)MOVD, (V)MOVQ, (V)MOVDDUP, (V)MOVLPD, (V)MOVLPS, (V)MOVHPD, (V)MOVHPS, (V)MOVSD, (V)MOVSS, (V)PEXTRB, (V)PEXTRD, (V)PEXTRW, (V)PEXTRQ, (V)PINSRB, (V)PINSRD, (V)PINSRW, (V)PINSRQ, (V)RCPSS, (V)RSQRTSS, (V)PMOVSX/ZX

Exception Class	Instruction
Type 6	VEXTRACTF128, VPERMILPD, VPERMILPS, VPERM2F128, VBROADCASTSS, VBROADCASTSD, VBROADCASTF128, VINSERTF128, VMASKMOVPS**, VMASKMOVPD**
Type 7	(V)MOVLHPS, (V)MOVHLPS, (V)MOVMSKPD, (V)MOVMSKPS, (V)PMOVMSKB, (V)PSLLDQ, (V)PSRLDQ, (V)PSLLW, (V)PSLLD, (V)PSLLQ, (V)PSRAW, (V)PSRAD, (V)PSRLW, (V)PSRLD, (V)PSRLQ
Type 8	VZEROALL, VZEROUPPER
Type 9	VLDMXCSR*, VSTMXCSR

(*) - Additional exception restrictions are present - see the Instruction description for details

(**) - Instruction behavior on alignment check reporting with mask bits of less than all 1s are the same as with mask bits of all 1s, i.e. no alignment checks are performed.

Table 2-14 classifies exception behaviors for AVX instructions. Within each class of exception conditions that are listed in Table 2-17 through Table 2-23, certain subsets of AVX instructions may be subject to #UD exception depending on the encoded value of the VEX.L field. Table 2-16 provides supplemental information of AVX instructions that may be subject to #UD exception if encoded with incorrect values in the VEX.W or VEX.L field.

Table 2-15. #UD Exception and VEX.W=1 Encoding

Exception Class	#UD If VEX.W = 1 in all modes	#UD If VEX.W = 1 in non-64-bit modes
Type 1		
Type 2		
Type 3		
Type 4	VBLENDVPD, VBLENDVPS, VPBLENDVB, VTESTPD, VTESTPS	
Type 5		VPEXTRQ, VPINSRQ,
Type 6	VEXTRACTF128, VPERMILPD, VPERMILPS, VPERM2F128, VBROADCASTSS, VBROADCASTSD, VBROADCASTF128, VINSERTF128, VMASKMOVPS, VMASKMOVPD	
Type 7		
Type 8		
Type 9		

Table 2-16. #UD Exception and VEX.L Field Encoding

Exception Class	#UD If VEX.L = 0	#UD If VEX.L = 1
Type 1		VMOVNTDQA
Type 2		VDPPD
Type 3		
Type 4		VMASKMOVDQU, VMPSADBW, VPABSB/W/D, VPACKSSWB/DW, VPACKUSWB/DW, VPADDB/W/D, VPADDQ, VPADDSB/W, VPADDUSB/W, VPALIGNR, VPAND, VPANDN, VPAVGB/W, VPBLENDVB, VPBLENDW, VPCMP(E/I)STRI/M, VPCMPEQB/W/D/Q, VPCMPGTB/W/D/Q, VPHADDW/D, VPHADDSW, VPHMINPOSUW, VPHSUBD/W, VPHSUBSW, VPMADDWD, VPMADDUBSW, VPMAXSB/W/D, VPMAXUB/W/D, VPMINSB/W/D, VPMINUB/W/D, VPMULHUW, VPMULHRSW, VPMULHW/LW, VPMULLD, VPMULUDQ, VPMULDQ, VPOR, VPSADBW, VPSHUFB/D, VPSHUFHW/LW, VPSIGNB/W/D, VPSLLW/D/Q, VPSRAW/D, VPSRLW/D/Q, VPSUBB/W/D/Q, VPSUBSB/W, VPUNPCKHBW/WD/DQ, VPUNPCKHQDQ, VPUNPCKLBW/WD/DQ, VPUNPCKLQDQ, VPXOR
Type 5		VEXTRACTPS, VINSERTPS, VMOVD, VMOVQ, VMOVLPD, VMOVLPS, VMOVHPD, VMOVHPS, VPEXTRB, VPEXTRD, VPEXTRW, VPEXTRQ, VPINSRB, VPINSRD, VPINSRW, VPINSRQ, VPMOVSX/ZX
Type 6	VEXTRACTF128, VPERM2F128, VBROADCASTSD, VBROADCASTF128, VINSERTF128,	
Type 7		VMOVLHPS, VMOVHLPS, VPMOVMSKB, VPSLLDQ, VPSRLDQ, VPSLLW, VPSLLD, VPSLLQ, VPSRAW, VPSRAD, VPSRLW, VPSRLD, VPSRLQ
Type 8		
Type 9		VLDMXCSR, VSTMXCSR

2.4.1 Exceptions Type 1 (Aligned memory reference)

Table 2-17. Type 1 Class Exception Conditions

Exception	Real	Virtual 80x86	Protected and Compatibility	64-bit	Cause of Exception
Invalid Opcode, #UD	X	X			VEX prefix.
			X	X	VEX prefix: If XCR0[2:1] != '11b'. If CR4.OSXSAVE[bit 18]=0.
	X	X	X	X	Legacy SSE instruction: If CR0.EM[bit 2] = 1. If CR4.OSFXSR[bit 9] = 0.
	X	X	X	X	If preceded by a LOCK prefix (F0H).
			X	X	If any REX, F2, F3, or 66 prefixes precede a VEX prefix.
	X	X	X	X	If any corresponding CPUID feature flag is '0'.
Device Not Available, #NM	X	X	X	X	If CR0.TS[bit 3]=1.
Stack, SS(0)			X		For an illegal address in the SS segment.
				X	If a memory address referencing the SS segment is in a non-canonical form.
General Protection, #GP(0)			X	X	VEX.256: Memory operand is not 32-byte aligned. VEX.128: Memory operand is not 16-byte aligned.
	X	X	X	X	Legacy SSE: Memory operand is not 16-byte aligned.
			X		For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
				X	If the memory address is in a non-canonical form.
	X	X			If any part of the operand lies outside the effective address space from 0 to FFFFH.
Page Fault #PF(fault-code)		X	X	X	For a page fault.

2.4.2 Exceptions Type 2 (>=16 Byte Memory Reference, Unaligned)

Table 2-18. Type 2 Class Exception Conditions

Exception	Real	Virtual 8086	Protected and Compatibility	64-bit	Cause of Exception
Invalid Opcode, #UD	X	X			VEX prefix.
	X	X	X	X	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
			X	X	VEX prefix: If XCR0[2:1] != '11b'. If CR4.OSXSAVE[bit 18]=0.
	X	X	X	X	Legacy SSE instruction: If CR0.EM[bit 2] = 1. If CR4.OSFXSR[bit 9] = 0.
	X	X	X	X	If preceded by a LOCK prefix (F0H).
			X	X	If any REX, F2, F3, or 66 prefixes precede a VEX prefix.
	X	X	X	X	If any corresponding CPUID feature flag is '0'.
Device Not Available, #NM	X	X	X	X	If CR0.TS[bit 3]=1.
Stack, SS(0)			X		For an illegal address in the SS segment.
				X	If a memory address referencing the SS segment is in a non-canonical form.
General Protection, #GP(0)	X	X	X	X	Legacy SSE: Memory operand is not 16-byte aligned.
			X		For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
				X	If the memory address is in a non-canonical form.
	X	X			If any part of the operand lies outside the effective address space from 0 to FFFFH.
Page Fault #PF(fault-code)		X	X	X	For a page fault.
SIMD Floating-point Exception, #XM	X	X	X	X	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.

2.4.3 Exceptions Type 3 (<16 Byte memory argument)

Table 2-19. Type 3 Class Exception Conditions

Exception	Real	Virtual 80x86	Protected and Compatibility	64-bit	Cause of Exception
Invalid Opcode, #UD	X	X			VEX prefix.
	X	X	X	X	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
			X	X	VEX prefix: If XCR0[2:1] != '11b'. If CR4.OSXSAVE[bit 18]=0.
	X	X	X	X	Legacy SSE instruction: If CR0.EM[bit 2] = 1. If CR4.OSFXSR[bit 9] = 0.
	X	X	X	X	If preceded by a LOCK prefix (F0H).
			X	X	If any REX, F2, F3, or 66 prefixes precede a VEX prefix.
	X	X	X	X	If any corresponding CPUID feature flag is '0'.
Device Not Available, #NM	X	X	X	X	If CR0.TS[bit 3]=1.
Stack, SS(0)			X		For an illegal address in the SS segment.
				X	If a memory address referencing the SS segment is in a non-canonical form.
General Protection, #GP(0)			X		For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
				X	If the memory address is in a non-canonical form.
	X	X			If any part of the operand lies outside the effective address space from 0 to FFFFH.
Page Fault #PF(fault-code)		X	X	X	For a page fault.
Alignment Check #AC(0)		X	X	X	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
SIMD Floating-point Exception, #XM	X	X	X	X	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.

2.4.4 Exceptions Type 4 (>=16 Byte mem arg no alignment, no floating-point exceptions)

Table 2-20. Type 4 Class Exception Conditions

Exception	Real	Virtual 80x86	Protected and Compatibility	64-bit	Cause of Exception
Invalid Opcode, #UD	X	X			VEX prefix.
			X	X	VEX prefix: If XCR0[2:1] != '11b'. If CR4.OSXSAVE[bit 18]=0.
	X	X	X	X	Legacy SSE instruction: If CR0.EM[bit 2] = 1. If CR4.OSFXSR[bit 9] = 0.
	X	X	X	X	If preceded by a LOCK prefix (F0H).
			X	X	If any REX, F2, F3, or 66 prefixes precede a VEX prefix.
	X	X	X	X	If any corresponding CPUID feature flag is '0'.
Device Not Available, #NM	X	X	X	X	If CR0.TS[bit 3]=1.
Stack, SS(0)			X		For an illegal address in the SS segment.
				X	If a memory address referencing the SS segment is in a non-canonical form.
General Protection, #GP(0)	X	X	X	X	Legacy SSE: Memory operand is not 16-byte aligned.
			X		For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
				X	If the memory address is in a non-canonical form.
	X	X			If any part of the operand lies outside the effective address space from 0 to FFFFH.
Page Fault #PF(fault-code)		X	X	X	For a page fault.

2.4.5 Exceptions Type 5 (<16 Byte mem arg and no FP exceptions)

Table 2-21. Type 5 Class Exception Conditions

Exception	Real	Virtual 80x86	Protected and Compatibility	64-bit	Cause of Exception
Invalid Opcode, #UD	X	X			VEX prefix.
			X	X	VEX prefix: If XCR0[2:1] != '11b'. If CR4.OSXSAVE[bit 18]=0.
	X	X	X	X	Legacy SSE instruction: If CR0.EM[bit 2] = 1. If CR4.OSFXSR[bit 9] = 0.
	X	X	X	X	If preceded by a LOCK prefix (F0H).
			X	X	If any REX, F2, F3, or 66 prefixes precede a VEX prefix.
	X	X	X	X	If any corresponding CPUID feature flag is '0'.
Device Not Available, #NM	X	X	X	X	If CR0.TS[bit 3]=1.
Stack, SS(0)			X		For an illegal address in the SS segment.
				X	If a memory address referencing the SS segment is in a non-canonical form.
General Protection, #GP(0)			X		For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
				X	If the memory address is in a non-canonical form.
	X	X			If any part of the operand lies outside the effective address space from 0 to FFFFH.
Page Fault #PF(fault-code)		X	X	X	For a page fault.
Alignment Check #AC(0)		X	X	X	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

2.4.6 Exceptions Type 6 (VEX-Encoded Instructions Without Legacy SSE Analogues)

Note: At present, the AVX instructions in this category do not generate floating-point exceptions.

Table 2-22. Type 6 Class Exception Conditions

Exception	Real	Virtual 80x86	Protected and Compatibility	64-bit	Cause of Exception
Invalid Opcode, #UD	X	X			VEX prefix.
			X	X	If XCR0[2:1] != '11b'. If CR4.OSXSAVE[bit 18]=0.
			X	X	If preceded by a LOCK prefix (F0H).
			X	X	If any REX, F2, F3, or 66 prefixes precede a VEX prefix.
			X	X	If any corresponding CPUID feature flag is '0'.
Device Not Available, #NM			X	X	If CR0.TS[bit 3]=1.
Stack, SS(0)			X		For an illegal address in the SS segment.
				X	If a memory address referencing the SS segment is in a non-canonical form.
General Protection, #GP(0)			X		For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
				X	If the memory address is in a non-canonical form.
Page Fault #PF(fault-code)			X	X	For a page fault.
Alignment Check #AC(0)			X	X	For 4 or 8 byte memory references if alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

2.4.7 Exceptions Type 7 (No FP exceptions, no memory arg)

Table 2-23. Type 7 Class Exception Conditions

Exception	Real	Virtual 80x86	Protected and Compatibility	64-bit	Cause of Exception
Invalid Opcode, #UD	X	X			VEX prefix.
			X	X	VEX prefix: If XCR0[2:1] != '11b'. If CR4.OSXSAVE[bit 18]=0.
	X	X	X	X	Legacy SSE instruction: If CR0.EM[bit 2] = 1. If CR4.OSFXSR[bit 9] = 0.
	X	X	X	X	If preceded by a LOCK prefix (F0H).
			X	X	If any REX, F2, F3, or 66 prefixes precede a VEX prefix.
	X	X	X	X	If any corresponding CPUID feature flag is '0'.
Device Not Available, #NM			X	X	If CR0.TS[bit 3]=1.

2.4.8 Exceptions Type 8 (AVX and no memory argument)

Table 2-24. Type 8 Class Exception Conditions

Exception	Real	Virtual 80x86	Protected and Compatibility	64-bit	Cause of Exception
Invalid Opcode, #UD	X	X			Always in Real or Virtual 80x86 mode.
			X	X	If XCR0[2:1] != '11b'. If CR4.OSXSAVE[bit 18]=0. If CPUID.01H.ECX.AVX[bit 28]=0. If VEX.vvvv != 1111B.
	X	X	X	X	If proceeded by a LOCK prefix (F0H).
Device Not Available, #NM			X	X	If CR0.TS[bit 3]=1.

2.4.9 Exception Type 9 (Intel AVX)

Table 2-25. Type 9 Class Exception Conditions

Exception	Real	Virtual 80x86	Protected and Compatibility	64-bit	Cause of Exception
Invalid Opcode, #UD	X	X			Always in Real or Virtual 80x86 mode.
			X	X	If CR4.OSXSAVE[bit 18]=0. If CR0.EM[bit 2] = 1. If CPUID.01H.ECX.AVX[bit 28]=0. If VEX.L = 1.
			X	X	If preceded by a LOCK prefix (F0H).
			X	X	If any REX, F2, F3, or 66 prefixes precede a VEX prefix.
Device Not Available, #NM			X	X	If CR0.TS[bit 3]=1.
Stack, SS(0)			X		For an illegal address in the SS segment.
				X	If a memory address referencing the SS segment is in a non-canonical form.
General Protection, #GP(0)			X		For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
				X	If the memory address is in a non-canonical form.
Page Fault #PF(fault-code)			X	X	For a page fault.
Alignment Check #AC(0)			X	X	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

CHAPTER 3
INSTRUCTION SET REFERENCE, A-M

This chapter describes the instruction set for the Intel 64 and IA-32 architectures (A-M) in IA-32e, protected, Virtual-8086, and real modes of operation. The set includes general-purpose, x87 FPU, MMX, SSE/SSE2/SSE3/SSSE3/SSE4, AESNI/PCLMULQDQ, AVX, and system instructions. See also Chapter 4, "Instruction Set Reference, N-Z," in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 2B*.

For each instruction, each operand combination is described. A description of the instruction and its operand, an operational description, a description of the effect of the instructions on flags in the EFLAGS register, and a summary of exceptions that can be generated are also provided.

3.1 INTERPRETING THE INSTRUCTION REFERENCE PAGES

This section describes the format of information contained in the instruction reference pages in this chapter. It explains notational conventions and abbreviations used in these sections.

3.1.1 Instruction Format

The following is an example of the format used for each instruction description in this chapter. The heading below introduces the example. The table below provides an example summary table.

CMC—Complement Carry Flag [this is an example]

Opcode	Instruction	Op/En	64/32-bit Mode	CPUID Feature Flag	Description
F5	CMC	A	V/V	NA	Complement carry flag.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

3.1.1.1 Opcode Column in the Instruction Summary Table (Instructions without VEX prefix)

The “Opcode” column in the table above shows the object code produced for each form of the instruction. When possible, codes are given as hexadecimal bytes in the same order in which they appear in memory. Definitions of entries other than hexadecimal bytes are as follows:

- **REX.W** — Indicates the use of a REX prefix that affects operand size or instruction semantics. The ordering of the REX prefix and other optional/mandatory instruction prefixes are discussed Chapter 2. Note that REX prefixes that promote legacy instructions to 64-bit behavior are not listed explicitly in the opcode column.
- **/digit** — A digit between 0 and 7 indicates that the ModR/M byte of the instruction uses only the r/m (register or memory) operand. The reg field contains the digit that provides an extension to the instruction's opcode.
- **/r —** Indicates that the ModR/M byte of the instruction contains a register operand and an r/m operand.
- **cb, cw, cd, cp, co, ct —** A 1-byte (cb), 2-byte (cw), 4-byte (cd), 6-byte (cp), 8-byte (co) or 10-byte (ct) value following the opcode. This value is used to specify a code offset and possibly a new value for the code segment register.
- **ib, iw, id, io —** A 1-byte (ib), 2-byte (iw), 4-byte (id) or 8-byte (io) immediate operand to the instruction that follows the opcode, ModR/M bytes or scale-indexing bytes. The opcode determines if the operand is a signed value. All words, doublewords and quadwords are given with the low-order byte first.
- **+rb, +rw, +rd, +ro —** A register code, from 0 through 7, added to the hexadecimal byte given at the left of the plus sign to form a single opcode byte. See Table 3-1 for the codes. The +ro columns in the table are applicable only in 64-bit mode.
- +**i —** A number used in floating-point instructions when one of the operands is ST(i) from the FPU register stack. The number i (which can range from 0 to 7) is added to the hexadecimal byte given at the left of the plus sign to form a single opcode byte.

Table 3-1. Register Codes Associated With +rb, +rw, +rd, +ro

byte register			word register			dword register			quadword register (64-Bit Mode only)		
Register	**REX.B**	**Reg Field**	**Register**	**REX.B**	**Reg Field**	**Register**	**REX.B**	**Reg Field**	**Register**	**REX.B**	**Reg Field**
AL	None	0	AX	None	0	EAX	None	0	RAX	None	0
CL	None	1	CX	None	1	ECX	None	1	RCX	None	1
DL	None	2	DX	None	2	EDX	None	2	RDX	None	2
BL	None	3	BX	None	3	EBX	None	3	RBX	None	3
AH	Not encodable (N.E.)	4	SP	None	4	ESP	None	4	N/A	N/A	N/A
CH	N.E.	5	BP	None	5	EBP	None	5	N/A	N/A	N/A
DH	N.E.	6	SI	None	6	ESI	None	6	N/A	N/A	N/A
BH	N.E.	7	DI	None	7	EDI	None	7	N/A	N/A	N/A
SPL	Yes	4	SP	None	4	ESP	None	4	RSP	None	4
BPL	Yes	5	BP	None	5	EBP	None	5	RBP	None	5
SIL	Yes	6	SI	None	6	ESI	None	6	RSI	None	6
DIL	Yes	7	DI	None	7	EDI	None	7	RDI	None	7
Registers R8 - R15 (see below): Available in 64-Bit Mode Only											
R8L	Yes	0	R8W	Yes	0	R8D	Yes	0	R8	Yes	0
R9L	Yes	1	R9W	Yes	1	R9D	Yes	1	R9	Yes	1
R10L	Yes	2	R10W	Yes	2	R10D	Yes	2	R10	Yes	2
R11L	Yes	3	R11W	Yes	3	R11D	Yes	3	R11	Yes	3
R12L	Yes	4	R12W	Yes	4	R12D	Yes	4	R12	Yes	4
R13L	Yes	5	R13W	Yes	5	R13D	Yes	5	R13	Yes	5
R14L	Yes	6	R14W	Yes	6	R14D	Yes	6	R14	Yes	6
R15L	Yes	7	R15W	Yes	7	R15D	Yes	7	R15	Yes	7

3.1.1.2 Opcode Column in the Instruction Summary Table (Instructions with VEX prefix)

In the Instruction Summary Table, the Opcode column presents each instruction encoded using the VEX prefix in following form (including the modR/M byte if applicable, the immediate byte if applicable):

VEX.[NDS].[128,256].[66,F2,F3].0F/0F3A/0F38.[W0,W1] opcode [/r] [/ib,/is4]

- **VEX:** indicates the presence of the VEX prefix is required. The VEX prefix can be encoded using the three-byte form (the first byte is C4H), or using the two-byte form (the first byte is C5H). The two-byte form of VEX only applies to those instructions that do not require the following fields to be encoded: VEX.mmmmm, VEX.W, VEX.X, VEX.B. Refer to Section 2.3 for more detail on the VEX prefix.

  The encoding of various sub-fields of the VEX prefix is described using the following notations:

  — **NDS, NDD, DDS:** specifies that VEX.vvvv field is valid for the encoding of a register operand:

    - VEX.NDS: VEX.vvvv encodes the first source register in an instruction syntax where the content of source registers will be preserved.
    - VEX.NDD: VEX.vvvv encodes the destination register that cannot be encoded by ModR/M:reg field.
    - VEX.DDS: VEX.vvvv encodes the second source register in a three-operand instruction syntax where the content of first source register will be overwritten by the result.
    - If none of NDS, NDD, and DDS is present, VEX.vvvv must be 1111b (i.e. VEX.vvvv does not encode an operand). The VEX.vvvv field can be encoded using either the 2-byte or 3-byte form of the VEX prefix.

  — **128,256:** VEX.L field can be 0 (denoted by VEX.128 or VEX.LZ) or 1 (denoted by VEX.256). The VEX.L field can be encoded using either the 2-byte or 3-byte form of the VEX prefix. The presence of the notation VEX.256 or VEX.128 in the opcode column should be interpreted as follows:

    - If VEX.256 is present in the opcode column: The semantics of the instruction must be encoded with VEX.L = 1. An attempt to encode this instruction with VEX.L= 0 can result in one of two situations: (a) if VEX.128 version is defined, the processor will behave according to the defined VEX.128 behavior; (b) an #UD occurs if there is no VEX.128 version defined.
    - If VEX.128 is present in the opcode column but there is no VEX.256 version defined for the same opcode byte: Two situations apply: (a) For VEX-encoded, 128-bit SIMD integer instructions, software must encode the instruction with VEX.L = 0. The processor will treat the opcode byte encoded with VEX.L= 1 by causing an #UD exception; (b) For VEX-

encoded, 128-bit packed floating-point instructions, software must encode the instruction with VEX.L = 0. The processor will treat the opcode byte encoded with VEX.L= 1 by causing an #UD exception (e.g. VMOVLPS).

- If VEX.LIG is present in the opcode column: The VEX.L value is ignored. This generally applies to VEX-encoded scalar SIMD floating-point instructions. Scalar SIMD floating-point instruction can be distinguished from the mnemonic of the instruction. Generally, the last two letters of the instruction mnemonic would be either "SS", "SD", or "SI" for SIMD floating-point conversion instructions.
- If VEX.LZ is present in the opcode column: The VEX.L must be encoded to be 0B, an #UD occurs if VEX.L is not zero.

— **66,F2,F3:** The presence or absence of these values map to the VEX.pp field encodings. If absent, this corresponds to VEX.pp=00B. If present, the corresponding VEX.pp value affects the "opcode" byte in the same way as if a SIMD prefix (66H, F2H or F3H) does to the ensuing opcode byte. Thus a non-zero encoding of VEX.pp may be considered as an implied 66H/F2H/F3H prefix. The VEX.pp field may be encoded using either the 2-byte or 3-byte form of the VEX prefix.

— **0F,0F3A,0F38:** The presence maps to a valid encoding of the VEX.mmmmm field. Only three encoded values of VEX.mmmmm are defined as valid, corresponding to the escape byte sequence of 0FH, 0F3AH and 0F38H. The effect of a valid VEX.mmmmm encoding on the ensuing opcode byte is same as if the corresponding escape byte sequence on the ensuing opcode byte for non-VEX encoded instructions. Thus a valid encoding of VEX.mmmmm may be consider as an implies escape byte sequence of either 0FH, 0F3AH or 0F38H. The VEX.mmmmm field must be encoded using the 3-byte form of VEX prefix.

— **0F,0F3A,0F38 and 2-byte/3-byte VEX.** The presence of 0F3A and 0F38 in the opcode column implies that opcode can only be encoded by the three-byte form of VEX. The presence of 0F in the opcode column does not preclude the opcode to be encoded by the two-byte of VEX if the semantics of the opcode does not require any subfield of VEX not present in the two-byte form of the VEX prefix.

— **W0:** VEX.W=0.

— **W1:** VEX.W=1.

— The presence of W0/W1 in the opcode column applies to two situations: (a) it is treated as an extended opcode bit, (b) the instruction semantics support an operand size promotion to 64-bit of a general-purpose register operand or a 32-bit memory operand. The presence of W1 in the opcode column implies the opcode must be encoded using the 3-byte form of the VEX prefix. The presence of W0 in the opcode column does not preclude the opcode to be encoded using the C5H form of the VEX prefix, if the semantics of the opcode does not require other VEX subfields not present in the two-byte form of the VEX prefix. Please see Section 2.3 on the subfield definitions within VEX.

- **WIG:** can use C5H form (if not requiring VEX.mmmmm) or VEX.W value is ignored in the C4H form of VEX prefix.
- If WIG is present, the instruction may be encoded using either the two-byte form or the three-byte form of VEX. When encoding the instruction using the three-byte form of VEX, the value of VEX.W is ignored.

- **opcode:** Instruction opcode.
- **/is4:** An 8-bit immediate byte is present containing a source register specifier in imm[7:4] and instruction-specific payload in imm[3:0].
- In general, the encoding o f VEX.R, VEX.X, VEX.B field are not shown explicitly in the opcode column. The encoding scheme of VEX.R, VEX.X, VEX.B fields must follow the rules defined in Section 2.3.

3.1.1.3 Instruction Column in the Opcode Summary Table

The "Instruction" column gives the syntax of the instruction statement as it would appear in an ASM386 program. The following is a list of the symbols used to represent operands in the instruction statements:

- **rel8** — A relative address in the range from 128 bytes before the end of the instruction to 127 bytes after the end of the instruction.
- **rel16, rel32, rel64** — A relative address within the same code segment as the instruction assembled. The rel16 symbol applies to instructions with an operand-size attribute of 16 bits; the rel32 symbol applies to instructions with an operand-size attribute of 32 bits; the rel64 symbol applies to instructions with an operand-size attribute of 64 bits.
- **ptr16:16, ptr16:32 and ptr16:64** — A far pointer, typically to a code segment different from that of the instruction. The notation *16:16* indicates that the value of the pointer has two parts. The value to the left of the colon is a 16-bit selector or value destined for the code segment register. The value to the right corresponds to the offset within the destination segment. The ptr16:16 symbol is used when the instruction's operand-size attribute is 16 bits; the ptr16:32 symbol is used when the operand-size attribute is 32 bits; the ptr16:64 symbol is used when the operand-size attribute is 64 bits.
- **r8** — One of the byte general-purpose registers: AL, CL, DL, BL, AH, CH, DH, BH, BPL, SPL, DIL and SIL; or one of the byte registers (R8L - R15L) available when using REX.R and 64-bit mode.
- **r16** — One of the word general-purpose registers: AX, CX, DX, BX, SP, BP, SI, DI; or one of the word registers (R8-R15) available when using REX.R and 64-bit mode.
- **r32** — One of the doubleword general-purpose registers: EAX, ECX, EDX, EBX, ESP, EBP, ESI, EDI; or one of the doubleword registers (R8D - R15D) available when using REX.R in 64-bit mode.

- **r64** — One of the quadword general-purpose registers: RAX, RBX, RCX, RDX, RDI, RSI, RBP, RSP, R8–R15. These are available when using REX.R and 64-bit mode.
- **imm8** — An immediate byte value. The imm8 symbol is a signed number between –128 and +127 inclusive. For instructions in which imm8 is combined with a word or doubleword operand, the immediate value is sign-extended to form a word or doubleword. The upper byte of the word is filled with the topmost bit of the immediate value.
- **imm16** — An immediate word value used for instructions whose operand-size attribute is 16 bits. This is a number between –32,768 and +32,767 inclusive.
- **imm32** — An immediate doubleword value used for instructions whose operand-size attribute is 32 bits. It allows the use of a number between +2,147,483,647 and –2,147,483,648 inclusive.
- **imm64** — An immediate quadword value used for instructions whose operand-size attribute is 64 bits. The value allows the use of a number between +9,223,372,036,854,775,807 and –9,223,372,036,854,775,808 inclusive.
- **r/m8** — A byte operand that is either the contents of a byte general-purpose register (AL, CL, DL, BL, AH, CH, DH, BH, BPL, SPL, DIL and SIL) or a byte from memory. Byte registers R8L - R15L are available using REX.R in 64-bit mode.
- **r/m16** — A word general-purpose register or memory operand used for instructions whose operand-size attribute is 16 bits. The word general-purpose registers are: AX, CX, DX, BX, SP, BP, SI, DI. The contents of memory are found at the address provided by the effective address computation. Word registers R8W - R15W are available using REX.R in 64-bit mode.
- **r/m32** — A doubleword general-purpose register or memory operand used for instructions whose operand-size attribute is 32 bits. The doubleword general-purpose registers are: EAX, ECX, EDX, EBX, ESP, EBP, ESI, EDI. The contents of memory are found at the address provided by the effective address computation. Doubleword registers R8D - R15D are available when using REX.R in 64-bit mode.
- **r/m64** — A quadword general-purpose register or memory operand used for instructions whose operand-size attribute is 64 bits when using REX.W. Quadword general-purpose registers are: RAX, RBX, RCX, RDX, RDI, RSI, RBP, RSP, R8–R15; these are available only in 64-bit mode. The contents of memory are found at the address provided by the effective address computation.
- **m** — A 16-, 32- or 64-bit operand in memory.
- **m8** — A byte operand in memory, usually expressed as a variable or array name, but pointed to by the DS:(E)SI or ES:(E)DI registers. In 64-bit mode, it is pointed to by the RSI or RDI registers.
- **m16** — A word operand in memory, usually expressed as a variable or array name, but pointed to by the DS:(E)SI or ES:(E)DI registers. This nomenclature is used only with the string instructions.

- **m32** — A doubleword operand in memory, usually expressed as a variable or array name, but pointed to by the DS:(E)SI or ES:(E)DI registers. This nomenclature is used only with the string instructions.
- **m64** — A memory quadword operand in memory.
- **m128** — A memory double quadword operand in memory.
- **m16:16, m16:32 & m16:64** — A memory operand containing a far pointer composed of two numbers. The number to the left of the colon corresponds to the pointer's segment selector. The number to the right corresponds to its offset.
- **m16&32, m16&16, m32&32, m16&64** — A memory operand consisting of data item pairs whose sizes are indicated on the left and the right side of the ampersand. All memory addressing modes are allowed. The m16&16 and m32&32 operands are used by the BOUND instruction to provide an operand containing an upper and lower bounds for array indices. The m16&32 operand is used by LIDT and LGDT to provide a word with which to load the limit field, and a doubleword with which to load the base field of the corresponding GDTR and IDTR registers. The m16&64 operand is used by LIDT and LGDT in 64-bit mode to provide a word with which to load the limit field, and a quadword with which to load the base field of the corresponding GDTR and IDTR registers.
- **moffs8, moffs16, moffs32, moffs64** — A simple memory variable (memory offset) of type byte, word, or doubleword used by some variants of the MOV instruction. The actual address is given by a simple offset relative to the segment base. No ModR/M byte is used in the instruction. The number shown with moffs indicates its size, which is determined by the address-size attribute of the instruction.
- **Sreg** — A segment register. The segment register bit assignments are ES = 0, CS = 1, SS = 2, DS = 3, FS = 4, and GS = 5.
- **m32fp, m64fp, m80fp** — A single-precision, double-precision, and double extended-precision (respectively) floating-point operand in memory. These symbols designate floating-point values that are used as operands for x87 FPU floating-point instructions.
- **m16int, m32int, m64int** — A word, doubleword, and quadword integer (respectively) operand in memory. These symbols designate integers that are used as operands for x87 FPU integer instructions.
- **ST or ST(0)** — The top element of the FPU register stack.
- **ST(i)** — The i^{th} element from the top of the FPU register stack ($i \leftarrow 0$ through 7).
- **mm** — An MMX register. The 64-bit MMX registers are: MM0 through MM7.
- **mm/m32** — The low order 32 bits of an MMX register or a 32-bit memory operand. The 64-bit MMX registers are: MM0 through MM7. The contents of memory are found at the address provided by the effective address computation.
- **mm/m64** — An MMX register or a 64-bit memory operand. The 64-bit MMX registers are: MM0 through MM7. The contents of memory are found at the address provided by the effective address computation.

- **xmm** — An XMM register. The 128-bit XMM registers are: XMM0 through XMM7; XMM8 through XMM15 are available using REX.R in 64-bit mode.
- **xmm/m32**— An XMM register or a 32-bit memory operand. The 128-bit XMM registers are XMM0 through XMM7; XMM8 through XMM15 are available using REX.R in 64-bit mode. The contents of memory are found at the address provided by the effective address computation.
- **xmm/m64** — An XMM register or a 64-bit memory operand. The 128-bit SIMD floating-point registers are XMM0 through XMM7; XMM8 through XMM15 are available using REX.R in 64-bit mode. The contents of memory are found at the address provided by the effective address computation.
- **xmm/m128** — An XMM register or a 128-bit memory operand. The 128-bit XMM registers are XMM0 through XMM7; XMM8 through XMM15 are available using REX.R in 64-bit mode. The contents of memory are found at the address provided by the effective address computation.
- **<XMM0>**— indicates implied use of the XMM0 register.

  When there is ambiguity, xmm1 indicates the first source operand using an XMM register and xmm2 the second source operand using an XMM register.

  Some instructions use the XMM0 register as the third source operand, indicated by <XMM0>. The use of the third XMM register operand is implicit in the instruction encoding and does not affect the ModR/M encoding.
- **ymm** — a YMM register. The 256-bit YMM registers are: YMM0 through YMM7**;** YMM8 through YMM15 are available in 64-bit mode.
- **m256** — A 32-byte operand in memory. This nomenclature is used only with AVX instructions.
- **ymm/m256** — a YMM register or 256-bit memory operand.
- **<YMM0>**— indicates use of the YMM0 register as an implicit argument.
- **SRC1** — Denotes the first source operand in the instruction syntax of an instruction encoded with the VEX prefix and having two or more source operands.
- **SRC2** — Denotes the second source operand in the instruction syntax of an instruction encoded with the VEX prefix and having two or more source operands.
- **SRC3** — Denotes the third source operand in the instruction syntax of an instruction encoded with the VEX prefix and having three source operands.
- **SRC** — The source in a AVX single-source instruction or the source in a Legacy SSE instruction.
- **DST** — the destination in a AVX instruction. In Legacy SSE instructions can be either the destination, first source, or both. This field is encoded by reg_field.

3.1.1.4 Operand Encoding Column in the Instruction Summary Table

The “operand encoding” column is abbreviated as Op/En in the Instruction Summary table heading. Instruction operand encoding information is provided for each assembly instruction syntax using a letter to cross reference to a row entry in the operand encoding definition table that follows the instruction summary table. The operand encoding table in each instruction reference page lists each instruction operand (according to each instruction syntax and operand ordering shown in the instruction column) relative to the ModRM byte, VEX.vvvv field or additional operand encoding placement.

NOTES

- The letters in the Op/En column of an instruction apply ONLY to the encoding definition table immediately following the instruction summary table.
- In the encoding definition table, the letter 'r' within a pair of parenthesis denotes the content of the operand will be read by the processor. The letter 'w' within a pair of parenthesis denotes the content of the operand will be updated by the processor.

3.1.1.5 64/32-bit Mode Column in the Instruction Summary Table

The "64/32-bit Mode" column indicates whether the opcode sequence is supported in (a) 64-bit mode or (b) the Compatibility mode and other IA-32 modes that apply in conjunction with the CPUID feature flag associated specific instruction extensions.

The 64-bit mode support is to the left of the 'slash' and has the following notation:

- **V** — Supported.
- **I** — Not supported.
- **N.E.** — Indicates an instruction syntax is not encodable in 64-bit mode (it may represent part of a sequence of valid instructions in other modes).
- **N.P.** — Indicates the REX prefix does not affect the legacy instruction in 64-bit mode.
- **N.I.** — Indicates the opcode is treated as a new instruction in 64-bit mode.
- **N.S.** — Indicates an instruction syntax that requires an address override prefix in 64-bit mode and is not supported. Using an address override prefix in 64-bit mode may result in model-specific execution behavior.

The Compatibility/Legacy Mode support is to the right of the 'slash' and has the following notation:

• **V —** Supported.

• **I —** Not supported.

• **N.E. —** Indicates an Intel 64 instruction mnemonics/syntax that is not encodable; the opcode sequence is not applicable as an individual instruction in compatibility mode or IA-32 mode. The opcode may represent a valid sequence of legacy IA-32 instructions.

3.1.1.6 CPUID Support Column in the Instruction Summary Table

The fourth column holds abbreviated CPUID feature flags (e.g. appropriate bit in CPUID.1.ECX, CPUID.1.EDX for SSE/SSE2/SSE3/SSSE3/SSE4.1/SSE4.2/AES-NI/PCLMULQDQ/AVX support) that indicate processor support for the instruction. If the corresponding flag is '0', the instruction will #UD.

3.1.1.7 Description Column in the Instruction Summary Table

The "Description" column briefly explains forms of the instruction.

3.1.1.8 Description Section

Each instruction is then described by number of information sections. The "Description" section describes the purpose of the instructions and required operands in more detail.

Summary of terms that may be used in the description section:

- **Legacy SSE:** Refers to SSE, SSE2, SSE3, SSSE3, SSE4, AESNI, PCLMULQDQ and any future instruction sets referencing XMM registers and encoded without a VEX prefix.
- **VEX.vvvv**. The VEX bitfield specifying a source or destination register (in 1's complement form).
- **rm_field**: shorthand for the ModR/M *r/m* field and any REX.B
- **reg_field:** shorthand for the ModR/M *reg* field and any REX.R

3.1.1.9 Operation Section

The "Operation" section contains an algorithm description (frequently written in pseudo-code) for the instruction. Algorithms are composed of the following elements:

- Comments are enclosed within the symbol pairs "(*" and "*)".
- Compound statements are enclosed in keywords, such as: IF, THEN, ELSE and FI for an if statement; DO and OD for a do statement; or CASE... OF for a case statement.
- A register name implies the contents of the register. A register name enclosed in brackets implies the contents of the location whose address is contained in that register. For example, ES:[DI] indicates the contents of the location whose ES segment relative address is in register DI. [SI] indicates the contents of the address contained in register SI relative to the SI register's default segment (DS) or the overridden segment.
- Parentheses around the "E" in a general-purpose register name, such as (E)SI, indicates that the offset is read from the SI register if the address-size attribute is 16, from the ESI register if the address-size attribute is 32. Parentheses

around the "R" in a general-purpose register name, (R)SI, in the presence of a 64-bit register definition such as (R)SI, indicates that the offset is read from the 64-bit RSI register if the address-size attribute is 64.

- Brackets are used for memory operands where they mean that the contents of the memory location is a segment-relative offset. For example, [SRC] indicates that the content of the source operand is a segment-relative offset.
- A ← B indicates that the value of B is assigned to A.
- The symbols =, ≠, >, <, ≥, and ≤ are relational operators used to compare two values: meaning equal, not equal, greater or equal, less or equal, respectively. A relational expression such as A ← B is TRUE if the value of A is equal to B; otherwise it is FALSE.
- The expression "« COUNT" and "» COUNT" indicates that the destination operand should be shifted left or right by the number of bits indicated by the count operand.

The following identifiers are used in the algorithmic descriptions:

- **OperandSize and AddressSize** — The OperandSize identifier represents the operand-size attribute of the instruction, which is 16, 32 or 64-bits. The AddressSize identifier represents the address-size attribute, which is 16, 32 or 64-bits. For example, the following pseudo-code indicates that the operand-size attribute depends on the form of the MOV instruction used.

```
IF Instruction ← MOVW
    THEN OperandSize = 16;
ELSE
    IF Instruction ← MOVD
        THEN OperandSize = 32;
    ELSE
        IF Instruction ← MOVQ
            THEN OperandSize = 64;
        FI;
    FI;
FI;
```

  See "Operand-Size and Address-Size Attributes" in Chapter 3 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for guidelines on how these attributes are determined.

- **StackAddrSize** — Represents the stack address-size attribute associated with the instruction, which has a value of 16, 32 or 64-bits. See "Address-Size Attribute for Stack" in Chapter 6, "Procedure Calls, Interrupts, and Exceptions," of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*.
- **SRC** — Represents the source operand.
- **DEST** — Represents the destination operand.
- **VLMAX** — The maximum vector register width pertaining to the instruction. This is not the vector-length encoding in the instruction's prefix but is instead

determined by the current value of XCR0. For existing processors, VLMAX is 256 whenever XCR0.YMM[bit 2] is 1. Future processors may defined new bits in XCR0 whose setting may imply other values for VLMAX.

VLMAX Definition

XCR0 Component	VLMAX
XCR0.YMM	256

The following functions are used in the algorithmic descriptions:

- **ZeroExtend(value)** — Returns a value zero-extended to the operand-size attribute of the instruction. For example, if the operand-size attribute is 32, zero extending a byte value of –10 converts the byte from F6H to a doubleword value of 000000F6H. If the value passed to the ZeroExtend function and the operand-size attribute are the same size, ZeroExtend returns the value unaltered.
- **SignExtend(value)** — Returns a value sign-extended to the operand-size attribute of the instruction. For example, if the operand-size attribute is 32, sign extending a byte containing the value –10 converts the byte from F6H to a doubleword value of FFFFFFF6H. If the value passed to the SignExtend function and the operand-size attribute are the same size, SignExtend returns the value unaltered.
- **SaturateSignedWordToSignedByte** — Converts a signed 16-bit value to a signed 8-bit value. If the signed 16-bit value is less than –128, it is represented by the saturated value -128 (80H); if it is greater than 127, it is represented by the saturated value 127 (7FH).
- **SaturateSignedDwordToSignedWord** — Converts a signed 32-bit value to a signed 16-bit value. If the signed 32-bit value is less than –32768, it is represented by the saturated value –32768 (8000H); if it is greater than 32767, it is represented by the saturated value 32767 (7FFFH).
- **SaturateSignedWordToUnsignedByte** — Converts a signed 16-bit value to an unsigned 8-bit value. If the signed 16-bit value is less than zero, it is represented by the saturated value zero (00H); if it is greater than 255, it is represented by the saturated value 255 (FFH).
- **SaturateToSignedByte** — Represents the result of an operation as a signed 8-bit value. If the result is less than –128, it is represented by the saturated value –128 (80H); if it is greater than 127, it is represented by the saturated value 127 (7FH).
- **SaturateToSignedWord** — Represents the result of an operation as a signed 16-bit value. If the result is less than –32768, it is represented by the saturated value –32768 (8000H); if it is greater than 32767, it is represented by the saturated value 32767 (7FFFH).
- **SaturateToUnsignedByte** — Represents the result of an operation as a signed 8-bit value. If the result is less than zero it is represented by the saturated value

zero (00H); if it is greater than 255, it is represented by the saturated value 255 (FFH).

- **SaturateToUnsignedWord** — Represents the result of an operation as a signed 16-bit value. If the result is less than zero it is represented by the saturated value zero (00H); if it is greater than 65535, it is represented by the saturated value 65535 (FFFFH).
- **LowOrderWord(DEST * SRC)** — Multiplies a word operand by a word operand and stores the least significant word of the doubleword result in the destination operand.
- **HighOrderWord(DEST * SRC)** — Multiplies a word operand by a word operand and stores the most significant word of the doubleword result in the destination operand.
- **Push(value)** — Pushes a value onto the stack. The number of bytes pushed is determined by the operand-size attribute of the instruction. See the "Operation" subsection of the "PUSH—Push Word, Doubleword or Quadword Onto the Stack" section in Chapter 4 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 2B*.
- **Pop()** removes the value from the top of the stack and returns it. The statement EAX ← Pop(); assigns to EAX the 32-bit value from the top of the stack. Pop will return either a word, a doubleword or a quadword depending on the operand-size attribute. See the "Operation" subsection in the "POP—Pop a Value from the Stack" section of Chapter 4 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 2B*.
- **PopRegisterStack** — Marks the FPU ST(0) register as empty and increments the FPU register stack pointer (TOP) by 1.
- **Switch-Tasks** — Performs a task switch.
- **Bit(BitBase, BitOffset)** — Returns the value of a bit within a bit string. The bit string is a sequence of bits in memory or a register. Bits are numbered from low-order to high-order within registers and within memory bytes. If the BitBase is a register, the BitOffset can be in the range 0 to [15, 31, 63] depending on the mode and register size. See Figure 3-1: the function Bit[RAX, 21] is illustrated.

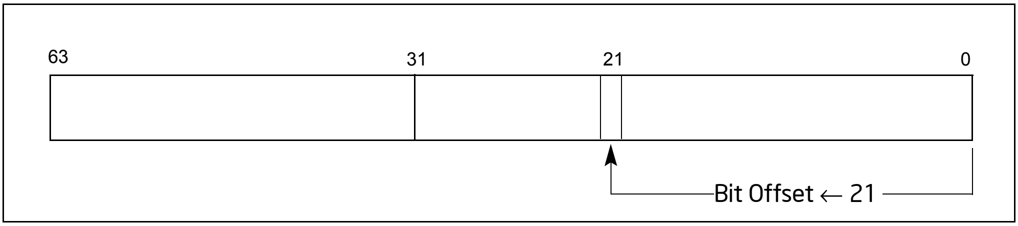


Figure 3-1. Bit Offset for BIT[RAX, 21]

If BitBase is a memory address, the BitOffset can range has different ranges depending on the operand size (see Table 3-2).

Table 3-2. Range of Bit Positions Specified by Bit Offset Operands

Operand Size	Immediate BitOffset	Register BitOffset
16	0 to 15	-2^{15} to $2^{15}-1$
32	0 to 31	-2^{31} to $2^{31}-1$
64	0 to 63	-2^{63} to $2^{63}-1$

The addressed bit is numbered (Offset MOD 8) within the byte at address (BitBase + (BitOffset DIV 8)) where DIV is signed division with rounding towards negative infinity and MOD returns a positive number (see Figure 3-2).

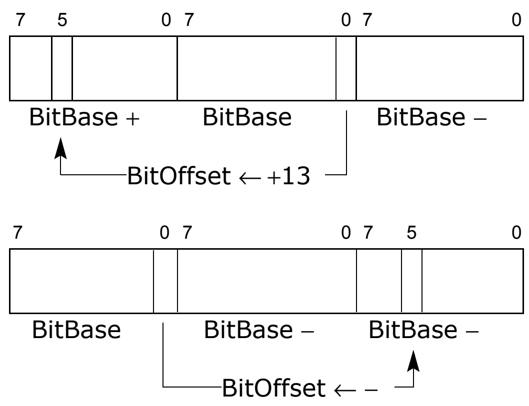


Figure 3-2. Memory Bit Indexing

3.1.1.10 Intel® C/C++ Compiler Intrinsics Equivalents Section

The Intel C/C++ compiler intrinsics equivalents are special C/C++ coding extensions that allow using the syntax of C function calls and C variables instead of hardware registers. Using these intrinsics frees programmers from having to manage registers and assembly programming. Further, the compiler optimizes the instruction scheduling so that executable run faster.

The following sections discuss the intrinsics API and the MMX technology and SIMD floating-point intrinsics. Each intrinsic equivalent is listed with the instruction description. There may be additional intrinsics that do not have an instruction equivalent. It is strongly recommended that the reader reference the compiler documentation for the complete list of supported intrinsics.

See Appendix C, “Intel® C/C++ Compiler Intrinsics and Functional Equivalents,” in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 2B*, for more information on using intrinsics.

Intrinsics API

The benefit of coding with MMX technology intrinsics and the SSE/SSE2/SSE3 intrinsics is that you can use the syntax of C function calls and C variables instead of hardware registers. This frees you from managing registers and programming assembly. Further, the compiler optimizes the instruction scheduling so that your executable runs faster. For each computational and data manipulation instruction in the new instruction set, there is a corresponding C intrinsic that implements it directly. The intrinsics allow you to specify the underlying implementation (instruction selection) of an algorithm yet leave instruction scheduling and register allocation to the compiler.

MMX™ Technology Intrinsics

The MMX technology intrinsics are based on a __m64 data type that represents the specific contents of an MMX technology register. You can specify values in bytes, short integers, 32-bit values, or a 64-bit object. The __m64 data type, however, is not a basic ANSI C data type, and therefore you must observe the following usage restrictions:

- Use __m64 data only on the left-hand side of an assignment, as a return value, or as a parameter. You cannot use it with other arithmetic expressions (“+”, “>>”, and so on).
- Use __m64 objects in aggregates, such as unions to access the byte elements and structures; the address of an __m64 object may be taken.
- Use __m64 data only with the MMX technology intrinsics described in this manual and Intel® C/C++ compiler documentation.
- See:
  - http://www.intel.com/support/performancetools/
  - Appendix C, “Intel® C/C++ Compiler Intrinsics and Functional Equivalents,” in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 2B*, for more information on using intrinsics.
  - SSE/SSE2/SSE3 Intrinsics
  - SSE/SSE2/SSE3 intrinsics all make use of the XMM registers of the Pentium III, Pentium 4, and Intel Xeon processors. There are three data types supported by these intrinsics: __m128, __m128d, and __m128i.

- The __m128 data type is used to represent the contents of an XMM register used by an SSE intrinsic. This is either four packed single-precision floating-point values or a scalar single-precision floating-point value.
- The __m128d data type holds two packed double-precision floating-point values or a scalar double-precision floating-point value.
- The __m128i data type can hold sixteen byte, eight word, or four doubleword, or two quadword integer values.

The compiler aligns __m128, __m128d, and __m128i local and global data to 16-byte boundaries on the stack. To align integer, float, or double arrays, use the declspec statement as described in Intel C/C++ compiler documentation. See http://www.intel.com/support/performancetools/.

The __m128, __m128d, and __m128i data types are not basic ANSI C data types and therefore some restrictions are placed on its usage:

- Use __m128, __m128d, and __m128i only on the left-hand side of an assignment, as a return value, or as a parameter. Do not use it in other arithmetic expressions such as "+" and ">>."
- Do not initialize __m128, __m128d, and __m128i with literals; there is no way to express 128-bit constants.
- Use __m128, __m128d, and __m128i objects in aggregates, such as unions (for example, to access the float elements) and structures. The address of these objects may be taken.
- Use __m128, __m128d, and __m128i data only with the intrinsics described in this user's guide. See Appendix C, "Intel® C/C++ Compiler Intrinsics and Functional Equivalents," in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 2B*, for more information on using intrinsics.

The compiler aligns __m128, __m128d, and __m128i local data to 16-byte boundaries on the stack. Global __m128 data is also aligned on 16-byte boundaries. (To align float arrays, you can use the alignment declspec described in the following section.) Because the new instruction set treats the SIMD floating-point registers in the same way whether you are using packed or scalar data, there is no __m32 data type to represent scalar data as you might expect. For scalar operations, you should use the __m128 objects and the "scalar" forms of the intrinsics; the compiler and the processor implement these operations with 32-bit memory references.

The suffixes ps and ss are used to denote "packed single" and "scalar single" precision operations. The packed floats are represented in right-to-left order, with the lowest word (right-most) being used for scalar operations: [z, y, x, w]. To explain how memory storage reflects this, consider the following example.

The operation:

```
float a[4] ← { 1.0, 2.0, 3.0, 4.0 };
__m128 t ← _mm_load_ps(a);
```

Produces the same result as follows:

```
__m128 t ← _mm_set_ps(4.0, 3.0, 2.0, 1.0);
```

In other words:

```
t ← [ 4.0, 3.0, 2.0, 1.0 ]
```

Where the "scalar" element is 1.0.

Some intrinsics are "composites" because they require more than one instruction to implement them. You should be familiar with the hardware features provided by the SSE, SSE2, SSE3, and MMX technology when writing programs with the intrinsics.

Keep the following important issues in mind:

- Certain intrinsics, such as _mm_loadr_ps and _mm_cmpgt_ss, are not directly supported by the instruction set. While these intrinsics are convenient programming aids, be mindful of their implementation cost.
- Data loaded or stored as __m128 objects must generally be 16-byte-aligned.
- Some intrinsics require that their argument be immediates, that is, constant integers (literals), due to the nature of the instruction.
- The result of arithmetic operations acting on two NaN (Not a Number) arguments is undefined. Therefore, floating-point operations using NaN arguments may not match the expected behavior of the corresponding assembly instructions.

For a more detailed description of each intrinsic and additional information related to its usage, refer to Intel C/C++ compiler documentation. See:

— http://www.intel.com/support/performancetools/

— Appendix C, "Intel® C/C++ Compiler Intrinsics and Functional Equivalents," in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 2B*, for more information on using intrinsics.

3.1.1.11 Flags Affected Section

The "Flags Affected" section lists the flags in the EFLAGS register that are affected by the instruction. When a flag is cleared, it is equal to 0; when it is set, it is equal to 1. The arithmetic and logical instructions usually assign values to the status flags in a uniform manner (see Appendix A, "EFLAGS Cross-Reference," in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*). Non-conventional assignments are described in the "Operation" section. The values of flags listed as **undefined** may be changed by the instruction in an indeterminate manner. Flags that are not listed are unchanged by the instruction.

3.1.1.12 FPU Flags Affected Section

The floating-point instructions have an "FPU Flags Affected" section that describes how each instruction can affect the four condition code flags of the FPU status word.

3.1.1.13 Protected Mode Exceptions Section

The “Protected Mode Exceptions” section lists the exceptions that can occur when the instruction is executed in protected mode and the reasons for the exceptions. Each exception is given a mnemonic that consists of a pound sign (#) followed by two letters and an optional error code in parentheses. For example, #GP(0) denotes a general protection exception with an error code of 0. Table 3-3 associates each two-letter mnemonic with the corresponding interrupt vector number and exception name. See Chapter 6, “Interrupt and Exception Handling,” in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 3A*, for a detailed description of the exceptions.

Application programmers should consult the documentation provided with their operating systems to determine the actions taken when exceptions occur.

Table 3-3. Intel 64 and IA-32 General Exceptions

Vector No.	Name	Source	Protected Mode[1]	Real Address Mode	Virtual 8086 Mode
0	#DE—Divide Error	DIV and IDIV instructions.	Yes	Yes	Yes
1	#DB—Debug	Any code or data reference.	Yes	Yes	Yes
3	#BP—Breakpoint	INT 3 instruction.	Yes	Yes	Yes
4	#OF—Overflow	INTO instruction.	Yes	Yes	Yes
5	#BR—BOUND Range Exceeded	BOUND instruction.	Yes	Yes	Yes
6	#UD—Invalid Opcode (Undefined Opcode)	UD2 instruction or reserved opcode.	Yes	Yes	Yes
7	#NM—Device Not Available (No Math Coprocessor)	Floating-point or WAIT/FWAIT instruction.	Yes	Yes	Yes
8	#DF—Double Fault	Any instruction that can generate an exception, an NMI, or an INTR.	Yes	Yes	Yes
10	#TS—Invalid TSS	Task switch or TSS access.	Yes	Reserved	Yes
11	#NP—Segment Not Present	Loading segment registers or accessing system segments.	Yes	Reserved	Yes
12	#SS—Stack Segment Fault	Stack operations and SS register loads.	Yes	Yes	Yes
13	#GP—General Protection[2]	Any memory reference and other protection checks.	Yes	Yes	Yes

Table 3-3. Intel 64 and IA-32 General Exceptions (Contd.)

Vector No.	Name	Source	Protected Mode[1]	Real Address Mode	Virtual 8086 Mode
14	#PF—Page Fault	Any memory reference.	Yes	Reserved	Yes
16	#MF—Floating-Point Error (Math Fault)	Floating-point or WAIT/FWAIT instruction.	Yes	Yes	Yes
17	#AC—Alignment Check	Any data reference in memory.	Yes	Reserved	Yes
18	#MC—Machine Check	Model dependent machine check errors.	Yes	Yes	Yes
19	#XM—SIMD Floating-Point Numeric Error	SSE/SSE2/SSE3 floating-point instructions.	Yes	Yes	Yes

NOTES:

1. Apply to protected mode, compatibility mode, and 64-bit mode.
2. In the real-address mode, vector 13 is the segment overrun exception.

3.1.1.14 Real-Address Mode Exceptions Section

The “Real-Address Mode Exceptions” section lists the exceptions that can occur when the instruction is executed in real-address mode (see Table 3-3).

3.1.1.15 Virtual-8086 Mode Exceptions Section

The “Virtual-8086 Mode Exceptions” section lists the exceptions that can occur when the instruction is executed in virtual-8086 mode (see Table 3-3).

3.1.1.16 Floating-Point Exceptions Section

The “Floating-Point Exceptions” section lists exceptions that can occur when an x87 FPU floating-point instruction is executed. All of these exception conditions result in a floating-point error exception (#MF, vector number 16) being generated. Table 3-4 associates a one- or two-letter mnemonic with the corresponding exception name. See “Floating-Point Exception Conditions” in Chapter 8 of the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*, for a detailed description of these exceptions.

Table 3-4. x87 FPU Floating-Point Exceptions

Mnemonic	Name	Source

Table 3-4. x87 FPU Floating-Point Exceptions

#IS #IA	Floating-point invalid operation: - Stack overflow or underflow - Invalid arithmetic operation	 - x87 FPU stack overflow or underflow - Invalid FPU arithmetic operation
#Z	Floating-point divide-by-zero	Divide-by-zero
#D	Floating-point denormal operand	Source operand that is a denormal number
#O	Floating-point numeric overflow	Overflow in result
#U	Floating-point numeric underflow	Underflow in result
#P	Floating-point inexact result (precision)	Inexact result (precision)

3.1.1.17 SIMD Floating-Point Exceptions Section

The "SIMD Floating-Point Exceptions" section lists exceptions that can occur when an SSE/SSE2/SSE3 floating-point instruction is executed. All of these exception conditions result in a SIMD floating-point error exception (#XM, vector number 19) being generated. Table 3-5 associates a one-letter mnemonic with the corresponding exception name. For a detailed description of these exceptions, refer to "SSE and SSE2 Exceptions", in Chapter 11 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*.

Table 3-5. SIMD Floating-Point Exceptions

Mnemonic	Name	Source
#I	Floating-point invalid operation	Invalid arithmetic operation or source operand
#Z	Floating-point divide-by-zero	Divide-by-zero
#D	Floating-point denormal operand	Source operand that is a denormal number
#O	Floating-point numeric overflow	Overflow in result
#U	Floating-point numeric underflow	Underflow in result
#P	Floating-point inexact result	Inexact result (precision)

3.1.1.18 Compatibility Mode Exceptions Section

This section lists exceptions that occur within compatibility mode.

3.1.1.19 64-Bit Mode Exceptions Section

This section lists exceptions that occur within 64-bit mode.

3.2 INSTRUCTIONS (A-M)

The remainder of this chapter provides descriptions of Intel 64 and IA-32 instructions (A-M). See also: Chapter 4, “Instruction Set Reference, N-Z,” in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 2B*.

AAA—ASCII Adjust After Addition

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
37	AAA	A	Invalid	Valid	ASCII adjust AL after addition.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Adjusts the sum of two unpacked BCD values to create an unpacked BCD result. The AL register is the implied source and destination operand for this instruction. The AAA instruction is only useful when it follows an ADD instruction that adds (binary addition) two unpacked BCD values and stores a byte result in the AL register. The AAA instruction then adjusts the contents of the AL register to contain the correct 1-digit unpacked BCD result.

If the addition produces a decimal carry, the AH register increments by 1, and the CF and AF flags are set. If there was no decimal carry, the CF and AF flags are cleared and the AH register is unchanged. In either case, bits 4 through 7 of the AL register are set to 0.

This instruction executes as described in compatibility mode and legacy mode. It is not valid in 64-bit mode.

Operation

```
IF 64-Bit Mode
    THEN
        #UD;
    ELSE
        IF ((AL AND 0FH) > 9) or (AF = 1)
            THEN
                AL ← AL + 6;
                AH ← AH + 1;
                AF ← 1;
                CF ← 1;
                AL ← AL AND 0FH;
            ELSE
                AF ← 0;
                CF ← 0;
                AL ← AL AND 0FH;
        FI;
```

FI;

Flags Affected

The AF and CF flags are set to 1 if the adjustment results in a decimal carry; otherwise they are set to 0. The OF, SF, ZF, and PF flags are undefined.

Protected Mode Exceptions

#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as protected mode.

Compatibility Mode Exceptions

Same exceptions as protected mode.

64-Bit Mode Exceptions

#UD	If in 64-bit mode.

AAD—ASCII Adjust AX Before Division

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
D5 0A	AAD	A	Invalid	Valid	ASCII adjust AX before division.
D5 *ib*	(No mnemonic)	A	Invalid	Valid	Adjust AX before division to number base *imm8*.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Adjusts two unpacked BCD digits (the least-significant digit in the AL register and the most-significant digit in the AH register) so that a division operation performed on the result will yield a correct unpacked BCD value. The AAD instruction is only useful when it precedes a DIV instruction that divides (binary division) the adjusted value in the AX register by an unpacked BCD value.

The AAD instruction sets the value in the AL register to (AL + (10 * AH)), and then clears the AH register to 00H. The value in the AX register is then equal to the binary equivalent of the original unpacked two-digit (base 10) number in registers AH and AL.

The generalized version of this instruction allows adjustment of two unpacked digits of any number base (see the "Operation" section below), by setting the *imm8* byte to the selected number base (for example, 08H for octal, 0AH for decimal, or 0CH for base 12 numbers). The AAD mnemonic is interpreted by all assemblers to mean adjust ASCII (base 10) values. To adjust values in another number base, the instruction must be hand coded in machine code (D5 *imm8*).

This instruction executes as described in compatibility mode and legacy mode. It is not valid in 64-bit mode.

Operation

```
IF 64-Bit Mode
    THEN
        #UD;
    ELSE
        tempAL ← AL;
        tempAH ← AH;
        AL ← (tempAL + (tempAH ∗ imm8)) AND FFH;
        (* imm8 is set to 0AH for the AAD mnemonic.*)
```

```
        AH ← 0;
FI;
```

The immediate value (*imm8*) is taken from the second byte of the instruction.

Flags Affected

The SF, ZF, and PF flags are set according to the resulting binary value in the AL register; the OF, AF, and CF flags are undefined.

Protected Mode Exceptions

#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as protected mode.

Compatibility Mode Exceptions

Same exceptions as protected mode.

64-Bit Mode Exceptions

#UD	If in 64-bit mode.

AAM—ASCII Adjust AX After Multiply

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
D4 0A	AAM	A	Invalid	Valid	ASCII adjust AX after multiply.
D4 *ib*	(No mnemonic)	A	Invalid	Valid	Adjust AX after multiply to number base *imm8.*

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Adjusts the result of the multiplication of two unpacked BCD values to create a pair of unpacked (base 10) BCD values. The AX register is the implied source and destination operand for this instruction. The AAM instruction is only useful when it follows an MUL instruction that multiplies (binary multiplication) two unpacked BCD values and stores a word result in the AX register. The AAM instruction then adjusts the contents of the AX register to contain the correct 2-digit unpacked (base 10) BCD result.

The generalized version of this instruction allows adjustment of the contents of the AX to create two unpacked digits of any number base (see the “Operation” section below). Here, the *imm8* byte is set to the selected number base (for example, 08H for octal, 0AH for decimal, or 0CH for base 12 numbers). The AAM mnemonic is interpreted by all assemblers to mean adjust to ASCII (base 10) values. To adjust to values in another number base, the instruction must be hand coded in machine code (D4 *imm8*).

This instruction executes as described in compatibility mode and legacy mode. It is not valid in 64-bit mode.

Operation

```
IF 64-Bit Mode
    THEN
        #UD;
    ELSE
        tempAL ← AL;
        AH ← tempAL / imm8; (* imm8 is set to 0AH for the AAM mnemonic *)
        AL ← tempAL MOD imm8;
FI;
```

The immediate value (*imm8*) is taken from the second byte of the instruction.

Flags Affected

The SF, ZF, and PF flags are set according to the resulting binary value in the AL register. The OF, AF, and CF flags are undefined.

Protected Mode Exceptions

#DE	If an immediate value of 0 is used.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as protected mode.

Compatibility Mode Exceptions

Same exceptions as protected mode.

64-Bit Mode Exceptions

#UD	If in 64-bit mode.

AAS—ASCII Adjust AL After Subtraction

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
3F	AAS	A	Invalid	Valid	ASCII adjust AL after subtraction.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Adjusts the result of the subtraction of two unpacked BCD values to create a unpacked BCD result. The AL register is the implied source and destination operand for this instruction. The AAS instruction is only useful when it follows a SUB instruction that subtracts (binary subtraction) one unpacked BCD value from another and stores a byte result in the AL register. The AAA instruction then adjusts the contents of the AL register to contain the correct 1-digit unpacked BCD result.

If the subtraction produced a decimal carry, the AH register decrements by 1, and the CF and AF flags are set. If no decimal carry occurred, the CF and AF flags are cleared, and the AH register is unchanged. In either case, the AL register is left with its top four bits set to 0.

This instruction executes as described in compatibility mode and legacy mode. It is not valid in 64-bit mode.

Operation

```
IF 64-bit mode
    THEN
        #UD;
    ELSE
        IF ((AL AND 0FH) > 9) or (AF = 1)
            THEN
                AL ← AL - 6;
                AH ← AH - 1;
                AF ← 1;
                CF ← 1;
                AL ← AL AND 0FH;
            ELSE
                CF ← 0;
                AF ← 0;
                AL ← AL AND 0FH;
```

```
    FI;
FI;
```

Flags Affected

The AF and CF flags are set to 1 if there is a decimal borrow; otherwise, they are cleared to 0. The OF, SF, ZF, and PF flags are undefined.

Protected Mode Exceptions

#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as protected mode.

Compatibility Mode Exceptions

Same exceptions as protected mode.

64-Bit Mode Exceptions

#UD	If in 64-bit mode.

ADC—Add with Carry

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
14 *ib*	ADC AL, *imm8*	C	Valid	Valid	Add with carry *imm8* to AL.
15 *iw*	ADC AX, *imm16*	C	Valid	Valid	Add with carry *imm16* to AX.
15 *id*	ADC EAX, *imm32*	C	Valid	Valid	Add with carry *imm32* to EAX.
REX.W + 15 *id*	ADC RAX, *imm32*	C	Valid	N.E.	Add with carry *imm32 sign extended to 64-bits* to RAX.
80 */2 ib*	ADC *r/m8, imm8*	B	Valid	Valid	Add with carry *imm8* to *r/m8.*
REX + 80 */2 ib*	ADC *r/m8**, *imm8*	B	Valid	N.E.	Add with carry *imm8* to *r/m8.*
81 */2 iw*	ADC *r/m16, imm16*	B	Valid	Valid	Add with carry *imm16* to *r/m16.*
81 */2 id*	ADC *r/m32, imm32*	B	Valid	Valid	Add with CF *imm32* to *r/m32.*
REX.W + 81 */2 id*	ADC *r/m64, imm32*	B	Valid	N.E.	Add with CF *imm32* sign extended to 64-bits to *r/m64.*
83 */2 ib*	ADC *r/m16, imm8*	B	Valid	Valid	Add with CF sign-extended *imm8* to *r/m16.*
83 */2 ib*	ADC *r/m32, imm8*	B	Valid	Valid	Add with CF sign-extended *imm8* into *r/m32.*
REX.W + 83 */2 ib*	ADC *r/m64, imm8*	B	Valid	N.E.	Add with CF sign-extended *imm8* into *r/m64.*
10 */r*	ADC *r/m8, r8*	A	Valid	Valid	Add with carry byte register to *r/m8.*
REX + 10 */r*	ADC *r/m8**, *r8**	A	Valid	N.E.	Add with carry byte register to *r/m64.*
11 */r*	ADC *r/m16, r16*	A	Valid	Valid	Add with carry *r16* to *r/m16.*
11 */r*	ADC *r/m32, r32*	A	Valid	Valid	Add with CF *r32* to *r/m32.*
REX.W + 11 */r*	ADC *r/m64, r64*	A	Valid	N.E.	Add with CF *r64* to *r/m64.*
12 */r*	ADC *r8, r/m8*	A	Valid	Valid	Add with carry *r/m8* to byte register.
REX + 12 */r*	ADC *r8**, *r/m8**	A	Valid	N.E.	Add with carry *r/m64* to byte register.

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
13 /*r*	ADC *r16, r/m16*	A	Valid	Valid	Add with carry *r/m16* to *r16.*
13 /*r*	ADC *r32, r/m32*	A	Valid	Valid	Add with CF *r/m32* to *r32.*
REX.W + 13 /*r*	ADC *r64, r/m64*	A	Valid	N.E.	Add with CF *r/m64* to *r64.*

NOTES:

*In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r, w)	ModRM:reg (r)	NA	NA
B	ModRM:r/m (r, w)	imm8	NA	NA
C	AL/AX/EAX/RAX	imm8	NA	NA

Description

Adds the destination operand (first operand), the source operand (second operand), and the carry (CF) flag and stores the result in the destination operand. The destination operand can be a register or a memory location; the source operand can be an immediate, a register, or a memory location. (However, two memory operands cannot be used in one instruction.) The state of the CF flag represents a carry from a previous addition. When an immediate value is used as an operand, it is sign-extended to the length of the destination operand format.

The ADC instruction does not distinguish between signed or unsigned operands. Instead, the processor evaluates the result for both data types and sets the OF and CF flags to indicate a carry in the signed or unsigned result, respectively. The SF flag indicates the sign of the signed result.

The ADC instruction is usually executed as part of a multibyte or multiword addition in which an ADD instruction is followed by an ADC instruction.

This instruction can be used with a LOCK prefix to allow the instruction to be executed atomically.

In 64-bit mode, the instruction's default operation size is 32 bits. Using a REX prefix in the form of REX.R permits access to additional registers (R8-R15). Using a REX prefix in the form of REX.W promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
DEST ← DEST + SRC + CF;
```

Flags Affected

The OF, SF, ZF, AF, CF, and PF flags are set according to the result.

Protected Mode Exceptions

#GP(0)	If the destination is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

ADD—Add

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
04 *ib*	ADD AL, *imm8*	C	Valid	Valid	Add *imm8* to AL.
05 *iw*	ADD AX, *imm16*	C	Valid	Valid	Add *imm16* to AX.
05 *id*	ADD EAX, *imm32*	C	Valid	Valid	Add *imm32* to EAX.
REX.W + 05 *id*	ADD RAX, *imm32*	C	Valid	N.E.	Add *imm32 sign-extended to 64-bits* to RAX.
80 /0 *ib*	ADD *r/m8, imm8*	B	Valid	Valid	Add *imm8* to *r/m8.*
REX + 80 /0 *ib*	ADD *r/m8**, *imm8*	B	Valid	N.E.	Add *sign*-extended *imm8* to *r/m64.*
81 /0 *iw*	ADD *r/m16, imm16*	B	Valid	Valid	Add *imm16* to *r/m16.*
81 /0 *id*	ADD *r/m32, imm32*	B	Valid	Valid	Add *imm32* to *r/m32*.
REX.W + 81 /0 *id*	ADD *r/m64, imm32*	B	Valid	N.E.	Add *imm32 sign-extended to 64-bits* to *r/m64*.
83 /0 *ib*	ADD *r/m16, imm8*	B	Valid	Valid	Add *sign*-extended *imm8* to *r/m16.*
83 /0 *ib*	ADD *r/m32, imm8*	B	Valid	Valid	Add *sign*-extended *imm8* to *r/m32.*
REX.W + 83 /0 *ib*	ADD *r/m64, imm8*	B	Valid	N.E.	Add *sign*-extended *imm8* to *r/m64.*
00 /*r*	ADD *r/m8, r8*	A	Valid	Valid	Add *r8* to *r/m8.*
REX + 00 /*r*	ADD *r/m8**, *r8**	A	Valid	N.E.	Add *r8* to *r/m8.*
01 /*r*	ADD *r/m16, r16*	A	Valid	Valid	Add *r16* to *r/m16.*
01 /*r*	ADD *r/m32, r32*	A	Valid	Valid	Add r32 to *r/m32.*
REX.W + 01 /*r*	ADD *r/m64, r64*	A	Valid	N.E.	Add r64 to *r/m64.*
02 /*r*	ADD *r8, r/m8*	A	Valid	Valid	Add *r/m8* to *r8.*
REX + 02 /*r*	ADD *r8**, *r/m8**	A	Valid	N.E.	Add *r/m8* to *r8.*
03 /*r*	ADD *r16, r/m16*	A	Valid	Valid	Add *r/m16* to *r16.*
03 /*r*	ADD *r32, r/m32*	A	Valid	Valid	Add *r/m32* to *r32.*
REX.W + 03 /*r*	ADD *r64, r/m64*	A	Valid	N.E.	Add *r/m64* to *r64.*

NOTES:

*In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (r, w)	imm8	NA	NA
C	AL/AX/EAX/RAX	imm8	NA	NA

Description

Adds the destination operand (first operand) and the source operand (second operand) and then stores the result in the destination operand. The destination operand can be a register or a memory location; the source operand can be an immediate, a register, or a memory location. (However, two memory operands cannot be used in one instruction.) When an immediate value is used as an operand, it is sign-extended to the length of the destination operand format.

The ADD instruction performs integer addition. It evaluates the result for both signed and unsigned integer operands and sets the OF and CF flags to indicate a carry (overflow) in the signed or unsigned result, respectively. The SF flag indicates the sign of the signed result.

This instruction can be used with a LOCK prefix to allow the instruction to be executed atomically.

In 64-bit mode, the instruction’s default operation size is 32 bits. Using a REX prefix in the form of REX.R permits access to additional registers (R8-R15). Using a REX a REX prefix in the form of REX.W promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
DEST ← DEST + SRC;
```

Flags Affected

The OF, SF, ZF, AF, CF, and PF flags are set according to the result.

Protected Mode Exceptions

#GP(0)	If the destination is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.

#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

ADDPD—Add Packed Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 58 /r ADDPD *xmm1, xmm2/m128*	A	V/V	SSE2	Add packed double-precision floating-point values from *xmm2/m128* to *xmm1*.
VEX.NDS.128.66.0F.WIG 58 /r VADDPD xmm1,xmm2, xmm3/m128	B	V/V	AVX	Add packed double-precision floating-point values from xmm3/mem to xmm2 and stores result in xmm1.
VEX.NDS.256.66.0F.WIG 58 /r VADDPD ymm1, ymm2, ymm3/m256	B	V/V	AVX	Add packed double-precision floating-point values from ymm3/mem to ymm2 and stores result in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs a SIMD add of the two packed double-precision floating-point values from the source operand (second operand) and the destination operand (first operand), and stores the packed double-precision floating-point results in the destination operand.

In 64-bit mode, using a REX prefix in the form of REX.R permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified. See Chapter 11 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for an overview of SIMD double-precision floating-point operation.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

ADDPD (128-bit Legacy SSE version)

DEST[63:0] ← DEST[63:0] + SRC[63:0];
DEST[127:64] ← DEST[127:64] + SRC[127:64];
DEST[VLMAX-1:128] (Unmodified)

VADDPD (VEX.128 encoded version)

DEST[63:0] ← SRC1[63:0] + SRC2[63:0]
DEST[127:64] ← SRC1[127:64] + SRC2[127:64]
DEST[VLMAX-1:128] ← 0

VADDPD (VEX.256 encoded version)

DEST[63:0] ← SRC1[63:0] + SRC2[63:0]
DEST[127:64] ← SRC1[127:64] + SRC2[127:64]
DEST[191:128] ← SRC1[191:128] + SRC2[191:128]
DEST[255:192] ← SRC1[255:192] + SRC2[255:192]

.Intel C/C++ Compiler Intrinsic Equivalent

ADDPD __m128d _mm_add_pd (__m128d a, __m128d b)

VADDPD __m256d _mm256_add_pd (__m256d a, __m256d b)

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 2.

ADDPS—Add Packed Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 58 /r ADDPS *xmm1*, *xmm2/m128*	A	V/V	SSE	Add packed single-precision floating-point values from *xmm2/m128* to *xmm1* and stores result in xmm1.
VEX.NDS.128.0F.WIG 58 /r VADDPS xmm1,xmm2, xmm3/m128	B	V/V	AVX	Add packed single-precision floating-point values from xmm3/mem to xmm2 and stores result in xmm1.
VEX.NDS.256.0F.WIG 58 /r VADDPS ymm1, ymm2, ymm3/m256	B	V/V	AVX	Add packed single-precision floating-point values from ymm3/mem to ymm2 and stores result in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r))	NA

Description

Performs a SIMD add of the four packed single-precision floating-point values from the source operand (second operand) and the destination operand (first operand), and stores the packed single-precision floating-point results in the destination operand.

In 64-bit mode, using a REX prefix in the form of REX.R permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified. See Chapter 10 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for an overview of SIMD single-precision floating-point operation.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

ADDPS (128-bit Legacy SSE version)

DEST[31:0] ← DEST[31:0] + SRC[31:0];
DEST[63:32] ← DEST[63:32] + SRC[63:32];
DEST[95:64] ← DEST[95:64] + SRC[95:64];
DEST[127:96] ← DEST[127:96] + SRC[127:96];
DEST[VLMAX-1:128] (Unmodified)

VADDPS (VEX.128 encoded version)

DEST[31:0] ← SRC1[31:0] + SRC2[31:0]
DEST[63:32] ← SRC1[63:32] + SRC2[63:32]
DEST[95:64] ← SRC1[95:64] + SRC2[95:64]
DEST[127:96] ← SRC1[127:96] + SRC2[127:96]
DEST[VLMAX-1:128] ← 0

VADDPS (VEX.256 encoded version)

DEST[31:0] ← SRC1[31:0] + SRC2[31:0]
DEST[63:32] ← SRC1[63:32] + SRC2[63:32]
DEST[95:64] ← SRC1[95:64] + SRC2[95:64]
DEST[127:96] ← SRC1[127:96] + SRC2[127:96]
DEST[159:128] ← SRC1[159:128] + SRC2[159:128]
DEST[191:160]← SRC1[191:160] + SRC2[191:160]
DEST[223:192] ← SRC1[223:192] + SRC2[223:192]
DEST[255:224] ← SRC1[255:224] + SRC2[255:224]

Intel C/C++ Compiler Intrinsic Equivalent

ADDPS __m128 _mm_add_ps(__m128 a, __m128 b)

VADDPS __m256 _mm256_add_ps (__m256 a, __m256 b)

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 2.

ADDSD—Add Scalar Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 58 /r ADDSD *xmm1*, *xmm2/m64*	A	V/V	SSE2	Add the low double-precision floating-point value from *xmm2/m64* to *xmm1*.
VEX.NDS.LIG.F2.0F.WIG 58 /r VADDSD xmm1, xmm2, xmm3/m64	B	V/V	AVX	Add the low double-precision floating-point value from xmm3/mem to xmm2 and store the result in xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r))	NA

Description

Adds the low double-precision floating-point values from the source operand (second operand) and the destination operand (first operand), and stores the double-precision floating-point result in the destination operand.

The source operand can be an XMM register or a 64-bit memory location. The destination operand is an XMM register. See Chapter 11 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for an overview of a scalar double-precision floating-point operation.

In 64-bit mode, using a REX prefix in the form of REX.R permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: Bits (VLMAX-1:64) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (127:64) of the XMM register destination are copied from corresponding bits in the first source operand. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

ADDSD (128-bit Legacy SSE version)

DEST[63:0] ← DEST[63:0] + SRC[63:0]
DEST[VLMAX-1:64] (Unmodified)

VADDSD (VEX.128 encoded version)

DEST[63:0] ← SRC1[63:0] + SRC2[63:0]
DEST[127:64] ← SRC1[127:64]
DEST[VLMAX-1:128] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

ADDSD __m128d _mm_add_sd (m128d a, m128d b)

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 3.

ADDSS—Add Scalar Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 58 /r ADDSS *xmm1, xmm2/m32*	A	V/V	SSE	Add the low single-precision floating-point value from *xmm2/m32* to *xmm1*.
VEX.NDS.LIG.F3.0F.WIG 58 /r VADDSS xmm1,xmm2, xmm3/m32	B	V/V	AVX	Add the low single-precision floating-point value from xmm3/mem to xmm2 and store the result in xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Adds the low single-precision floating-point values from the source operand (second operand) and the destination operand (first operand), and stores the single-precision floating-point result in the destination operand.

The source operand can be an XMM register or a 32-bit memory location. The destination operand is an XMM register. See Chapter 10 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for an overview of a scalar single-precision floating-point operation.

In 64-bit mode, using a REX prefix in the form of REX.R permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: Bits (VLMAX-1:32) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (127:32) of the XMM register destination are copied from corresponding bits in the first source operand. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

ADDSS DEST, SRC (128-bit Legacy SSE version)

```
DEST[31:0] ← DEST[31:0] + SRC[31:0];
DEST[VLMAX-1:32] (Unmodified)
```

VADDSS DEST, SRC1, SRC2 (VEX.128 encoded version)

```
DEST[31:0] ← SRC1[31:0] + SRC2[31:0]
DEST[127:32] ← SRC1[127:32]
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

ADDSS __m128 _mm_add_ss(__m128 a, __m128 b)

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 3.

ADDSUBPD—Packed Double-FP Add/Subtract

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F D0 /r ADDSUBPD *xmm1*, *xmm2/m128*	A	V/V	SSE3	Add/subtract double-precision floating-point values from *xmm2/m128* to *xmm1*.
VEX.NDS.128.66.0F.WIG D0 /r VADDSUBPD xmm1, xmm2, xmm3/m128	B	V/V	AVX	Add/subtract packed double-precision floating-point values from xmm3/mem to xmm2 and stores result in xmm1.
VEX.NDS.256.66.0F.WIG D0 /r VADDSUBPD ymm1, ymm2, ymm3/m256	B	V/V	AVX	Add / subtract packed double-precision floating-point values from ymm3/mem to ymm2 and stores result in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Adds odd-numbered double-precision floating-point values of the first source operand (second operand) with the corresponding double-precision floating-point values from the second source operand (third operand); stores the result in the odd-numbered values of the destination operand (first operand). Subtracts the even-numbered double-precision floating-point values from the second source operand from the corresponding double-precision floating values in the first source operand; stores the result into the even-numbered values of the destination operand.

In 64-bit mode, using a REX prefix in the form of REX.R permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified. See Figure 3-3.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

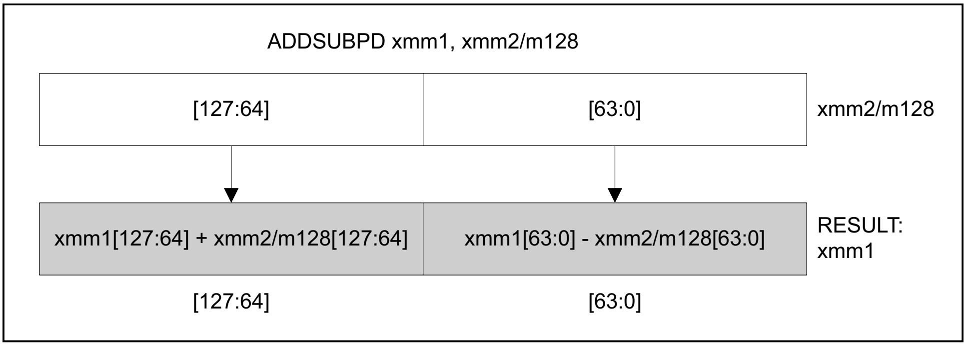


Figure 3-3. ADDSUBPD—Packed Double-FP Add/Subtract

Operation

ADDSUBPD (128-bit Legacy SSE version)

DEST[63:0] ← DEST[63:0] - SRC[63:0]
DEST[127:64] ← DEST[127:64] + SRC[127:64]
DEST[VLMAX-1:128] (Unmodified)

VADDSUBPD (VEX.128 encoded version)

DEST[63:0] ← SRC1[63:0] - SRC2[63:0]
DEST[127:64] ← SRC1[127:64] + SRC2[127:64]
DEST[VLMAX-1:128] ← 0

VADDSUBPD (VEX.256 encoded version)

DEST[63:0] ← SRC1[63:0] - SRC2[63:0]
DEST[127:64] ← SRC1[127:64] + SRC2[127:64]
DEST[191:128] ← SRC1[191:128] - SRC2[191:128]
DEST[255:192] ← SRC1[255:192] + SRC2[255:192]

Intel C/C++ Compiler Intrinsic Equivalent

ADDSUBPD __m128d _mm_addsub_pd(__m128d a, __m128d b)

VADDSUBPD __m256d _mm256_addsub_pd (__m256d a, __m256d b)

Exceptions

When the source operand is a memory operand, it must be aligned on a 16-byte boundary or a general-protection exception (#GP) will be generated.

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 2.

ADDSUBPS—Packed Single-FP Add/Subtract

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F D0 /r ADDSUBPS *xmm1*, *xmm2/m128*	A	V/V	SSE3	Add/subtract single-precision floating-point values from *xmm2/m128* to *xmm1*.
VEX.NDS.128.F2.0F.WIG D0 /r VADDSUBPS xmm1, xmm2, xmm3/m128	B	V/V	AVX	Add/subtract single-precision floating-point values from xmm3/mem to xmm2 and stores result in xmm1.
VEX.NDS.256.F2.0F.WIG D0 /r VADDSUBPS ymm1, ymm2, ymm3/m256	B	V/V	AVX	Add / subtract single-precision floating-point values from ymm3/mem to ymm2 and stores result in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Adds odd-numbered single-precision floating-point values of the first source operand (second operand) with the corresponding single-precision floating-point values from the second source operand (third operand); stores the result in the odd-numbered values of the destination operand (first operand). Subtracts the even-numbered single-precision floating-point values from the second source operand from the corresponding single-precision floating values in the first source operand; stores the result into the even-numbered values of the destination operand.

In 64-bit mode, using a REX prefix in the form of REX.R permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified. See Figure 3-4.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

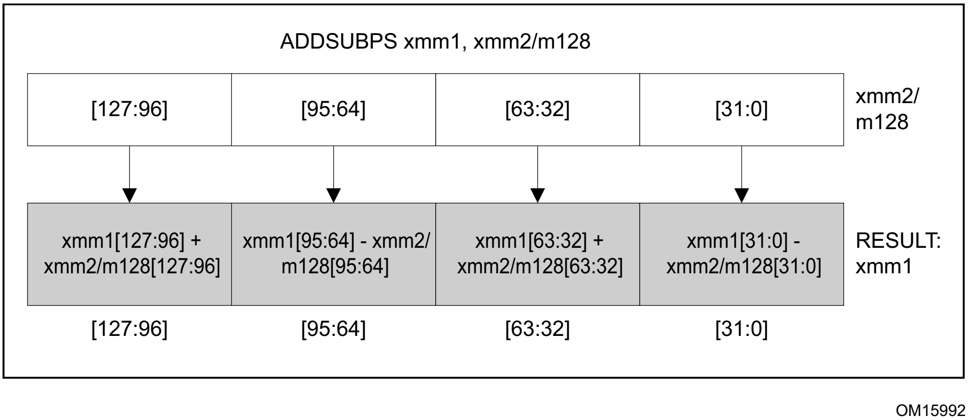


Figure 3-4. ADDSUBPS—Packed Single-FP Add/Subtract

Operation

ADDSUBPS (128-bit Legacy SSE version)

DEST[31:0] ← DEST[31:0] - SRC[31:0]
DEST[63:32] ← DEST[63:32] + SRC[63:32]
DEST[95:64] ← DEST[95:64] - SRC[95:64]
DEST[127:96] ← DEST[127:96] + SRC[127:96]
DEST[VLMAX-1:128] (Unmodified)

VADDSUBPS (VEX.128 encoded version)

DEST[31:0] ← SRC1[31:0] - SRC2[31:0]
DEST[63:32] ← SRC1[63:32] + SRC2[63:32]
DEST[95:64] ← SRC1[95:64] - SRC2[95:64]
DEST[127:96] ← SRC1[127:96] + SRC2[127:96]
DEST[VLMAX-1:128] ← 0

VADDSUBPS (VEX.256 encoded version)

DEST[31:0] ← SRC1[31:0] - SRC2[31:0]
DEST[63:32] ← SRC1[63:32] + SRC2[63:32]
DEST[95:64] ← SRC1[95:64] - SRC2[95:64]

DEST[127:96] ← SRC1[127:96] + SRC2[127:96]
DEST[159:128] ← SRC1[159:128] - SRC2[159:128]
DEST[191:160]← SRC1[191:160] + SRC2[191:160]
DEST[223:192] ← SRC1[223:192] - SRC2[223:192]
DEST[255:224] ← SRC1[255:224] + SRC2[255:224].

Intel C/C++ Compiler Intrinsic Equivalent

ADDSUBPS __m128 _mm_addsub_ps(__m128 a, __m128 b)

VADDSUBPS __m256 _mm256_addsub_ps (__m256 a, __m256 b)

Exceptions

When the source operand is a memory operand, the operand must be aligned on a 16-byte boundary or a general-protection exception (#GP) will be generated.

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 2.

AESDEC—Perform One Round of an AES Decryption Flow

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 38 DE /r AESDEC xmm1, xmm2/m128	A	V/V	AES	Perform one round of an AES decryption flow, using the Equivalent Inverse Cipher, operating on a 128-bit data (state) from xmm1 with a 128-bit round key from xmm2/m128.
VEX.NDS.128.66.0F38.WIG DE /r VAESDEC xmm1, xmm2, xmm3/m128	B	V/V	Both AES and AVX flags	Perform one round of an AES decryption flow, using the Equivalent Inverse Cipher, operating on a 128-bit data (state) from xmm2 with a 128-bit round key from xmm3/m128; store the result in xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand2	Operand3	Operand4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

This instruction performs a single round of the AES decryption flow using the Equivalent Inverse Cipher, with the round key from the second source operand, operating on a 128-bit data (state) from the first source operand, and store the result in the destination operand.

Use the AESDEC instruction for all but the last decryption round. For the last decryption round, use the AESDECCLAST instruction.

128-bit Legacy SSE version: The first source operand and the destination operand are the same and must be an XMM register. The second source operand can be an XMM register or a 128-bit memory location. Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: The first source operand and the destination operand are XMM registers. The second source operand can be an XMM register or a 128-bit memory location. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

AESDEC

```
STATE ← SRC1;
RoundKey ← SRC2;
STATE ← InvShiftRows( STATE );
STATE ← InvSubBytes( STATE );
STATE ← InvMixColumns( STATE );
DEST[127:0] ← STATE XOR RoundKey;
DEST[VLMAX-1:128] (Unmodified)
```

VAESDEC

```
STATE ← SRC1;
RoundKey ← SRC2;
STATE ← InvShiftRows( STATE );
STATE ← InvSubBytes( STATE );
STATE ← InvMixColumns( STATE );
DEST[127:0] ← STATE XOR RoundKey;
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

(V)AESDEC __m128i _mm_aesdec (__m128i, __m128i)

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 4.

AESDECLAST—Perform Last Round of an AES Decryption Flow

Opcode	Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 38 DF /r AESDECLAST xmm1, xmm2/m128		A	V/V	AES	Perform the last round of an AES decryption flow, using the Equivalent Inverse Cipher, operating on a 128-bit data (state) from xmm1 with a 128-bit round key from xmm2/m128.
VEX.NDS.128.66.0F38.WIG DF /r VAESDECLAST xmm1, xmm2, xmm3/m128		B	V/V	Both AES and AVX flags	Perform the last round of an AES decryption flow, using the Equivalent Inverse Cipher, operating on a 128-bit data (state) from xmm2 with a 128-bit round key from xmm3/m128; store the result in xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand2	Operand3	Operand4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

This instruction performs the last round of the AES decryption flow using the Equivalent Inverse Cipher, with the round key from the second source operand, operating on a 128-bit data (state) from the first source operand, and store the result in the destination operand.

128-bit Legacy SSE version: The first source operand and the destination operand are the same and must be an XMM register. The second source operand can be an XMM register or a 128-bit memory location. Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: The first source operand and the destination operand are XMM registers. The second source operand can be an XMM register or a 128-bit memory location. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

AESDECLAST

```
STATE ← SRC1;
RoundKey ← SRC2;
STATE ← InvShiftRows( STATE );
STATE ← InvSubBytes( STATE );
DEST[127:0] ← STATE XOR RoundKey;
DEST[VLMAX-1:128] (Unmodified)
```

VAESDECLAST

```
STATE ← SRC1;
RoundKey ← SRC2;
STATE ← InvShiftRows( STATE );
STATE ← InvSubBytes( STATE );
DEST[127:0] ← STATE XOR RoundKey;
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

(V)AESDECLAST __m128i _mm_aesdeclast (__m128i, __m128i)

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 4.

AESENC—Perform One Round of an AES Encryption Flow

Opcode	Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 38 DC /r	AESENC xmm1, xmm2/m128	A	V/V	AES	Perform one round of an AES encryption flow, operating on a 128-bit data (state) from xmm1 with a 128-bit round key from xmm2/m128.
VEX.NDS.128.66.0F38.WIG DC /r	VAESENC xmm1, xmm2, xmm3/m128	B	V/V	Both AES and AVX flags	Perform one round of an AES encryption flow, operating on a 128-bit data (state) from xmm2 with a 128-bit round key from the xmm3/m128; store the result in xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand2	Operand3	Operand4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

This instruction performs a single round of an AES encryption flow using a round key from the second source operand, operating on 128-bit data (state) from the first source operand, and store the result in the destination operand.

Use the AESENC instruction for all but the last encryption rounds. For the last encryption round, use the AESENCCLAST instruction.

128-bit Legacy SSE version: The first source operand and the destination operand are the same and must be an XMM register. The second source operand can be an XMM register or a 128-bit memory location. Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: The first source operand and the destination operand are XMM registers. The second source operand can be an XMM register or a 128-bit memory location. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

AESENC

```
STATE ← SRC1;
RoundKey ← SRC2;
STATE ← ShiftRows( STATE );
STATE ← SubBytes( STATE );
STATE ← MixColumns( STATE );
DEST[127:0] ← STATE XOR RoundKey;
DEST[VLMAX-1:128] (Unmodified)
```

VAESENC

```
STATE ← SRC1;
RoundKey ← SRC2;
STATE ← ShiftRows( STATE );
STATE ← SubBytes( STATE );
STATE ← MixColumns( STATE );
DEST[127:0] ← STATE XOR RoundKey;
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

(V)AESENC __m128i _mm_aesenc (__m128i, __m128i)

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 4.

AESENCLAST—Perform Last Round of an AES Encryption Flow

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 38 DD /r AESENCLAST xmm1, xmm2/m128	A	V/V	AES	Perform the last round of an AES encryption flow, operating on a 128-bit data (state) from xmm1 with a 128-bit round key from xmm2/m128.
VEX.NDS.128.66.0F38.WIG DD /r VAESENCLAST xmm1, xmm2, xmm3/m128	B	V/V	Both AES and AVX flags	Perform the last round of an AES encryption flow, operating on a 128-bit data (state) from xmm2 with a 128 bit round key from xmm3/m128; store the result in xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand2	Operand3	Operand4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

This instruction performs the last round of an AES encryption flow using a round key from the second source operand, operating on 128-bit data (state) from the first source operand, and store the result in the destination operand.

128-bit Legacy SSE version: The first source operand and the destination operand are the same and must be an XMM register. The second source operand can be an XMM register or a 128-bit memory location. Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: The first source operand and the destination operand are XMM registers. The second source operand can be an XMM register or a 128-bit memory location. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

AESENCLAST

```
STATE ← SRC1;
RoundKey ← SRC2;
```

STATE ← ShiftRows(STATE);
STATE ← SubBytes(STATE);
DEST[127:0] ← STATE XOR RoundKey;
DEST[VLMAX-1:128] (Unmodified)

VAESENCLAST
STATE ← SRC1;
RoundKey ← SRC2;
STATE ← ShiftRows(STATE);
STATE ← SubBytes(STATE);
DEST[127:0] ← STATE XOR RoundKey;
DEST[VLMAX-1:128] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

(V)AESENCLAST __m128i _mm_aesenclast (__m128i, __m128i)

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 4.

AESIMC—Perform the AES InvMixColumn Transformation

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 38 DB /r AESIMC xmm1, xmm2/m128	A	V/V	AES	Perform the InvMixColumn transformation on a 128-bit round key from xmm2/m128 and store the result in xmm1.
VEX.128.66.0F38.WIG DB /r VAESIMC xmm1, xmm2/m128	A	V/V	Both AES and AVX flags	Perform the InvMixColumn transformation on a 128-bit round key from xmm2/m128 and store the result in xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand2	Operand3	Operand4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Perform the InvMixColumns transformation on the source operand and store the result in the destination operand. The destination operand is an XMM register. The source operand can be an XMM register or a 128-bit memory location.

Note: the AESIMC instruction should be applied to the expanded AES round keys (except for the first and last round key) in order to prepare them for decryption using the "Equivalent Inverse Cipher" (defined in FIPS 197).

128-bit Legacy SSE version: Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

AESIMC

DEST[127:0] ← InvMixColumns(SRC);
DEST[VLMAX-1:128] (Unmodified)

VAESIMC

DEST[127:0] ← InvMixColumns(SRC);
DEST[VLMAX-1:128] ← 0;

Intel C/C++ Compiler Intrinsic Equivalent

(V)AESIMC __m128i _mm_aesimc (__m128i)

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 4.

AESKEYGENASSIST—AES Round Key Generation Assist

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 3A DF /r ib AESKEYGENASSIST xmm1, xmm2/m128, imm8	A	V/V	AES	Assist in AES round key generation using an 8 bits Round Constant (RCON) specified in the immediate byte, operating on 128 bits of data specified in xmm2/m128 and stores the result in xmm1.
VEX.128.66.0F3A.WIG DF /r ib VAESKEYGENASSIST xmm1, xmm2/m128, imm8	A	V/V	Both AES and AVX flags	Assist in AES round key generation using 8 bits Round Constant (RCON) specified in the immediate byte, operating on 128 bits of data specified in xmm2/m128 and stores the result in xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand2	Operand3	Operand4
A	ModRM:reg (w)	ModRM:r/m (r)	imm8	NA

Description

Assist in expanding the AES cipher key, by computing steps towards generating a round key for encryption, using 128-bit data specified in the source operand and an 8-bit round constant specified as an immediate, store the result in the destination operand.

The destination operand is an XMM register. The source operand can be an XMM register or a 128-bit memory location.

128-bit Legacy SSE version:Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

AESKEYGENASSIST

X3[31:0] ← SRC [127: 96];
X2[31:0] ← SRC [95: 64];

```
X1[31:0] ← SRC [63: 32];
X0[31:0] ← SRC [31: 0];
RCON[31:0] ← ZeroExtend(Imm8[7:0]);
DEST[31:0] ← SubWord(X1);
DEST[63:32 ] ← RotWord( SubWord(X1) ) XOR RCON;
DEST[95:64] ← SubWord(X3);
DEST[127:96] ← RotWord( SubWord(X3) ) XOR RCON;
DEST[VLMAX-1:128] (Unmodified)
```

VAESKEYGENASSIST

```
X3[31:0] ← SRC [127: 96];
X2[31:0] ← SRC [95: 64];
X1[31:0] ← SRC [63: 32];
X0[31:0] ← SRC [31: 0];
RCON[31:0] ← ZeroExtend(Imm8[7:0]);
DEST[31:0] ← SubWord(X1);
DEST[63:32 ] ← RotWord( SubWord(X1) ) XOR RCON;
DEST[95:64] ← SubWord(X3);
DEST[127:96] ← RotWord( SubWord(X3) ) XOR RCON;
DEST[VLMAX-1:128] ← 0;
```

Intel C/C++ Compiler Intrinsic Equivalent

(V)AESKEYGENASSIST __m128i _mm_aesimc (__m128i, const int)

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 4.

AND—Logical AND

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
24 *ib*	AND AL, *imm8*	C	Valid	Valid	AL AND *imm8.*
25 *iw*	AND AX, *imm16*	C	Valid	Valid	AX AND i*mm16.*
25 *id*	AND EAX, *imm32*	C	Valid	Valid	EAX AND *imm32.*
REX.W + 25 *id*	AND RAX, *imm32*	C	Valid	N.E.	RAX AND *imm32 sign-extended to 64-bits.*
80 /4 *ib*	AND *r/m8, imm8*	B	Valid	Valid	*r/m8* AND *imm8.*
REX + 80 /4 *ib*	AND *r/m8**, *imm8*	B	Valid	N.E.	*r/m8* AND *imm8.*
81 /4 *iw*	AND *r/m16, imm16*	B	Valid	Valid	*r/m16* AND *imm16.*
81 /4 *id*	AND *r/m32, imm32*	B	Valid	Valid	*r/m32* AND *imm32.*
REX.W + 81 /4 *id*	AND *r/m64, imm32*	B	Valid	N.E.	*r/m64* AND *imm32 sign extended to 64-bits.*
83 /4 *ib*	AND *r/m16, imm8*	B	Valid	Valid	*r/m16* AND *imm8 (sign-extended).*
83 /4 *ib*	AND *r/m32, imm8*	B	Valid	Valid	*r/m32* AND *imm8 (sign-extended).*
REX.W + 83 /4 *ib*	AND *r/m64, imm8*	B	Valid	N.E.	*r/m64* AND *imm8 (sign-extended).*
20 /*r*	AND *r/m8, r8*	A	Valid	Valid	*r/m8* AND *r8.*
REX + 20 /*r*	AND *r/m8**, *r8**	A	Valid	N.E.	*r/m64* AND *r8 (sign-extended).*
21 /*r*	AND *r/m16, r16*	A	Valid	Valid	*r/m16* AND *r16.*
21 /*r*	AND *r/m32, r32*	A	Valid	Valid	*r/m32* AND *r32.*
REX.W + 21 /*r*	AND *r/m64, r64*	A	Valid	N.E.	*r/m64* AND *r32.*
22 /*r*	AND *r8, r/m8*	A	Valid	Valid	*r8* AND *r/m8.*
REX + 22 /*r*	AND *r8**, *r/m8**	A	Valid	N.E.	*r/m64* AND *r8 (sign-extended).*
23 /*r*	AND *r16, r/m16*	A	Valid	Valid	*r16* AND *r/m16.*
23 /*r*	AND *r32, r/m32*	A	Valid	Valid	*r32* AND *r/m32.*
REX.W + 23 /*r*	AND *r64, r/m64*	A	Valid	N.E.	*r64* AND *r/m64.*

NOTES:

*In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (r, w)	imm8	NA	NA
C	AL/AX/EAX/RAX	imm8	NA	NA

Description

Performs a bitwise AND operation on the destination (first) and source (second) operands and stores the result in the destination operand location. The source operand can be an immediate, a register, or a memory location; the destination operand can be a register or a memory location. (However, two memory operands cannot be used in one instruction.) Each bit of the result is set to 1 if both corresponding bits of the first and second operands are 1; otherwise, it is set to 0.

This instruction can be used with a LOCK prefix to allow the it to be executed atomically.

In 64-bit mode, the instruction's default operation size is 32 bits. Using a REX prefix in the form of REX.R permits access to additional registers (R8-R15). Using a REX prefix in the form of REX.W promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
DEST ← DEST AND SRC;
```

Flags Affected

The OF and CF flags are cleared; the SF, ZF, and PF flags are set according to the result. The state of the AF flag is undefined.

Protected Mode Exceptions

#GP(0)	If the destination operand points to a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

#UD	If the LOCK prefix is used but the destination is not a memory operand.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

ANDPD—Bitwise Logical AND of Packed Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 54 /r ANDPD *xmm1*, *xmm2/m128*	A	V/V	SSE2	Return the bitwise logical AND of packed double-precision floating-point values in *xmm1* and *xmm2/m128*.
VEX.NDS.128.66.0F.WIG 54 /r VANDPD xmm1, xmm2, xmm3/m128	B	V/V	AVX	Return the bitwise logical AND of packed double-precision floating-point values in xmm2 and xmm3/mem.
VEX.NDS.256.66.0F.WIG 54 /r VANDPD ymm1, ymm2, ymm3/m256	B	V/V	AVX	Return the bitwise logical AND of packed double-precision floating-point values in ymm2 and ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs a bitwise logical AND of the two packed double-precision floating-point values from the source operand (second operand) and the destination operand (first operand), and stores the result in the destination operand.

In 64-bit mode, using a REX prefix in the form of REX.R permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

ANDPD (128-bit Legacy SSE version)

DEST[63:0] ← DEST[63:0] BITWISE AND SRC[63:0]
DEST[127:64] ← DEST[127:64] BITWISE AND SRC[127:64]
DEST[VLMAX-1:128] (Unmodified)

VANDPD (VEX.128 encoded version)

DEST[63:0] ← SRC1[63:0] BITWISE AND SRC2[63:0]
DEST[127:64] ← SRC1[127:64] BITWISE AND SRC2[127:64]
DEST[VLMAX-1:128] ← 0

VANDPD (VEX.256 encoded version)

DEST[63:0] ← SRC1[63:0] BITWISE AND SRC2[63:0]
DEST[127:64] ← SRC1[127:64] BITWISE AND SRC2[127:64]
DEST[191:128] ← SRC1[191:128] BITWISE AND SRC2[191:128]
DEST[255:192] ← SRC1[255:192] BITWISE AND SRC2[255:192]

Intel C/C++ Compiler Intrinsic Equivalent

ANDPD __m128d _mm_and_pd(__m128d a, __m128d b)

VANDPD __m256d _mm256_and_pd (__m256d a, __m256d b)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 4.

ANDPS—Bitwise Logical AND of Packed Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 54 /r ANDPS *xmm1*, *xmm2/m128*	A	V/V	SSE	Bitwise logical AND of *xmm2/m128* and *xmm1*.
VEX.NDS.128.0F.WIG 54 /r VANDPS xmm1,xmm2, xmm3/m128	B	V/V	AVX	Return the bitwise logical AND of packed single-precision floating-point values in xmm2 and xmm3/mem.
VEX.NDS.256.0F.WIG 54 /r VANDPS ymm1, ymm2, ymm3/m256	B	V/V	AVX	Return the bitwise logical AND of packed single-precision floating-point values in ymm2 and ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs a bitwise logical AND of the four or eight packed single-precision floating-point values from the first source operand and the second source operand, and stores the result in the destination operand.

In 64-bit mode, using a REX prefix in the form of REX.R permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

ANDPS (128-bit Legacy SSE version)
DEST[31:0] ← DEST[31:0] BITWISE AND SRC[31:0]
DEST[63:32] ← DEST[63:32] BITWISE AND SRC[63:32]
DEST[95:64] ← DEST[95:64] BITWISE AND SRC[95:64]
DEST[127:96] ← DEST[127:96] BITWISE AND SRC[127:96]
DEST[VLMAX-1:128] (Unmodified)

VANDPS (VEX.128 encoded version)
DEST[31:0] ← SRC1[31:0] BITWISE AND SRC2[31:0]
DEST[63:32] ← SRC1[63:32] BITWISE AND SRC2[63:32]
DEST[95:64] ← SRC1[95:64] BITWISE AND SRC2[95:64]
DEST[127:96] ← SRC1[127:96] BITWISE AND SRC2[127:96]
DEST[VLMAX-1:128] ← 0

VANDPS (VEX.256 encoded version)
DEST[31:0] ← SRC1[31:0] BITWISE AND SRC2[31:0]
DEST[63:32] ← SRC1[63:32] BITWISE AND SRC2[63:32]
DEST[95:64] ← SRC1[95:64] BITWISE AND SRC2[95:64]
DEST[127:96] ← SRC1[127:96] BITWISE AND SRC2[127:96]
DEST[159:128] ← SRC1[159:128] BITWISE AND SRC2[159:128]
DEST[191:160]← SRC1[191:160] BITWISE AND SRC2[191:160]
DEST[223:192] ← SRC1[223:192] BITWISE AND SRC2[223:192]
DEST[255:224] ← SRC1[255:224] BITWISE AND SRC2[255:224].

Intel C/C++ Compiler Intrinsic Equivalent

ANDPS __m128 _mm_and_ps(__m128 a, __m128 b)

VANDPS __m256 _mm256_and_ps (__m256 a, __m256 b)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 4.

ANDNPD—Bitwise Logical AND NOT of Packed Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 55 /r ANDNPD *xmm1*, *xmm2/m128*	A	V/V	SSE2	Bitwise logical AND NOT of *xmm2/m128* and *xmm1*.
VEX.NDS.128.66.0F.WIG 55 /r VANDNPD xmm1, xmm2, xmm3/m128	B	V/V	AVX	Return the bitwise logical AND NOT of packed double-precision floating-point values in xmm2 and xmm3/mem.
VEX.NDS.256.66.0F.WIG 55/r VANDNPD ymm1, ymm2, ymm3/m256	B	V/V	AVX	Return the bitwise logical AND NOT of packed double-precision floating-point values in ymm2 and ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs a bitwise logical AND NOT of the two or four packed double-precision floating-point values from the first source operand and the second source operand, and stores the result in the destination operand.

In 64-bit mode, using a REX prefix in the form of REX.R permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

ANDNPD (128-bit Legacy SSE version)
DEST[63:0] ← (NOT(DEST[63:0])) BITWISE AND SRC[63:0]
DEST[127:64] ← (NOT(DEST[127:64])) BITWISE AND SRC[127:64]
DEST[VLMAX-1:128] (Unmodified)

VANDNPD (VEX.128 encoded version)
DEST[63:0] ← (NOT(SRC1[63:0])) BITWISE AND SRC2[63:0]
DEST[127:64] ← (NOT(SRC1[127:64])) BITWISE AND SRC2[127:64]
DEST[VLMAX-1:128] ← 0

VANDNPD (VEX.256 encoded version)
DEST[63:0] ← (NOT(SRC1[63:0])) BITWISE AND SRC2[63:0]
DEST[127:64] ← (NOT(SRC1[127:64])) BITWISE AND SRC2[127:64]
DEST[191:128] ← (NOT(SRC1[191:128])) BITWISE AND SRC2[191:128]
DEST[255:192] ← (NOT(SRC1[255:192])) BITWISE AND SRC2[255:192]

Intel C/C++ Compiler Intrinsic Equivalent

ANDNPD __m128d _mm_andnot_pd(__m128d a, __m128d b)

VANDNPD __m256d _mm256_andnot_pd (__m256d a, __m256d b)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 4.

ANDNPS—Bitwise Logical AND NOT of Packed Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 55 /r ANDNPS *xmm1, xmm2/m128*	A	V/V	SSE	Bitwise logical AND NOT of *xmm2/m128* and *xmm1*.
VEX.NDS.128.0F.WIG 55 /r VANDNPS xmm1, xmm2, xmm3/m128	B	V/V	AVX	Return the bitwise logical AND NOT of packed single-precision floating-point values in xmm2 and xmm3/mem.
VEX.NDS.256.0F.WIG 55 /r VANDNPS ymm1, ymm2, ymm3/m256	B	V/V	AVX	Return the bitwise logical AND NOT of packed single-precision floating-point values in ymm2 and ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Inverts the bits of the four packed single-precision floating-point values in the destination operand (first operand), performs a bitwise logical AND of the four packed single-precision floating-point values in the source operand (second operand) and the temporary inverted result, and stores the result in the destination operand.

In 64-bit mode, using a REX prefix in the form of REX.R permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

ANDNPS (128-bit Legacy SSE version)

DEST[31:0] ← (NOT(DEST[31:0])) BITWISE AND SRC[31:0]
DEST[63:32] ← (NOT(DEST[63:32])) BITWISE AND SRC[63:32]
DEST[95:64] ← (NOT(DEST[95:64])) BITWISE AND SRC[95:64]
DEST[127:96] ← (NOT(DEST[127:96])) BITWISE AND SRC[127:96]
DEST[VLMAX-1:128] (Unmodified)

VANDNPS (VEX.128 encoded version)

DEST[31:0] ← (NOT(SRC1[31:0])) BITWISE AND SRC2[31:0]
DEST[63:32] ← (NOT(SRC1[63:32])) BITWISE AND SRC2[63:32]
DEST[95:64] ← (NOT(SRC1[95:64])) BITWISE AND SRC2[95:64]
DEST[127:96] ← (NOT(SRC1[127:96])) BITWISE AND SRC2[127:96]
DEST[VLMAX-1:128] ← 0

VANDNPS (VEX.256 encoded version)

DEST[31:0] ← (NOT(SRC1[31:0])) BITWISE AND SRC2[31:0]
DEST[63:32] ← (NOT(SRC1[63:32])) BITWISE AND SRC2[63:32]
DEST[95:64] ← (NOT(SRC1[95:64])) BITWISE AND SRC2[95:64]
DEST[127:96] ← (NOT(SRC1[127:96])) BITWISE AND SRC2[127:96]
DEST[159:128] ← (NOT(SRC1[159:128])) BITWISE AND SRC2[159:128]
DEST[191:160]← (NOT(SRC1[191:160])) BITWISE AND SRC2[191:160]
DEST[223:192] ← (NOT(SRC1[223:192])) BITWISE AND SRC2[223:192]
DEST[255:224] ← (NOT(SRC1[255:224])) BITWISE AND SRC2[255:224].

Intel C/C++ Compiler Intrinsic Equivalent

ANDNPS __m128 _mm_andnot_ps(__m128 a, __m128 b)

VANDNPS __m256 _mm256_andnot_ps (__m256 a, __m256 b)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 4.

ARPL—Adjust RPL Field of Segment Selector

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
63 */r*	ARPL *r/m16, r16*	A	N. E.	Valid	Adjust RPL of *r/m16* to not less than RPL of *r16.*

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

Compares the RPL fields of two segment selectors. The first operand (the destination operand) contains one segment selector and the second operand (source operand) contains the other. (The RPL field is located in bits 0 and 1 of each operand.) If the RPL field of the destination operand is less than the RPL field of the source operand, the ZF flag is set and the RPL field of the destination operand is increased to match that of the source operand. Otherwise, the ZF flag is cleared and no change is made to the destination operand. (The destination operand can be a word register or a memory location; the source operand must be a word register.)

The ARPL instruction is provided for use by operating-system procedures (however, it can also be used by applications). It is generally used to adjust the RPL of a segment selector that has been passed to the operating system by an application program to match the privilege level of the application program. Here the segment selector passed to the operating system is placed in the destination operand and segment selector for the application program's code segment is placed in the source operand. (The RPL field in the source operand represents the privilege level of the application program.) Execution of the ARPL instruction then ensures that the RPL of the segment selector received by the operating system is no lower (does not have a higher privilege) than the privilege level of the application program (the segment selector for the application program's code segment can be read from the stack following a procedure call).

This instruction executes as described in compatibility mode and legacy mode. It is not encodable in 64-bit mode.

See "Checking Caller Access Privileges" in Chapter 3, "Protected-Mode Memory Management," of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*, for more information about the use of this instruction.

Operation

```
IF 64-BIT MODE
    THEN
        See MOVSXD;
    ELSE
        IF DEST[RPL) < SRC[RPL)
            THEN
                ZF ← 1;
                DEST[RPL) ← SRC[RPL);
            ELSE
                ZF ← 0;
        FI;
FI;
```

Flags Affected

The ZF flag is set to 1 if the RPL field of the destination operand is less than that of the source operand; otherwise, it is set to 0.

Protected Mode Exceptions

#GP(0)	If the destination is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#UD	The ARPL instruction is not recognized in real-address mode.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#UD	The ARPL instruction is not recognized in virtual-8086 mode.
	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Not applicable.

BLENDPD — Blend Packed Double Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 3A 0D /r ib BLENDPD *xmm1*, *xmm2/m128, imm8*	A	V/V	SSE4_1	Select packed DP-FP values from *xmm1* and *xmm2/m128* from mask specified in imm8 and store the values into *xmm1*.
VEX.NDS.128.66.0F3A.WIG 0D /r ib VBLENDPD xmm1, xmm2, xmm3/m128, imm8	B	V/V	AVX	Select packed double-precision floating-point Values from xmm2 and xmm3/m128 from mask in imm8 and store the values in xmm1.
VEX.NDS.256.66.0F3A.WIG 0D /r ib VBLENDPD ymm1, ymm2, ymm3/m256, imm8	B	V/V	AVX	Select packed double-precision floating-point Values from ymm2 and ymm3/m256 from mask in imm8 and store the values in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	imm8	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	imm8[3:0]

Description

Double-precision floating-point values from the second source operand (third operand) are conditionally merged with values from the first source operand (second operand) and written to the destination operand (first operand). The immediate bits [3:0] determine whether the corresponding double-precision floating-point value in the destination is copied from the second source or first source. If a bit in the mask, corresponding to a word, is “1", then the double-precision floating-point value in the second source operand is copied, else the value in the first source operand is copied.

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register. The second source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

BLENDPD (128-bit Legacy SSE version)

```
IF (IMM8[0] = 0)THEN DEST[63:0] ← DEST[63:0]
        ELSE DEST [63:0] ← SRC[63:0] FI
IF (IMM8[1] = 0) THEN DEST[127:64] ← DEST[127:64]
        ELSE DEST [127:64] ← SRC[127:64] FI
DEST[VLMAX-1:128] (Unmodified)
```

VBLENDPD (VEX.128 encoded version)

```
IF (IMM8[0] = 0)THEN DEST[63:0] ← SRC1[63:0]
        ELSE DEST [63:0] ← SRC2[63:0] FI
IF (IMM8[1] = 0) THEN DEST[127:64] ← SRC1[127:64]
        ELSE DEST [127:64] ← SRC2[127:64] FI
DEST[VLMAX-1:128] ← 0
```

VBLENDPD (VEX.256 encoded version)

```
IF (IMM8[0] = 0)THEN DEST[63:0] ← SRC1[63:0]
        ELSE DEST [63:0] ← SRC2[63:0] FI
IF (IMM8[1] = 0) THEN DEST[127:64] ← SRC1[127:64]
        ELSE DEST [127:64] ← SRC2[127:64] FI
IF (IMM8[2] = 0) THEN DEST[191:128] ← SRC1[191:128]
        ELSE DEST [191:128] ← SRC2[191:128] FI
IF (IMM8[3] = 0) THEN DEST[255:192] ← SRC1[255:192]
        ELSE DEST [255:192] ← SRC2[255:192] FI
```

Intel C/C++ Compiler Intrinsic Equivalent

BLENDPD __m128d _mm_blend_pd (__m128d v1, __m128d v2, const int mask);

VBLENDPD __m256d _mm256_blend_pd (__m256d a, __m256d b, const int mask);

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 4.

BLENDPS — Blend Packed Single Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 3A 0C /r ib BLENDPS *xmm1, xmm2/m128, imm8*	A	V/V	SSE4_1	Select packed single precision floating-point values from *xmm1* and *xmm2/m128* from mask specified in *imm8* and store the values into *xmm1.*
VEX.NDS.128.66.0F3A.WIG 0C /r ib VBLENDPS xmm1, xmm2, xmm3/m128, imm8	B	V/V	AVX	Select packed single-precision floating-point values from xmm2 and xmm3/m128 from mask in imm8 and store the values in xmm1.
VEX.NDS.256.66.0F3A.WIG 0C /r ib VBLENDPS ymm1, ymm2, ymm3/m256, imm8	B	V/V	AVX	Select packed single-precision floating-point values from ymm2 and ymm3/m256 from mask in imm8 and store the values in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	imm8	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	imm8

Description

Packed single-precision floating-point values from the second source operand (third operand) are conditionally merged with values from the first source operand (second operand) and written to the destination operand (first operand). The immediate bits [7:0] determine whether the corresponding single precision floating-point value in the destination is copied from the second source or first source. If a bit in the mask, corresponding to a word, is "1", then the single-precision floating-point value in the second source operand is copied, else the value in the first source operand is copied.

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: The first source operand an XMM register. The second source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

BLENDPS (128-bit Legacy SSE version)

```
IF (IMM8[0] = 0) THEN DEST[31:0] ←DEST[31:0]
        ELSE DEST [31:0] ← SRC[31:0] FI
IF (IMM8[1] = 0) THEN DEST[63:32] ← DEST[63:32]
        ELSE DEST [63:32] ← SRC[63:32] FI
IF (IMM8[2] = 0) THEN DEST[95:64] ← DEST[95:64]
        ELSE DEST [95:64] ← SRC[95:64] FI
IF (IMM8[3] = 0) THEN DEST[127:96] ← DEST[127:96]
        ELSE DEST [127:96] ← SRC[127:96] FI
DEST[VLMAX-1:128] (Unmodified)
```

VBLENDPS (VEX.128 encoded version)

```
IF (IMM8[0] = 0) THEN DEST[31:0] ←SRC1[31:0]
        ELSE DEST [31:0] ← SRC2[31:0] FI
IF (IMM8[1] = 0) THEN DEST[63:32] ← SRC1[63:32]
        ELSE DEST [63:32] ← SRC2[63:32] FI
IF (IMM8[2] = 0) THEN DEST[95:64] ← SRC1[95:64]
        ELSE DEST [95:64] ← SRC2[95:64] FI
IF (IMM8[3] = 0) THEN DEST[127:96] ← SRC1[127:96]
        ELSE DEST [127:96] ← SRC2[127:96] FI
DEST[VLMAX-1:128] ← 0
```

VBLENDPS (VEX.256 encoded version)

```
IF (IMM8[0] = 0) THEN DEST[31:0] ←SRC1[31:0]
        ELSE DEST [31:0] ← SRC2[31:0] FI
IF (IMM8[1] = 0) THEN DEST[63:32] ← SRC1[63:32]
        ELSE DEST [63:32] ← SRC2[63:32] FI
IF (IMM8[2] = 0) THEN DEST[95:64] ← SRC1[95:64]
        ELSE DEST [95:64] ← SRC2[95:64] FI
IF (IMM8[3] = 0) THEN DEST[127:96] ← SRC1[127:96]
        ELSE DEST [127:96] ← SRC2[127:96] FI
IF (IMM8[4] = 0) THEN DEST[159:128] ← SRC1[159:128]
        ELSE DEST [159:128] ← SRC2[159:128] FI
IF (IMM8[5] = 0) THEN DEST[191:160] ← SRC1[191:160]
        ELSE DEST [191:160] ← SRC2[191:160] FI
```

```
IF (IMM8[6] = 0) THEN DEST[223:192] ← SRC1[223:192]
        ELSE DEST [223:192] ← SRC2[223:192] FI
IF (IMM8[7] = 0) THEN DEST[255:224] ← SRC1[255:224]
        ELSE DEST [255:224] ← SRC2[255:224] FI.
```

Intel C/C++ Compiler Intrinsic Equivalent

BLENDPS __m128 _mm_blend_ps (__m128 v1, __m128 v2, const int mask);

VBLENDPS __m256 _mm256_blend_ps (__m256 a, __m256 b, const int mask);

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 4.

BLENDVPD — Variable Blend Packed Double Precision Floating-Point Values

Opcode/ Instruction	**Op/ En**	**64/32-bit Mode**	**CPUID Feature Flag**	**Description**
66 0F 38 15 /r BLENDVPD *xmm1, xmm2/m128 , <XMM0>*	A	V/V	SSE4_1	Select packed DP FP values from *xmm1* and *xmm2* from mask specified in *XMM0* and store the values in *xmm1*.
VEX.NDS.128.66.0F3A.W0 4B /r /is4 VBLENDVPD xmm1, xmm2, xmm3/m128, xmm4	B	V/V	AVX	Conditionally copy double-precision floating-point values from xmm2 or xmm3/m128 to xmm1, based on mask bits in the mask operand, xmm4.
VEX.NDS.256.66.0F3A.W0 4B /r /is4 VBLENDVPD ymm1, ymm2, ymm3/m256, ymm4	B	V/V	AVX	Conditionally copy double-precision floating-point values from ymm2 or ymm3/m256 to ymm1, based on mask bits in the mask operand, ymm4.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	implicit XMM0	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	imm8[7:4]

Description

Conditionally copy each quadword data element of double-precision floating-point value from the second source operand and the first source operand depending on mask bits defined in the mask register operand. The mask bits are the most significant bit in each quadword element of the mask register.

Each quadword element of the destination operand is copied from:

- the corresponding quadword element in the second source operand, If a mask bit is “1"; or
- the corresponding quadword element in the first source operand, If a mask bit is “0"

The register assignment of the implicit mask operand for BLENDVPD is defined to be the architectural register XMM0.

128-bit Legacy SSE version: The first source operand and the destination operand is the same. Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged. The mask register operand is implicitly defined to be the architectural register XMM0. An attempt to execute BLENDVPD with a VEX prefix will cause #UD.

VEX.128 encoded version: The first source operand and the destination operand are XMM registers. The second source operand is an XMM register or 128-bit memory location. The mask operand is the third source register, and encoded in bits[7:4] of the immediate byte(imm8). The bits[3:0] of imm8 are ignored. In 32-bit mode, imm8[7] is ignored. The upper bits (VLMAX-1:128) of the corresponding YMM register (destination register) are zeroed. VEX.W must be 0, otherwise, the instruction will #UD.

VEX.256 encoded version: The first source operand and destination operand are YMM registers. The second source operand can be a YMM register or a 256-bit memory location. The mask operand is the third source register, and encoded in bits[7:4] of the immediate byte(imm8). The bits[3:0] of imm8 are ignored. In 32-bit mode, imm8[7] is ignored. VEX.W must be 0, otherwise, the instruction will #UD.

VBLENDVPD permits the mask to be any XMM or YMM register. In contrast, BLENDVPD treats XMM0 implicitly as the mask and do not support non-destructive destination operation.

Operation

BLENDVPD (128-bit Legacy SSE version)

```
MASK ← XMM0
IF (MASK[63] = 0) THEN DEST[63:0] ← DEST[63:0]
        ELSE DEST [63:0] ← SRC[63:0] FI
IF (MASK[127] = 0) THEN DEST[127:64] ← DEST[127:64]
        ELSE DEST [127:64] ← SRC[127:64] FI
DEST[VLMAX-1:128] (Unmodified)
```

VBLENDVPD (VEX.128 encoded version)

```
MASK ← SRC3
IF (MASK[63] = 0) THEN DEST[63:0] ← SRC1[63:0]
        ELSE DEST [63:0] ← SRC2[63:0] FI
IF (MASK[127] = 0) THEN DEST[127:64] ← SRC1[127:64]
        ELSE DEST [127:64] ← SRC2[127:64] FI
DEST[VLMAX-1:128] ← 0
```

VBLENDVPD (VEX.256 encoded version)

```
MASK ← SRC3
IF (MASK[63] = 0) THEN DEST[63:0] ← SRC1[63:0]
        ELSE DEST [63:0] ← SRC2[63:0] FI
IF (MASK[127] = 0) THEN DEST[127:64] ← SRC1[127:64]
        ELSE DEST [127:64] ← SRC2[127:64] FI
```

```
IF (MASK[191] = 0) THEN DEST[191:128] ← SRC1[191:128]
        ELSE DEST [191:128] ← SRC2[191:128] FI
IF (MASK[255] = 0) THEN DEST[255:192] ← SRC1[255:192]
        ELSE DEST [255:192] ← SRC2[255:192] FI
```

Intel C/C++ Compiler Intrinsic Equivalent

BLENDVPD __m128d _mm_blendv_pd(__m128d v1, __m128d v2, __m128d v3);

VBLENDVPD __m128 _mm_blendv_pd (__m128d a, __m128d b, __m128d mask);

VBLENDVPD __m256 _mm256_blendv_pd (__m256d a, __m256d b, __m256d mask);

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 4; additionally

#UD	If VEX.W = 1.

BLENDVPS — Variable Blend Packed Single Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 38 14 /r BLENDVPS *xmm1, xmm2/m128,* *<XMM0>*	A	V/V	SSE4_1	Select packed single precision floating-point values from *xmm1* and *xmm2/m128* from mask specified in *XMM0* and store the values into *xmm1*.
VEX.NDS.128.66.0F3A.W0 4A /r /is4 VBLENDVPS xmm1, xmm2, xmm3/m128, xmm4	B	V/V	AVX	Conditionally copy single-precision floating-point values from xmm2 or xmm3/m128 to xmm1, based on mask bits in the specified mask operand, xmm4.
VEX.NDS.256.66.0F3A.W0 4A /r /is4 VBLENDVPS ymm1, ymm2, ymm3/m256, ymm4	B	V/V	AVX	Conditionally copy single-precision floating-point values from ymm2 or ymm3/m256 to ymm1, based on mask bits in the specified mask register, ymm4.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	implicit XMM0	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	imm8[7:4]

Description

Conditionally copy each dword data element of single-precision floating-point value from the second source operand and the first source operand depending on mask bits defined in the mask register operand. The mask bits are the most significant bit in each dword element of the mask register.

Each quadword element of the destination operand is copied from:

- the corresponding dword element in the second source operand, If a mask bit is "1"; or
- the corresponding dword element in the first source operand, If a mask bit is "0"

The register assignment of the implicit mask operand for BLENDVPS is defined to be the architectural register XMM0.

128-bit Legacy SSE version: The first source operand and the destination operand is the same. Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged. The mask register operand is implicitly defined to be the architectural register XMM0. An attempt to execute BLENDVPS with a VEX prefix will cause #UD.

VEX.128 encoded version: The first source operand and the destination operand are XMM registers. The second source operand is an XMM register or 128-bit memory location. The mask operand is the third source register, and encoded in bits[7:4] of the immediate byte(imm8). The bits[3:0] of imm8 are ignored. In 32-bit mode, imm8[7] is ignored. The upper bits (VLMAX-1:128) of the corresponding YMM register (destination register) are zeroed. VEX.W must be 0, otherwise, the instruction will #UD.

VEX.256 encoded version: The first source operand and destination operand are YMM registers. The second source operand can be a YMM register or a 256-bit memory location. The mask operand is the third source register, and encoded in bits[7:4] of the immediate byte(imm8). The bits[3:0] of imm8 are ignored. In 32-bit mode, imm8[7] is ignored. VEX.W must be 0, otherwise, the instruction will #UD.

VBLENDVPS permits the mask to be any XMM or YMM register. In contrast, BLENDVPS treats XMM0 implicitly as the mask and do not support non-destructive destination operation.

Operation

BLENDVPS (128-bit Legacy SSE version)

```
MASK ← XMM0
IF (MASK[31] = 0) THEN DEST[31:0] ← DEST[31:0]
        ELSE DEST [31:0] ← SRC[31:0] FI
IF (MASK[63] = 0) THEN DEST[63:32] ← DEST[63:32]
        ELSE DEST [63:32] ← SRC[63:32] FI
IF (MASK[95] = 0) THEN DEST[95:64] ← DEST[95:64]
        ELSE DEST [95:64] ← SRC[95:64] FI
IF (MASK[127] = 0) THEN DEST[127:96] ← DEST[127:96]
        ELSE DEST [127:96] ← SRC[127:96] FI
DEST[VLMAX-1:128] (Unmodified)
```

VBLENDVPS (VEX.128 encoded version)

```
MASK ← SRC3
IF (MASK[31] = 0) THEN DEST[31:0] ← SRC1[31:0]
        ELSE DEST [31:0] ← SRC2[31:0] FI
IF (MASK[63] = 0) THEN DEST[63:32] ← SRC1[63:32]
        ELSE DEST [63:32] ← SRC2[63:32] FI
IF (MASK[95] = 0) THEN DEST[95:64] ← SRC1[95:64]
        ELSE DEST [95:64] ← SRC2[95:64] FI
```

```
IF (MASK[127] = 0) THEN DEST[127:96] ← SRC1[127:96]
        ELSE DEST [127:96] ← SRC2[127:96] FI
DEST[VLMAX-1:128] ← 0
```

VBLENDVPS (VEX.256 encoded version)

```
MASK ← SRC3
IF (MASK[31] = 0) THEN DEST[31:0] ← SRC1[31:0]
        ELSE DEST [31:0] ← SRC2[31:0] FI
IF (MASK[63] = 0) THEN DEST[63:32] ← SRC1[63:32]
        ELSE DEST [63:32] ← SRC2[63:32] FI
IF (MASK[95] = 0) THEN DEST[95:64] ← SRC1[95:64]
        ELSE DEST [95:64] ← SRC2[95:64] FI
IF (MASK[127] = 0) THEN DEST[127:96] ← SRC1[127:96]
        ELSE DEST [127:96] ← SRC2[127:96] FI
IF (MASK[159] = 0) THEN DEST[159:128] ← SRC1[159:128]
        ELSE DEST [159:128] ← SRC2[159:128] FI
IF (MASK[191] = 0) THEN DEST[191:160] ← SRC1[191:160]
        ELSE DEST [191:160] ← SRC2[191:160] FI
IF (MASK[223] = 0) THEN DEST[223:192] ← SRC1[223:192]
        ELSE DEST [223:192] ← SRC2[223:192] FI
IF (MASK[255] = 0) THEN DEST[255:224] ← SRC1[255:224]
        ELSE DEST [255:224] ← SRC2[255:224] FI
```

Intel C/C++ Compiler Intrinsic Equivalent

BLENDVPS __m128 _mm_blendv_ps(__m128 v1, __m128 v2, __m128 v3);

VBLENDVPS __m128 _mm_blendv_ps (__m128 a, __m128 b, __m128 mask);

VBLENDVPS __m256 _mm256_blendv_ps (__m256 a, __m256 b, __m256 mask);

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 4; additionally

#UD	If VEX.W = 1.

BOUND—Check Array Index Against Bounds

Opcode	Instruction	Op/En	64-bit Mode	Compat/Leg Mode	Description
62 /*r*	BOUND *r16, m16&16*	A	Invalid	Valid	Check if *r16* (array index) is within bounds specified by *m16&16.*
62 /*r*	BOUND *r32, m32&32*	A	Invalid	Valid	Check if *r32* (array index) is within bounds specified by *m16&16.*

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r)	ModRM:r/m (r)	NA	NA

Description

BOUND determines if the first operand (array index) is within the bounds of an array specified the second operand (bounds operand). The array index is a signed integer located in a register. The bounds operand is a memory location that contains a pair of signed doubleword-integers (when the operand-size attribute is 32) or a pair of signed word-integers (when the operand-size attribute is 16). The first doubleword (or word) is the lower bound of the array and the second doubleword (or word) is the upper bound of the array. The array index must be greater than or equal to the lower bound and less than or equal to the upper bound plus the operand size in bytes. If the index is not within bounds, a BOUND range exceeded exception (#BR) is signaled. When this exception is generated, the saved return instruction pointer points to the BOUND instruction.

The bounds limit data structure (two words or doublewords containing the lower and upper limits of the array) is usually placed just before the array itself, making the limits addressable via a constant offset from the beginning of the array. Because the address of the array already will be present in a register, this practice avoids extra bus cycles to obtain the effective address of the array bounds.

This instruction executes as described in compatibility mode and legacy mode. It is not valid in 64-bit mode.

Operation

```
IF 64bit Mode
    THEN
        #UD;
    ELSE
        IF (ArrayIndex < LowerBound OR ArrayIndex > UpperBound)
        (* Below lower bound or above upper bound *)
```

```
        THEN #BR; FI;
FI;
```

Flags Affected

None.

Protected Mode Exceptions

#BR	If the bounds test fails.
#UD	If second operand is not a memory location.
	If the LOCK prefix is used.
#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

Real-Address Mode Exceptions

#BR	If the bounds test fails.
#UD	If second operand is not a memory location.
	If the LOCK prefix is used.
#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.

Virtual-8086 Mode Exceptions

#BR	If the bounds test fails.
#UD	If second operand is not a memory location.
	If the LOCK prefix is used.
#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#UD	If in 64-bit mode.

VBROADCAST—Load with Broadcast

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
VEX.128.66.0F38.W0 18 /r VBROADCASTSS xmm1, m32	A	I/V	AVX	Broadcast single-precision floating-point element in mem to four locations in xmm1.
VEX.256.66.0F38.W0 18 /r VBROADCASTSS ymm1, m32	A	V/V	AVX	Broadcast single-precision floating-point element in mem to eight locations in ymm1.
VEX.256.66.0F38.W0 19 /r VBROADCASTSD ymm1, m64	A	V/V	AVX	Broadcast double-precision floating-point element in mem to four locations in ymm1.
VEX.256.66.0F38.W0 1A /r VBROADCASTF128 ymm1, m128	A	V/V	AVX	Broadcast 128 bits of floating-point data in mem to low and high 128-bits in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Load floating point values from the source operand (second operand) and broadcast to all elements of the destination operand (first operand).

The destination operand is a YMM register. The source operand is either a 32-bit, 64-bit, or 128-bit memory location. Register source encodings are reserved and will #UD.

VBROADCASTSD and VBROADCASTF128 are only supported as 256-bit wide versions. VBROADCASTSS is supported in both 128-bit and 256-bit wide versions.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

If VBROADCASTSD or VBROADCASTF128 is encoded with VEX.L= 0, an attempt to execute the instruction encoded with VEX.L= 0 will cause an #UD exception.

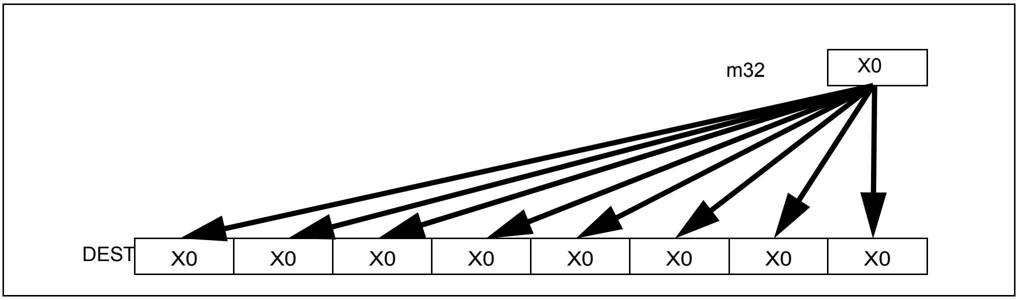


Figure 3-5. VBROADCASTSS Operation (VEX.256 encoded version)

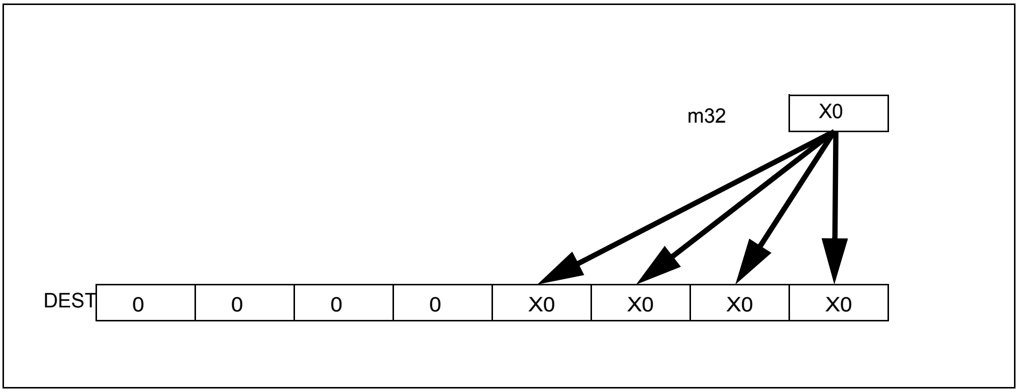


Figure 3-6. VBROADCASTSS Operation (128-bit version)

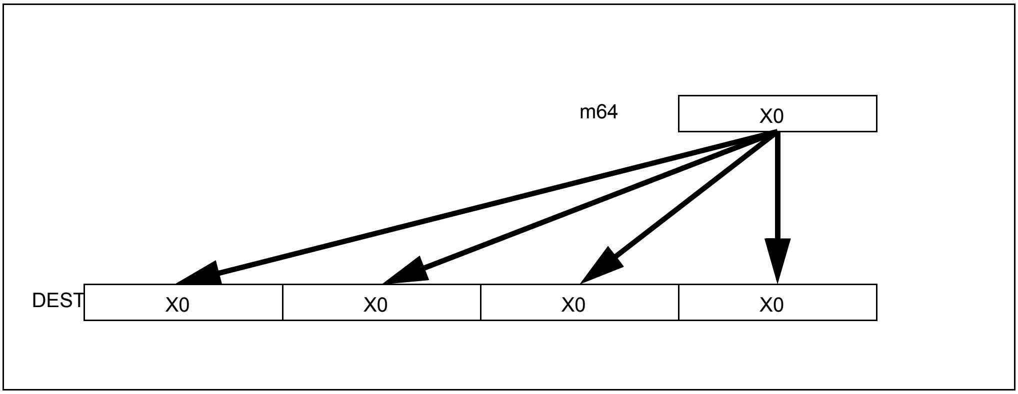


Figure 3-7. VBROADCASTSD Operation

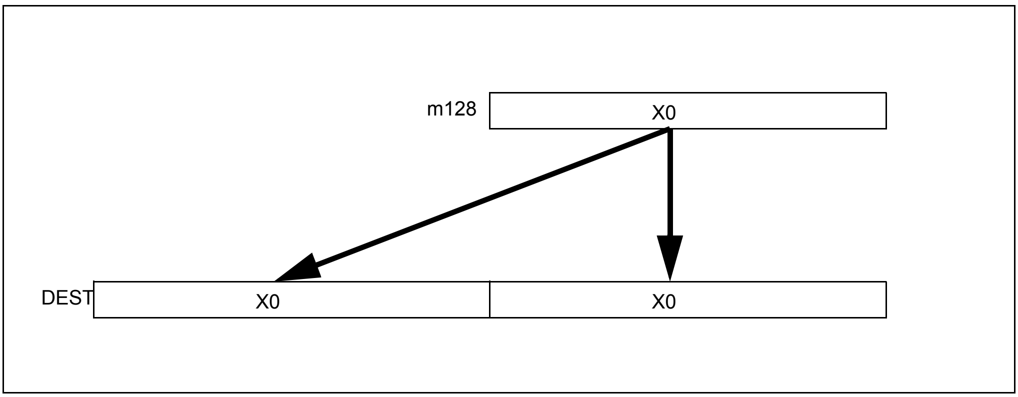


Figure 3-8. VBROADCASTF128 Operation

Operation

VBROADCASTSS (128 bit version)

temp ← SRC[31:0]
DEST[31:0] ← temp
DEST[63:32] ← temp
DEST[95:64] ← temp
DEST[127:96] ← temp
DEST[VLMAX-1:128] ← 0

VBROADCASTSS (VEX.256 encoded version)

```
temp ← SRC[31:0]
DEST[31:0] ← temp
DEST[63:32] ← temp
DEST[95:64] ← temp
DEST[127:96] ← temp
DEST[159:128] ← temp
DEST[191:160] ← temp
DEST[223:192] ← temp
DEST[255:224] ← temp
```

VBROADCASTSD (VEX.256 encoded version)

```
temp ← SRC[63:0]
DEST[63:0] ← temp
DEST[127:64] ← temp
DEST[191:128] ← temp
DEST[255:192] ← temp
```

VBROADCASTF128

```
temp ← SRC[127:0]
DEST[127:0] ← temp
DEST[VLMAX-1:128] ← temp
```

Intel C/C++ Compiler Intrinsic Equivalent

VBROADCASTSS __m128 _mm_broadcast_ss(float *a);

VBROADCASTSS __m256 _mm256_broadcast_ss(float *a);

VBROADCASTSD __m256d _mm256_broadcast_sd(double *a);

VBROADCASTF128 __m256 _mm256_broadcast_ps(__m128 * a);

VBROADCASTF128 __m256d _mm256_broadcast_pd(__m128d * a);

Flags Affected

None.

Other Exceptions

See Exceptions Type 6; additionally

#UD	If VEX.L = 0 for VBROADCASTSD
	If VEX.L = 0 for VBROADCASTF128
	If VEX.W = 1.

BSF—Bit Scan Forward

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
0F BC /r	BSF *r16, r/m16*	A	Valid	Valid	Bit scan forward on *r/m16.*
0F BC /r	BSF *r32, r/m32*	A	Valid	Valid	Bit scan forward on *r/m32.*
REX.W + 0F BC	BSF *r64, r/m64*	A	Valid	N.E.	Bit scan forward on *r/m64.*

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Searches the source operand (second operand) for the least significant set bit (1 bit). If a least significant 1 bit is found, its bit index is stored in the destination operand (first operand). The source operand can be a register or a memory location; the destination operand is a register. The bit index is an unsigned offset from bit 0 of the source operand. If the content of the source operand is 0, the content of the destination operand is undefined.

In 64-bit mode, the instruction's default operation size is 32 bits. Using a REX prefix in the form of REX.R permits access to additional registers (R8-R15). Using a REX prefix in the form of REX.W promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
IF SRC = 0
    THEN
        ZF ← 1;
        DEST is undefined;
    ELSE
        ZF ← 0;
        temp ← 0;
        WHILE Bit(SRC, temp) = 0
        DO
            temp ← temp + 1;
            DEST ← temp;
        OD;
FI;
```

Flags Affected

The ZF flag is set to 1 if all the source operand is 0; otherwise, the ZF flag is cleared. The CF, OF, SF, AF, and PF, flags are undefined.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.

#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

BSR—Bit Scan Reverse

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
0F BD /r	BSR *r16, r/m16*	A	Valid	Valid	Bit scan reverse on *r/m16.*
0F BD /r	BSR *r32, r/m32*	A	Valid	Valid	Bit scan reverse on *r/m32.*
REX.W + 0F BD	BSR *r64, r/m64*	A	Valid	N.E.	Bit scan reverse on *r/m64.*

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Searches the source operand (second operand) for the most significant set bit (1 bit). If a most significant 1 bit is found, its bit index is stored in the destination operand (first operand). The source operand can be a register or a memory location; the destination operand is a register. The bit index is an unsigned offset from bit 0 of the source operand. If the content source operand is 0, the content of the destination operand is undefined.

In 64-bit mode, the instruction's default operation size is 32 bits. Using a REX prefix in the form of REX.R permits access to additional registers (R8-R15). Using a REX prefix in the form of REX.W promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
IF SRC = 0
    THEN
        ZF ← 1;
        DEST is undefined;
    ELSE
        ZF ← 0;
        temp ← OperandSize – 1;
        WHILE Bit(SRC, temp) = 0
        DO
            temp ← temp - 1;
            DEST ← temp;
        OD;
FI;
```

Flags Affected

The ZF flag is set to 1 if all the source operand is 0; otherwise, the ZF flag is cleared. The CF, OF, SF, AF, and PF, flags are undefined.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.

#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

BSWAP—Byte Swap

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
0F C8+*rd*	BSWAP *r32*	A	Valid*	Valid	Reverses the byte order of a 32-bit register.
REX.W + 0F C8+*rd*	BSWAP *r64*	A	Valid	N.E.	Reverses the byte order of a 64-bit register.

NOTES:

* See IA-32 Architecture Compatibility section below.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	reg (r, w)	NA	NA	NA

Description

Reverses the byte order of a 32-bit or 64-bit (destination) register. This instruction is provided for converting little-endian values to big-endian format and vice versa. To swap bytes in a word value (16-bit register), use the XCHG instruction. When the BSWAP instruction references a 16-bit register, the result is undefined.

In 64-bit mode, the instruction’s default operation size is 32 bits. Using a REX prefix in the form of REX.R permits access to additional registers (R8-R15). Using a REX prefix in the form of REX.W promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

IA-32 Architecture Legacy Compatibility

The BSWAP instruction is not supported on IA-32 processors earlier than the Intel486™ processor family. For compatibility with this instruction, software should include functionally equivalent code for execution on Intel processors earlier than the Intel486 processor family.

Operation

```
TEMP ← DEST
IF 64-bit mode AND OperandSize = 64
    THEN
        DEST[7:0] ← TEMP[63:56];
        DEST[15:8] ← TEMP[55:48];
        DEST[23:16] ← TEMP[47:40];
        DEST[31:24] ← TEMP[39:32];
        DEST[39:32] ← TEMP[31:24];
```

```
        DEST[47:40] ← TEMP[23:16];
        DEST[55:48] ← TEMP[15:8];
        DEST[63:56] ← TEMP[7:0];
    ELSE
        DEST[7:0] ← TEMP[31:24];
        DEST[15:8] ← TEMP[23:16];
        DEST[23:16] ← TEMP[15:8];
        DEST[31:24] ← TEMP[7:0];
FI;
```

Flags Affected

None.

Exceptions (All Operating Modes)

#UD	If the LOCK prefix is used.

BT—Bit Test

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
0F A3	BT *r/m16, r16*	A	Valid	Valid	Store selected bit in CF flag.
0F A3	BT *r/m32, r32*	A	Valid	Valid	Store selected bit in CF flag.
REX.W + 0F A3	BT *r/m64, r64*	A	Valid	N.E.	Store selected bit in CF flag.
0F BA /4 *ib*	BT *r/m16, imm8*	B	Valid	Valid	Store selected bit in CF flag.
0F BA /4 *ib*	BT *r/m32, imm8*	B	Valid	Valid	Store selected bit in CF flag.
REX.W + 0F BA /4 *ib*	BT *r/m64, imm8*	B	Valid	N.E.	Store selected bit in CF flag.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r)	ModRM:reg (r)	NA	NA
B	ModRM:r/m (r)	imm8	NA	NA

Description

Selects the bit in a bit string (specified with the first operand, called the bit base) at the bit-position designated by the bit offset (specified by the second operand) and stores the value of the bit in the CF flag. The bit base operand can be a register or a memory location; the bit offset operand can be a register or an immediate value:

- If the bit base operand specifies a register, the instruction takes the modulo 16, 32, or 64 of the bit offset operand (modulo size depends on the mode and register size; 64-bit operands are available only in 64-bit mode).
- If the bit base operand specifies a memory location, the operand represents the address of the byte in memory that contains the bit base (bit 0 of the specified byte) of the bit string. The range of the bit position that can be referenced by the offset operand depends on the operand size.

See also: **Bit(BitBase, BitOffset)** on page 3-14.

Some assemblers support immediate bit offsets larger than 31 by using the immediate bit offset field in combination with the displacement field of the memory operand. In this case, the low-order 3 or 5 bits (3 for 16-bit operands, 5 for 32-bit operands) of the immediate bit offset are stored in the immediate bit offset field, and the high-order bits are shifted and combined with the byte displacement in the addressing mode by the assembler. The processor will ignore the high order bits if they are not zero.

When accessing a bit in memory, the processor may access 4 bytes starting from the memory address for a 32-bit operand size, using by the following relationship:

Effective Address + (4 ∗ (BitOffset DIV 32))

Or, it may access 2 bytes starting from the memory address for a 16-bit operand, using this relationship:

Effective Address + (2 ∗ (BitOffset DIV 16))

It may do so even when only a single byte needs to be accessed to reach the given bit. When using this bit addressing mechanism, software should avoid referencing areas of memory close to address space holes. In particular, it should avoid references to memory-mapped I/O registers. Instead, software should use the MOV instructions to load from or store to these addresses, and use the register form of these instructions to manipulate the data.

In 64-bit mode, the instruction's default operation size is 32 bits. Using a REX prefix in the form of REX.R permits access to additional registers (R8-R15). Using a REX prefix in the form of REX.W promotes operation to 64 bit operands. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
CF ← Bit(BitBase, BitOffset);
```

Flags Affected

The CF flag contains the value of the selected bit. The ZF flag is unaffected. The OF, SF, AF, and PF flags are undefined.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

BTC—Bit Test and Complement

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
0F BB	BTC *r/m16, r16*	A	Valid	Valid	Store selected bit in CF flag and complement.
0F BB	BTC *r/m32, r32*	A	Valid	Valid	Store selected bit in CF flag and complement.
REX.W + 0F BB	BTC *r/m64, r64*	A	Valid	N.E.	Store selected bit in CF flag and complement.
0F BA /7 *ib*	BTC *r/m16, imm8*	B	Valid	Valid	Store selected bit in CF flag and complement.
0F BA /7 *ib*	BTC *r/m32, imm8*	B	Valid	Valid	Store selected bit in CF flag and complement.
REX.W + 0F BA /7 *ib*	BTC *r/m64, imm8*	B	Valid	N.E.	Store selected bit in CF flag and complement.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r, w)	ModRM:reg (r)	NA	NA
B	ModRM:r/m (r, w)	imm8	NA	NA

Description

Selects the bit in a bit string (specified with the first operand, called the bit base) at the bit-position designated by the bit offset operand (second operand), stores the value of the bit in the CF flag, and complements the selected bit in the bit string. The bit base operand can be a register or a memory location; the bit offset operand can be a register or an immediate value:

- If the bit base operand specifies a register, the instruction takes the modulo 16, 32, or 64 of the bit offset operand (modulo size depends on the mode and register size; 64-bit operands are available only in 64-bit mode). This allows any bit position to be selected.
- If the bit base operand specifies a memory location, the operand represents the address of the byte in memory that contains the bit base (bit 0 of the specified byte) of the bit string. The range of the bit position that can be referenced by the offset operand depends on the operand size.

See also: **Bit(BitBase, BitOffset)** on page 3-14.

Some assemblers support immediate bit offsets larger than 31 by using the immediate bit offset field in combination with the displacement field of the memory operand. See “BT—Bit Test” in this chapter for more information on this addressing mechanism.

This instruction can be used with a LOCK prefix to allow the instruction to be executed atomically.

In 64-bit mode, the instruction's default operation size is 32 bits. Using a REX prefix in the form of REX.R permits access to additional registers (R8-R15). Using a REX prefix in the form of REX.W promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
CF ← Bit(BitBase, BitOffset);
Bit(BitBase, BitOffset) ← NOT Bit(BitBase, BitOffset);
```

Flags Affected

The CF flag contains the value of the selected bit before it is complemented. The ZF flag is unaffected. The OF, SF, AF, and PF flags are undefined.

Protected Mode Exceptions

#GP(0)	If the destination operand points to a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.

#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

BTR—Bit Test and Reset

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
0F B3	BTR *r/m16, r16*	A	Valid	Valid	Store selected bit in CF flag and clear.
0F B3	BTR *r/m32, r32*	A	Valid	Valid	Store selected bit in CF flag and clear.
REX.W + 0F B3	BTR *r/m64, r64*	A	Valid	N.E.	Store selected bit in CF flag and clear.
0F BA /6 *ib*	BTR *r/m16, imm8*	B	Valid	Valid	Store selected bit in CF flag and clear.
0F BA /6 *ib*	BTR *r/m32, imm8*	B	Valid	Valid	Store selected bit in CF flag and clear.
REX.W + 0F BA /6 *ib*	BTR *r/m64, imm8*	B	Valid	N.E.	Store selected bit in CF flag and clear.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r, w)	ModRM:reg (r)	NA	NA
B	ModRM:r/m (r, w)	imm8	NA	NA

Description

Selects the bit in a bit string (specified with the first operand, called the bit base) at the bit-position designated by the bit offset operand (second operand), stores the value of the bit in the CF flag, and clears the selected bit in the bit string to 0. The bit base operand can be a register or a memory location; the bit offset operand can be a register or an immediate value:

- If the bit base operand specifies a register, the instruction takes the modulo 16, 32, or 64 of the bit offset operand (modulo size depends on the mode and register size; 64-bit operands are available only in 64-bit mode). This allows any bit position to be selected.
- If the bit base operand specifies a memory location, the operand represents the address of the byte in memory that contains the bit base (bit 0 of the specified byte) of the bit string. The range of the bit position that can be referenced by the offset operand depends on the operand size.

See also: **Bit(BitBase, BitOffset)** on page 3-14.

Some assemblers support immediate bit offsets larger than 31 by using the immediate bit offset field in combination with the displacement field of the memory operand. See “BT—Bit Test” in this chapter for more information on this addressing mechanism.

This instruction can be used with a LOCK prefix to allow the instruction to be executed atomically.

In 64-bit mode, the instruction's default operation size is 32 bits. Using a REX prefix in the form of REX.R permits access to additional registers (R8-R15). Using a REX prefix in the form of REX.W promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
CF ← Bit(BitBase, BitOffset);
Bit(BitBase, BitOffset) ← 0;
```

Flags Affected

The CF flag contains the value of the selected bit before it is cleared. The ZF flag is unaffected. The OF, SF, AF, and PF flags are undefined.

Protected Mode Exceptions

#GP(0)	If the destination operand points to a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.

#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

BTS—Bit Test and Set

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
0F AB	BTS *r/m16, r16*	A	Valid	Valid	Store selected bit in CF flag and set.
0F AB	BTS *r/m32, r32*	A	Valid	Valid	Store selected bit in CF flag and set.
REX.W + 0F AB	BTS *r/m64, r64*	A	Valid	N.E.	Store selected bit in CF flag and set.
0F BA /5 *ib*	BTS *r/m16, imm8*	B	Valid	Valid	Store selected bit in CF flag and set.
0F BA /5 *ib*	BTS *r/m32, imm8*	B	Valid	Valid	Store selected bit in CF flag and set.
REX.W + 0F BA /5 *ib*	BTS *r/m64, imm8*	B	Valid	N.E.	Store selected bit in CF flag and set.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r, w)	ModRM:reg (r)	NA	NA
B	ModRM:r/m (r, w)	imm8	NA	NA

Description

Selects the bit in a bit string (specified with the first operand, called the bit base) at the bit-position designated by the bit offset operand (second operand), stores the value of the bit in the CF flag, and sets the selected bit in the bit string to 1. The bit base operand can be a register or a memory location; the bit offset operand can be a register or an immediate value:

- If the bit base operand specifies a register, the instruction takes the modulo 16, 32, or 64 of the bit offset operand (modulo size depends on the mode and register size; 64-bit operands are available only in 64-bit mode). This allows any bit position to be selected.
- If the bit base operand specifies a memory location, the operand represents the address of the byte in memory that contains the bit base (bit 0 of the specified byte) of the bit string. The range of the bit position that can be referenced by the offset operand depends on the operand size.

See also: **Bit(BitBase, BitOffset)** on page 3-14.

Some assemblers support immediate bit offsets larger than 31 by using the immediate bit offset field in combination with the displacement field of the memory operand. See “BT—Bit Test” in this chapter for more information on this addressing mechanism.

This instruction can be used with a LOCK prefix to allow the instruction to be executed atomically.

In 64-bit mode, the instruction's default operation size is 32 bits. Using a REX prefix in the form of REX.R permits access to additional registers (R8-R15). Using a REX prefix in the form of REX.W promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
CF ← Bit(BitBase, BitOffset);
Bit(BitBase, BitOffset) ← 1;
```

Flags Affected

The CF flag contains the value of the selected bit before it is set. The ZF flag is unaffected. The OF, SF, AF, and PF flags are undefined.

Protected Mode Exceptions

#GP(0)	If the destination operand points to a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Virtual-8086 Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.

#SS	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

CALL—Call Procedure

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
E8 *cw*	CALL *rel16*	B	N.S.	Valid	Call near, relative, displacement relative to next instruction.
E8 *cd*	CALL *rel32*	B	Valid	Valid	Call near, relative, displacement relative to next instruction. 32-bit displacement sign extended to 64-bits in 64-bit mode.
FF /2	CALL *r/m16*	B	N.E.	Valid	Call near, absolute indirect, address given in *r/m16.*
FF /2	CALL *r/m32*	B	N.E.	Valid	Call near, absolute indirect, address given in *r/m32.*
FF /2	CALL *r/m64*	B	Valid	N.E.	Call near, absolute indirect, address given in *r/m*64.
9A *cd*	CALL *ptr16:16*	A	Invalid	Valid	Call far, absolute, address given in operand.
9A *cp*	CALL *ptr16:32*	A	Invalid	Valid	Call far, absolute, address given in operand.
FF /3	CALL *m16:16*	B	Valid	Valid	Call far, absolute indirect address given in *m16:16.* In 32-bit mode: if selector points to a gate, then RIP = 32-bit zero extended displacement taken from gate; else RIP = zero extended 16-bit offset from far pointer referenced in the instruction.
FF /3	CALL *m16:32*	B	Valid	Valid	In 64-bit mode: If selector points to a gate, then RIP = 64-bit displacement taken from gate; else RIP = zero extended 32-bit offset from far pointer referenced in the instruction.

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
REX.W + FF /3	CALL *m16:64*	B	Valid	N.E.	In 64-bit mode: If selector points to a gate, then RIP = 64-bit displacement taken from gate; else RIP = 64-bit offset from far pointer referenced in the instruction.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	Offset	NA	NA	NA
B	ModRM:r/m (r)	NA	NA	NA

Description

Saves procedure linking information on the stack and branches to the called procedure specified using the target operand. The target operand specifies the address of the first instruction in the called procedure. The operand can be an immediate value, a general-purpose register, or a memory location.

This instruction can be used to execute four types of calls:

- **Near Call** — A call to a procedure in the current code segment (the segment currently pointed to by the CS register), sometimes referred to as an intra-segment call.
- **Far Call** — A call to a procedure located in a different segment than the current code segment, sometimes referred to as an inter-segment call.
- **Inter-privilege-level far call** — A far call to a procedure in a segment at a different privilege level than that of the currently executing program or procedure.
- **Task switch** — A call to a procedure located in a different task.

The latter two call types (inter-privilege-level call and task switch) can only be executed in protected mode. See “Calling Procedures Using Call and RET” in Chapter 6 of the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*, for additional information on near, far, and inter-privilege-level calls. See Chapter 7, “Task Management,” in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 3A*, for information on performing task switches with the CALL instruction.

Near Call. When executing a near call, the processor pushes the value of the EIP register (which contains the offset of the instruction following the CALL instruction) on the stack (for use later as a return-instruction pointer). The processor then branches to the address in the current code segment specified by the target operand. The target operand specifies either an absolute offset in the code segment (an offset from the base of the code segment) or a relative offset (a signed displacement relative to the current value of the instruction pointer in the EIP register; this value points to the instruction following the CALL instruction). The CS register is not changed on near calls.

For a near call absolute, an absolute offset is specified indirectly in a general-purpose register or a memory location (*r/m16*, *r/m32, or r/m64*). The operand-size attribute determines the size of the target operand (16, 32 or 64 bits). When in 64-bit mode, the operand size for near call (and all near branches) is forced to 64-bits. Absolute offsets are loaded directly into the EIP(RIP) register. If the operand size attribute is 16, the upper two bytes of the EIP register are cleared, resulting in a maximum instruction pointer size of 16 bits. When accessing an absolute offset indirectly using the stack pointer [ESP] as the base register, the base value used is the value of the ESP before the instruction executes.

A relative offset (*rel16* or *rel32*) is generally specified as a label in assembly code. But at the machine code level, it is encoded as a signed, 16- or 32-bit immediate value. This value is added to the value in the EIP(RIP) register. In 64-bit mode the relative offset is always a 32-bit immediate value which is sign extended to 64-bits before it is added to the value in the RIP register for the target calculation. As with absolute offsets, the operand-size attribute determines the size of the target operand (16, 32, or 64 bits). In 64-bit mode the target operand will always be 64-bits because the operand size is forced to 64-bits for near branches.

Far Calls in Real-Address or Virtual-8086 Mode. When executing a far call in real-address or virtual-8086 mode, the processor pushes the current value of both the CS and EIP registers on the stack for use as a return-instruction pointer. The processor then performs a “far branch” to the code segment and offset specified with the target operand for the called procedure. The target operand specifies an absolute far address either directly with a pointer (*ptr16:16* or *ptr16:32*) or indirectly with a memory location (*m16:16* or *m16:32*). With the pointer method, the segment and offset of the called procedure is encoded in the instruction using a 4-byte (16-bit operand size) or 6-byte (32-bit operand size) far address immediate. With the indirect method, the target operand specifies a memory location that contains a 4-byte (16-bit operand size) or 6-byte (32-bit operand size) far address. The operand-size attribute determines the size of the offset (16 or 32 bits) in the far address. The far address is loaded directly into the CS and EIP registers. If the operand-size attribute is 16, the upper two bytes of the EIP register are cleared.

Far Calls in Protected Mode. When the processor is operating in protected mode, the CALL instruction can be used to perform the following types of far calls:

- Far call to the same privilege level
- Far call to a different privilege level (inter-privilege level call)
- Task switch (far call to another task)

In protected mode, the processor always uses the segment selector part of the far address to access the corresponding descriptor in the GDT or LDT. The descriptor type (code segment, call gate, task gate, or TSS) and access rights determine the type of call operation to be performed.

If the selected descriptor is for a code segment, a far call to a code segment at the same privilege level is performed. (If the selected code segment is at a different privilege level and the code segment is non-conforming, a general-protection exception is generated.) A far call to the same privilege level in protected mode is very similar to one carried out in real-address or virtual-8086 mode. The target operand specifies an absolute far address either directly with a pointer (*ptr16:16* or *ptr16:32*) or indirectly with a memory location (*m16:16* or *m16:32*). The operand- size attribute determines the size of the offset (16 or 32 bits) in the far address. The new code segment selector and its descriptor are loaded into CS register; the offset from the instruction is loaded into the EIP register.

A call gate (described in the next paragraph) can also be used to perform a far call to a code segment at the same privilege level. Using this mechanism provides an extra level of indirection and is the preferred method of making calls between 16-bit and 32-bit code segments.

When executing an inter-privilege-level far call, the code segment for the procedure being called must be accessed through a call gate. The segment selector specified by the target operand identifies the call gate. The target operand can specify the call gate segment selector either directly with a pointer (*ptr16:16* or *ptr16:32*) or indirectly with a memory location (*m16:16* or *m16:32*). The processor obtains the segment selector for the new code segment and the new instruction pointer (offset) from the call gate descriptor. (The offset from the target operand is ignored when a call gate is used.)

On inter-privilege-level calls, the processor switches to the stack for the privilege level of the called procedure. The segment selector for the new stack segment is specified in the TSS for the currently running task. The branch to the new code segment occurs after the stack switch. (Note that when using a call gate to perform a far call to a segment at the same privilege level, no stack switch occurs.) On the new stack, the processor pushes the segment selector and stack pointer for the calling procedure's stack, an optional set of parameters from the calling procedures stack, and the segment selector and instruction pointer for the calling procedure's code segment. (A value in the call gate descriptor determines how many parameters to copy to the new stack.) Finally, the processor branches to the address of the procedure being called within the new code segment.

Executing a task switch with the CALL instruction is similar to executing a call through a call gate. The target operand specifies the segment selector of the task gate for the new task activated by the switch (the offset in the target operand is ignored). The task gate in turn points to the TSS for the new task, which contains the segment selectors for the task's code and stack segments. Note that the TSS also contains the EIP value for the next instruction that was to be executed before the calling task was suspended. This instruction pointer value is loaded into the EIP register to re-start the calling task.

The CALL instruction can also specify the segment selector of the TSS directly, which eliminates the indirection of the task gate. See Chapter 7, "Task Management," in the

Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A, for information on the mechanics of a task switch.

When you execute at task switch with a CALL instruction, the nested task flag (NT) is set in the EFLAGS register and the new TSS's previous task link field is loaded with the old task's TSS selector. Code is expected to suspend this nested task by executing an IRET instruction which, because the NT flag is set, automatically uses the previous task link to return to the calling task. (See "Task Linking" in Chapter 7 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*, for information on nested tasks.) Switching tasks with the CALL instruction differs in this regard from JMP instruction. JMP does not set the NT flag and therefore does not expect an IRET instruction to suspend the task.

Mixing 16-Bit and 32-Bit Calls. When making far calls between 16-bit and 32-bit code segments, use a call gate. If the far call is from a 32-bit code segment to a 16-bit code segment, the call should be made from the first 64 KBytes of the 32-bit code segment. This is because the operand-size attribute of the instruction is set to 16, so only a 16-bit return address offset can be saved. Also, the call should be made using a 16-bit call gate so that 16-bit values can be pushed on the stack. See Chapter 18, "Mixing 16-Bit and 32-Bit Code," in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*, for more information.

Far Calls in Compatibility Mode. When the processor is operating in compatibility mode, the CALL instruction can be used to perform the following types of far calls:

- Far call to the same privilege level, remaining in compatibility mode
- Far call to the same privilege level, transitioning to 64-bit mode
- Far call to a different privilege level (inter-privilege level call), transitioning to 64-bit mode

Note that a CALL instruction can not be used to cause a task switch in compatibility mode since task switches are not supported in IA-32e mode.

In compatibility mode, the processor always uses the segment selector part of the far address to access the corresponding descriptor in the GDT or LDT. The descriptor type (code segment, call gate) and access rights determine the type of call operation to be performed.

If the selected descriptor is for a code segment, a far call to a code segment at the same privilege level is performed. (If the selected code segment is at a different privilege level and the code segment is non-conforming, a general-protection exception is generated.) A far call to the same privilege level in compatibility mode is very similar to one carried out in protected mode. The target operand specifies an absolute far address either directly with a pointer (*ptr16:16* or *ptr16:32*) or indirectly with a memory location (*m16:16* or *m16:32*). The operand-size attribute determines the size of the offset (16 or 32 bits) in the far address. The new code segment selector and its descriptor are loaded into CS register and the offset from the instruction is loaded into the EIP register. The difference is that 64-bit mode may be entered. This specified by the L bit in the new code segment descriptor.

Note that a 64-bit call gate (described in the next paragraph) can also be used to perform a far call to a code segment at the same privilege level. However, using this mechanism requires that the target code segment descriptor have the L bit set, causing an entry to 64-bit mode.

When executing an inter-privilege-level far call, the code segment for the procedure being called must be accessed through a 64-bit call gate. The segment selector specified by the target operand identifies the call gate. The target operand can specify the call gate segment selector either directly with a pointer (*ptr16:16* or *ptr16:32*) or indirectly with a memory location (*m16:16* or *m16:32*). The processor obtains the segment selector for the new code segment and the new instruction pointer (offset) from the 16-byte call gate descriptor. (The offset from the target operand is ignored when a call gate is used.)

On inter-privilege-level calls, the processor switches to the stack for the privilege level of the called procedure. The segment selector for the new stack segment is set to NULL. The new stack pointer is specified in the TSS for the currently running task. The branch to the new code segment occurs after the stack switch. (Note that when using a call gate to perform a far call to a segment at the same privilege level, an implicit stack switch occurs as a result of entering 64-bit mode. The SS selector is unchanged, but stack segment accesses use a segment base of 0x0, the limit is ignored, and the default stack size is 64-bits. The full value of RSP is used for the offset, of which the upper 32-bits are undefined.) On the new stack, the processor pushes the segment selector and stack pointer for the calling procedure's stack and the segment selector and instruction pointer for the calling procedure's code segment. (Parameter copy is not supported in IA-32e mode.) Finally, the processor branches to the address of the procedure being called within the new code segment.

Near/(Far) Calls in 64-bit Mode. When the processor is operating in 64-bit mode, the CALL instruction can be used to perform the following types of far calls:

- Far call to the same privilege level, transitioning to compatibility mode
- Far call to the same privilege level, remaining in 64-bit mode
- Far call to a different privilege level (inter-privilege level call), remaining in 64-bit mode

Note that in this mode the CALL instruction can not be used to cause a task switch in 64-bit mode since task switches are not supported in IA-32e mode.

In 64-bit mode, the processor always uses the segment selector part of the far address to access the corresponding descriptor in the GDT or LDT. The descriptor type (code segment, call gate) and access rights determine the type of call operation to be performed.

If the selected descriptor is for a code segment, a far call to a code segment at the same privilege level is performed. (If the selected code segment is at a different privilege level and the code segment is non-conforming, a general-protection exception is generated.) A far call to the same privilege level in 64-bit mode is very similar to one carried out in compatibility mode. The target operand specifies an absolute far address indirectly with a memory location (*m16:16, m16:32* or *m16:64*). The form of CALL with a direct specification of absolute far address is not defined in 64-bit mode. The operand-size attribute determines the size of the offset (16, 32, or 64 bits) in the far address. The new code segment selector and its descriptor are loaded into the CS register; the offset from the instruction is loaded into the EIP register. The new code segment may specify entry either into compatibility or 64-bit mode, based on the L bit value.

A 64-bit call gate (described in the next paragraph) can also be used to perform a far call to a code segment at the same privilege level. However, using this mechanism requires that the target code segment descriptor have the L bit set.

When executing an inter-privilege-level far call, the code segment for the procedure being called must be accessed through a 64-bit call gate. The segment selector specified by the target operand identifies the call gate. The target operand can only specify the call gate segment selector indirectly with a memory location (*m16:16, m16:32* or *m16:64*). The processor obtains the segment selector for the new code segment and the new instruction pointer (offset) from the 16-byte call gate descriptor. (The offset from the target operand is ignored when a call gate is used.)

On inter-privilege-level calls, the processor switches to the stack for the privilege level of the called procedure. The segment selector for the new stack segment is set to NULL. The new stack pointer is specified in the TSS for the currently running task. The branch to the new code segment occurs after the stack switch.

Note that when using a call gate to perform a far call to a segment at the same privilege level, an implicit stack switch occurs as a result of entering 64-bit mode. The SS selector is unchanged, but stack segment accesses use a segment base of 0x0, the limit is ignored, and the default stack size is 64-bits. (The full value of RSP is used for the offset.) On the new stack, the processor pushes the segment selector and stack pointer for the calling procedure's stack and the segment selector and instruction pointer for the calling procedure's code segment. (Parameter copy is not supported in IA-32e mode.) Finally, the processor branches to the address of the procedure being called within the new code segment.

Operation

```
IF near call
   THEN IF near relative call
        THEN
             IF OperandSize = 64
                  THEN
                       tempDEST ← SignExtend(DEST); (* DEST is rel32 *)
                       tempRIP ← RIP + tempDEST;
                       IF stack not large enough for a 8-byte return address
                            THEN #SS(0); FI;
                       Push(RIP);
                       RIP ← tempRIP;
             FI;
             IF OperandSize = 32
```

```
            THEN
                tempEIP ← EIP + DEST; (* DEST is rel32 *)
                IF tempEIP is not within code segment limit THEN #GP(0); FI;
                IF stack not large enough for a 4-byte return address
                    THEN #SS(0); FI;
                Push(EIP);
                EIP ← tempEIP;
        FI;
        IF OperandSize = 16
            THEN
                tempEIP ← (EIP + DEST) AND 0000FFFFH; (* DEST is rel16 *)
                IF tempEIP is not within code segment limit THEN #GP(0); FI;
                IF stack not large enough for a 2-byte return address
                    THEN #SS(0); FI;
                Push(IP);
                EIP ← tempEIP;
        FI;
    ELSE (* Near absolute call *)
        IF OperandSize = 64
            THEN
                tempRIP ← DEST; (* DEST is r/m64 *)
                IF stack not large enough for a 8-byte return address
                    THEN #SS(0); FI;
                Push(RIP);
                RIP ← tempRIP;
        FI;
        IF OperandSize = 32
            THEN
                tempEIP ← DEST; (* DEST is r/m32 *)
                IF tempEIP is not within code segment limit THEN #GP(0); FI;
                IF stack not large enough for a 4-byte return address
                    THEN #SS(0); FI;
                Push(EIP);
                EIP ← tempEIP;
        FI;
        IF OperandSize = 16
            THEN
                tempEIP ← DEST AND 0000FFFFH; (* DEST is r/m16 *)
                IF tempEIP is not within code segment limit THEN #GP(0); FI;
                IF stack not large enough for a 2-byte return address
                    THEN #SS(0); FI;
                Push(IP);
                EIP ← tempEIP;
```

```
            FI;
    FI;rel/abs
FI; near

IF far call and (PE = 0 or (PE = 1 and VM = 1)) (* Real-address or virtual-8086 mode *)
    THEN
        IF OperandSize = 32
            THEN
                IF stack not large enough for a 6-byte return address
                    THEN #SS(0); FI;
                IF DEST[31:16] is not zero THEN #GP(0); FI;
                Push(CS); (* Padded with 16 high-order bits *)
                Push(EIP);
                CS ← DEST[47:32]; (* DEST is ptr16:32 or [m16:32] *)
                EIP ← DEST[31:0]; (* DEST is ptr16:32 or [m16:32] *)
            ELSE (* OperandSize = 16 *)
                IF stack not large enough for a 4-byte return address
                    THEN #SS(0); FI;
                Push(CS);
                Push(IP);
                CS ← DEST[31:16]; (* DEST is ptr16:16 or [m16:16] *)
                EIP ← DEST[15:0]; (* DEST is ptr16:16 or [m16:16]; clear upper 16 bits *)
        FI;
FI;

IF far call and (PE = 1 and VM = 0) (* Protected mode or IA-32e Mode, not virtual-8086 mode*)
    THEN
        IF segment selector in target operand NULL
            THEN #GP(0); FI;
        IF segment selector index not within descriptor table limits
            THEN #GP(new code segment selector); FI;
        Read type and access rights of selected segment descriptor;
        IF IA32_EFER.LMA = 0
            THEN
                IF segment type is not a conforming or nonconforming code segment, call
                gate, task gate, or TSS
                    THEN #GP(segment selector); FI;
            ELSE
                IF segment type is not a conforming or nonconforming code segment or
                64-bit call gate,
                    THEN #GP(segment selector); FI;
        FI;
        Depending on type and access rights:
```

```
                GO TO CONFORMING-CODE-SEGMENT;
                GO TO NONCONFORMING-CODE-SEGMENT;
                GO TO CALL-GATE;
                GO TO TASK-GATE;
                GO TO TASK-STATE-SEGMENT;
FI;

CONFORMING-CODE-SEGMENT:
    IF L bit = 1 and D bit = 1 and IA32_EFER.LMA = 1
        THEN GP(new code segment selector); FI;
    IF DPL > CPL
        THEN #GP(new code segment selector); FI;
    IF segment not present
        THEN #NP(new code segment selector); FI;
    IF stack not large enough for return address
        THEN #SS(0); FI;
    tempEIP ← DEST(Offset);
    IF OperandSize = 16
        THEN
            tempEIP ← tempEIP AND 0000FFFFH; FI; (* Clear upper 16 bits *)
    IF (EFER.LMA = 0 or target mode = Compatibility mode) and (tempEIP outside new code
    segment limit)
        THEN #GP(0); FI;
    IF tempEIP is non-canonical
        THEN #GP(0); FI;
    IF OperandSize = 32
        THEN
            Push(CS); (* Padded with 16 high-order bits *)
            Push(EIP);
            CS ← DEST(CodeSegmentSelector);
            (* Segment descriptor information also loaded *)
            CS(RPL) ← CPL;
            EIP ← tempEIP;
        ELSE
            IF OperandSize = 16
                THEN
                    Push(CS);
                    Push(IP);
                    CS ← DEST(CodeSegmentSelector);
                    (* Segment descriptor information also loaded *)
                    CS(RPL) ← CPL;
                    EIP ← tempEIP;
                ELSE (* OperandSize = 64 *)
```

```
                    Push(CS); (* Padded with 48 high-order bits *)
                    Push(RIP);
                    CS ← DEST(CodeSegmentSelector);
                    (* Segment descriptor information also loaded *)
                    CS(RPL) ← CPL;
                    RIP ← tempEIP;
            FI;
    FI;
END;

NONCONFORMING-CODE-SEGMENT:
    IF L-Bit = 1 and D-BIT = 1 and IA32_EFER.LMA = 1
        THEN GP(new code segment selector); FI;
    IF (RPL > CPL) or (DPL ≠ CPL)
        THEN #GP(new code segment selector); FI;
    IF segment not present
        THEN #NP(new code segment selector); FI;
    IF stack not large enough for return address
        THEN #SS(0); FI;
    tempEIP ← DEST(Offset);
    IF OperandSize = 16
        THEN tempEIP ← tempEIP AND 0000FFFFH; FI; (* Clear upper 16 bits *)
    IF (EFER.LMA = 0 or target mode = Compatibility mode) and (tempEIP outside new code
    segment limit)
        THEN #GP(0); FI;
    IF tempEIP is non-canonical
        THEN #GP(0); FI;
    IF OperandSize = 32
        THEN
            Push(CS); (* Padded with 16 high-order bits *)
            Push(EIP);
            CS ← DEST(CodeSegmentSelector);
            (* Segment descriptor information also loaded *)
            CS(RPL) ← CPL;
            EIP ← tempEIP;
        ELSE
            IF OperandSize = 16
                THEN
                    Push(CS);
                    Push(IP);
                    CS ← DEST(CodeSegmentSelector);
                    (* Segment descriptor information also loaded *)
                    CS(RPL) ← CPL;
                    EIP ← tempEIP;
```

```
                ELSE (* OperandSize = 64 *)
                    Push(CS); (* Padded with 48 high-order bits *)
                    Push(RIP);
                    CS ← DEST(CodeSegmentSelector);
                    (* Segment descriptor information also loaded *)
                    CS(RPL) ← CPL;
                    RIP ← tempEIP;
            FI;
    FI;
END;

CALL-GATE:
    IF call gate (DPL < CPL) or (RPL > DPL)
        THEN #GP(call-gate selector); FI;
    IF call gate not present
        THEN #NP(call-gate selector); FI;
    IF call-gate code-segment selector is NULL
        THEN #GP(0); FI;
    IF call-gate code-segment selector index is outside descriptor table limits
        THEN #GP(call-gate code-segment selector); FI;
    Read call-gate code-segment descriptor;
    IF call-gate code-segment descriptor does not indicate a code segment
    or call-gate code-segment descriptor DPL > CPL
        THEN #GP(call-gate code-segment selector); FI;
    IF IA32_EFER.LMA = 1 AND (call-gate code-segment descriptor is
    not a 64-bit code segment or call-gate code-segment descriptor has both L-bit and D-bit set)
        THEN #GP(call-gate code-segment selector); FI;
    IF call-gate code segment not present
        THEN #NP(call-gate code-segment selector); FI;
    IF call-gate code segment is non-conforming and DPL < CPL
        THEN go to MORE-PRIVILEGE;
        ELSE go to SAME-PRIVILEGE;
    FI;
END;

MORE-PRIVILEGE:
    IF current TSS is 32-bit
        THEN
            TSSstackAddress ← (new code-segment DPL ∗ 8) + 4;
            IF (TSSstackAddress + 5) > current TSS limit
                THEN #TS(current TSS selector); FI;
            NewSS ← 2 bytes loaded from (TSS base + TSSstackAddress + 4);
            NewESP ← 4 bytes loaded from (TSS base + TSSstackAddress);
        ELSE
```

```
            IF current TSS is 16-bit
                THEN
                    TSSstackAddress ← (new code-segment DPL ∗ 4) + 2
                    IF (TSSstackAddress + 3) > current TSS limit
                        THEN #TS(current TSS selector); FI;
                    NewSS ← 2 bytes loaded from (TSS base + TSSstackAddress + 2);
                    NewESP ← 2 bytes loaded from (TSS base + TSSstackAddress);
                ELSE (* current TSS is 64-bit *)
                    TSSstackAddress ← (new code-segment DPL ∗ 8) + 4;
                    IF (TSSstackAddress + 7) > current TSS limit
                        THEN #TS(current TSS selector); FI;
                    NewSS ← new code-segment DPL; (* NULL selector with RPL = new CPL *)
                    NewRSP ← 8 bytes loaded from (current TSS base + TSSstackAddress);
            FI;
    FI;
    IF IA32_EFER.LMA = 0 and NewSS is NULL
        THEN #TS(NewSS); FI;
    Read new code-segment descriptor and new stack-segment descriptor;
    IF IA32_EFER.LMA = 0 and (NewSS RPL ≠ new code-segment DPL
    or new stack-segment DPL ≠ new code-segment DPL or new stack segment is not a
    writable data segment)
        THEN #TS(NewSS); FI
    IF IA32_EFER.LMA = 0 and new stack segment not present
        THEN #SS(NewSS); FI;
    IF CallGateSize = 32
        THEN
            IF new stack does not have room for parameters plus 16 bytes
                THEN #SS(NewSS); FI;
            IF CallGate(InstructionPointer) not within new code-segment limit
                THEN #GP(0); FI;
            SS ← newSS; (* Segment descriptor information also loaded *)
            ESP ← newESP;
            CS:EIP ← CallGate(CS:InstructionPointer);
            (* Segment descriptor information also loaded *)
            Push(oldSS:oldESP); (* From calling procedure *)
            temp ← parameter count from call gate, masked to 5 bits;
            Push(parameters from calling procedure’s stack, temp)
            Push(oldCS:oldEIP); (* Return address to calling procedure *)
        ELSE
            IF CallGateSize = 16
                THEN
                    IF new stack does not have room for parameters plus 8 bytes
                        THEN #SS(NewSS); FI;
                    IF (CallGate(InstructionPointer) AND FFFFH) not in new code-segment limit
```

```
                        THEN #GP(0); FI;
                    SS ← newSS; (* Segment descriptor information also loaded *)
                    ESP ← newESP;
                    CS:IP ← CallGate(CS:InstructionPointer);
                    (* Segment descriptor information also loaded *)
                    Push(oldSS:oldESP); (* From calling procedure *)
                    temp ← parameter count from call gate, masked to 5 bits;
                    Push(parameters from calling procedure's stack, temp)
                    Push(oldCS:oldEIP); (* Return address to calling procedure *)
                ELSE (* CallGateSize = 64 *)
                    IF pushing 32 bytes on the stack would use a non-canonical address
                        THEN #SS(NewSS); FI;
                    IF (CallGate(InstructionPointer) is non-canonical)
                        THEN #GP(0); FI;
                    SS ← NewSS; (* NewSS is NULL)
                    RSP ← NewESP;
                    CS:IP ← CallGate(CS:InstructionPointer);
                    (* Segment descriptor information also loaded *)
                    Push(oldSS:oldESP); (* From calling procedure *)
                    Push(oldCS:oldEIP); (* Return address to calling procedure *)
            FI;
    FI;
    CPL ← CodeSegment(DPL)
    CS(RPL) ← CPL
END;

SAME-PRIVILEGE:
    IF CallGateSize = 32
        THEN
            IF stack does not have room for 8 bytes
                THEN #SS(0); FI;
            IF CallGate(InstructionPointer) not within code segment limit
                THEN #GP(0); FI;
            CS:EIP ← CallGate(CS:EIP) (* Segment descriptor information also loaded *)
            Push(oldCS:oldEIP); (* Return address to calling procedure *)
        ELSE
            If CallGateSize = 16
                THEN
                    IF stack does not have room for 4 bytes
                        THEN #SS(0); FI;
                    IF CallGate(InstructionPointer) not within code segment limit
                        THEN #GP(0); FI;
                    CS:IP ← CallGate(CS:instruction pointer);
```

```
                        (* Segment descriptor information also loaded *)
                        Push(oldCS:oldIP); (* Return address to calling procedure *)
                    ELSE (* CallGateSize = 64)
                        IF pushing 16 bytes on the stack touches non-canonical addresses
                            THEN #SS(0); FI;
                        IF RIP non-canonical
                            THEN #GP(0); FI;
                        CS:IP ← CallGate(CS:instruction pointer);
                        (* Segment descriptor information also loaded *)
                        Push(oldCS:oldIP); (* Return address to calling procedure *)
                FI;
    FI;
    CS(RPL) ← CPL
END;

TASK-GATE:
    IF task gate DPL < CPL or RPL
        THEN #GP(task gate selector); FI;
    IF task gate not present
        THEN #NP(task gate selector); FI;
    Read the TSS segment selector in the task-gate descriptor;
    IF TSS segment selector local/global bit is set to local
    or index not within GDT limits
        THEN #GP(TSS selector); FI;
    Access TSS descriptor in GDT;
    IF TSS descriptor specifies that the TSS is busy (low-order 5 bits set to 00001)
        THEN #GP(TSS selector); FI;
    IF TSS not present
        THEN #NP(TSS selector); FI;
    SWITCH-TASKS (with nesting) to TSS;
    IF EIP not within code segment limit
        THEN #GP(0); FI;
END;

TASK-STATE-SEGMENT:
    IF TSS DPL < CPL or RPL
    or TSS descriptor indicates TSS not available
        THEN #GP(TSS selector); FI;
    IF TSS is not present
        THEN #NP(TSS selector); FI;
    SWITCH-TASKS (with nesting) to TSS;
    IF EIP not within code segment limit
        THEN #GP(0); FI;
```

```
END;
```

Flags Affected

All flags are affected if a task switch occurs; no flags are affected if a task switch does not occur.

Protected Mode Exceptions

#GP(0)	If the target offset in destination operand is beyond the new code segment limit.
	If the segment selector in the destination operand is NULL.
	If the code segment selector in the gate is NULL.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#GP(selector)	If a code segment or gate or TSS selector index is outside descriptor table limits.
	If the segment descriptor pointed to by the segment selector in the destination operand is not for a conforming-code segment, nonconforming-code segment, call gate, task gate, or task state segment.
	If the DPL for a nonconforming-code segment is not equal to the CPL or the RPL for the segment's segment selector is greater than the CPL.
	If the DPL for a conforming-code segment is greater than the CPL.
	If the DPL from a call-gate, task-gate, or TSS segment descriptor is less than the CPL or than the RPL of the call-gate, task-gate, or TSS's segment selector.
	If the segment descriptor for a segment selector from a call gate does not indicate it is a code segment.
	If the segment selector from a call gate is beyond the descriptor table limits.
	If the DPL for a code-segment obtained from a call gate is greater than the CPL.
	If the segment selector for a TSS has its local/global bit set for local.
	If a TSS segment descriptor specifies that the TSS is busy or not available.

#SS(0)	If pushing the return address, parameters, or stack segment pointer onto the stack exceeds the bounds of the stack segment, when no stack switch occurs.
	If a memory operand effective address is outside the SS segment limit.
#SS(selector)	If pushing the return address, parameters, or stack segment pointer onto the stack exceeds the bounds of the stack segment, when a stack switch occurs.
	If the SS register is being loaded as part of a stack switch and the segment pointed to is marked not present.
	If stack segment does not have room for the return address, parameters, or stack segment pointer, when stack switch occurs.
#NP(selector)	If a code segment, data segment, stack segment, call gate, task gate, or TSS is not present.
#TS(selector)	If the new stack segment selector and ESP are beyond the end of the TSS.
	If the new stack segment selector is NULL.
	If the RPL of the new stack segment selector in the TSS is not equal to the DPL of the code segment being accessed.
	If DPL of the stack segment descriptor for the new stack segment is not equal to the DPL of the code segment descriptor.
	If the new stack segment is not a writable data segment.
	If segment-selector index for stack segment is outside descriptor table limits.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the target offset is beyond the code segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the target offset is beyond the code segment limit.
#PF(fault-code)	If a page fault occurs.

#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

#GP(selector)	If a memory address accessed by the selector is in non-canonical space.
#GP(0)	If the target offset in the destination operand is non-canonical.

64-Bit Mode Exceptions

#GP(0)	If a memory address is non-canonical.
	If target offset in destination operand is non-canonical.
	If the segment selector in the destination operand is NULL.
	If the code segment selector in the 64-bit gate is NULL.
#GP(selector)	If code segment or 64-bit call gate is outside descriptor table limits.
	If code segment or 64-bit call gate overlaps non-canonical space.
	If the segment descriptor pointed to by the segment selector in the destination operand is not for a conforming-code segment, nonconforming-code segment, or 64-bit call gate.
	If the segment descriptor pointed to by the segment selector in the destination operand is a code segment and has both the D-bit and the L- bit set.
	If the DPL for a nonconforming-code segment is not equal to the CPL, or the RPL for the segment’s segment selector is greater than the CPL.
	If the DPL for a conforming-code segment is greater than the CPL.
	If the DPL from a 64-bit call-gate is less than the CPL or than the RPL of the 64-bit call-gate.
	If the upper type field of a 64-bit call gate is not 0x0.
	If the segment selector from a 64-bit call gate is beyond the descriptor table limits.
	If the DPL for a code-segment obtained from a 64-bit call gate is greater than the CPL.
	If the code segment descriptor pointed to by the selector in the 64-bit gate doesn't have the L-bit set and the D-bit clear.
	If the segment descriptor for a segment selector from the 64-bit call gate does not indicate it is a code segment.

#SS(0)	If pushing the return offset or CS selector onto the stack exceeds the bounds of the stack segment when no stack switch occurs.
	If a memory operand effective address is outside the SS segment limit.
	If the stack address is in a non-canonical form.
#SS(selector)	If pushing the old values of SS selector, stack pointer, EFLAGS, CS selector, offset, or error code onto the stack violates the canonical boundary when a stack switch occurs.
#NP(selector)	If a code segment or 64-bit call gate is not present.
#TS(selector)	If the load of the new RSP exceeds the limit of the TSS.
#UD	(64-bit mode only) If a far call is direct to an absolute address in memory.
	If the LOCK prefix is used.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

CBW/CWDE/CDQE—Convert Byte to Word/Convert Word to Doubleword/Convert Doubleword to Quadword

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
98	CBW	A	Valid	Valid	AX ← sign-extend of AL.
98	CWDE	A	Valid	Valid	EAX ← sign-extend of AX.
REX.W + 98	CDQE	A	Valid	N.E.	RAX ← sign-extend of EAX.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Double the size of the source operand by means of sign extension. The CBW (convert byte to word) instruction copies the sign (bit 7) in the source operand into every bit in the AH register. The CWDE (convert word to doubleword) instruction copies the sign (bit 15) of the word in the AX register into the high 16 bits of the EAX register.

CBW and CWDE reference the same opcode. The CBW instruction is intended for use when the operand-size attribute is 16; CWDE is intended for use when the operand-size attribute is 32. Some assemblers may force the operand size. Others may treat these two mnemonics as synonyms (CBW/CWDE) and use the setting of the operand-size attribute to determine the size of values to be converted.

In 64-bit mode, the default operation size is the size of the destination register. Use of the REX.W prefix promotes this instruction (CDQE when promoted) to operate on 64-bit operands. In which case, CDQE copies the sign (bit 31) of the doubleword in the EAX register into the high 32 bits of RAX.

Operation

```
IF OperandSize = 16 (* Instruction = CBW *)
    THEN
        AX ← SignExtend(AL);
    ELSE IF (OperandSize = 32, Instruction = CWDE)
        EAX ← SignExtend(AX); FI;
    ELSE (* 64-Bit Mode, OperandSize = 64, Instruction = CDQE*)
        RAX ← SignExtend(EAX);
FI;
```

Flags Affected

None.

Exceptions (All Operating Modes)

#UD	If the LOCK prefix is used.

CLC—Clear Carry Flag

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
F8	CLC	A	Valid	Valid	Clear CF flag.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Clears the CF flag in the EFLAGS register. Operation is the same in all non-64-bit modes and 64-bit mode.

Operation

CF $\leftarrow$ 0;

Flags Affected

The CF flag is set to 0. The OF, ZF, SF, AF, and PF flags are unaffected.

Exceptions (All Operating Modes)

#UD	If the LOCK prefix is used.

CLD—Clear Direction Flag

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
FC	CLD	A	Valid	Valid	Clear DF flag.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Clears the DF flag in the EFLAGS register. When the DF flag is set to 0, string operations increment the index registers (ESI and/or EDI). Operation is the same in all non-64-bit modes and 64-bit mode.

Operation

DF ← 0;

Flags Affected

The DF flag is set to 0. The CF, OF, ZF, SF, AF, and PF flags are unaffected.

Exceptions (All Operating Modes)

#UD	If the LOCK prefix is used.

CLFLUSH—Flush Cache Line

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
0F AE /7	CLFLUSH *m8*	A	Valid	Valid	Flushes cache line containing *m8*.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (w)	NA	NA	NA

Description

Invalidates the cache line that contains the linear address specified with the source operand from all levels of the processor cache hierarchy (data and instruction). The invalidation is broadcast throughout the cache coherence domain. If, at any level of the cache hierarchy, the line is inconsistent with memory (dirty) it is written to memory before invalidation. The source operand is a byte memory location.

The availability of CLFLUSH is indicated by the presence of the CPUID feature flag CLFSH (bit 19 of the EDX register, see "CPUID—CPU Identification" in this chapter). The aligned cache line size affected is also indicated with the CPUID instruction (bits 8 through 15 of the EBX register when the initial value in the EAX register is 1).

The memory attribute of the page containing the affected line has no effect on the behavior of this instruction. It should be noted that processors are free to speculatively fetch and cache data from system memory regions assigned a memory-type allowing for speculative reads (such as, the WB, WC, and WT memory types). PREFETCH*h* instructions can be used to provide the processor with hints for this speculative behavior. Because this speculative fetching can occur at any time and is not tied to instruction execution, the CLFLUSH instruction is not ordered with respect to PREFETCH*h* instructions or any of the speculative fetching mechanisms (that is, data can be speculatively loaded into a cache line just before, during, or after the execution of a CLFLUSH instruction that references the cache line).

CLFLUSH is only ordered by the MFENCE instruction. It is not guaranteed to be ordered by any other fencing or serializing instructions or by another CLFLUSH instruction. For example, software can use an MFENCE instruction to ensure that previous stores are included in the write-back.

The CLFLUSH instruction can be used at all privilege levels and is subject to all permission checking and faults associated with a byte load (and in addition, a CLFLUSH instruction is allowed to flush a linear address in an execute-only segment). Like a load, the CLFLUSH instruction sets the A bit but not the D bit in the page tables.

The CLFLUSH instruction was introduced with the SSE2 extensions; however, because it has its own CPUID feature flag, it can be implemented in IA-32 processors that do not include the SSE2 extensions. Also, detecting the presence of the SSE2 extensions with the CPUID instruction does not guarantee that the CLFLUSH instruction is implemented in the processor.

CLFLUSH operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
Flush_Cache_Line(SRC);
```

Intel C/C++ Compiler Intrinsic Equivalents

CLFLUSH void _mm_clflush(void const *p)

Protected Mode Exceptions

#GP(0)	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
#SS(0)	For an illegal address in the SS segment.
#PF(fault-code)	For a page fault.
#UD	If CPUID.01H:EDX.CLFSH[bit 19] = 0.
	If the LOCK prefix is used.

Real-Address Mode Exceptions

GP	If any part of the operand lies outside the effective address space from 0 to FFFFH.
#UD	If CPUID.01H:EDX.CLFSH[bit 19] = 0.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	For a page fault.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	For a page fault.
#UD	If CPUID.01H:EDX.CLFSH[bit 19] = 0.
	If the LOCK prefix is used.

CLI — Clear Interrupt Flag

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
FA	CLI	A	Valid	Valid	Clear interrupt flag; interrupts disabled when interrupt flag cleared.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

If protected-mode virtual interrupts are not enabled, CLI clears the IF flag in the EFLAGS register. No other flags are affected. Clearing the IF flag causes the processor to ignore maskable external interrupts. The IF flag and the CLI and STI instruction have no affect on the generation of exceptions and NMI interrupts.

When protected-mode virtual interrupts are enabled, CPL is 3, and IOPL is less than 3; CLI clears the VIF flag in the EFLAGS register, leaving IF unaffected. Table 3-6 indicates the action of the CLI instruction depending on the processor operating mode and the CPL/IOPL of the running program or procedure.

CLI operation is the same in non-64-bit modes and 64-bit mode.

Table 3-6. Decision Table for CLI Results

PE	VM	IOPL	CPL	PVI	VIP	VME	CLI Result
0	X	X	X	X	X	X	**IF = 0**
1	0	≥ CPL	X	X	X	X	**IF = 0**
1	0	< CPL	3	1	X	X	**VIF = 0**
1	0	< CPL	< 3	X	X	X	**GP Fault**
1	0	< CPL	X	0	X	X	**GP Fault**
1	1	3	X	X	X	X	**IF = 0**
1	1	< 3	X	X	X	1	**VIF = 0**
1	1	< 3	X	X	X	0	**GP Fault**

NOTES:

* X = This setting has no impact.

Operation

IF PE = 0

```
    THEN
        IF ← 0; (* Reset Interrupt Flag *)
    ELSE
        IF VM = 0;
            THEN
                IF IOPL ≥ CPL
                    THEN
                        IF ← 0; (* Reset Interrupt Flag *)
                ELSE
                    IF ((IOPL < CPL) and (CPL = 3) and (PVI = 1))
                        THEN
                            VIF ← 0; (* Reset Virtual Interrupt Flag *)
                        ELSE
                            #GP(0);
                    FI;
                FI;
            ELSE (* VM = 1 *)
                IF IOPL = 3
                    THEN
                        IF ← 0; (* Reset Interrupt Flag *)
                    ELSE
                        IF (IOPL < 3) AND (VME = 1)
                            THEN
                                VIF ← 0; (* Reset Virtual Interrupt Flag *)
                            ELSE
                                #GP(0);
                        FI;
                FI;
        FI;
FI;
```

Flags Affected

If protected-mode virtual interrupts are not enabled, IF is set to 0 if the CPL is equal to or less than the IOPL; otherwise, it is not affected. The other flags in the EFLAGS register are unaffected.

When protected-mode virtual interrupts are enabled, CPL is 3, and IOPL is less than 3; CLI clears the VIF flag in the EFLAGS register, leaving IF unaffected.

Protected Mode Exceptions

#GP(0)	If the CPL is greater (has less privilege) than the IOPL of the current program or procedure.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If the CPL is greater (has less privilege) than the IOPL of the current program or procedure.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(0)	If the CPL is greater (has less privilege) than the IOPL of the current program or procedure.
#UD	If the LOCK prefix is used.

CLTS—Clear Task-Switched Flag in CR0

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
0F 06	CLTS	A	Valid	Valid	Clears TS flag in CR0.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Clears the task-switched (TS) flag in the CR0 register. This instruction is intended for use in operating-system procedures. It is a privileged instruction that can only be executed at a CPL of 0. It is allowed to be executed in real-address mode to allow initialization for protected mode.

The processor sets the TS flag every time a task switch occurs. The flag is used to synchronize the saving of FPU context in multitasking applications. See the description of the TS flag in the section titled "Control Registers" in Chapter 2 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*, for more information about this flag.

CLTS operation is the same in non-64-bit modes and 64-bit mode.

See Chapter 22, "VMX Non-Root Operation," of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3B*, for more information about the behavior of this instruction in VMX non-root operation.

Operation

CR0.TS[bit 3] ← 0;

Flags Affected

The TS flag in CR0 register is cleared.

Protected Mode Exceptions

#GP(0)	If the current privilege level is not 0.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	CLTS is not recognized in virtual-8086 mode.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(0)	If the CPL is greater than 0.
#UD	If the LOCK prefix is used.

CMC—Complement Carry Flag

Opcode	Instruction	Op/ En	64-bit Mode	Compat/ Leg Mode	Description
F5	CMC	A	Valid	Valid	Complement CF flag.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Complements the CF flag in the EFLAGS register. CMC operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
EFLAGS.CF[bit 0]← NOT EFLAGS.CF[bit 0];
```

Flags Affected

The CF flag contains the complement of its original value. The OF, ZF, SF, AF, and PF flags are unaffected.

Exceptions (All Operating Modes)

#UD	If the LOCK prefix is used.

CMOV*cc*—Conditional Move

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 47 */r*	CMOVA *r16, r/m16*	A	Valid	Valid	Move if above (CF=0 and ZF=0).
0F 47 */r*	CMOVA *r32, r/m32*	A	Valid	Valid	Move if above (CF=0 and ZF=0).
REX.W + 0F 47 */r*	CMOVA *r64, r/m64*	A	Valid	N.E.	Move if above (CF=0 and ZF=0).
0F 43 */r*	CMOVAE *r16, r/m16*	A	Valid	Valid	Move if above or equal (CF=0).
0F 43 */r*	CMOVAE *r32, r/m32*	A	Valid	Valid	Move if above or equal (CF=0).
REX.W + 0F 43 */r*	CMOVAE *r64, r/m64*	A	Valid	N.E.	Move if above or equal (CF=0).
0F 42 */r*	CMOVB *r16, r/m16*	A	Valid	Valid	Move if below (CF=1).
0F 42 */r*	CMOVB *r32, r/m32*	A	Valid	Valid	Move if below (CF=1).
REX.W + 0F 42 */r*	CMOVB *r64, r/m64*	A	Valid	N.E.	Move if below (CF=1).
0F 46 */r*	CMOVBE *r16, r/m16*	A	Valid	Valid	Move if below or equal (CF=1 or ZF=1).
0F 46 */r*	CMOVBE *r32, r/m32*	A	Valid	Valid	Move if below or equal (CF=1 or ZF=1).
REX.W + 0F 46 */r*	CMOVBE *r64, r/m64*	A	Valid	N.E.	Move if below or equal (CF=1 or ZF=1).
0F 42 */r*	CMOVC *r16, r/m16*	A	Valid	Valid	Move if carry (CF=1).
0F 42 */r*	CMOVC *r32, r/m32*	A	Valid	Valid	Move if carry (CF=1).
REX.W + 0F 42 */r*	CMOVC *r64, r/m64*	A	Valid	N.E.	Move if carry (CF=1).
0F 44 */r*	CMOVE *r16, r/m16*	A	Valid	Valid	Move if equal (ZF=1).
0F 44 */r*	CMOVE *r32, r/m32*	A	Valid	Valid	Move if equal (ZF=1).
REX.W + 0F 44 */r*	CMOVE *r64, r/m64*	A	Valid	N.E.	Move if equal (ZF=1).
0F 4F */r*	CMOVG *r16, r/m16*	A	Valid	Valid	Move if greater (ZF=0 and SF=OF).
0F 4F */r*	CMOVG *r32, r/m32*	A	Valid	Valid	Move if greater (ZF=0 and SF=OF).
REX.W + 0F 4F */r*	CMOVG *r64, r/m64*	A	V/N.E.	NA	Move if greater (ZF=0 and SF=OF).

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 4D /r	CMOVGE *r16, r/m16*	A	Valid	Valid	Move if greater or equal (SF=OF).
0F 4D /r	CMOVGE *r32, r/m32*	A	Valid	Valid	Move if greater or equal (SF=OF).
REX.W + 0F 4D /r	CMOVGE *r64, r/m64*	A	Valid	N.E.	Move if greater or equal (SF=OF).
0F 4C /r	CMOVL *r16, r/m16*	A	Valid	Valid	Move if less (SF≠ OF).
0F 4C /r	CMOVL *r32, r/m32*	A	Valid	Valid	Move if less (SF≠ OF).
REX.W + 0F 4C /r	CMOVL *r64, r/m64*	A	Valid	N.E.	Move if less (SF≠ OF).
0F 4E /r	CMOVLE *r16, r/m16*	A	Valid	Valid	Move if less or equal (ZF=1 or SF≠ OF).
0F 4E /r	CMOVLE *r32, r/m32*	A	Valid	Valid	Move if less or equal (ZF=1 or SF≠ OF).
REX.W + 0F 4E /r	CMOVLE *r64, r/m64*	A	Valid	N.E.	Move if less or equal (ZF=1 or SF≠ OF).
0F 46 /r	CMOVNA *r16, r/m16*	A	Valid	Valid	Move if not above (CF=1 or ZF=1).
0F 46 /r	CMOVNA *r32, r/m32*	A	Valid	Valid	Move if not above (CF=1 or ZF=1).
REX.W + 0F 46 /r	CMOVNA *r64, r/m64*	A	Valid	N.E.	Move if not above (CF=1 or ZF=1).
0F 42 /r	CMOVNAE *r16, r/m16*	A	Valid	Valid	Move if not above or equal (CF=1).
0F 42 /r	CMOVNAE *r32, r/m32*	A	Valid	Valid	Move if not above or equal (CF=1).
REX.W + 0F 42 /r	CMOVNAE *r64, r/m64*	A	Valid	N.E.	Move if not above or equal (CF=1).
0F 43 /r	CMOVNB *r16, r/m16*	A	Valid	Valid	Move if not below (CF=0).
0F 43 /r	CMOVNB *r32, r/m32*	A	Valid	Valid	Move if not below (CF=0).
REX.W + 0F 43 /r	CMOVNB *r64, r/m64*	A	Valid	N.E.	Move if not below (CF=0).
0F 47 /r	CMOVNBE *r16, r/m16*	A	Valid	Valid	Move if not below or equal (CF=0 and ZF=0).
0F 47 /r	CMOVNBE *r32, r/m32*	A	Valid	Valid	Move if not below or equal (CF=0 and ZF=0).

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
REX.W + 0F 47 /r	CMOVNBE *r64, r/m64*	A	Valid	N.E.	Move if not below or equal (CF=0 and ZF=0).
0F 43 /r	CMOVNC *r16, r/m16*	A	Valid	Valid	Move if not carry (CF=0).
0F 43 /r	CMOVNC *r32, r/m32*	A	Valid	Valid	Move if not carry (CF=0).
REX.W + 0F 43 /r	CMOVNC *r64, r/m64*	A	Valid	N.E.	Move if not carry (CF=0).
0F 45 /r	CMOVNE *r16, r/m16*	A	Valid	Valid	Move if not equal (ZF=0).
0F 45 /r	CMOVNE *r32, r/m32*	A	Valid	Valid	Move if not equal (ZF=0).
REX.W + 0F 45 /r	CMOVNE *r64, r/m64*	A	Valid	N.E.	Move if not equal (ZF=0).
0F 4E /r	CMOVNG *r16, r/m16*	A	Valid	Valid	Move if not greater (ZF=1 or SF≠ OF).
0F 4E /r	CMOVNG *r32, r/m32*	A	Valid	Valid	Move if not greater (ZF=1 or SF≠ OF).
REX.W + 0F 4E /r	CMOVNG *r64, r/m64*	A	Valid	N.E.	Move if not greater (ZF=1 or SF≠ OF).
0F 4C /r	CMOVNGE *r16, r/m16*	A	Valid	Valid	Move if not greater or equal (SF≠ OF).
0F 4C /r	CMOVNGE *r32, r/m32*	A	Valid	Valid	Move if not greater or equal (SF≠ OF).
REX.W + 0F 4C /r	CMOVNGE *r64, r/m64*	A	Valid	N.E.	Move if not greater or equal (SF≠ OF).
0F 4D /r	CMOVNL *r16, r/m16*	A	Valid	Valid	Move if not less (SF=OF).
0F 4D /r	CMOVNL *r32, r/m32*	A	Valid	Valid	Move if not less (SF=OF).
REX.W + 0F 4D /r	CMOVNL *r64, r/m64*	A	Valid	N.E.	Move if not less (SF=OF).
0F 4F /r	CMOVNLE *r16, r/m16*	A	Valid	Valid	Move if not less or equal (ZF=0 and SF=OF).
0F 4F /r	CMOVNLE *r32, r/m32*	A	Valid	Valid	Move if not less or equal (ZF=0 and SF=OF).
REX.W + 0F 4F /r	CMOVNLE *r64, r/m64*	A	Valid	N.E.	Move if not less or equal (ZF=0 and SF=OF).
0F 41 /r	CMOVNO *r16, r/m16*	A	Valid	Valid	Move if not overflow (OF=0).
0F 41 /r	CMOVNO *r32, r/m32*	A	Valid	Valid	Move if not overflow (OF=0).

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
REX.W + 0F 41 /r	CMOVNO *r64, r/m64*	A	Valid	N.E.	Move if not overflow (OF=0).
0F 4B /r	CMOVNP *r16, r/m16*	A	Valid	Valid	Move if not parity (PF=0).
0F 4B /r	CMOVNP *r32, r/m32*	A	Valid	Valid	Move if not parity (PF=0).
REX.W + 0F 4B /r	CMOVNP *r64, r/m64*	A	Valid	N.E.	Move if not parity (PF=0).
0F 49 /r	CMOVNS *r16, r/m16*	A	Valid	Valid	Move if not sign (SF=0).
0F 49 /r	CMOVNS *r32, r/m32*	A	Valid	Valid	Move if not sign (SF=0).
REX.W + 0F 49 /r	CMOVNS *r64, r/m64*	A	Valid	N.E.	Move if not sign (SF=0).
0F 45 /r	CMOVNZ *r16, r/m16*	A	Valid	Valid	Move if not zero (ZF=0).
0F 45 /r	CMOVNZ *r32, r/m32*	A	Valid	Valid	Move if not zero (ZF=0).
REX.W + 0F 45 /r	CMOVNZ *r64, r/m64*	A	Valid	N.E.	Move if not zero (ZF=0).
0F 40 /r	CMOVO *r16, r/m16*	A	Valid	Valid	Move if overflow (OF=1).
0F 40 /r	CMOVO *r32, r/m32*	A	Valid	Valid	Move if overflow (OF=1).
REX.W + 0F 40 /r	CMOVO *r64, r/m64*	A	Valid	N.E.	Move if overflow (OF=1).
0F 4A /r	CMOVP *r16, r/m16*	A	Valid	Valid	Move if parity (PF=1).
0F 4A /r	CMOVP *r32, r/m32*	A	Valid	Valid	Move if parity (PF=1).
REX.W + 0F 4A /r	CMOVP *r64, r/m64*	A	Valid	N.E.	Move if parity (PF=1).
0F 4A /r	CMOVPE *r16, r/m16*	A	Valid	Valid	Move if parity even (PF=1).
0F 4A /r	CMOVPE *r32, r/m32*	A	Valid	Valid	Move if parity even (PF=1).
REX.W + 0F 4A /r	CMOVPE *r64, r/m64*	A	Valid	N.E.	Move if parity even (PF=1).
0F 4B /r	CMOVPO *r16, r/m16*	A	Valid	Valid	Move if parity odd (PF=0).
0F 4B /r	CMOVPO *r32, r/m32*	A	Valid	Valid	Move if parity odd (PF=0).
REX.W + 0F 4B /r	CMOVPO *r64, r/m64*	A	Valid	N.E.	Move if parity odd (PF=0).
0F 48 /r	CMOVS *r16, r/m16*	A	Valid	Valid	Move if sign (SF=1).
0F 48 /r	CMOVS *r32, r/m32*	A	Valid	Valid	Move if sign (SF=1).
REX.W + 0F 48 /r	CMOVS *r64, r/m64*	A	Valid	N.E.	Move if sign (SF=1).
0F 44 /r	CMOVZ *r16, r/m16*	A	Valid	Valid	Move if zero (ZF=1).

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 44 */r*	CMOVZ *r32, r/m32*	A	Valid	Valid	Move if zero (ZF=1).
REX.W + 0F 44 */r*	CMOVZ *r64, r/m64*	A	Valid	N.E.	Move if zero (ZF=1).

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA

Description

The CMOV*cc* instructions check the state of one or more of the status flags in the EFLAGS register (CF, OF, PF, SF, and ZF) and perform a move operation if the flags are in a specified state (or condition). A condition code (*cc*) is associated with each instruction to indicate the condition being tested for. If the condition is not satisfied, a move is not performed and execution continues with the instruction following the CMOV*cc* instruction.

These instructions can move 16-bit, 32-bit or 64-bit values from memory to a general-purpose register or from one general-purpose register to another. Conditional moves of 8-bit register operands are not supported.

The condition for each CMOV*cc* mnemonic is given in the description column of the above table. The terms "less" and "greater" are used for comparisons of signed integers and the terms "above" and "below" are used for unsigned integers.

Because a particular state of the status flags can sometimes be interpreted in two ways, two mnemonics are defined for some opcodes. For example, the CMOVA (conditional move if above) instruction and the CMOVNBE (conditional move if not below or equal) instruction are alternate mnemonics for the opcode 0F 47H.

The CMOV*cc* instructions were introduced in P6 family processors; however, these instructions may not be supported by all IA-32 processors. Software can determine if the CMOV*cc* instructions are supported by checking the processor's feature information with the CPUID instruction (see "CPUID—CPU Identification" in this chapter).

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

temp ← SRC

IF condition TRUE

```
    THEN
        DEST ← temp;
    FI;
ELSE
    IF (OperandSize = 32 and IA-32e mode active)
        THEN
            DEST[63:32] ← 0;
    FI;
FI;
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

CMP—Compare Two Operands

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
3C *ib*	CMP AL, *imm8*	D	Valid	Valid	Compare *imm8* with AL.
3D *iw*	CMP AX, *imm16*	D	Valid	Valid	Compare *imm16* with AX.
3D *id*	CMP EAX, *imm32*	D	Valid	Valid	Compare *imm32* with EAX.
REX.W + 3D *id*	CMP RAX, *imm32*	D	Valid	N.E.	Compare *imm32 sign-extended to 64-bits* with RAX.
80 /7 *ib*	CMP *r/m8, imm8*	C	Valid	Valid	Compare *imm8* with *r/m8.*
REX + 80 /7 *ib*	CMP *r/m8**, *imm8*	C	Valid	N.E.	Compare *imm8* with *r/m8.*
81 /7 *iw*	CMP *r/m16, imm16*	C	Valid	Valid	Compare *imm16* with *r/m16.*
81 /7 *id*	CMP *r/m32, imm32*	C	Valid	Valid	Compare *imm32* with *r/m32.*
REX.W + 81 /7 *id*	CMP *r/m64, imm32*	C	Valid	N.E.	Compare *imm32 sign-extended to 64-bits* with *r/m64.*
83 /7 *ib*	CMP *r/m16, imm8*	C	Valid	Valid	Compare *imm8* with *r/m16.*
83 /7 *ib*	CMP *r/m32, imm8*	C	Valid	Valid	Compare *imm8* with *r/m32.*
REX.W + 83 /7 *ib*	CMP *r/m64, imm8*	C	Valid	N.E.	Compare *imm8* with *r/m64.*
38 /*r*	CMP *r/m8, r8*	B	Valid	Valid	Compare *r8* with *r/m8.*
REX + 38 /*r*	CMP *r/m8**, *r8**	B	Valid	N.E.	Compare *r8* with *r/m8.*
39 /*r*	CMP *r/m16, r16*	B	Valid	Valid	Compare *r16* with *r/m16.*
39 /*r*	CMP *r/m32, r32*	B	Valid	Valid	Compare *r32* with *r/m32.*
REX.W + 39 /*r*	CMP *r/m64,r64*	B	Valid	N.E.	Compare *r64* with *r/m64.*
3A /*r*	CMP *r8, r/m8*	A	Valid	Valid	Compare *r/m8* with *r8.*
REX + 3A /*r*	CMP *r8**, *r/m8**	A	Valid	N.E.	Compare *r/m8 with r8.*
3B /*r*	CMP *r16, r/m16*	A	Valid	Valid	Compare *r/m16* with *r16.*
3B /*r*	CMP *r32, r/m32*	A	Valid	Valid	Compare *r/m32* with *r32.*
REX.W + 3B /*r*	CMP *r64, r/m64*	A	Valid	N.E.	Compare *r/m64* with *r64.*

NOTES:

* In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (r, w)	ModRM:reg (w)	NA	NA
C	ModRM:r/m (r, w)	imm8	NA	NA
D	AL/AX/EAX/RAX	imm8	NA	NA

Description

Compares the first source operand with the second source operand and sets the status flags in the EFLAGS register according to the results. The comparison is performed by subtracting the second operand from the first operand and then setting the status flags in the same manner as the SUB instruction. When an immediate value is used as an operand, it is sign-extended to the length of the first operand.

The condition codes used by the J*cc*, CMOV*cc*, and SET*cc* instructions are based on the results of a CMP instruction. Appendix B, "EFLAGS Condition Codes," in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, shows the relationship of the status flags and the condition codes.

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
temp ← SRC1 − SignExtend(SRC2);
ModifyStatusFlags; (* Modify status flags in the same manner as the SUB instruction*)
```

Flags Affected

The CF, OF, SF, ZF, AF, and PF flags are set according to the result.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

CMPPD—Compare Packed Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F C2 /*r* ib CMPPD *xmm1*, *xmm2/m128*, *imm8*	A	V/V	SSE2	Compare packed double-precision floating-point values in *xmm2/m128* and *xmm1* using imm8 as comparison predicate.
VEX.NDS.128.66.0F.WIG C2 /r ib VCMPPD xmm1, xmm2, xmm3/m128, imm8	B	V/V	AVX	Compare packed double-precision floating-point values in xmm3/m128 and xmm2 using bits 4:0 of imm8 as a comparison predicate.
VEX.NDS.256.66.0F.WIG C2 /r ib VCMPPD ymm1, ymm2, ymm3/m256, imm8	B	V/V	AVX	Compare packed double-precision floating-point values in ymm3/m256 and ymm2 using bits 4:0 of imm8 as a comparison predicate.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	imm8	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs a SIMD compare of the packed double-precision floating-point values in the source operand (second operand) and the destination operand (first operand) and returns the results of the comparison to the destination operand. The comparison predicate operand (third operand) specifies the type of comparison performed on each of the pairs of packed values. The result of each comparison is a quadword mask of all 1s (comparison true) or all 0s (comparison false).

128-bit Legacy SSE version: The first source and destination operand (first operand) is an XMM register. The second source operand (second operand) can be an XMM register or 128-bit memory location. The comparison predicate operand is an 8-bit immediate, bits 2:0 of the immediate define the type of comparison to be performed (see Table 3-7). Bits 7:3 of the immediate is reserved. Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged. Two comparisons are performed with results written to bits 127:0 of the destination operand.

Table 3-7. Comparison Predicate for CMPPD and CMPPS Instructions

Predi-cate	imm8 Encod-ing	Description	Relation where: A Is 1st Operand B Is 2nd Operand	Emulation	Result if NaN Operand	QNaN Oper-and Signals Invalid
EQ	000B	Equal	A = B		False	No
LT	001B	Less-than	A < B		False	Yes
LE	010B	Less-than-or-equal	A ≤ B		False	Yes
		Greater than	A > B	Swap Operands, Use LT	False	Yes
		Greater-than-or-equal	A ≥ B	Swap Operands, Use LE	False	Yes
UNORD	011B	Unordered	A, B = Unordered		True	No
NEQ	100B	Not-equal	A ≠ B		True	No
NLT	101B	Not-less-than	NOT(A < B)		True	Yes
NLE	110B	Not-less-than-or-equal	NOT(A ≤ B)		True	Yes
		Not-greater-than	NOT(A > B)	Swap Operands, Use NLT	True	Yes
		Not-greater-than-or-equal	NOT(A ≥ B)	Swap Operands, Use NLE	True	Yes
ORD	111B	Ordered	A , B = Ordered		False	No

The unordered relationship is true when at least one of the two source operands being compared is a NaN; the ordered relationship is true when neither source operand is a NaN.

A subsequent computational instruction that uses the mask result in the destination operand as an input operand will not generate an exception, because a mask of all 0s corresponds to a floating-point value of +0.0 and a mask of all 1s corresponds to a QNaN.

Note that the processors with "CPUID.1H:ECX.AVX =0" do not implement the greater-than, greater-than-or-equal, not-greater-than, and not-greater-than-or-equal relations. These comparisons can be made either by using the inverse relationship (that is, use the "not-less-than-or-equal" to make a "greater-than" comparison) or by using software emulation. When using software emulation, the program must swap the operands (copying registers when necessary to protect the data that will now be in the destination), and then perform the compare using a different predicate. The predicate to be used for these emulations is listed in Table 3-7 under the heading Emulation.

Compilers and assemblers may implement the following two-operand pseudo-ops in addition to the three-operand CMPPD instruction, for processors with "CPUID.1H:ECX.AVX =0". See Table 3-8. Compiler should treat reserved Imm8 values as illegal syntax.

Table 3-8. Pseudo-Op and CMPPD Implementation

Pseudo-Op	CMPPD Implementation
CMPEQPD *xmm1, xmm2*	CMPPD *xmm1, xmm2, 0*
CMPLTPD *xmm1, xmm2*	CMPPD *xmm1, xmm2, 1*
CMPLEPD *xmm1, xmm2*	CMPPD *xmm1, xmm2, 2*
CMPUNORDPD *xmm1, xmm2*	CMPPD *xmm1, xmm2, 3*
CMPNEQPD *xmm1, xmm2*	CMPPD *xmm1, xmm2, 4*
CMPNLTPD *xmm1, xmm2*	CMPPD *xmm1, xmm2, 5*
CMPNLEPD *xmm1, xmm2*	CMPPD *xmm1, xmm2, 6*
CMPORDPD *xmm1, xmm2*	CMPPD *xmm1, xmm2, 7*

The greater-than relations that the processor does not implement, require more than one instruction to emulate in software and therefore should not be implemented as pseudo-ops. (For these, the programmer should reverse the operands of the corresponding less than relations and use move instructions to ensure that the mask is moved to the correct destination register and that the source operand is left intact.)

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Enhanced Comparison Predicate for VEX-Encoded VCMPPD

VEX.128 encoded version: The first source operand (second operand) is an XMM register. The second source operand (third operand) can be an XMM register or a 128-bit memory location. Bits (VLMAX-1:128) of the destination YMM register are zeroed. Two comparisons are performed with results written to bits 127:0 of the destination operand.

VEX.256 encoded version: The first source operand (second operand) is a YMM register. The second source operand (third operand) can be a YMM register or a 256-bit memory location. The destination operand (first operand) is a YMM register. Four comparisons are performed with results written to the destination operand.

The comparison predicate operand is an 8-bit immediate:

- For instructions encoded using the VEX prefix, bits 4:0 define the type of comparison to be performed (see Table 3-9). Bits 5 through 7 of the immediate are reserved.

Table 3-9. Comparison Predicate for VCMPPD and VCMPPS Instructions

Predicate	imm8 Value	Description	Result: A Is 1st Operand, B Is 2nd Operand				Signals #IA on QNAN
			A >B	A < B	A = B	Unordered[1]	
EQ_OQ (EQ)	0H	Equal (ordered, non-signaling)	False	False	True	False	No
LT_OS (LT)	1H	Less-than (ordered, signaling)	False	True	False	False	Yes
LE_OS (LE)	2H	Less-than-or-equal (ordered, signaling)	False	True	True	False	Yes
UNORD_Q (UNORD)	3H	Unordered (non-signaling)	False	False	False	True	No
NEQ_UQ (NEQ)	4H	Not-equal (unordered, non-signaling)	True	True	False	True	No
NLT_US (NLT)	5H	Not-less-than (unordered, signaling)	True	False	True	True	Yes
NLE_US (NLE)	6H	Not-less-than-or-equal (unordered, signaling)	True	False	False	True	Yes
ORD_Q (ORD)	7H	Ordered (non-signaling)	True	True	True	False	No
EQ_UQ	8H	Equal (unordered, non-signaling)	False	False	True	True	No
NGE_US (NGE)	9H	Not-greater-than-or-equal (unordered, signaling)	False	True	False	True	Yes
NGT_US (NGT)	AH	Not-greater-than (unordered, signal-ing)	False	True	True	True	Yes
FALSE_OQ(FALSE)	BH	False (ordered, non-signaling)	False	False	False	False	No
NEQ_OQ	CH	Not-equal (ordered, non-signaling)	True	True	False	False	No
GE_OS (GE)	DH	Greater-than-or-equal (ordered, sig-naling)	True	False	True	False	Yes

Table 3-9. Comparison Predicate for VCMPPD and VCMPPS Instructions (Contd.)

Predicate	imm8 Value	Description	Result: A Is 1st Operand, B Is 2nd Operand				Signals #IA on QNAN
			A >B	A < B	A = B	Unordered[1]	
GT_OS (GT)	EH	Greater-than (ordered, signaling)	True	False	False	False	Yes
TRUE_UQ(TRUE)	FH	True (unordered, non-signaling)	True	True	True	True	No
EQ_OS	10H	Equal (ordered, signaling)	False	False	True	False	Yes
LT_OQ	11H	Less-than (ordered, nonsignaling)	False	True	False	False	No
LE_OQ	12H	Less-than-or-equal (ordered, nonsignaling)	False	True	True	False	No
UNORD_S	13H	Unordered (signaling)	False	False	False	True	Yes
NEQ_US	14H	Not-equal (unordered, signaling)	True	True	False	True	Yes
NLT_UQ	15H	Not-less-than (unordered, nonsignaling)	True	False	True	True	No
NLE_UQ	16H	Not-less-than-or-equal (unordered, nonsignaling)	True	False	False	True	No
ORD_S	17H	Ordered (signaling)	True	True	True	False	Yes
EQ_US	18H	Equal (unordered, signaling)	False	False	True	True	Yes
NGE_UQ	19H	Not-greater-than-or-equal (unordered, nonsignaling)	False	True	False	True	No
NGT_UQ	1AH	Not-greater-than (unordered, nonsignaling)	False	True	True	True	No
FALSE_OS	1BH	False (ordered, signaling)	False	False	False	False	Yes
NEQ_OS	1CH	Not-equal (ordered, signaling)	True	True	False	False	Yes

Table 3-9. Comparison Predicate for VCMPPD and VCMPPS Instructions (Contd.)

Predicate	imm8 Value	Description	Result: A Is 1st Operand, B Is 2nd Operand				Signals #IA on QNAN
			A >B	A < B	A = B	Unordered[1]	
GE_OQ	1DH	Greater-than-or-equal (ordered, non-signaling)	True	False	True	False	No
GT_OQ	1EH	Greater-than (ordered, nonsignaling)	True	False	False	False	No
TRUE_US	1FH	True (unordered, signaling)	True	True	True	True	Yes

NOTES:

1. If either operand A or B is a NAN.

Processors with "CPUID.1H:ECX.AVX =1" implement the full complement of 32 predicates shown in Table 3-9, software emulation is no longer needed. Compilers and assemblers may implement the following three-operand pseudo-ops in addition to the four-operand VCMPPD instruction. See Table 3-10, where the notations of reg1 reg2, and reg3 represent either XMM registers or YMM registers. Compiler should treat reserved Imm8 values as illegal syntax. Alternately, intrinsics can map the pseudo-ops to pre-defined constants to support a simpler intrinsic interface.

Table 3-10. Pseudo-Op and VCMPPD Implementation

Pseudo-Op	CMPPD Implementation
VCMPEQPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 0*
VCMPLTPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 1*
VCMPLEPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 2*
VCMPUNORDPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 3*
VCMPNEQPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 4*
VCMPNLTPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 5*
VCMPNLEPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 6*
VCMPORDPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 7*
VCMPEQ_UQPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 8*
VCMPNGEPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 9*
VCMPNGTPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 0AH*
VCMPFALSEPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 0BH*
VCMPNEQ_OQPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 0CH*

Table 3-10. Pseudo-Op and VCMPPD Implementation

Pseudo-Op	CMPPD Implementation
VCMPGEPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 0DH*
VCMPGTPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 0EH*
VCMPTRUEPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 0FH*
VCMPEQ_OSPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 10H*
VCMPLT_OQPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 11H*
VCMPLE_OQPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 12H*
VCMPUNORD_SPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 13H*
VCMPNEQ_USPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 14H*
VCMPNLT_UQPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 15H*
VCMPNLE_UQPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 16H*
VCMPORD_SPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 17H*
VCMPEQ_USPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 18H*
VCMPNGE_UQPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 19H*
VCMPNGT_UQPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 1AH*
VCMPFALSE_OSPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 1BH*
VCMPNEQ_OSPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 1CH*
VCMPGE_OQPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 1DH*
VCMPGT_OQPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 1EH*
VCMPTRUE_USPD *reg1, reg2, reg3*	VCMPPD *reg1, reg2, reg3, 1FH*

Operation

```
CASE (COMPARISON PREDICATE) OF
    0: OP3 ← EQ_OQ; OP5 ← EQ_OQ;
    1: OP3 ← LT_OS; OP5 ← LT_OS;
    2: OP3 ← LE_OS; OP5 ← LE_OS;
    3: OP3 ← UNORD_Q; OP5 ← UNORD_Q;
    4: OP3 ← NEQ_UQ; OP5 ← NEQ_UQ;
    5: OP3 ← NLT_US; OP5 ← NLT_US;
    6: OP3 ← NLE_US; OP5 ← NLE_US;
    7: OP3 ← ORD_Q; OP5 ← ORD_Q;
    8: OP5 ← EQ_UQ;
    9: OP5 ← NGE_US;
```

```
    10: OP5 ← NGT_US;
    11: OP5 ← FALSE_OQ;
    12: OP5 ← NEQ_OQ;
    13: OP5 ← GE_OS;
    14: OP5 ← GT_OS;
    15: OP5 ← TRUE_UQ;
    16: OP5 ← EQ_OS;
    17: OP5 ← LT_OQ;
    18: OP5 ← LE_OQ;
    19: OP5 ← UNORD_S;
    20: OP5 ← NEQ_US;
    21: OP5 ← NLT_UQ;
    22: OP5 ← NLE_UQ;
    23: OP5 ← ORD_S;
    24: OP5 ← EQ_US;
    25: OP5 ← NGE_UQ;
    26: OP5 ← NGT_UQ;
    27: OP5 ← FALSE_OS;
    28: OP5 ← NEQ_OS;
    29: OP5 ← GE_OQ;
    30: OP5 ← GT_OQ;
    31: OP5 ← TRUE_US;
    DEFAULT: Reserved;
```

CMPPD (128-bit Legacy SSE version)

```
CMP0 ← SRC1[63:0] OP3 SRC2[63:0];
CMP1 ← SRC1[127:64] OP3 SRC2[127:64];
IF CMP0 = TRUE
    THEN DEST[63:0] ← FFFFFFFFFFFFFFFFH;
    ELSE DEST[63:0] ← 0000000000000000H; FI;
IF CMP1 = TRUE
    THEN DEST[127:64] ← FFFFFFFFFFFFFFFFH;
    ELSE DEST[127:64] ← 0000000000000000H; FI;
DEST[VLMAX-1:128] (Unmodified)
```

VCMPPD (VEX.128 encoded version)

```
CMP0 ← SRC1[63:0] OP5 SRC2[63:0];
CMP1 ← SRC1[127:64] OP5 SRC2[127:64];
IF CMP0 = TRUE
    THEN DEST[63:0] ← FFFFFFFFFFFFFFFFH;
    ELSE DEST[63:0] ← 0000000000000000H; FI;
IF CMP1 = TRUE
    THEN DEST[127:64] ← FFFFFFFFFFFFFFFFH;
```

```
    ELSE DEST[127:64] ← 0000000000000000H; FI;
DEST[VLMAX-1:128] ← 0
```

VCMPPD (VEX.256 encoded version)

```
CMP0 ← SRC1[63:0] OP5 SRC2[63:0];
CMP1 ← SRC1[127:64] OP5 SRC2[127:64];
CMP2 ← SRC1[191:128] OP5 SRC2[191:128];
CMP3 ← SRC1[255:192] OP5 SRC2[255:192];
IF CMP0 = TRUE
    THEN DEST[63:0] ← FFFFFFFFFFFFFFFFH;
    ELSE DEST[63:0] ← 0000000000000000H; FI;
IF CMP1 = TRUE
    THEN DEST[127:64] ← FFFFFFFFFFFFFFFFH;
    ELSE DEST[127:64] ← 0000000000000000H; FI;
IF CMP2 = TRUE
    THEN DEST[191:128] ← FFFFFFFFFFFFFFFFH;
    ELSE DEST[191:128] ← 0000000000000000H; FI;
IF CMP3 = TRUE
    THEN DEST[255:192] ← FFFFFFFFFFFFFFFFH;
    ELSE DEST[255:192] ← 0000000000000000H; FI;
```

Intel C/C++ Compiler Intrinsic Equivalents

CMPPD for equality __m128d _mm_cmpeq_pd(__m128d a, __m128d b)

CMPPD for less-than__m128d _mm_cmplt_pd(__m128d a, __m128d b)

CMPPD for less-than-or-equal__m128d _mm_cmple_pd(__m128d a, __m128d b)

CMPPD for greater-than__m128d _mm_cmpgt_pd(__m128d a, __m128d b)

CMPPD for greater-than-or-equal__m128d _mm_cmpge_pd(__m128d a, __m128d b)

CMPPD for inequality __m128d _mm_cmpneq_pd(__m128d a, __m128d b)

CMPPD for not-less-than __m128d _mm_cmpnlt_pd(__m128d a, __m128d b)

CMPPD for not-greater-than __m128d _mm_cmpngt_pd(__m128d a, __m128d b)

CMPPD for not-greater-than-or-equal__m128d _mm_cmpnge_pd(__m128d a, __m128d b)

CMPPD for ordered __m128d _mm_cmpord_pd(__m128d a, __m128d b)

CMPPD for unordered__m128d _mm_cmpunord_pd(__m128d a, __m128d b)

CMPPD for not-less-than-or-equal__m128d _mm_cmpnle_pd(__m128d a, __m128d b)

VCMPPD __m256 _mm256_cmp_pd(__m256 a, __m256 b, const int imm)

VCMPPD __m128 _mm_cmp_pd(__m128 a, __m128 b, const int imm)

SIMD Floating-Point Exceptions

Invalid if SNaN operand and invalid if QNaN and predicate as listed in above table, Denormal.

Other Exceptions

See Exceptions Type 2.

CMPPS—Compare Packed Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F C2 /*r* ib CMPPS *xmm1, xmm2/m128, imm8*	A	V/V	SSE	Compare packed single-precision floating-point values in *xmm2/mem* and *xmm1* using *imm8* as comparison predicate.
VEX.NDS.128.0F.WIG C2 /r ib VCMPPS xmm1, xmm2, xmm3/m128, imm8	B	V/V	AVX	Compare packed single-precision floating-point values in xmm3/m128 and xmm2 using bits 4:0 of imm8 as a comparison predicate.
VEX.NDS.256.0F.WIG C2 /r ib VCMPPS ymm1, ymm2, ymm3/m256, imm8	B	V/V	AVX	Compare packed single-precision floating-point values in ymm3/m256 and ymm2 using bits 2:0 of imm8 as a comparison predicate.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	imm8	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs a SIMD compare of the packed single-precision floating-point values in the source operand (second operand) and the destination operand (first operand) and returns the results of the comparison to the destination operand. The comparison predicate operand (third operand) specifies the type of comparison performed on each of the pairs of packed values. The result of each comparison is a doubleword mask of all 1s (comparison true) or all 0s (comparison false).

128-bit Legacy SSE version: The first source and destination operand (first operand) is an XMM register. The second source operand (second operand) can be an XMM register or 128-bit memory location. The comparison predicate operand is an 8-bit immediate, bits 2:0 of the immediate define the type of comparison to be performed (see Table 3-7). Bits 7:3 of the immediate is reserved. Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged. Four comparisons are performed with results written to bits 127:0 of the destination operand.

The unordered relationship is true when at least one of the two source operands being compared is a NaN; the ordered relationship is true when neither source operand is a NaN.

A subsequent computational instruction that uses the mask result in the destination operand as an input operand will not generate a fault, because a mask of all 0s corresponds to a floating-point value of +0.0 and a mask of all 1s corresponds to a QNaN.

Note that processors with "CPUID.1H:ECX.AVX =0" do not implement the "greater-than", "greater-than-or-equal", "not-greater than", and "not-greater-than-or-equal relations" predicates. These comparisons can be made either by using the inverse relationship (that is, use the "not-less-than-or-equal" to make a "greater-than" comparison) or by using software emulation. When using software emulation, the program must swap the operands (copying registers when necessary to protect the data that will now be in the destination), and then perform the compare using a different predicate. The predicate to be used for these emulations is listed in Table 3-7 under the heading Emulation.

Compilers and assemblers may implement the following two-operand pseudo-ops in addition to the three-operand CMPPS instruction, for processors with "CPUID.1H:ECX.AVX =0". See Table 3-11. Compiler should treat reserved Imm8 values as illegal syntax.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Table 3-11. Pseudo-Ops and CMPPS

Pseudo-Op	Implementation
CMPEQPS *xmm1, xmm2*	CMPPS *xmm1, xmm2, 0*
CMPLTPS *xmm1, xmm2*	CMPPS *xmm1, xmm2, 1*
CMPLEPS *xmm1, xmm2*	CMPPS *xmm1, xmm2, 2*
CMPUNORDPS *xmm1, xmm2*	CMPPS *xmm1, xmm2, 3*
CMPNEQPS *xmm1, xmm2*	CMPPS *xmm1, xmm2, 4*
CMPNLTPS *xmm1, xmm2*	CMPPS *xmm1, xmm2, 5*
CMPNLEPS *xmm1, xmm2*	CMPPS *xmm1, xmm2, 6*
CMPORDPS *xmm1, xmm2*	CMPPS *xmm1, xmm2, 7*

The greater-than relations not implemented by processor require more than one instruction to emulate in software and therefore should not be implemented as pseudo-ops. (For these, the programmer should reverse the operands of the corresponding less than relations and use move instructions to ensure that the mask is moved to the correct destination register and that the source operand is left intact.)

Enhanced Comparison Predicate for VEX-Encoded VCMPPS

VEX.128 encoded version: The first source operand (second operand) is an XMM register. The second source operand (third operand) can be an XMM register or a 128-bit memory location. Bits (VLMAX-1:128) of the destination YMM register are zeroed. Four comparisons are performed with results written to bits 127:0 of the destination operand.

VEX.256 encoded version: The first source operand (second operand) is a YMM register. The second source operand (third operand) can be a YMM register or a 256-bit memory location. The destination operand (first operand) is a YMM register. Eight comparisons are performed with results written to the destination operand.

The comparison predicate operand is an 8-bit immediate:

- For instructions encoded using the VEX prefix, bits 4:0 define the type of comparison to be performed (see Table 3-9). Bits 5 through 7 of the immediate are reserved.

Processors with "CPUID.1H:ECX.AVX =1" implement the full complement of 32 predicates shown in Table 3-9, software emulation is no longer needed. Compilers and assemblers may implement the following three-operand pseudo-ops in addition to the four-operand VCMPPS instruction. See Table 3-12, where the notation of reg1 and reg2 represent either XMM registers or YMM registers. Compiler should treat reserved Imm8 values as illegal syntax. Alternately, intrinsics can map the pseudo-ops to pre-defined constants to support a simpler intrinsic interface.

Table 3-12. Pseudo-Op and VCMPPS Implementation

Pseudo-Op	CMPPS Implementation
VCMPEQPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 0*
VCMPLTPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 1*
VCMPLEPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 2*
VCMPUNORDPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 3*
VCMPNEQPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 4*
VCMPNLTPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 5*
VCMPNLEPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 6*
VCMPORDPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 7*
VCMPEQ_UQPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 8*
VCMPNGEPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 9*
VCMPNGTPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 0AH*
VCMPFALSEPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 0BH*
VCMPNEQ_OQPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 0CH*
VCMPGEPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 0DH*
VCMPGTPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 0EH*

Table 3-12. Pseudo-Op and VCMPPS Implementation

Pseudo-Op	CMPPS Implementation
VCMPTRUEPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 0FH*
VCMPEQ_OSPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 10H*
VCMPLT_OQPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 11H*
VCMPLE_OQPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 12H*
VCMPUNORD_SPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 13H*
VCMPNEQ_USPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 14H*
VCMPNLT_UQPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 15H*
VCMPNLE_UQPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 16H*
VCMPORD_SPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 17H*
VCMPEQ_USPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 18H*
VCMPNGE_UQPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 19H*
VCMPNGT_UQPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 1AH*
VCMPFALSE_OSPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 1BH*
VCMPNEQ_OSPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 1CH*
VCMPGE_OQPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 1DH*
VCMPGT_OQPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 1EH*
VCMPTRUE_USPS *reg1, reg2, reg3*	VCMPPS *reg1, reg2, reg3, 1FH*

Operation

```
CASE (COMPARISON PREDICATE) OF
    0: OP3 ← EQ_OQ; OP5 ← EQ_OQ;
    1: OP3 ← LT_OS; OP5 ← LT_OS;
    2: OP3 ← LE_OS; OP5 ← LE_OS;
    3: OP3 ← UNORD_Q; OP5 ← UNORD_Q;
    4: OP3 ← NEQ_UQ; OP5 ← NEQ_UQ;
    5: OP3 ← NLT_US; OP5 ← NLT_US;
    6: OP3 ← NLE_US; OP5 ← NLE_US;
    7: OP3 ← ORD_Q; OP5 ← ORD_Q;
    8: OP5 ← EQ_UQ;
    9: OP5 ← NGE_US;
    10: OP5 ← NGT_US;
    11: OP5 ← FALSE_OQ;
    12: OP5 ← NEQ_OQ;
    13: OP5 ← GE_OS;
    14: OP5 ← GT_OS;
```

```
    15: OP5 ← TRUE_UQ;
    16: OP5 ← EQ_OS;
    17: OP5 ← LT_OQ;
    18: OP5 ← LE_OQ;
    19: OP5 ← UNORD_S;
    20: OP5 ← NEQ_US;
    21: OP5 ← NLT_UQ;
    22: OP5 ← NLE_UQ;
    23: OP5 ← ORD_S;
    24: OP5 ← EQ_US;
    25: OP5 ← NGE_UQ;
    26: OP5 ← NGT_UQ;
    27: OP5 ← FALSE_OS;
    28: OP5 ← NEQ_OS;
    29: OP5 ← GE_OQ;
    30: OP5 ← GT_OQ;
    31: OP5 ← TRUE_US;
    DEFAULT: Reserved
EASC;
```

CMPPS (128-bit Legacy SSE version)

```
CMP0 ← SRC1[31:0] OP3 SRC2[31:0];
CMP1 ← SRC1[63:32] OP3 SRC2[63:32];
CMP2 ← SRC1[95:64] OP3 SRC2[95:64];
CMP3 ← SRC1[127:96] OP3 SRC2[127:96];
IF CMP0 = TRUE
    THEN DEST[31:0] ←FFFFFFFFH;
    ELSE DEST[31:0] ← 000000000H; FI;
IF CMP1 = TRUE
    THEN DEST[63:32] ← FFFFFFFFH;
    ELSE DEST[63:32] ← 000000000H; FI;
IF CMP2 = TRUE
    THEN DEST[95:64] ← FFFFFFFFH;
    ELSE DEST[95:64] ← 000000000H; FI;
IF CMP3 = TRUE
    THEN DEST[127:96] ← FFFFFFFFH;
    ELSE DEST[127:96] ←000000000H; FI;
DEST[VLMAX-1:128] (Unmodified)
```

VCMPPS (VEX.128 encoded version)

```
CMP0 ← SRC1[31:0] OP5 SRC2[31:0];
CMP1 ← SRC1[63:32] OP5 SRC2[63:32];
CMP2 ← SRC1[95:64] OP5 SRC2[95:64];
```

```
CMP3 ← SRC1[127:96] OP5 SRC2[127:96];
IF CMP0 = TRUE
    THEN DEST[31:0] ←FFFFFFFFH;
    ELSE DEST[31:0] ← 000000000H; FI;
IF CMP1 = TRUE
    THEN DEST[63:32] ← FFFFFFFFH;
    ELSE DEST[63:32] ← 000000000H; FI;
IF CMP2 = TRUE
    THEN DEST[95:64] ← FFFFFFFFH;
    ELSE DEST[95:64] ← 000000000H; FI;
IF CMP3 = TRUE
    THEN DEST[127:96] ← FFFFFFFFH;
    ELSE DEST[127:96] ←000000000H; FI;
DEST[VLMAX-1:128] ← 0
```

VCMPPS (VEX.256 encoded version)

```
CMP0 ← SRC1[31:0] OP5 SRC2[31:0];
CMP1 ← SRC1[63:32] OP5 SRC2[63:32];
CMP2 ← SRC1[95:64] OP5 SRC2[95:64];
CMP3 ← SRC1[127:96] OP5 SRC2[127:96];
CMP4 ← SRC1[159:128] OP5 SRC2[159:128];
CMP5 ← SRC1[191:160] OP5 SRC2[191:160];
CMP6 ← SRC1[223:192] OP5 SRC2[223:192];
CMP7 ← SRC1[255:224] OP5 SRC2[255:224];
IF CMP0 = TRUE
    THEN DEST[31:0] ←FFFFFFFFH;
    ELSE DEST[31:0] ← 000000000H; FI;
IF CMP1 = TRUE
    THEN DEST[63:32] ← FFFFFFFFH;
    ELSE DEST[63:32] ←000000000H; FI;
IF CMP2 = TRUE
    THEN DEST[95:64] ← FFFFFFFFH;
    ELSE DEST[95:64] ← 000000000H; FI;
IF CMP3 = TRUE
    THEN DEST[127:96] ← FFFFFFFFH;
    ELSE DEST[127:96] ← 000000000H; FI;
IF CMP4 = TRUE
    THEN DEST[159:128] ← FFFFFFFFH;
    ELSE DEST[159:128] ← 000000000H; FI;
IF CMP5 = TRUE
    THEN DEST[191:160] ← FFFFFFFFH;
    ELSE DEST[191:160] ← 000000000H; FI;
IF CMP6 = TRUE
```

```
    THEN DEST[223:192] ← FFFFFFFFH;
    ELSE DEST[223:192] ←000000000H; FI;
IF CMP7 = TRUE
    THEN DEST[255:224] ← FFFFFFFFH;
    ELSE DEST[255:224] ← 000000000H; FI;
```

Intel C/C++ Compiler Intrinsic Equivalents

CMPPS for equality __m128 _mm_cmpeq_ps(__m128 a, __m128 b)

CMPPS for less-than__m128 _mm_cmplt_ps(__m128 a, __m128 b)

CMPPS for less-than-or-equal__m128 _mm_cmple_ps(__m128 a, __m128 b)

CMPPS for greater-than __m128 _mm_cmpgt_ps(__m128 a, __m128 b)

CMPPS for greater-than-or-equal__m128 _mm_cmpge_ps(__m128 a, __m128 b)

CMPPS for inequality __m128 _mm_cmpneq_ps(__m128 a, __m128 b)

CMPPS for not-less-than __m128 _mm_cmpnlt_ps(__m128 a, __m128 b)

CMPPS for not-greater-than __m128 _mm_cmpngt_ps(__m128 a, __m128 b)

CMPPS for not-greater-than-or-equal__m128 _mm_cmpnge_ps(__m128 a, __m128 b)

CMPPS for ordered__m128 _mm_cmpord_ps(__m128 a, __m128 b)

CMPPS for unordered__m128 _mm_cmpunord_ps(__m128 a, __m128 b)

CMPPS for not-less-than-or-equal__m128 _mm_cmpnle_ps(__m128 a, __m128 b)

VCMPPS __m256 _mm256_cmp_ps(__m256 a, __m256 b, const int imm)

VCMPPS __m128 _mm_cmp_ps(__m128 a, __m128 b, const int imm)

SIMD Floating-Point Exceptions

Invalid if SNaN operand and invalid if QNaN and predicate as listed in above table, Denormal.

Other Exceptions

See Exceptions Type 2.

CMPS/CMPSB/CMPSW/CMPSD/CMPSQ—Compare String Operands

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
A6	CMPS *m8, m8*	A	Valid	Valid	For legacy mode, compare byte at address DS:(E)SI with byte at address ES:(E)DI; For 64-bit mode compare byte at address (R|E)SI to byte at address (R|E)DI. The status flags are set accordingly.
A7	CMPS *m16, m16*	A	Valid	Valid	For legacy mode, compare word at address DS:(E)SI with word at address ES:(E)DI; For 64-bit mode compare word at address (R|E)SI with word at address (R|E)DI. The status flags are set accordingly.
A7	CMPS *m32, m32*	A	Valid	Valid	For legacy mode, compare dword at address DS:(E)SI at dword at address ES:(E)DI; For 64-bit mode compare dword at address (R|E)SI at dword at address (R|E)DI. The status flags are set accordingly.
REX.W + A7	CMPS *m64, m64*	A	Valid	N.E.	Compares quadword at address (R|E)SI with quadword at address (R|E)DI and sets the status flags accordingly.
A6	CMPSB	A	Valid	Valid	For legacy mode, compare byte at address DS:(E)SI with byte at address ES:(E)DI; For 64-bit mode compare byte at address (R|E)SI with byte at address (R|E)DI. The status flags are set accordingly.

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
A7	CMPSW	A	Valid	Valid	For legacy mode, compare word at address DS:(E)SI with word at address ES:(E)DI; For 64-bit mode compare word at address (R|E)SI with word at address (R|E)DI. The status flags are set accordingly.
A7	CMPSD	A	Valid	Valid	For legacy mode, compare dword at address DS:(E)SI with dword at address ES:(E)DI; For 64-bit mode compare dword at address (R|E)SI with dword at address (R|E)DI. The status flags are set accordingly.
REX.W + A7	CMPSQ	A	Valid	N.E.	Compares quadword at address (R|E)SI with quadword at address (R|E)DI and sets the status flags accordingly.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Compares the byte, word, doubleword, or quadword specified with the first source operand with the byte, word, doubleword, or quadword specified with the second source operand and sets the status flags in the EFLAGS register according to the results.

Both source operands are located in memory. The address of the first source operand is read from DS:SI, DS:ESI or RSI (depending on the address-size attribute of the instruction is 16, 32, or 64, respectively). The address of the second source operand is read from ES:DI, ES:EDI or RDI (again depending on the address-size attribute of the instruction is 16, 32, or 64). The DS segment may be overridden with a segment override prefix, but the ES segment cannot be overridden.

At the assembly-code level, two forms of this instruction are allowed: the “explicit-operands” form and the “no-operands” form. The explicit-operands form (specified with the CMPS mnemonic) allows the two source operands to be specified explicitly.

Here, the source operands should be symbols that indicate the size and location of the source values. This explicit-operand form is provided to allow documentation. However, note that the documentation provided by this form can be misleading. That is, the source operand symbols must specify the correct type (size) of the operands (bytes, words, or doublewords, quadwords), but they do not have to specify the correct location. Locations of the source operands are always specified by the DS:(E)SI (or RSI) and ES:(E)DI (or RDI) registers, which must be loaded correctly before the compare string instruction is executed.

The no-operands form provides "short forms" of the byte, word, and doubleword versions of the CMPS instructions. Here also the DS:(E)SI (or RSI) and ES:(E)DI (or RDI) registers are assumed by the processor to specify the location of the source operands. The size of the source operands is selected with the mnemonic: CMPSB (byte comparison), CMPSW (word comparison), CMPSD (doubleword comparison), or CMPSQ (quadword comparison using REX.W).

After the comparison, the (E/R)SI and (E/R)DI registers increment or decrement automatically according to the setting of the DF flag in the EFLAGS register. (If the DF flag is 0, the (E/R)SI and (E/R)DI register increment; if the DF flag is 1, the registers decrement.) The registers increment or decrement by 1 for byte operations, by 2 for word operations, 4 for doubleword operations. If operand size is 64, RSI and RDI registers increment by 8 for quadword operations.

The CMPS, CMPSB, CMPSW, CMPSD, and CMPSQ instructions can be preceded by the REP prefix for block comparisons. More often, however, these instructions will be used in a LOOP construct that takes some action based on the setting of the status flags before the next comparison is made. See "REP/REPE/REPZ /REPNE/REPNZ—Repeat String Operation Prefix" in Chapter 4 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 2B*, for a description of the REP prefix.

In 64-bit mode, the instruction's default address size is 64 bits, 32 bit address size is supported using the prefix 67H. Use of the REX.W prefix promotes doubleword operation to 64 bits (see CMPSQ). See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
temp ← SRC1 - SRC2;
SetStatusFlags(temp);

IF (64-Bit Mode)
    THEN
        IF (Byte comparison)
        THEN IF DF = 0
            THEN
                (R|E)SI ← (R|E)SI + 1;
                (R|E)DI ← (R|E)DI + 1;
            ELSE
```

```
                (R|E)SI ← (R|E)SI – 1;
                (R|E)DI ← (R|E)DI – 1;
        FI;
    ELSE IF (Word comparison)
        THEN IF DF = 0
            THEN
                (R|E)SI ← (R|E)SI + 2;
                (R|E)DI ← (R|E)DI + 2;
            ELSE
                (R|E)SI ← (R|E)SI – 2;
                (R|E)DI ← (R|E)DI – 2;
            FI;
    ELSE IF (Doubleword comparison)
        THEN IF DF = 0
            THEN
                (R|E)SI ← (R|E)SI + 4;
                (R|E)DI ← (R|E)DI + 4;
            ELSE
                (R|E)SI ← (R|E)SI – 4;
                (R|E)DI ← (R|E)DI – 4;
            FI;
    ELSE (* Quadword comparison *)
        THEN IF DF = 0
            (R|E)SI ← (R|E)SI + 8;
            (R|E)DI ← (R|E)DI + 8;
        ELSE
            (R|E)SI ← (R|E)SI – 8;
            (R|E)DI ← (R|E)DI – 8;
        FI;
    FI;
ELSE (* Non-64-bit Mode *)
    IF (byte comparison)
    THEN IF DF = 0
        THEN
            (E)SI ← (E)SI + 1;
            (E)DI ← (E)DI + 1;
        ELSE
            (E)SI ← (E)SI – 1;
            (E)DI ← (E)DI – 1;
        FI;
    ELSE IF (Word comparison)
        THEN IF DF = 0
            (E)SI ← (E)SI + 2;
```

```
                (E)DI ← (E)DI + 2;
            ELSE
                (E)SI ← (E)SI – 2;
                (E)DI ← (E)DI – 2;
            FI;
        ELSE (* Doubleword comparison *)
            THEN IF DF = 0
                (E)SI ← (E)SI + 4;
                (E)DI ← (E)DI + 4;
            ELSE
                (E)SI ← (E)SI – 4;
                (E)DI ← (E)DI – 4;
            FI;
        FI;
FI;
```

Flags Affected

The CF, OF, SF, ZF, AF, and PF flags are set according to the temporary result of the comparison.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

CMPSD—Compare Scalar Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F C2 /*r* ib CMPSD *xmm1*, *xmm2/m64*, *imm8*	A	V/V	SSE2	Compare low double-precision floating-point value in *xmm2/m64* and *xmm1* using *imm8* as comparison predicate.
VEX.NDS.LIG.F2.0F.WIG C2 /r ib VCMPSD xmm1, xmm2, xmm3/m64, imm8	B	V/V	AVX	Compare low double precision floating-point value in xmm3/m64 and xmm2 using bits 4:0 of imm8 as comparison predicate.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	imm8	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Compares the low double-precision floating-point values in the source operand (second operand) and the destination operand (first operand) and returns the results of the comparison to the destination operand. The comparison predicate operand (third operand) specifies the type of comparison performed. The comparison result is a quadword mask of all 1s (comparison true) or all 0s (comparison false).

128-bit Legacy SSE version: The first source and destination operand (first operand) is an XMM register. The second source operand (second operand) can be an XMM register or 64-bit memory location. The comparison predicate operand is an 8-bit immediate, bits 2:0 of the immediate define the type of comparison to be performed (see Table 3-7). Bits 7:3 of the immediate is reserved. Bits (VLMAX-1:64) of the corresponding YMM destination register remain unchanged.

The unordered relationship is true when at least one of the two source operands being compared is a NaN; the ordered relationship is true when neither source operand is a NaN.

A subsequent computational instruction that uses the mask result in the destination operand as an input operand will not generate a fault, because a mask of all 0s corresponds to a floating-point value of +0.0 and a mask of all 1s corresponds to a QNaN.

Note that processors with “CPUID.1H:ECX.AVX =0” do not implement the “greater-than”, “greater-than-or-equal”, “not-greater than”, and “not-greater-than-or-equal relations" predicates. These comparisons can be made either by using the inverse relationship (that is, use the "not-less-than-or-equal" to make a "greater-than" comparison) or by using software emulation. When using software emulation, the program must swap the operands (copying registers when necessary to protect the data that will now be in the destination operand), and then perform the compare using a different predicate. The predicate to be used for these emulations is listed in Table 3-7 under the heading Emulation.

Compilers and assemblers may implement the following two-operand pseudo-ops in addition to the three-operand CMPSD instruction, for processors with "CPUID.1H:ECX.AVX =0". See Table 3-13. Compiler should treat reserved Imm8 values as illegal syntax.

Table 3-13. Pseudo-Ops and CMPSD

Pseudo-Op	Implementation
CMPEQSD *xmm1, xmm2*	CMPSD *xmm1,xmm2, 0*
CMPLTSD *xmm1, xmm2*	CMPSD *xmm1,xmm2, 1*
CMPLESD *xmm1, xmm2*	CMPSD *xmm1,xmm2, 2*
CMPUNORDSD *xmm1, xmm2*	CMPSD *xmm1,xmm2, 3*
CMPNEQSD *xmm1, xmm2*	CMPSD *xmm1,xmm2, 4*
CMPNLTSD *xmm1, xmm2*	CMPSD *xmm1,xmm2, 5*
CMPNLESD *xmm1, xmm2*	CMPSD *xmm1,xmm2, 6*
CMPORDSD *xmm1, xmm2*	CMPSD *xmm1,xmm2, 7*

The greater-than relations not implemented in the processor require more than one instruction to emulate in software and therefore should not be implemented as pseudo-ops. (For these, the programmer should reverse the operands of the corresponding less than relations and use move instructions to ensure that the mask is moved to the correct destination register and that the source operand is left intact.)

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Enhanced Comparison Predicate for VEX-Encoded VCMPSD

VEX.128 encoded version: The first source operand (second operand) is an XMM register. The second source operand (third operand) can be an XMM register or a 64-bit memory location. Bits (VLMAX-1:128) of the destination YMM register are zeroed. The comparison predicate operand is an 8-bit immediate:

- For instructions encoded using the VEX prefix, bits 4:0 define the type of comparison to be performed (see Table 3-9). Bits 5 through 7 of the immediate are reserved.

Processors with "CPUID.1H:ECX.AVX =1" implement the full complement of 32 predicates shown in Table 3-9, software emulation is no longer needed. Compilers and assemblers may implement the following three-operand pseudo-ops in addition to the four-operand VCMPSD instruction. See Table 3-14, where the notations of reg1 reg2, and reg3 represent either XMM registers or YMM registers. Compiler should treat reserved Imm8 values as illegal syntax. Alternately, intrinsics can map the pseudo-ops to pre-defined constants to support a simpler intrinsic interface.

Table 3-14. Pseudo-Op and VCMPSD Implementation

Pseudo-Op	CMPSD Implementation
VCMPEQSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 0*
VCMPLTSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 1*
VCMPLESD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 2*
VCMPUNORDSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 3*
VCMPNEQSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 4*
VCMPNLTSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 5*
VCMPNLESD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 6*
VCMPORDSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 7*
VCMPEQ_UQSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 8*
VCMPNGESD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 9*
VCMPNGTSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 0AH*
VCMPFALSESD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 0BH*
VCMPNEQ_OQSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 0CH*
VCMPGESD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 0DH*
VCMPGTSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 0EH*
VCMPTRUESD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 0FH*
VCMPEQ_OSSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 10H*
VCMPLT_OQSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 11H*
VCMPLE_OQSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 12H*
VCMPUNORD_SSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 13H*
VCMPNEQ_USSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 14H*
VCMPNLT_UQSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 15H*
VCMPNLE_UQSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 16H*
VCMPORD_SSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 17H*
VCMPEQ_USSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 18H*
VCMPNGE_UQSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 19H*
VCMPNGT_UQSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 1AH*

Table 3-14. Pseudo-Op and VCMPSD Implementation

Pseudo-Op	CMPSD Implementation
VCMPFALSE_OSSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 1BH*
VCMPNEQ_OSSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 1CH*
VCMPGE_OQSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 1DH*
VCMPGT_OQSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 1EH*
VCMPTRUE_USSD *reg1, reg2, reg3*	VCMPSD *reg1, reg2, reg3, 1FH*

Operation

```
CASE (COMPARISON PREDICATE) OF
    0: OP3 ← EQ_OQ; OP5 ← EQ_OQ;
    1: OP3 ← LT_OS; OP5 ← LT_OS;
    2: OP3 ← LE_OS; OP5 ← LE_OS;
    3: OP3 ← UNORD_Q; OP5 ← UNORD_Q;
    4: OP3 ← NEQ_UQ; OP5 ← NEQ_UQ;
    5: OP3 ← NLT_US; OP5 ← NLT_US;
    6: OP3 ← NLE_US; OP5 ← NLE_US;
    7: OP3 ← ORD_Q; OP5 ← ORD_Q;
    8: OP5 ← EQ_UQ;
    9: OP5 ← NGE_US;
    10: OP5 ← NGT_US;
    11: OP5 ← FALSE_OQ;
    12: OP5 ← NEQ_OQ;
    13: OP5 ← GE_OS;
    14: OP5 ← GT_OS;
    15: OP5 ← TRUE_UQ;
    16: OP5 ← EQ_OS;
    17: OP5 ← LT_OQ;
    18: OP5 ← LE_OQ;
    19: OP5 ← UNORD_S;
    20: OP5 ← NEQ_US;
    21: OP5 ← NLT_UQ;
    22: OP5 ← NLE_UQ;
    23: OP5 ← ORD_S;
    24: OP5 ← EQ_US;
    25: OP5 ← NGE_UQ;
    26: OP5 ← NGT_UQ;
    27: OP5 ← FALSE_OS;
    28: OP5 ← NEQ_OS;
```

```
    29: OP5 ← GE_OQ;
    30: OP5 ← GT_OQ;
    31: OP5 ← TRUE_US;
    DEFAULT: Reserved
ESAC;
```

CMPSD (128-bit Legacy SSE version)

```
CMP0 ← DEST[63:0] OP3 SRC[63:0];
IF CMP0 = TRUE
THEN DEST[63:0] ← FFFFFFFFFFFFFFFFH;
ELSE DEST[63:0] ← 0000000000000000H; FI;
DEST[VLMAX-1:64] (Unmodified)
```

VCMPSD (VEX.128 encoded version)

```
CMP0 ← SRC1[63:0] OP5 SRC2[63:0];
IF CMP0 = TRUE
THEN DEST[63:0] ← FFFFFFFFFFFFFFFFH;
ELSE DEST[63:0] ← 0000000000000000H; FI;
DEST[127:64] ← SRC1[127:64]
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalents

CMPSD for equality__m128d _mm_cmpeq_sd(__m128d a, __m128d b)

CMPSD for less-than__m128d _mm_cmplt_sd(__m128d a, __m128d b)

CMPSD for less-than-or-equal__m128d _mm_cmple_sd(__m128d a, __m128d b)

CMPSD for greater-than__m128d _mm_cmpgt_sd(__m128d a, __m128d b)

CMPSD for greater-than-or-equal__m128d _mm_cmpge_sd(__m128d a, __m128d b)

CMPSD for inequality__m128d _mm_cmpneq_sd(__m128d a, __m128d b)

CMPSD for not-less-than__m128d _mm_cmpnlt_sd(__m128d a, __m128d b)

CMPSD for not-greater-than__m128d _mm_cmpngt_sd(__m128d a, __m128d b)

CMPSD for not-greater-than-or-equal__m128d _mm_cmpnge_sd(__m128d a, __m128d b)

CMPSD for ordered__m128d _mm_cmpord_sd(__m128d a, __m128d b)

CMPSD for unordered__m128d _mm_cmpunord_sd(__m128d a, __m128d b)

CMPSD for not-less-than-or-equal__m128d _mm_cmpnle_sd(__m128d a, __m128d b)

VCMPSD __m128 _mm_cmp_sd(__m128 a, __m128 b, const int imm)

SIMD Floating-Point Exceptions

Invalid if SNaN operand, Invalid if QNaN and predicate as listed in above table, Denormal.

Other Exceptions

See Exceptions Type 3.

CMPSS—Compare Scalar Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F C2 /r ib CMPSS *xmm1, xmm2/m32, imm8*	A	V/V	SSE	Compare low single-precision floating-point value in *xmm2/m32* and *xmm1* using *imm8* as comparison predicate.
VEX.NDS.LIG.F3.0F.WIG C2 /r ib VCMPSS xmm1, xmm2, xmm3/m32, imm8	B	V/V	AVX	Compare low single precision floating-point value in xmm3/m32 and xmm2 using bits 4:0 of imm8 as comparison predicate.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	imm8	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Compares the low single-precision floating-point values in the source operand (second operand) and the destination operand (first operand) and returns the results of the comparison to the destination operand. The comparison predicate operand (third operand) specifies the type of comparison performed. The comparison result is a doubleword mask of all 1s (comparison true) or all 0s (comparison false).

128-bit Legacy SSE version: The first source and destination operand (first operand) is an XMM register. The second source operand (second operand) can be an XMM register or 64-bit memory location. The comparison predicate operand is an 8-bit immediate, bits 2:0 of the immediate define the type of comparison to be performed (see Table 3-7). Bits 7:3 of the immediate is reserved. Bits (VLMAX-1:32) of the corresponding YMM destination register remain unchanged.

The unordered relationship is true when at least one of the two source operands being compared is a NaN; the ordered relationship is true when neither source operand is a NaN

A subsequent computational instruction that uses the mask result in the destination operand as an input operand will not generate a fault, since a mask of all 0s corresponds to a floating-point value of +0.0 and a mask of all 1s corresponds to a QNaN.

Note that processors with “CPUID.1H:ECX.AVX =0” do not implement the “greater-than”, “greater-than-or-equal”, “not-greater than”, and “not-greater-than-or-equal relations" predicates. These comparisons can be made either by using the inverse relationship (that is, use the "not-less-than-or-equal" to make a "greater-than" comparison) or by using software emulation. When using software emulation, the program must swap the operands (copying registers when necessary to protect the data that will now be in the destination operand), and then perform the compare using a different predicate. The predicate to be used for these emulations is listed in Table 3-7 under the heading Emulation.

Compilers and assemblers may implement the following two-operand pseudo-ops in addition to the three-operand CMPSS instruction, for processors with "CPUID.1H:ECX.AVX =0". See Table 3-15. Compiler should treat reserved Imm8 values as illegal syntax.

Table 3-15. Pseudo-Ops and CMPSS

Pseudo-Op	CMPSS Implementation
CMPEQSS *xmm1, xmm2*	CMPSS *xmm1, xmm2, 0*
CMPLTSS *xmm1, xmm2*	CMPSS *xmm1, xmm2, 1*
CMPLESS *xmm1, xmm2*	CMPSS *xmm1, xmm2, 2*
CMPUNORDSS *xmm1, xmm2*	CMPSS *xmm1, xmm2, 3*
CMPNEQSS *xmm1, xmm2*	CMPSS *xmm1, xmm2, 4*
CMPNLTSS *xmm1, xmm2*	CMPSS *xmm1, xmm2, 5*
CMPNLESS *xmm1, xmm2*	CMPSS *xmm1, xmm2, 6*
CMPORDSS *xmm1, xmm2*	CMPSS *xmm1, xmm2, 7*

The greater-than relations not implemented in the processor require more than one instruction to emulate in software and therefore should not be implemented as pseudo-ops. (For these, the programmer should reverse the operands of the corresponding less than relations and use move instructions to ensure that the mask is moved to the correct destination register and that the source operand is left intact.)

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Enhanced Comparison Predicate for VEX-Encoded VCMPSD

VEX.128 encoded version: The first source operand (second operand) is an XMM register. The second source operand (third operand) can be an XMM register or a 32-bit memory location. Bits (VLMAX-1:128) of the destination YMM register are zeroed. The comparison predicate operand is an 8-bit immediate:

- For instructions encoded using the VEX prefix, bits 4:0 define the type of comparison to be performed (see Table 3-9). Bits 5 through 7 of the immediate are reserved.

Processors with "CPUID.1H:ECX.AVX =1" implement the full complement of 32 predicates shown in Table 3-9, software emulation is no longer needed. Compilers and assemblers may implement the following three-operand pseudo-ops in addition to the four-operand VCMPSS instruction. See Table 3-16, where the notations of reg1 reg2, and reg3 represent either XMM registers or YMM registers. Compiler should treat reserved Imm8 values as illegal syntax. Alternately, intrinsics can map the pseudo-ops to pre-defined constants to support a simpler intrinsic interface.

Table 3-16. Pseudo-Op and VCMPSS Implementation

Pseudo-Op	CMPSS Implementation
VCMPEQSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 0*
VCMPLTSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 1*
VCMPLESS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 2*
VCMPUNORDSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 3*
VCMPNEQSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 4*
VCMPNLTSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 5*
VCMPNLESS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 6*
VCMPORDSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 7*
VCMPEQ_UQSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 8*
VCMPNGESS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 9*
VCMPNGTSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 0AH*
VCMPFALSESS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 0BH*
VCMPNEQ_OQSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 0CH*
VCMPGESS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 0DH*
VCMPGTSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 0EH*
VCMPTRUESS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 0FH*
VCMPEQ_OSSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 10H*
VCMPLT_OQSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 11H*
VCMPLE_OQSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 12H*
VCMPUNORD_SSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 13H*
VCMPNEQ_USSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 14H*
VCMPNLT_UQSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 15H*
VCMPNLE_UQSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 16H*
VCMPORD_SSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 17H*
VCMPEQ_USSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 18H*
VCMPNGE_UQSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 19H*
VCMPNGT_UQSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 1AH*

Table 3-16. Pseudo-Op and VCMPSS Implementation

Pseudo-Op	CMPSS Implementation
VCMPFALSE_OSSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 1BH*
VCMPNEQ_OSSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 1CH*
VCMPGE_OQSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 1DH*
VCMPGT_OQSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 1EH*
VCMPTRUE_USSS *reg1, reg2, reg3*	VCMPSS *reg1, reg2, reg3, 1FH*

Operation

```
CASE (COMPARISON PREDICATE) OF
    0: OP3 ← EQ_OQ; OP5 ← EQ_OQ;
    1: OP3 ← LT_OS; OP5 ← LT_OS;
    2: OP3 ← LE_OS; OP5 ← LE_OS;
    3: OP3 ← UNORD_Q; OP5 ← UNORD_Q;
    4: OP3 ← NEQ_UQ; OP5 ← NEQ_UQ;
    5: OP3 ← NLT_US; OP5 ← NLT_US;
    6: OP3 ← NLE_US; OP5 ← NLE_US;
    7: OP3 ← ORD_Q; OP5 ← ORD_Q;
    8: OP5 ← EQ_UQ;
    9: OP5 ← NGE_US;
    10: OP5 ← NGT_US;
    11: OP5 ← FALSE_OQ;
    12: OP5 ← NEQ_OQ;
    13: OP5 ← GE_OS;
    14: OP5 ← GT_OS;
    15: OP5 ← TRUE_UQ;
    16: OP5 ← EQ_OS;
    17: OP5 ← LT_OQ;
    18: OP5 ← LE_OQ;
    19: OP5 ← UNORD_S;
    20: OP5 ← NEQ_US;
    21: OP5 ← NLT_UQ;
    22: OP5 ← NLE_UQ;
    23: OP5 ← ORD_S;
    24: OP5 ← EQ_US;
    25: OP5 ← NGE_UQ;
    26: OP5 ← NGT_UQ;
    27: OP5 ← FALSE_OS;
    28: OP5 ← NEQ_OS;
    29: OP5 ← GE_OQ;
    30: OP5 ← GT_OQ;
```

```
    31: OP5 ← TRUE_US;
    DEFAULT: Reserved
ESAC;
```

CMPSS (128-bit Legacy SSE version)

```
CMP0 ← DEST[31:0] OP3 SRC[31:0];
IF CMP0 = TRUE
THEN DEST[31:0] ← FFFFFFFFH;
ELSE DEST[31:0] ← 00000000H; FI;
DEST[VLMAX-1:32] (Unmodified)
```

VCMPSS (VEX.128 encoded version)

```
CMP0 ← SRC1[31:0] OP5 SRC2[31:0];
IF CMP0 = TRUE
THEN DEST[31:0] ← FFFFFFFFH;
ELSE DEST[31:0] ← 00000000H; FI;
DEST[127:32] ← SRC1[127:32]
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalents

CMPSS for equality__m128 _mm_cmpeq_ss(__m128 a, __m128 b)

CMPSS for less-than__m128 _mm_cmplt_ss(__m128 a, __m128 b)

CMPSS for less-than-or-equal__m128 _mm_cmple_ss(__m128 a, __m128 b)

CMPSS for greater-than__m128 _mm_cmpgt_ss(__m128 a, __m128 b)

CMPSS for greater-than-or-equal__m128 _mm_cmpge_ss(__m128 a, __m128 b)

CMPSS for inequality__m128 _mm_cmpneq_ss(__m128 a, __m128 b)

CMPSS for not-less-than__m128 _mm_cmpnlt_ss(__m128 a, __m128 b)

CMPSS for not-greater-than__m128 _mm_cmpngt_ss(__m128 a, __m128 b)

CMPSS for not-greater-than-or-equal__m128 _mm_cmpnge_ss(__m128 a, __m128 b)

CMPSS for ordered__m128 _mm_cmpord_ss(__m128 a, __m128 b)

CMPSS for unordered__m128 _mm_cmpunord_ss(__m128 a, __m128 b)

CMPSS for not-less-than-or-equal__m128 _mm_cmpnle_ss(__m128 a, __m128 b)

VCMPSS __m128 _mm_cmp_ss(__m128 a, __m128 b, const int imm)

SIMD Floating-Point Exceptions

Invalid if SNaN operand, Invalid if QNaN and predicate as listed in above table, Denormal.

Other Exceptions

See Exceptions Type 3.

CMPXCHG—Compare and Exchange

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F B0/*r*	CMPXCHG *r/m8, r8*	A	Valid	Valid*	Compare AL with *r/m8*. If equal, ZF is set and *r8* is loaded into *r/m8*. Else, clear ZF and load *r/m8* into AL.
REX + 0F B0/*r*	CMPXCHG *r/m8***,*r8*	A	Valid	N.E.	Compare AL with *r/m8*. If equal, ZF is set and *r8* is loaded into *r/m8*. Else, clear ZF and load *r/m8* into AL.
0F B1/*r*	CMPXCHG *r/m16, r16*	A	Valid	Valid*	Compare AX with *r/m16*. If equal, ZF is set and *r16* is loaded into *r/m16*. Else, clear ZF and load *r/m16* into AX.
0F B1/*r*	CMPXCHG *r/m32, r32*	A	Valid	Valid*	Compare EAX with *r/m32*. If equal, ZF is set and *r32* is loaded into *r/m32*. Else, clear ZF and load *r/m32* into EAX.
REX.W + 0F B1/*r*	CMPXCHG *r/m64, r64*	A	Valid	N.E.	Compare RAX with *r/m64*. If equal, ZF is set and *r64* is loaded into *r/m64*. Else, clear ZF and load *r/m64* into RAX.

NOTES:

* See the IA-32 Architecture Compatibility section below.

** In 64-bit mode, *r/m8* can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r, w)	ModRM:reg (r)	NA	NA

Description

Compares the value in the AL, AX, EAX, or RAX register with the first operand (destination operand). If the two values are equal, the second operand (source operand) is loaded into the destination operand. Otherwise, the destination operand is loaded into the AL, AX, EAX or RAX register. RAX register is available only in 64-bit mode.

This instruction can be used with a LOCK prefix to allow the instruction to be executed atomically. To simplify the interface to the processor's bus, the destination operand receives a write cycle without regard to the result of the comparison. The destination operand is written back if the comparison fails; otherwise, the source operand is written into the destination. (The processor never produces a locked read without also producing a locked write.)

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

IA-32 Architecture Compatibility

This instruction is not supported on Intel processors earlier than the Intel486 processors.

Operation

```
(* Accumulator = AL, AX, EAX, or RAX depending on whether a byte, word, doubleword, or
quadword comparison is being performed *)

IF accumulator = DEST
    THEN
        ZF ← 1;
        DEST ← SRC;
    ELSE
        ZF ← 0;
        accumulator ← DEST;
FI;
```

Flags Affected

The ZF flag is set if the values in the destination operand and register AL, AX, or EAX are equal; otherwise it is cleared. The CF, PF, AF, SF, and OF flags are set according to the results of the comparison operation.

Protected Mode Exceptions

#GP(0)	If the destination is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.

#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

CMPXCHG8B/CMPXCHG16B—Compare and Exchange Bytes

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F C7 /1 *m64*	CMPXCHG8B *m64*	A	Valid	Valid*	Compare EDX:EAX with *m64*. If equal, set ZF and load ECX:EBX into *m64*. Else, clear ZF and load *m64* into EDX:EAX.
REX.W + 0F C7 /1 *m128*	CMPXCHG16B *m128*	A	Valid	N.E.	Compare RDX:RAX with *m128*. If equal, set ZF and load RCX:RBX into *m128*. Else, clear ZF and load *m128* into RDX:RAX.

NOTES:

*See IA-32 Architecture Compatibility section below.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r, w)	NA	NA	NA

Description

Compares the 64-bit value in EDX:EAX (or 128-bit value in RDX:RAX if operand size is 128 bits) with the operand (destination operand). If the values are equal, the 64-bit value in ECX:EBX (or 128-bit value in RCX:RBX) is stored in the destination operand. Otherwise, the value in the destination operand is loaded into EDX:EAX (or RDX:RAX). The destination operand is an 8-byte memory location (or 16-byte memory location if operand size is 128 bits). For the EDX:EAX and ECX:EBX register pairs, EDX and ECX contain the high-order 32 bits and EAX and EBX contain the low-order 32 bits of a 64-bit value. For the RDX:RAX and RCX:RBX register pairs, RDX and RCX contain the high-order 64 bits and RAX and RBX contain the low-order 64bits of a 128-bit value.

This instruction can be used with a LOCK prefix to allow the instruction to be executed atomically. To simplify the interface to the processor's bus, the destination operand receives a write cycle without regard to the result of the comparison. The destination operand is written back if the comparison fails; otherwise, the source operand is written into the destination. (The processor never produces a locked read without also producing a locked write.)

In 64-bit mode, default operation size is 64 bits. Use of the REX.W prefix promotes operation to 128 bits. Note that CMPXCHG16B requires that the destination (memory) operand be 16-byte aligned. See the summary chart at the beginning of this section for encoding data and limits. For information on the CPUID flag that indicates CMPXCHG16B, see page 3-216.

IA-32 Architecture Compatibility

This instruction encoding is not supported on Intel processors earlier than the Pentium processors.

Operation

```
IF (64-Bit Mode and OperandSize = 64)
    THEN
        IF (RDX:RAX = DEST)
        ZF ← 1;
            DEST ← RCX:RBX;
        ELSE
            ZF ← 0;
            RDX:RAX ← DEST;
        FI
    ELSE
        IF (EDX:EAX = DEST)
            ZF ← 1;
            DEST ← ECX:EBX;
        ELSE
            ZF ← 0;
            EDX:EAX ← DEST;
        FI;
FI;
```

Flags Affected

The ZF flag is set if the destination operand and EDX:EAX are equal; otherwise it is cleared. The CF, PF, AF, SF, and OF flags are unaffected.

Protected Mode Exceptions

#UD	If the destination is not a memory operand.
#GP(0)	If the destination is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

Real-Address Mode Exceptions

#UD	If the destination operand is not a memory location.
#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.

Virtual-8086 Mode Exceptions

#UD	If the destination operand is not a memory location.
#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
	If memory operand for CMPXCHG16B is not aligned on a 16-byte boundary.
	If CPUID.01H:ECX.CMPXCHG16B[bit 13] = 0.
#UD	If the destination operand is not a memory location.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

COMISD—Compare Scalar Ordered Double-Precision Floating-Point Values and Set EFLAGS

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 2F /*r* COMISD *xmm1*, *xmm2/m64*	A	V/V	SSE2	Compare low double-precision floating-point values in *xmm1* and *xmm2/mem64* and set the EFLAGS flags accordingly.
VEX.LIG.66.0F.WIG 2F /r VCOMISD xmm1, xmm2/m64	A	V/V	AVX	Compare low double precision floating-point values in xmm1 and xmm2/mem64 and set the EFLAGS flags accordingly.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r)	ModRM:r/m (r)	NA	NA

Description

Compares the double-precision floating-point values in the low quadwords of operand 1 (first operand) and operand 2 (second operand), and sets the ZF, PF, and CF flags in the EFLAGS register according to the result (unordered, greater than, less than, or equal). The OF, SF and AF flags in the EFLAGS register are set to 0. The unordered result is returned if either source operand is a NaN (QNaN or SNaN).

Operand 1 is an XMM register; operand 2 can be an XMM register or a 64 bit memory location.

The COMISD instruction differs from the UCOMISD instruction in that it signals a SIMD floating-point invalid operation exception (#I) when a source operand is either a QNaN or SNaN. The UCOMISD instruction signals an invalid numeric exception only if a source operand is an SNaN.

The EFLAGS register is not updated if an unmasked SIMD floating-point exception is generated.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b, otherwise instructions will #UD.

Operation

```
RESULT ← OrderedCompare(DEST[63:0] <> SRC[63:0]) {
(* Set EFLAGS *) CASE (RESULT) OF
    UNORDERED:          ZF,PF,CF ← 111;
    GREATER_THAN:       ZF,PF,CF ← 000;
    LESS_THAN:          ZF,PF,CF ← 001;
    EQUAL:              ZF,PF,CF ← 100;
ESAC;
OF, AF, SF ← 0; }
```

Intel C/C++ Compiler Intrinsic Equivalents

```
int _mm_comieq_sd (__m128d a, __m128d b)
int _mm_comilt_sd (__m128d a, __m128d b)
int _mm_comile_sd (__m128d a, __m128d b)
int _mm_comigt_sd (__m128d a, __m128d b)
int _mm_comige_sd (__m128d a, __m128d b)
int _mm_comineq_sd (__m128d a, __m128d b)
```

SIMD Floating-Point Exceptions

Invalid (if SNaN or QNaN operands), Denormal.

Other Exceptions

See Exceptions Type 3; additionally

#UD	If VEX.vvvv != 1111B.

COMISS—Compare Scalar Ordered Single-Precision Floating-Point Values and Set EFLAGS

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 2F /*r* COMISS *xmm1, xmm2/m32*	A	V/V	SSE	Compare low single-precision floating-point values in *xmm1* and *xmm2/mem32* and set the EFLAGS flags accordingly.
VEX.LIG.0F 2F.WIG /r VCOMISS xmm1, xmm2/m32	A	V/V	AVX	Compare low single precision floating-point values in xmm1 and xmm2/mem32 and set the EFLAGS flags accordingly.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r)	ModRM:r/m (r)	NA	NA

Description

Compares the single-precision floating-point values in the low doublewords of operand 1 (first operand) and operand 2 (second operand), and sets the ZF, PF, and CF flags in the EFLAGS register according to the result (unordered, greater than, less than, or equal). The OF, SF, and AF flags in the EFLAGS register are set to 0. The unordered result is returned if either source operand is a NaN (QNaN or SNaN).

Operand 1 is an XMM register; Operand 2 can be an XMM register or a 32 bit memory location.

The COMISS instruction differs from the UCOMISS instruction in that it signals a SIMD floating-point invalid operation exception (#I) when a source operand is either a QNaN or SNaN. The UCOMISS instruction signals an invalid numeric exception only if a source operand is an SNaN.

The EFLAGS register is not updated if an unmasked SIMD floating-point exception is generated.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b, otherwise instructions will #UD.

Operation

```
RESULT ← OrderedCompare(SRC1[31:0] <> SRC2[31:0]) {
(* Set EFLAGS *) CASE (RESULT) OF
    UNORDERED:          ZF,PF,CF ← 111;
    GREATER_THAN:       ZF,PF,CF ← 000;
    LESS_THAN:          ZF,PF,CF ← 001;
    EQUAL:              ZF,PF,CF ← 100;
ESAC;
OF,AF,SF ← 0; }
```

Intel C/C++ Compiler Intrinsic Equivalents

int _mm_comieq_ss (__m128 a, __m128 b)

int _mm_comilt_ss (__m128 a, __m128 b)

int _mm_comile_ss (__m128 a, __m128 b)

int _mm_comigt_ss (__m128 a, __m128 b)

int _mm_comige_ss (__m128 a, __m128 b)

int _mm_comineq_ss (__m128 a, __m128 b)

SIMD Floating-Point Exceptions

Invalid (if SNaN or QNaN operands), Denormal.

Other Exceptions

See Exceptions Type 3; additionally

#UD	If VEX.vvvv != 1111B.

CPUID—CPU Identification

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F A2	CPUID	A	Valid	Valid	Returns processor identification and feature information to the EAX, EBX, ECX, and EDX registers, as determined by input entered in EAX (in some cases, ECX as well).

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

The ID flag (bit 21) in the EFLAGS register indicates support for the CPUID instruction. If a software procedure can set and clear this flag, the processor executing the procedure supports the CPUID instruction. This instruction operates the same in non-64-bit modes and 64-bit mode.

CPUID returns processor identification and feature information in the EAX, EBX, ECX, and EDX registers.[1] The instruction's output is dependent on the contents of the EAX register upon execution (in some cases, ECX as well). For example, the following pseudocode loads EAX with 00H and causes CPUID to return a Maximum Return Value and the Vendor Identification String in the appropriate registers:

```
MOV EAX, 00H
CPUID
```

Table 3-17 shows information returned, depending on the initial value loaded into the EAX register. Table 3-18 shows the maximum CPUID input value recognized for each family of IA-32 processors on which CPUID is implemented.

Two types of information are returned: basic and extended function information. If a value entered for CPUID.EAX is higher than the maximum input value for basic or extended function for that processor then the data for the highest basic information leaf is returned. For example, using the Intel Core i7 processor, the following is true:

```
CPUID.EAX = 05H (* Returns MONITOR/MWAIT leaf. *)
CPUID.EAX = 0AH (* Returns Architectural Performance Monitoring leaf. *)
CPUID.EAX = 0BH (* Returns Extended Topology Enumeration leaf. *)
```

1. On Intel 64 processors, CPUID clears the high 32 bits of the RAX/RBX/RCX/RDX registers in all modes.

CPUID.EAX = 0CH (* INVALID: Returns the same information as CPUID.EAX = 0BH. *)
CPUID.EAX = 80000008H (* Returns linear/physical address size data. *)
CPUID.EAX = 8000000AH (* INVALID: Returns same information as CPUID.EAX = 0BH. *)

If a value entered for CPUID.EAX is less than or equal to the maximum input value and the leaf is not supported on that processor then 0 is returned in all the registers. For example, using the Intel Core i7 processor, the following is true:

CPUID.EAX = 07H (*Returns EAX=EBX=ECX=EDX=0. *)

When CPUID returns the highest basic leaf information as a result of an invalid input EAX value, any dependence on input ECX value in the basic leaf is honored.

CPUID can be executed at any privilege level to serialize instruction execution. Serializing instruction execution guarantees that any modifications to flags, registers, and memory for previous instructions are completed before the next instruction is fetched and executed.

See also:

"Serializing Instructions" in Chapter 8, "Multiple-Processor Management," in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*

"Caching Translation Information" in Chapter 4, "Paging," in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*.

Table 3-17. Information Returned by CPUID Instruction

Initial EAX Value	Information Provided about the Processor	
	Basic CPUID Information	
0H	EAX EBX ECX EDX	Maximum Input Value for Basic CPUID Information (see Table 3-18) "Genu" "ntel" "inel"
01H	EAX	Version Information: Type, Family, Model, and Stepping ID (see Figure 3-9)
	EBX	Bits 07-00: Brand Index Bits 15-08: CLFLUSH line size (Value ∗ 8 = cache line size in bytes) Bits 23-16: Maximum number of addressable IDs for logical processors in this physical package*. Bits 31-24: Initial APIC ID
	ECX EDX	Feature Information (see Figure 3-10 and Table 3-20) Feature Information (see Figure 3-11 and Table 3-21)
		NOTES: * The nearest power-of-2 integer that is not smaller than EBX[23:16] is the number of unique initial APIC IDs reserved for addressing different logical processors in a physical package.

Table 3-17. Information Returned by CPUID Instruction (Contd.)

Initial EAX Value	Information Provided about the Processor	
02H	EAX EBX ECX EDX	Cache and TLB Information (see Table 3-22) Cache and TLB Information Cache and TLB Information Cache and TLB Information
03H	EAX EBX ECX EDX	Reserved. Reserved. Bits 00-31 of 96 bit processor serial number. (Available in Pentium III processor only; otherwise, the value in this register is reserved.) Bits 32-63 of 96 bit processor serial number. (Available in Pentium III processor only; otherwise, the value in this register is reserved.) **NOTES:** Processor serial number (PSN) is not supported in the Pentium 4 processor or later. On all models, use the PSN flag (returned using CPUID) to check for PSN support before accessing the feature. See AP-485, *Intel Processor Identification and the CPUID Instruction* (Order Number 241618) for more information on PSN.
	CPUID leaves > 3 < 80000000 are visible only when IA32_MISC_ENABLE.BOOT_NT4[bit 22] = 0 (default).	
	Deterministic Cache Parameters Leaf	
04H		**NOTES:** Leaf 04H output depends on the initial value in ECX. See also: "INPUT EAX = 4: Returns Deterministic Cache Parameters for each level on page 3-227.
	EAX	Bits 04-00: Cache Type Field 0 = Null - No more caches 1 = Data Cache 2 = Instruction Cache 3 = Unified Cache 4-31 = Reserved Bits 07-05: Cache Level (starts at 1) Bit 08: Self Initializing cache level (does not need SW initialization) Bit 09: Fully Associative cache Bits 13-10: Reserved Bits 25-14: Maximum number of addressable IDs for logical processors sharing this cache*, ** Bits 31-26: Maximum number of addressable IDs for processor cores in the physical package*, ***, ****

Table 3-17. Information Returned by CPUID Instruction (Contd.)

Initial EAX Value	Information Provided about the Processor	
	EBX	Bits 11-00: L = System Coherency Line Size* Bits 21-12: P = Physical Line partitions* Bits 31-22: W = Ways of associativity*
	ECX	Bits 31-00: S = Number of Sets*
	EDX	Bit 0: Write-Back Invalidate/Invalidate 0 = WBINVD/INVD from threads sharing this cache acts upon lower level caches for threads sharing this cache. 1 = WBINVD/INVD is not guaranteed to act upon lower level caches of non-originating threads sharing this cache. Bit 1: Cache Inclusiveness 0 = Cache is not inclusive of lower cache levels. 1 = Cache is inclusive of lower cache levels. Bit 2: Complex Cache Indexing 0 = Direct mapped cache. 1 = A complex function is used to index the cache, potentially using all address bits. Bits 31-03: Reserved = 0
		NOTES: * Add one to the return value to get the result. ** The nearest power-of-2 integer that is not smaller than (1 + EAX[25:14]) is the number of unique initial APIC IDs reserved for addressing different logical processors sharing this cache *** The nearest power-of-2 integer that is not smaller than (1 + EAX[31:26]) is the number of unique Core_IDs reserved for addressing different processor cores in a physical package. Core ID is a subset of bits of the initial APIC ID. ****The returned value is constant for valid initial values in ECX. Valid ECX values start from 0.
	MONITOR/MWAIT Leaf	
05H	EAX	Bits 15-00: Smallest monitor-line size in bytes (default is processor's monitor granularity) Bits 31-16: Reserved = 0
	EBX	Bits 15-00: Largest monitor-line size in bytes (default is processor's monitor granularity) Bits 31-16: Reserved = 0

Table 3-17. Information Returned by CPUID Instruction (Contd.)

Initial EAX Value	Information Provided about the Processor	
	ECX	Bit 00: Enumeration of Monitor-Mwait extensions (beyond EAX and EBX registers) supported Bit 01: Supports treating interrupts as break-event for MWAIT, even when interrupts disabled Bits 31 - 02: Reserved
	EDX	Bits 03 - 00: Number of C0* sub C-states supported using MWAIT Bits 07 - 04: Number of C1* sub C-states supported using MWAIT Bits 11 - 08: Number of C2* sub C-states supported using MWAIT Bits 15 - 12: Number of C3* sub C-states supported using MWAIT Bits 19 - 16: Number of C4* sub C-states supported using MWAIT Bits 31 - 20: Reserved = 0 **NOTE:** * The definition of C0 through C4 states for MWAIT extension are processor-specific C-states, not ACPI C-states.
	Thermal and Power Management Leaf	
06H	EAX	Bit 00: Digital temperature sensor is supported if set Bit 01: Intel Turbo Boost Technology Available (see description of IA32_MISC_ENABLE[38]). Bit 02: ARAT. APIC-Timer-always-running feature is supported if set. Bit 03: Reserved Bit 04: PLN. Power limit notification controls are supported if set. Bit 05: ECMD. Clock modulation duty cycle extension is supported if set. Bit 06: PTM. Package thermal management is supported if set. Bits 31 - 07: Reserved
	EBX	Bits 03 - 00: Number of Interrupt Thresholds in Digital Thermal Sensor Bits 31 - 04: Reserved
	ECX	Bit 00: Hardware Coordination Feedback Capability (Presence of IA32_MPERF and IA32_APERF). The capability to provide a measure of delivered processor performance (since last reset of the counters), as a percentage of expected processor performance at frequency specified in CPUID Brand String Bits 02 - 01: Reserved = 0 Bit 03: The processor supports performance-energy bias preference if CPUID.06H:ECX.SETBH[bit 3] is set and it also implies the presence of a new architectural MSR called IA32_ENERGY_PERF_BIAS (1B0H) Bits 31 - 04: Reserved = 0
	EDX	Reserved = 0
	Direct Cache Access Information Leaf	

Table 3-17. Information Returned by CPUID Instruction (Contd.)

Initial EAX Value	Information Provided about the Processor	
09H	EAX	Value of bits [31:0] of IA32_PLATFORM_DCA_CAP MSR (address 1F8H)
	EBX	Reserved
	ECX	Reserved
	EDX	Reserved
	Architectural Performance Monitoring Leaf	
0AH	EAX	Bits 07 - 00: Version ID of architectural performance monitoring Bits 15- 08: Number of general-purpose performance monitoring counter per logical processor Bits 23 - 16: Bit width of general-purpose, performance monitoring counter Bits 31 - 24: Length of EBX bit vector to enumerate architectural performance monitoring events
	EBX	Bit 00: Core cycle event not available if 1 Bit 01: Instruction retired event not available if 1 Bit 02: Reference cycles event not available if 1 Bit 03: Last-level cache reference event not available if 1 Bit 04: Last-level cache misses event not available if 1 Bit 05: Branch instruction retired event not available if 1 Bit 06: Branch mispredict retired event not available if 1 Bits 31- 07: Reserved = 0
	ECX	Reserved = 0
	EDX	Bits 04 - 00: Number of fixed-function performance counters (if Version ID > 1) Bits 12- 05: Bit width of fixed-function performance counters (if Version ID > 1) Reserved = 0
	Extended Topology Enumeration Leaf	
0BH		**NOTES:** Most of Leaf 0BH output depends on the initial value in ECX. EDX output do not vary with initial value in ECX. ECX[7:0] output always reflect initial value in ECX. All other output value for an invalid initial value in ECX are 0. Leaf 0BH exists if EBX[15:0] is not zero.

Table 3-17. Information Returned by CPUID Instruction (Contd.)

Initial EAX Value	Information Provided about the Processor	
	EAX	Bits 04-00: Number of bits to shift right on x2APIC ID to get a unique topology ID of the next level type*. All logical processors with the same next level ID share current level. Bits 31-05: Reserved.
	EBX	Bits 15 - 00: Number of logical processors at this level type. The number reflects configuration as shipped by Intel**. Bits 31- 16: Reserved.
	ECX	Bits 07 - 00: Level number. Same value in ECX input Bits 15 - 08: Level type***. Bits 31 - 16:: Reserved.
	EDX	Bits 31- 00: x2APIC ID the current logical processor.
		NOTES: * Software should use this field (EAX[4:0]) to enumerate processor topology of the system. ** Software must not use EBX[15:0] to enumerate processor topology of the system. This value in this field (EBX[15:0]) is only intended for display/diagnostic purposes. The actual number of logical processors available to BIOS/OS/Applications may be different from the value of EBX[15:0], depending on software and platform hardware configurations. *** The value of the “level type” field is not related to level numbers in any way, higher “level type” values do not mean higher levels. Level type field has the following encoding: 0 : invalid 1 : SMT 2 : Core 3-255 : Reserved
Processor Extended State Enumeration Main Leaf (EAX = 0DH, ECX = 0)		
0DH		**NOTES:** Leaf 0DH main leaf (ECX = 0).
	EAX	Bits 31-00: Reports the valid bit fields of the lower 32 bits of XCR0. If a bit is 0, the corresponding bit field in XCR0 is reserved.
	EBX	Bits 31-00: Maximum size (bytes, from the beginning of the XSAVE/XRSTOR save area) required by enabled features in XCR0. May be different than ECX if some features at the end of the XSAVE save area are not enabled.

Table 3-17. Information Returned by CPUID Instruction (Contd.)

Initial EAX Value	Information Provided about the Processor	
	ECX	Bit 31-00: Maximum size (bytes, from the beginning of the XSAVE/XRSTOR save area) of the XSAVE/XRSTOR save area required by all supported features in the processor, i.e all the valid bit fields in XCR0.
	EDX	Bit 31-00: Reports the valid bit fields of the upper 32 bits of XCR0. If a bit is 0, the corresponding bit field in XCR0 is reserved.
	Processor Extended State Enumeration Sub-leaf (EAX = 0DH, ECX = 1)	
	EAX	Reserved
	EBX	Reserved
	ECX	Reserved
	EDX	Reserved
	Processor Extended State Enumeration Sub-leaves (EAX = 0DH, ECX = n, n > 1)	
0DH		**NOTES:** Leaf 0DH output depends on the initial value in ECX. If ECX contains an invalid sub leaf index, EAX/EBX/ECX/EDX return 0. Each valid sub-leaf index maps to a valid bit in the XCR0 register starting at bit position 2
	EAX	Bits 31-0: The size in bytes (from the offset specified in EBX) of the save area for an extended state feature associated with a valid sub-leaf index, *n*. This field reports 0 if the sub-leaf index, *n*, is invalid*.
	EBX	Bits 31-0: The offset in bytes of this extended state component's save area from the beginning of the XSAVE/XRSTOR area. This field reports 0 if the sub-leaf index, *n*, is invalid*.
	ECX	This field reports 0 if the sub-leaf index, *n*, is invalid*; otherwise it is reserved.
	EDX	This field reports 0 if the sub-leaf index, *n*, is invalid*; otherwise it is reserved.
	Unimplemented CPUID Leaf Functions	
40000000H - 4FFFFFFFH		Invalid. No existing or future CPU will return processor identification or feature information if the initial EAX value is in the range 40000000H to 4FFFFFFFH.
	Extended Function CPUID Information	
80000000H	EAX	Maximum Input Value for Extended Function CPUID Information (see Table 3-18).

Table 3-17. Information Returned by CPUID Instruction (Contd.)

Initial EAX Value	Information Provided about the Processor	
	EBX ECX EDX	Reserved Reserved Reserved
80000001H	EAX	Extended Processor Signature and Feature Bits.
	EBX	Reserved
	ECX	Bit 00: LAHF/SAHF available in 64-bit mode Bits 31-01 Reserved
	EDX	Bits 10-00: Reserved Bit 11: SYSCALL/SYSRET available (when in 64-bit mode) Bits 19-12: Reserved = 0 Bit 20: Execute Disable Bit available Bits 25-21: Reserved = 0 Bit 26: 1-GByte pages are available if 1 Bit 27: RDTSCP and IA32_TSC_AUX are available if 1 Bits 28: Reserved = 0 Bit 29: Intel® 64 Architecture available if 1 Bits 31-30: Reserved = 0
80000002H	EAX EBX ECX EDX	Processor Brand String Processor Brand String Continued Processor Brand String Continued Processor Brand String Continued
80000003H	EAX EBX ECX EDX	Processor Brand String Continued Processor Brand String Continued Processor Brand String Continued Processor Brand String Continued
80000004H	EAX EBX ECX EDX	Processor Brand String Continued Processor Brand String Continued Processor Brand String Continued Processor Brand String Continued
80000005H	EAX EBX ECX EDX	Reserved = 0 Reserved = 0 Reserved = 0 Reserved = 0
80000006H	EAX EBX	Reserved = 0 Reserved = 0

Table 3-17. Information Returned by CPUID Instruction (Contd.)

Initial EAX Value	Information Provided about the Processor	
	ECX	Bits 07-00: Cache Line size in bytes Bits 11-08: Reserved Bits 15-12: L2 Associativity field * Bits 31-16: Cache size in 1K units
	EDX	Reserved = 0
		NOTES: * L2 associativity field encodings: 00H - Disabled 01H - Direct mapped 02H - 2-way 04H - 4-way 06H - 8-way 08H - 16-way 0FH - Fully associative
80000007H	EAX	Reserved = 0
	EBX	Reserved = 0
	ECX	Reserved = 0
	EDX	Bits 07-00: Reserved = 0 Bit 08: Invariant TSC available if 1 Bits 31-09: Reserved = 0
80000008H	EAX	Linear/Physical Address size Bits 07-00: #Physical Address Bits* Bits 15-8: #Linear Address Bits Bits 31-16: Reserved = 0
	EBX	Reserved = 0
	ECX	Reserved = 0
	EDX	Reserved = 0
		NOTES: * If CPUID.80000008H:EAX[7:0] is supported, the maximum physical address number supported should come from this field.

INPUT EAX = 0: Returns CPUID's Highest Value for Basic Processor Information and the Vendor Identification String

When CPUID executes with EAX set to 0, the processor returns the highest value the CPUID recognizes for returning basic processor information. The value is returned in the EAX register (see Table 3-18) and is processor specific.

A vendor identification string is also returned in EBX, EDX, and ECX. For Intel processors, the string "GenuineIntel" and is expressed:

EBX ← 756e6547h (* "Genu", with G in the low eight bits of BL *)
EDX ← 49656e69h (* "ineI", with i in the low eight bits of DL *)
ECX ← 6c65746eh (* "ntel", with n in the low eight bits of CL *)

INPUT EAX = 80000000H: Returns CPUID's Highest Value for Extended Processor Information

When CPUID executes with EAX set to 80000000H, the processor returns the highest value the processor recognizes for returning extended processor information. The value is returned in the EAX register (see Table 3-18) and is processor specific.

Table 3-18. Highest CPUID Source Operand for Intel 64 and IA-32 Processors

Intel 64 or IA-32 Processors	Highest Value in EAX	
	Basic Information	Extended Function Information
Earlier Intel486 Processors	CPUID Not Implemented	CPUID Not Implemented
Later Intel486 Processors and Pentium Processors	01H	Not Implemented
Pentium Pro and Pentium II Processors, Intel® Celeron® Processors	02H	Not Implemented
Pentium III Processors	03H	Not Implemented
Pentium 4 Processors	02H	80000004H
Intel Xeon Processors	02H	80000004H
Pentium M Processor	02H	80000004H
Pentium 4 Processor supporting Hyper-Threading Technology	05H	80000008H
Pentium D Processor (8xx)	05H	80000008H
Pentium D Processor (9xx)	06H	80000008H
Intel Core Duo Processor	0AH	80000008H
Intel Core 2 Duo Processor	0AH	80000008H
Intel Xeon Processor 3000, 5100, 5200, 5300, 5400 Series	0AH	80000008H
Intel Core 2 Duo Processor 8000 Series	0DH	80000008H
Intel Xeon Processor 5200, 5400 Series	0AH	80000008H

Table 3-18. Highest CPUID Source Operand for Intel 64 and IA-32 Processors

Intel 64 or IA-32 Processors	Highest Value in EAX	
	Basic Information	Extended Function Information
Intel Atom Processor	0AH	80000008H
Intel Core i7 Processor	0BH	80000008H

IA32_BIOS_SIGN_ID Returns Microcode Update Signature

For processors that support the microcode update facility, the IA32_BIOS_SIGN_ID MSR is loaded with the update signature whenever CPUID executes. The signature is returned in the upper DWORD. For details, see Chapter 9 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*.

INPUT EAX = 1: Returns Model, Family, Stepping Information

When CPUID executes with EAX set to 1, version information is returned in EAX (see Figure 3-9). For example: model, family, and processor type for the Intel Xeon processor 5100 series is as follows:

- Model — 1111B
- Family — 0101B
- Processor Type — 00B

See Table 3-19 for available processor type values. Stepping IDs are provided as needed.

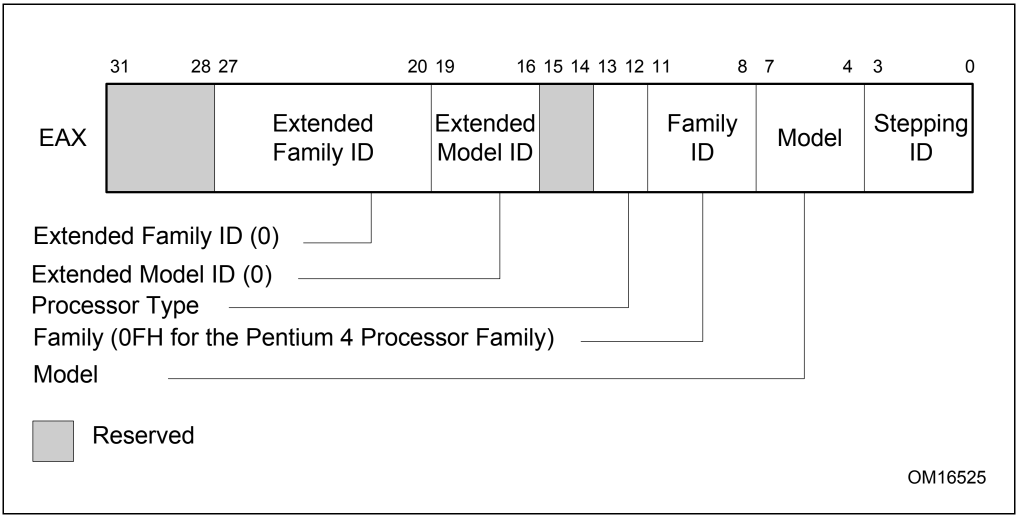


Figure 3-9. Version Information Returned by CPUID in EAX

Table 3-19. Processor Type Field

Type	Encoding
Original OEM Processor	00B
Intel OverDrive® Processor	01B
Dual processor (not applicable to Intel486 processors)	10B
Intel reserved	11B

NOTE

See Chapter 14 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for information on identifying earlier IA-32 processors.

The Extended Family ID needs to be examined only when the Family ID is 0FH. Integrate the fields into a display using the following rule:

```
IF Family_ID ≠ 0FH
    THEN DisplayFamily = Family_ID;
    ELSE DisplayFamily = Extended_Family_ID + Family_ID;
    (* Right justify and zero-extend 4-bit field. *)
FI;
(* Show DisplayFamily as HEX field. *)
```

The Extended Model ID needs to be examined only when the Family ID is 06H or 0FH. Integrate the field into a display using the following rule:

```
IF (Family_ID = 06H or Family_ID = 0FH)
    THEN DisplayModel = (Extended_Model_ID « 4) + Model_ID;
    (* Right justify and zero-extend 4-bit field; display Model_ID as HEX field.*)
    ELSE DisplayModel = Model_ID;
FI;
(* Show DisplayModel as HEX field. *)
```

INPUT EAX = 1: Returns Additional Information in EBX

When CPUID executes with EAX set to 1, additional information is returned to the EBX register:

- Brand index (low byte of EBX) — this number provides an entry into a brand string table that contains brand strings for IA-32 processors. More information about this field is provided later in this section.
- CLFLUSH instruction cache line size (second byte of EBX) — this number indicates the size of the cache line flushed with CLFLUSH instruction in 8-byte increments. This field was introduced in the Pentium 4 processor.

- Local APIC ID (high byte of EBX) — this number is the 8-bit ID that is assigned to the local APIC on the processor during power up. This field was introduced in the Pentium 4 processor.

INPUT EAX = 1: Returns Feature Information in ECX and EDX

When CPUID executes with EAX set to 1, feature information is returned in ECX and EDX.

- Figure 3-10 and Table 3-20 show encodings for ECX.
- Figure 3-11 and Table 3-21 show encodings for EDX.

For all feature flags, a 1 indicates that the feature is supported. Use Intel to properly interpret feature flags.

NOTE

Software must confirm that a processor feature is present using feature flags returned by CPUID prior to using the feature. Software should not depend on future offerings retaining all features.

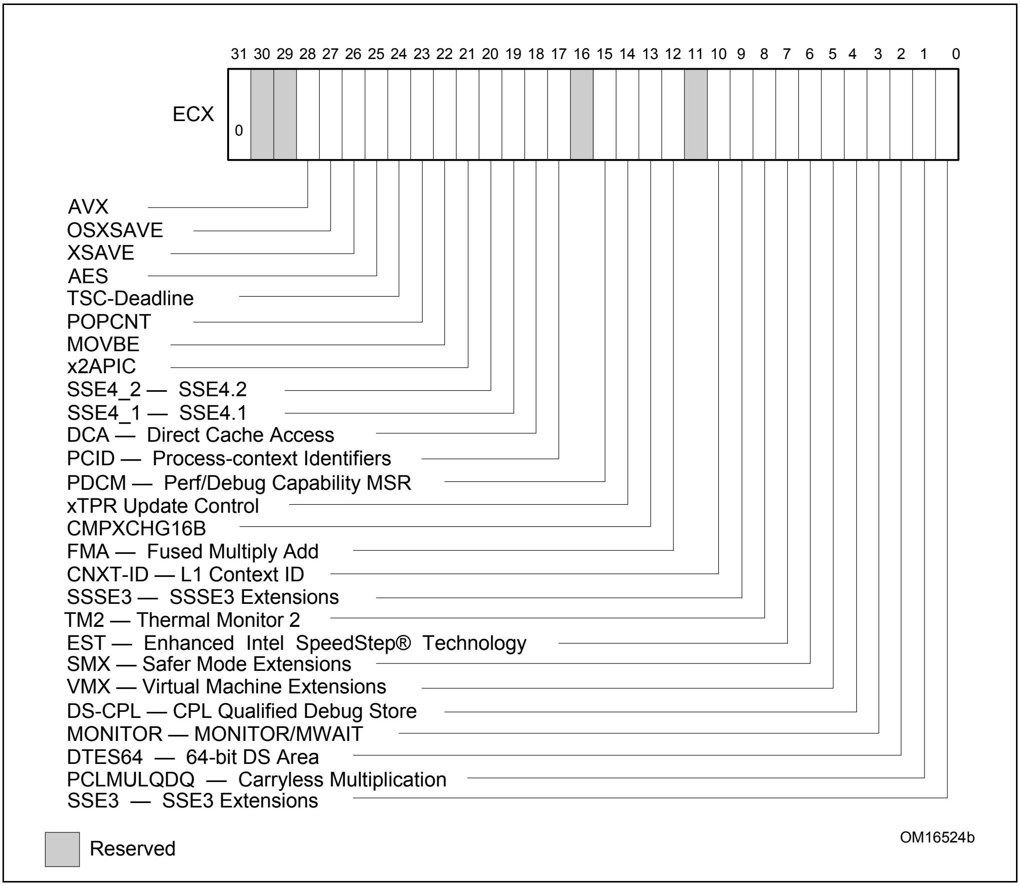


Figure 3-10. Feature Information Returned in the ECX Register

Table 3-20. Feature Information Returned in the ECX Register

Bit #	Mnemonic	Description
0	SSE3	**Streaming SIMD Extensions 3 (SSE3).** A value of 1 indicates the processor supports this technology.
1	PCLMULQDQ	**PCLMULQDQ.** A value of 1 indicates the processor supports the PCLMULQDQ instruction
2	DTES64	**64-bit DS Area.** A value of 1 indicates the processor supports DS area using 64-bit layout
3	MONITOR	**MONITOR/MWAIT.** A value of 1 indicates the processor supports this feature.

Table 3-20. Feature Information Returned in the ECX Register (Contd.)

Bit #	Mnemonic	Description
4	DS-CPL	**CPL Qualified Debug Store.** A value of 1 indicates the processor supports the extensions to the Debug Store feature to allow for branch message storage qualified by CPL.
5	VMX	**Virtual Machine Extensions.** A value of 1 indicates that the processor supports this technology
6	SMX	**Safer Mode Extensions.** A value of 1 indicates that the processor supports this technology. See Chapter 6, “Safer Mode Extensions Reference”.
7	EIST	**Enhanced Intel SpeedStep® technology.** A value of 1 indicates that the processor supports this technology.
8	TM2	**Thermal Monitor 2.** A value of 1 indicates whether the processor supports this technology.
9	SSSE3	A value of 1 indicates the presence of the Supplemental Streaming SIMD Extensions 3 (SSSE3). A value of 0 indicates the instruction extensions are not present in the processor
10	CNXT-ID	**L1 Context ID.** A value of 1 indicates the L1 data cache mode can be set to either adaptive mode or shared mode. A value of 0 indicates this feature is not supported. See definition of the IA32_MISC_ENABLE MSR Bit 24 (L1 Data Cache Context Mode) for details.
11	Reserved	Reserved
12	FMA	A value of 1 indicates the processor supports FMA extensions using YMM state.
13	CMPXCHG16B	**CMPXCHG16B Available.** A value of 1 indicates that the feature is available. See the “CMPXCHG8B/CMPXCHG16B—Compare and Exchange Bytes” section in this chapter for a description.
14	xTPR Update Control	**xTPR Update Control.** A value of 1 indicates that the processor supports changing IA32_MISC_ENABLE[bit 23].
15	PDCM	**Perfmon and Debug Capability:** A value of 1 indicates the processor supports the performance and debug feature indication MSR IA32_PERF_CAPABILITIES.
16	Reserved	Reserved
17	PCID	**Process-context identifiers.** A value of 1 indicates that the processor supports PCIDs and that software may set CR4.PCIDE to 1.
18	DCA	A value of 1 indicates the processor supports the ability to prefetch data from a memory mapped device.
19	SSE4.1	A value of 1 indicates that the processor supports SSE4.1.
20	SSE4.2	A value of 1 indicates that the processor supports SSE4.2.

Table 3-20. Feature Information Returned in the ECX Register (Contd.)

Bit #	Mnemonic	Description
21	x2APIC	A value of 1 indicates that the processor supports x2APIC feature.
22	MOVBE	A value of 1 indicates that the processor supports MOVBE instruction.
23	POPCNT	A value of 1 indicates that the processor supports the POPCNT instruction.
24	TSC-Deadline	A value of 1 indicates that the processor's local APIC timer supports one-shot operation using a TSC deadline value.
25	AESNI	A value of 1 indicates that the processor supports the AESNI instruction extensions.
26	XSAVE	A value of 1 indicates that the processor supports the XSAVE/XRSTOR processor extended states feature, the XSETBV/XGETBV instructions, and XCR0.
27	OSXSAVE	A value of 1 indicates that the OS has enabled XSETBV/XGETBV instructions to access XCR0, and support for processor extended state management using XSAVE/XRSTOR.
28	AVX	A value of 1 indicates the processor supports the AVX instruction extensions.
30 - 29	Reserved	Reserved
31	Not Used	Always returns 0

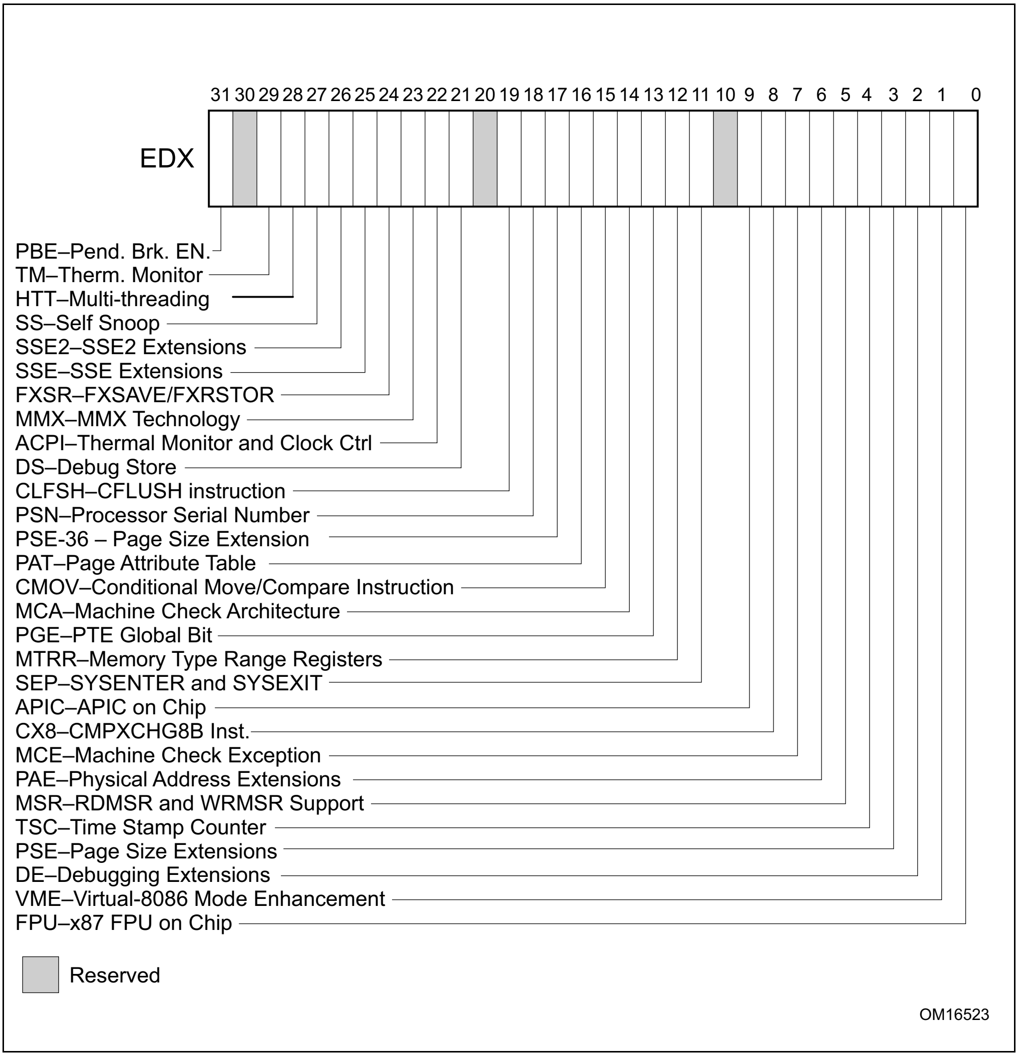


Figure 3-11. Feature Information Returned in the EDX Register

Table 3-21. More on Feature Information Returned in the EDX Register

Bit #	Mnemonic	Description
0	FPU	**Floating Point Unit On-Chip.** The processor contains an x87 FPU.
1	VME	**Virtual 8086 Mode Enhancements.** Virtual 8086 mode enhancements, including CR4.VME for controlling the feature, CR4.PVI for protected mode virtual interrupts, software interrupt indirection, expansion of the TSS with the software indirection bitmap, and EFLAGS.VIF and EFLAGS.VIP flags.
2	DE	**Debugging Extensions.** Support for I/O breakpoints, including CR4.DE for controlling the feature, and optional trapping of accesses to DR4 and DR5.
3	PSE	**Page Size Extension.** Large pages of size 4 MByte are supported, including CR4.PSE for controlling the feature, the defined dirty bit in PDE (Page Directory Entries), optional reserved bit trapping in CR3, PDEs, and PTEs.
4	TSC	**Time Stamp Counter.** The RDTSC instruction is supported, including CR4.TSD for controlling privilege.
5	MSR	**Model Specific Registers RDMSR and WRMSR Instructions.** The RDMSR and WRMSR instructions are supported. Some of the MSRs are implementation dependent.
6	PAE	**Physical Address Extension.** Physical addresses greater than 32 bits are supported: extended page table entry formats, an extra level in the page translation tables is defined, 2-MByte pages are supported instead of 4 Mbyte pages if PAE bit is 1.
7	MCE	**Machine Check Exception.** Exception 18 is defined for Machine Checks, including CR4.MCE for controlling the feature. This feature does not define the model-specific implementations of machine-check error logging, reporting, and processor shutdowns. Machine Check exception handlers may have to depend on processor version to do model specific processing of the exception, or test for the presence of the Machine Check feature.
8	CX8	**CMPXCHG8B Instruction.** The compare-and-exchange 8 bytes (64 bits) instruction is supported (implicitly locked and atomic).
9	APIC	**APIC On-Chip.** The processor contains an Advanced Programmable Interrupt Controller (APIC), responding to memory mapped commands in the physical address range FFFE0000H to FFFE0FFFH (by default - some processors permit the APIC to be relocated).
10	Reserved	Reserved
11	SEP	**SYSENTER and SYSEXIT Instructions.** The SYSENTER and SYSEXIT and associated MSRs are supported.
12	MTRR	**Memory Type Range Registers.** MTRRs are supported. The MTRRcap MSR contains feature bits that describe what memory types are supported, how many variable MTRRs are supported, and whether fixed MTRRs are supported.

Table 3-21. More on Feature Information Returned in the EDX Register (Contd.)

Bit #	Mnemonic	Description
13	PGE	**Page Global Bit.** The global bit is supported in paging-structure entries that map a page, indicating TLB entries that are common to different processes and need not be flushed. The CR4.PGE bit controls this feature.
14	MCA	**Machine Check Architecture.** The Machine Check Architecture, which provides a compatible mechanism for error reporting in P6 family, Pentium 4, Intel Xeon processors, and future processors, is supported. The MCG_CAP MSR contains feature bits describing how many banks of error reporting MSRs are supported.
15	CMOV	**Conditional Move Instructions.** The conditional move instruction CMOV is supported. In addition, if x87 FPU is present as indicated by the CPUID.FPU feature bit, then the FCOMI and FCMOV instructions are supported
16	PAT	**Page Attribute Table.** Page Attribute Table is supported. This feature augments the Memory Type Range Registers (MTRRs), allowing an operating system to specify attributes of memory accessed through a linear address on a 4KB granularity.
17	PSE-36	**36-Bit Page Size Extension.** 4-MByte pages addressing physical memory beyond 4 GBytes are supported with 32-bit paging. This feature indicates that upper bits of the physical address of a 4-MByte page are encoded in bits 20:13 of the page-directory entry. Such physical addresses are limited by MAXPHYADDR and may be up to 40 bits in size.
18	PSN	**Processor Serial Number.** The processor supports the 96-bit processor identification number feature and the feature is enabled.
19	CLFSH	**CLFLUSH Instruction.** CLFLUSH Instruction is supported.
20	Reserved	Reserved
21	DS	**Debug Store.** The processor supports the ability to write debug information into a memory resident buffer. This feature is used by the branch trace store (BTS) and precise event-based sampling (PEBS) facilities (see Chapter 20, "Introduction to Virtual-Machine Extensions," in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3B*).
22	ACPI	**Thermal Monitor and Software Controlled Clock Facilities.** The processor implements internal MSRs that allow processor temperature to be monitored and processor performance to be modulated in predefined duty cycles under software control.
23	MMX	**Intel MMX Technology.** The processor supports the Intel MMX technology.
24	FXSR	**FXSAVE and FXRSTOR Instructions.** The FXSAVE and FXRSTOR instructions are supported for fast save and restore of the floating point context. Presence of this bit also indicates that CR4.OSFXSR is available for an operating system to indicate that it supports the FXSAVE and FXRSTOR instructions.

Table 3-21. More on Feature Information Returned in the EDX Register (Contd.)

Bit #	Mnemonic	Description
25	SSE	**SSE.** The processor supports the SSE extensions.
26	SSE2	**SSE2.** The processor supports the SSE2 extensions.
27	SS	**Self Snoop.** The processor supports the management of conflicting memory types by performing a snoop of its own cache structure for transactions issued to the bus.
28	HTT	**Multi-Threading.** The physical processor package is capable of supporting more than one logical processor.
29	TM	**Thermal Monitor.** The processor implements the thermal monitor automatic thermal control circuitry (TCC).
30	Reserved	Reserved
31	PBE	**Pending Break Enable.** The processor supports the use of the FERR#/PBE# pin when the processor is in the stop-clock state (STPCLK# is asserted) to signal the processor that an interrupt is pending and that the processor should return to normal operation to handle the interrupt. Bit 10 (PBE enable) in the IA32_MISC_ENABLE MSR enables this capability.

INPUT EAX = 2: TLB/Cache/Prefetch Information Returned in EAX, EBX, ECX, EDX

When CPUID executes with EAX set to 2, the processor returns information about the processor's internal TLBs, cache and prefetch hardware in the EAX, EBX, ECX, and EDX registers. The information is reported in encoded form and fall into the following categories:

- The least-significant byte in register EAX (register AL) indicates the number of times the CPUID instruction must be executed with an input value of 2 to get a complete description of the processor's TLB/Cache/Prefetch hardware. The Intel Xeon processor 7400 series will return a 1.
- The most significant bit (bit 31) of each register indicates whether the register contains valid information (set to 0) or is reserved (set to 1).
- If a register contains valid information, the information is contained in 1 byte descriptors. There are four types of encoding values for the byte descriptor, the encoding type is noted in the second column of Table 3-22. Table 3-22 lists the encoding of these descriptors. Note that the order of descriptors in the EAX, EBX, ECX, and EDX registers is not defined; that is, specific bytes are not designated to contain descriptors for specific cache, prefetch, or TLB types. The descriptors may appear in any order. Note also a processor may report a general descriptor type (FFH) and not report any byte descriptor of "cache type" via CPUID leaf 2.

Table 3-22. Encoding of CPUID Leaf 2 Descriptors

Value	Type	Description
00H	General	Null descriptor, this byte contains no information
01H	TLB	Instruction TLB: 4 KByte pages, 4-way set associative, 32 entries
02H	TLB	Instruction TLB: 4 MByte pages, fully associative, 2 entries
03H	TLB	Data TLB: 4 KByte pages, 4-way set associative, 64 entries
04H	TLB	Data TLB: 4 MByte pages, 4-way set associative, 8 entries
05H	TLB	Data TLB1: 4 MByte pages, 4-way set associative, 32 entries
06H	Cache	1st-level instruction cache: 8 KBytes, 4-way set associative, 32 byte line size
08H	Cache	1st-level instruction cache: 16 KBytes, 4-way set associative, 32 byte line size
09H	Cache	1st-level instruction cache: 32KBytes, 4-way set associative, 64 byte line size
0AH	Cache	1st-level data cache: 8 KBytes, 2-way set associative, 32 byte line size
0BH	TLB	Instruction TLB: 4 MByte pages, 4-way set associative, 4 entries
0CH	Cache	1st-level data cache: 16 KBytes, 4-way set associative, 32 byte line size
0DH	Cache	1st-level data cache: 16 KBytes, 4-way set associative, 64 byte line size
0EH	Cache	1st-level data cache: 24 KBytes, 6-way set associative, 64 byte line size
21H	Cache	2nd-level cache: 256 KBytes, 8-way set associative, 64 byte line size
22H	Cache	3rd-level cache: 512 KBytes, 4-way set associative, 64 byte line size, 2 lines per sector
23H	Cache	3rd-level cache: 1 MBytes, 8-way set associative, 64 byte line size, 2 lines per sector
25H	Cache	3rd-level cache: 2 MBytes, 8-way set associative, 64 byte line size, 2 lines per sector
29H	Cache	3rd-level cache: 4 MBytes, 8-way set associative, 64 byte line size, 2 lines per sector
2CH	Cache	1st-level data cache: 32 KBytes, 8-way set associative, 64 byte line size
30H	Cache	1st-level instruction cache: 32 KBytes, 8-way set associative, 64 byte line size
40H	Cache	No 2nd-level cache or, if processor contains a valid 2nd-level cache, no 3rd-level cache
41H	Cache	2nd-level cache: 128 KBytes, 4-way set associative, 32 byte line size
42H	Cache	2nd-level cache: 256 KBytes, 4-way set associative, 32 byte line size
43H	Cache	2nd-level cache: 512 KBytes, 4-way set associative, 32 byte line size
44H	Cache	2nd-level cache: 1 MByte, 4-way set associative, 32 byte line size
45H	Cache	2nd-level cache: 2 MByte, 4-way set associative, 32 byte line size

Table 3-22. Encoding of CPUID Leaf 2 Descriptors (Contd.)

Value	Type	Description
46H	Cache	3rd-level cache: 4 MByte, 4-way set associative, 64 byte line size
47H	Cache	3rd-level cache: 8 MByte, 8-way set associative, 64 byte line size
48H	Cache	2nd-level cache: 3MByte, 12-way set associative, 64 byte line size
49H	Cache	3rd-level cache: 4MB, 16-way set associative, 64-byte line size (Intel Xeon processor MP, Family 0FH, Model 06H); 2nd-level cache: 4 MByte, 16-way set associative, 64 byte line size
4AH	Cache	3rd-level cache: 6MByte, 12-way set associative, 64 byte line size
4BH	Cache	3rd-level cache: 8MByte, 16-way set associative, 64 byte line size
4CH	Cache	3rd-level cache: 12MByte, 12-way set associative, 64 byte line size
4DH	Cache	3rd-level cache: 16MByte, 16-way set associative, 64 byte line size
4EH	Cache	2nd-level cache: 6MByte, 24-way set associative, 64 byte line size
4FH	TLB	Instruction TLB: 4 KByte pages, 32 entries
50H	TLB	Instruction TLB: 4 KByte and 2-MByte or 4-MByte pages, 64 entries
51H	TLB	Instruction TLB: 4 KByte and 2-MByte or 4-MByte pages, 128 entries
52H	TLB	Instruction TLB: 4 KByte and 2-MByte or 4-MByte pages, 256 entries
55H	TLB	Instruction TLB: 2-MByte or 4-MByte pages, fully associative, 7 entries
56H	TLB	Data TLB0: 4 MByte pages, 4-way set associative, 16 entries
57H	TLB	Data TLB0: 4 KByte pages, 4-way associative, 16 entries
59H	TLB	Data TLB0: 4 KByte pages, fully associative, 16 entries
5AH	TLB	Data TLB0: 2-MByte or 4 MByte pages, 4-way set associative, 32 entries
5BH	TLB	Data TLB: 4 KByte and 4 MByte pages, 64 entries
5CH	TLB	Data TLB: 4 KByte and 4 MByte pages,128 entries
5DH	TLB	Data TLB: 4 KByte and 4 MByte pages,256 entries
60H	Cache	1st-level data cache: 16 KByte, 8-way set associative, 64 byte line size
66H	Cache	1st-level data cache: 8 KByte, 4-way set associative, 64 byte line size
67H	Cache	1st-level data cache: 16 KByte, 4-way set associative, 64 byte line size
68H	Cache	1st-level data cache: 32 KByte, 4-way set associative, 64 byte line size
70H	Cache	Trace cache: 12 K-µop, 8-way set associative
71H	Cache	Trace cache: 16 K-µop, 8-way set associative
72H	Cache	Trace cache: 32 K-µop, 8-way set associative
76H	TLB	Instruction TLB: 2M/4M pages, fully associative, 8 entries
78H	Cache	2nd-level cache: 1 MByte, 4-way set associative, 64byte line size

Table 3-22. Encoding of CPUID Leaf 2 Descriptors (Contd.)

Value	Type	Description
79H	Cache	2nd-level cache: 128 KByte, 8-way set associative, 64 byte line size, 2 lines per sector
7AH	Cache	2nd-level cache: 256 KByte, 8-way set associative, 64 byte line size, 2 lines per sector
7BH	Cache	2nd-level cache: 512 KByte, 8-way set associative, 64 byte line size, 2 lines per sector
7CH	Cache	2nd-level cache: 1 MByte, 8-way set associative, 64 byte line size, 2 lines per sector
7DH	Cache	2nd-level cache: 2 MByte, 8-way set associative, 64byte line size
7FH	Cache	2nd-level cache: 512 KByte, 2-way set associative, 64-byte line size
80H	Cache	2nd-level cache: 512 KByte, 8-way set associative, 64-byte line size
82H	Cache	2nd-level cache: 256 KByte, 8-way set associative, 32 byte line size
83H	Cache	2nd-level cache: 512 KByte, 8-way set associative, 32 byte line size
84H	Cache	2nd-level cache: 1 MByte, 8-way set associative, 32 byte line size
85H	Cache	2nd-level cache: 2 MByte, 8-way set associative, 32 byte line size
86H	Cache	2nd-level cache: 512 KByte, 4-way set associative, 64 byte line size
87H	Cache	2nd-level cache: 1 MByte, 8-way set associative, 64 byte line size
B0H	TLB	Instruction TLB: 4 KByte pages, 4-way set associative, 128 entries
B1H	TLB	Instruction TLB: 2M pages, 4-way, 8 entries or 4M pages, 4-way, 4 entries
B2H	TLB	Instruction TLB: 4KByte pages, 4-way set associative, 64 entries
B3H	TLB	Data TLB: 4 KByte pages, 4-way set associative, 128 entries
B4H	TLB	Data TLB1: 4 KByte pages, 4-way associative, 256 entries
BAH	TLB	Data TLB1: 4 KByte pages, 4-way associative, 64 entries
C0H	TLB	Data TLB: 4 KByte and 4 MByte pages, 4-way associative, 8 entries
CAH	STLB	Shared 2nd-Level TLB: 4 KByte pages, 4-way associative, 512 entries
D0H	Cache	3rd-level cache: 512 KByte, 4-way set associative, 64 byte line size
D1H	Cache	3rd-level cache: 1 MByte, 4-way set associative, 64 byte line size
D2H	Cache	3rd-level cache: 2 MByte, 4-way set associative, 64 byte line size
D6H	Cache	3rd-level cache: 1 MByte, 8-way set associative, 64 byte line size
D7H	Cache	3rd-level cache: 2 MByte, 8-way set associative, 64 byte line size
D8H	Cache	3rd-level cache: 4 MByte, 8-way set associative, 64 byte line size
DCH	Cache	3rd-level cache: 1.5 MByte, 12-way set associative, 64 byte line size
DDH	Cache	3rd-level cache: 3 MByte, 12-way set associative, 64 byte line size
DEH	Cache	3rd-level cache: 6 MByte, 12-way set associative, 64 byte line size

Table 3-22. Encoding of CPUID Leaf 2 Descriptors (Contd.)

Value	Type	Description
E2H	Cache	3rd-level cache: 2 MByte, 16-way set associative, 64 byte line size
E3H	Cache	3rd-level cache: 4 MByte, 16-way set associative, 64 byte line size
E4H	Cache	3rd-level cache: 8 MByte, 16-way set associative, 64 byte line size
EAH	Cache	3rd-level cache: 12MByte, 24-way set associative, 64 byte line size
EBH	Cache	3rd-level cache: 18MByte, 24-way set associative, 64 byte line size
ECH	Cache	3rd-level cache: 24MByte, 24-way set associative, 64 byte line size
F0H	Prefetch	64-Byte prefetching
F1H	Prefetch	128-Byte prefetching
FFH	General	CPUID leaf 2 does not report cache descriptor information, use CPUID leaf 4 to query cache parameters

Example 3-1. Example of Cache and TLB Interpretation

The first member of the family of Pentium 4 processors returns the following information about caches and TLBs when the CPUID executes with an input value of 2:

```
EAX     66 5B 50 01H
EBX     0H
ECX     0H
EDX     00 7A 70 00H
```

Which means:

- The least-significant byte (byte 0) of register EAX is set to 01H. This indicates that CPUID needs to be executed once with an input value of 2 to retrieve complete information about caches and TLBs.
- The most-significant bit of all four registers (EAX, EBX, ECX, and EDX) is set to 0, indicating that each register contains valid 1-byte descriptors.
- Bytes 1, 2, and 3 of register EAX indicate that the processor has:
  - 50H - a 64-entry instruction TLB, for mapping 4-KByte and 2-MByte or 4-MByte pages.
  - 5BH - a 64-entry data TLB, for mapping 4-KByte and 4-MByte pages.
  - 66H - an 8-KByte 1st level data cache, 4-way set associative, with a 64-Byte cache line size.
- The descriptors in registers EBX and ECX are valid, but contain NULL descriptors.
- Bytes 0, 1, 2, and 3 of register EDX indicate that the processor has:
  - 00H - NULL descriptor.
  - 70H - Trace cache: 12 K-μop, 8-way set associative.

— 7AH - a 256-KByte 2nd level cache, 8-way set associative, with a sectored, 64-byte cache line size.

— 00H - NULL descriptor.

INPUT EAX = 04H: Returns Deterministic Cache Parameters for Each Level

When CPUID executes with EAX set to 04H and ECX contains an index value, the processor returns encoded data that describe a set of deterministic cache parameters (for the cache level associated with the input in ECX). Valid index values start from 0.

Software can enumerate the deterministic cache parameters for each level of the cache hierarchy starting with an index value of 0, until the parameters report the value associated with the cache type field is 0. The architecturally defined fields reported by deterministic cache parameters are documented in Table 3-17.

This Cache Size in Bytes

= (Ways + 1) * (Partitions + 1) * (Line_Size + 1) * (Sets + 1)

= (EBX[31:22] + 1) * (EBX[21:12] + 1) * (EBX[11:0] + 1) * (ECX + 1)

The CPUID leaf 04H also reports data that can be used to derive the topology of processor cores in a physical package. This information is constant for all valid index values. Software can query the raw data reported by executing CPUID with EAX=04H and ECX=0 and use it as part of the topology enumeration algorithm described in Chapter 8, “Multiple-Processor Management,” in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 3A*.

INPUT EAX = 05H: Returns MONITOR and MWAIT Features

When CPUID executes with EAX set to 05H, the processor returns information about features available to MONITOR/MWAIT instructions. The MONITOR instruction is used for address-range monitoring in conjunction with MWAIT instruction. The MWAIT instruction optionally provides additional extensions for advanced power management. See Table 3-17.

INPUT EAX = 06H: Returns Thermal and Power Management Features

When CPUID executes with EAX set to 06H, the processor returns information about thermal and power management features. See Table 3-17.

INPUT EAX = 09H: Returns Direct Cache Access Information

When CPUID executes with EAX set to 09H, the processor returns information about Direct Cache Access capabilities. See Table 3-17.

INPUT EAX = 0AH: Returns Architectural Performance Monitoring Features

When CPUID executes with EAX set to 0AH, the processor returns information about support for architectural performance monitoring capabilities. Architectural performance monitoring is supported if the version ID (see Table 3-17) is greater than Pn 0. See Table 3-17.

For each version of architectural performance monitoring capability, software must enumerate this leaf to discover the programming facilities and the architectural performance events available in the processor. The details are described in Chapter 20, “Introduction to Virtual-Machine Extensions,” in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 3B*.

INPUT EAX = 0BH: Returns Extended Topology Information

When CPUID executes with EAX set to 0BH, the processor returns information about extended topology enumeration data. Software must detect the presence of CPUID leaf 0BH by verifying (a) the highest leaf index supported by CPUID is >= 0BH, and (b) CPUID.0BH:EBX[15:0] reports a non-zero value. See Table 3-17.

INPUT EAX = 0DH: Returns Processor Extended States Enumeration Information

When CPUID executes with EAX set to 0DH and ECX = 0, the processor returns information about the bit-vector representation of all processor state extensions that are supported in the processor and storage size requirements of the XSAVE/XRSTOR area. See Table 3-17.

When CPUID executes with EAX set to 0DH and ECX = n (n > 1, and is a valid subleaf index), the processor returns information about the size and offset of each processor extended state save area within the XSAVE/XRSTOR area. See Table 3-17. Software can use the forward-extendable technique depicted below to query the valid sub-leaves and obtain size and offset information for each processor extended state save area:

```
For i = 2 to 62 // sub-leaf 1 is reserved
    IF (CPUID.(EAX=0DH, ECX=0):VECTOR[i] = 1 ) // VECTOR is the 64-bit value of EDX:EAX
        Execute CPUID.(EAX=0DH, ECX = i) to examine size and offset for sub-leaf i;
    FI;
```

METHODS FOR RETURNING BRANDING INFORMATION

Use the following techniques to access branding information:

1. Processor brand string method; this method also returns the processor’s maximum operating frequency
2. Processor brand index; this method uses a software supplied brand string table.

These two methods are discussed in the following sections. For methods that are available in early processors, see Section: “Identification of Earlier IA-32 Processors” in Chapter 14 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*.

The Processor Brand String Method

Figure 3-12 describes the algorithm used for detection of the brand string. Processor brand identification software should execute this algorithm on all Intel 64 and IA-32 processors.

This method (introduced with Pentium 4 processors) returns an ASCII brand identification string and the maximum operating frequency of the processor to the EAX, EBX, ECX, and EDX registers.

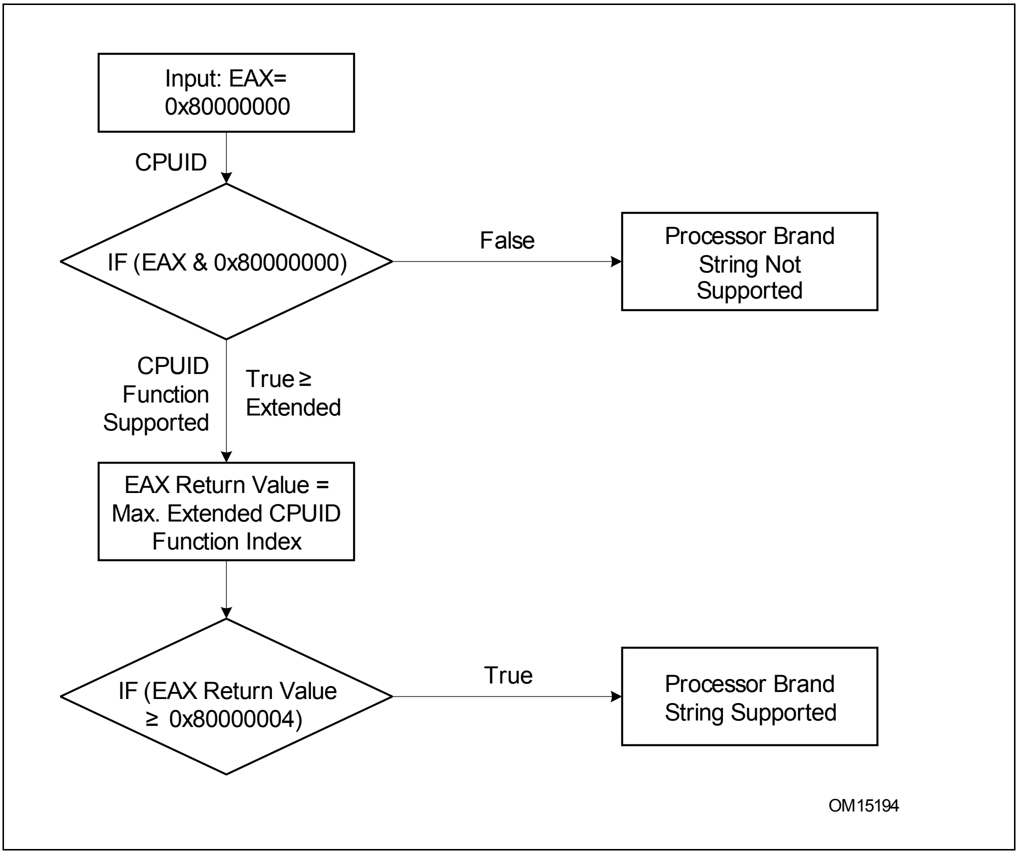


Figure 3-12. Determination of Support for the Processor Brand String

How Brand Strings Work

To use the brand string method, execute CPUID with EAX input of 8000002H through 80000004H. For each input value, CPUID returns 16 ASCII characters using EAX, EBX, ECX, and EDX. The returned string will be NULL-terminated.

Table 3-23 shows the brand string that is returned by the first processor in the Pentium 4 processor family.

Table 3-23. Processor Brand String Returned with Pentium 4 Processor

EAX Input Value	Return Values	ASCII Equivalent
80000002H	EAX = 20202020H	“ ”
	EBX = 20202020H	“ ”
	ECX = 20202020H	“ ”
	EDX = 6E492020H	“nI ”
80000003H	EAX = 286C6574H	“(let”
	EBX = 50202952H	“P)R”
	ECX = 69746E65H	“itne”
	EDX = 52286D75H	“R(mu”
80000004H	EAX = 20342029H	“ 4)”
	EBX = 20555043H	“ UPC”
	ECX = 30303531H	“0051”
	EDX = 007A484DH	“\0zHM”

Extracting the Maximum Processor Frequency from Brand Strings

Figure 3-13 provides an algorithm which software can use to extract the maximum processor operating frequency from the processor brand string.

NOTE

When a frequency is given in a brand string, it is the maximum qualified frequency of the processor, not the frequency at which the processor is currently running.

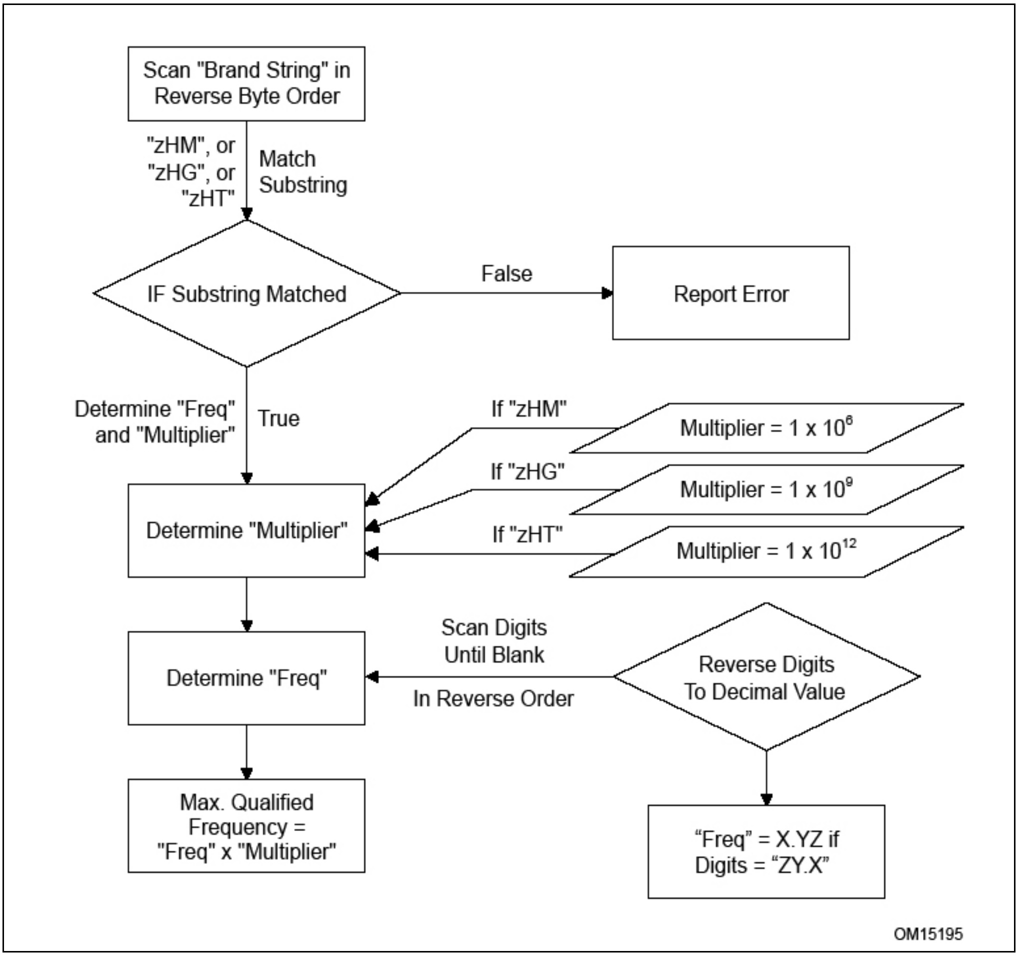


Figure 3-13. Algorithm for Extracting Maximum Processor Frequency

The Processor Brand Index Method

The brand index method (introduced with Pentium® III Xeon® processors) provides an entry point into a brand identification table that is maintained in memory by system software and is accessible from system- and user-level code. In this table, each brand index is associate with an ASCII brand identification string that identifies the official Intel family and model number of a processor.

When CPUID executes with EAX set to 1, the processor returns a brand index to the low byte in EBX. Software can then use this index to locate the brand identification string for the processor in the brand identification table. The first entry (brand index 0) in this table is reserved, allowing for backward compatibility with processors that do not support the brand identification feature. Starting with processor signature family ID = 0FH, model = 03H, brand index method is no longer supported. Use brand string method instead.

Table 3-24 shows brand indices that have identification strings associated with them.

Table 3-24. Mapping of Brand Indices; and Intel 64 and IA-32 Processor Brand Strings

Brand Index	Brand String
00H	This processor does not support the brand identification feature
01H	Intel(R) Celeron(R) processor[1]
02H	Intel(R) Pentium(R) III processor[1]
03H	Intel(R) Pentium(R) III Xeon(R) processor; If processor signature = 000006B1h, then Intel(R) Celeron(R) processor
04H	Intel(R) Pentium(R) III processor
06H	Mobile Intel(R) Pentium(R) III processor-M
07H	Mobile Intel(R) Celeron(R) processor[1]
08H	Intel(R) Pentium(R) 4 processor
09H	Intel(R) Pentium(R) 4 processor
0AH	Intel(R) Celeron(R) processor[1]
0BH	Intel(R) Xeon(R) processor; If processor signature = 00000F13h, then Intel(R) Xeon(R) processor MP
0CH	Intel(R) Xeon(R) processor MP
0EH	Mobile Intel(R) Pentium(R) 4 processor-M; If processor signature = 00000F13h, then Intel(R) Xeon(R) processor
0FH	Mobile Intel(R) Celeron(R) processor[1]
11H	Mobile Genuine Intel(R) processor
12H	Intel(R) Celeron(R) M processor
13H	Mobile Intel(R) Celeron(R) processor[1]
14H	Intel(R) Celeron(R) processor
15H	Mobile Genuine Intel(R) processor
16H	Intel(R) Pentium(R) M processor
17H	Mobile Intel(R) Celeron(R) processor[1]
18H – 0FFH	RESERVED

NOTES:

1. Indicates versions of these processors that were introduced after the Pentium III

IA-32 Architecture Compatibility

CPUID is not supported in early models of the Intel486 processor or in any IA-32 processor earlier than the Intel486 processor.

Operation

```
IA32_BIOS_SIGN_ID MSR ← Update with installed microcode revision number;

CASE (EAX) OF
    EAX = 0:
        EAX ← Highest basic function input value understood by CPUID;
        EBX ← Vendor identification string;
        EDX ← Vendor identification string;
        ECX ← Vendor identification string;
    BREAK;
    EAX = 1H:
        EAX[3:0] ← Stepping ID;
        EAX[7:4] ← Model;
        EAX[11:8] ← Family;
        EAX[13:12] ← Processor type;
        EAX[15:14] ← Reserved;
        EAX[19:16] ← Extended Model;
        EAX[27:20] ← Extended Family;
        EAX[31:28] ← Reserved;
        EBX[7:0] ← Brand Index; (* Reserved if the value is zero. *)
        EBX[15:8] ← CLFLUSH Line Size;
        EBX[16:23] ← Reserved; (* Number of threads enabled = 2 if MT enable fuse set. *)
        EBX[24:31] ← Initial APIC ID;
        ECX ← Feature flags; (* See Figure 3-10. *)
        EDX ← Feature flags; (* See Figure 3-11. *)
    BREAK;
    EAX = 2H:
        EAX ← Cache and TLB information;
        EBX ← Cache and TLB information;
        ECX ← Cache and TLB information;
        EDX ← Cache and TLB information;
    BREAK;
    EAX = 3H:
        EAX ← Reserved;
        EBX ← Reserved;
        ECX ← ProcessorSerialNumber[31:0];
        (* Pentium III processors only, otherwise reserved. *)
        EDX ← ProcessorSerialNumber[63:32];
        (* Pentium III processors only, otherwise reserved. *
```

```
BREAK
EAX = 4H:
    EAX ← Deterministic Cache Parameters Leaf; (* See Table 3-17. *)
    EBX ← Deterministic Cache Parameters Leaf;
    ECX ← Deterministic Cache Parameters Leaf;
    EDX ← Deterministic Cache Parameters Leaf;
BREAK;
EAX = 5H:
    EAX ← MONITOR/MWAIT Leaf; (* See Table 3-17. *)
    EBX ← MONITOR/MWAIT Leaf;
    ECX ← MONITOR/MWAIT Leaf;
    EDX ← MONITOR/MWAIT Leaf;
BREAK;
EAX = 6H:
    EAX ← Thermal and Power Management Leaf; (* See Table 3-17. *)
    EBX ← Thermal and Power Management Leaf;
    ECX ← Thermal and Power Management Leaf;
    EDX ← Thermal and Power Management Leaf;
BREAK;
EAX = 7H or 8H:
    EAX ← Reserved = 0;
    EBX ← Reserved = 0;
    ECX ← Reserved = 0;
    EDX ← Reserved = 0;
BREAK;
EAX = 9H:
    EAX ← Direct Cache Access Information Leaf; (* See Table 3-17. *)
    EBX ← Direct Cache Access Information Leaf;
    ECX ← Direct Cache Access Information Leaf;
    EDX ← Direct Cache Access Information Leaf;
BREAK;
EAX = AH:
    EAX ← Architectural Performance Monitoring Leaf; (* See Table 3-17. *)
    EBX ← Architectural Performance Monitoring Leaf;
    ECX ← Architectural Performance Monitoring Leaf;
    EDX ← Architectural Performance Monitoring Leaf;
    BREAK
EAX = BH:
    EAX ← Extended Topology Enumeration Leaf; (* See Table 3-17. *)
    EBX ← Extended Topology Enumeration Leaf;
    ECX ← Extended Topology Enumeration Leaf;
    EDX ← Extended Topology Enumeration Leaf;
BREAK;
```

```
    EAX = CH:
        EAX ← Reserved = 0;
        EBX ← Reserved = 0;
        ECX ← Reserved = 0;
        EDX ← Reserved = 0;
    BREAK;
    EAX = DH:
        EAX ← Processor Extended State Enumeration Leaf; (* See Table 3-17. *)
        EBX ← Processor Extended State Enumeration Leaf;
        ECX ← Processor Extended State Enumeration Leaf;
        EDX ← Processor Extended State Enumeration Leaf;
    BREAK;
BREAK;
    EAX = 80000000H:
        EAX ← Highest extended function input value understood by CPUID;
        EBX ← Reserved;
        ECX ← Reserved;
        EDX ← Reserved;
    BREAK;
    EAX = 80000001H:
        EAX ← Reserved;
        EBX ← Reserved;
        ECX ← Extended Feature Bits (* See Table 3-17.*);
        EDX ← Extended Feature Bits (* See Table 3-17. *);
    BREAK;
    EAX = 80000002H:
        EAX ← Processor Brand String;
        EBX ← Processor Brand String, continued;
        ECX ← Processor Brand String, continued;
        EDX ← Processor Brand String, continued;
    BREAK;
    EAX = 80000003H:
        EAX ← Processor Brand String, continued;
        EBX ← Processor Brand String, continued;
        ECX ← Processor Brand String, continued;
        EDX ← Processor Brand String, continued;
    BREAK;
    EAX = 80000004H:
        EAX ← Processor Brand String, continued;
        EBX ← Processor Brand String, continued;
        ECX ← Processor Brand String, continued;
        EDX ← Processor Brand String, continued;
    BREAK;
```

```
    EAX = 80000005H:
        EAX ← Reserved = 0;
        EBX ← Reserved = 0;
        ECX ← Reserved = 0;
        EDX ← Reserved = 0;
    BREAK;
    EAX = 80000006H:
        EAX ← Reserved = 0;
        EBX ← Reserved = 0;
        ECX ← Cache information;
        EDX ← Reserved = 0;
    BREAK;
    EAX = 80000007H:
        EAX ← Reserved = 0;
        EBX ← Reserved = 0;
        ECX ← Reserved = 0;
        EDX ← Reserved = Misc Feature Flags;
    BREAK;
    EAX = 80000008H:
        EAX ← Reserved = Physical Address Size Information;
        EBX ← Reserved = Virtual Address Size Information;
        ECX ← Reserved = 0;
        EDX ← Reserved = 0;
    BREAK;
    EAX >= 40000000H and EAX <= 4FFFFFFFH:
    DEFAULT: (* EAX = Value outside of recognized range for CPUID. *)
        (* If the highest basic information leaf data depend on ECX input value, ECX is honored.*)
        EAX ← Reserved; (* Information returned for highest basic information leaf. *)
        EBX ← Reserved; (* Information returned for highest basic information leaf. *)
        ECX ← Reserved; (* Information returned for highest basic information leaf. *)
        EDX ← Reserved; (* Information returned for highest basic information leaf. *)
    BREAK;
ESAC;
```

Flags Affected

None.

Exceptions (All Operating Modes)

#UD	If the LOCK prefix is used.
	In earlier IA-32 processors that do not support the CPUID instruction, execution of the instruction results in an invalid opcode (#UD) exception being generated.

CRC32 — Accumulate CRC32 Value

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
F2 0F 38 F0 */r*	CRC32 *r32, r/m8*	A	Valid	Valid	Accumulate CRC32 on *r/m8*.
F2 REX 0F 38 F0 */r*	CRC32 *r32, r/m8**	A	Valid	N.E.	Accumulate CRC32 on *r/m8*.
F2 0F 38 F1 */r*	CRC32 *r32, r/m16*	A	Valid	Valid	Accumulate CRC32 on *r/m16*.
F2 0F 38 F1 */r*	CRC32 *r32, r/m32*	A	Valid	Valid	Accumulate CRC32 on *r/m32.*
F2 REX.W 0F 38 F0 */r*	CRC32 *r64, r/m8*	A	Valid	N.E.	Accumulate CRC32 on *r/m8*.
F2 REX.W 0F 38 F1 */r*	CRC32 *r64, r/m64*	A	Valid	N.E.	Accumulate CRC32 on *r/m64.*

NOTES:

*In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA

Description

Starting with an initial value in the first operand (destination operand), accumulates a CRC32 (polynomial 0x11EDC6F41) value for the second operand (source operand) and stores the result in the destination operand. The source operand can be a register or a memory location. The destination operand must be an r32 or r64 register. If the destination is an r64 register, then the 32-bit result is stored in the least significant double word and 00000000H is stored in the most significant double word of the r64 register.

The initial value supplied in the destination operand is a double word integer stored in the r32 register or the least significant double word of the r64 register. To incrementally accumulate a CRC32 value, software retains the result of the previous CRC32 operation in the destination operand, then executes the CRC32 instruction again with new input data in the source operand. Data contained in the source operand is processed in reflected bit order. This means that the most significant bit of the source operand is treated as the least significant bit of the quotient, and so on, for all the bits of the source operand. Likewise, the result of the CRC operation is stored in the destination operand in reflected bit order. This means that the most significant bit of the resulting CRC (bit 31) is stored in the least significant bit of the destination operand (bit 0), and so on, for all the bits of the CRC.

Operation

Notes:

```
BIT_REFLECT64: DST[63-0] = SRC[0-63]
BIT_REFLECT32: DST[31-0] = SRC[0-31]
BIT_REFLECT16: DST[15-0] = SRC[0-15]
BIT_REFLECT8: DST[7-0] = SRC[0-7]
MOD2: Remainder from Polynomial division modulus 2
```

CRC32 instruction for 64-bit source operand and 64-bit destination operand:

```
TEMP1[63-0] ← BIT_REFLECT64 (SRC[63-0])
TEMP2[31-0] ← BIT_REFLECT32 (DEST[31-0])
TEMP3[95-0] ← TEMP1[63-0] « 32
TEMP4[95-0] ← TEMP2[31-0] « 64
TEMP5[95-0] ← TEMP3[95-0] XOR TEMP4[95-0]
TEMP6[31-0] ← TEMP5[95-0] MOD2 11EDC6F41H
DEST[31-0] ← BIT_REFLECT (TEMP6[31-0])
DEST[63-32] ← 00000000H
```

CRC32 instruction for 32-bit source operand and 32-bit destination operand:

```
TEMP1[31-0] ← BIT_REFLECT32 (SRC[31-0])
TEMP2[31-0] ← BIT_REFLECT32 (DEST[31-0])
TEMP3[63-0] ← TEMP1[31-0] « 32
TEMP4[63-0] ← TEMP2[31-0] « 32
TEMP5[63-0] ← TEMP3[63-0] XOR TEMP4[63-0]
TEMP6[31-0] ← TEMP5[63-0] MOD2 11EDC6F41H
DEST[31-0] ← BIT_REFLECT (TEMP6[31-0])
```

CRC32 instruction for 16-bit source operand and 32-bit destination operand:

```
TEMP1[15-0] ← BIT_REFLECT16 (SRC[15-0])
TEMP2[31-0] ← BIT_REFLECT32 (DEST[31-0])
TEMP3[47-0] ← TEMP1[15-0] « 32
TEMP4[47-0] ← TEMP2[31-0] « 16
TEMP5[47-0] ← TEMP3[47-0] XOR TEMP4[47-0]
TEMP6[31-0] ← TEMP5[47-0] MOD2 11EDC6F41H
DEST[31-0] ← BIT_REFLECT (TEMP6[31-0])
```

CRC32 instruction for 8-bit source operand and 64-bit destination operand:

```
TEMP1[7-0] ← BIT_REFLECT8(SRC[7-0])
TEMP2[31-0] ← BIT_REFLECT32 (DEST[31-0])
TEMP3[39-0] ← TEMP1[7-0] « 32
TEMP4[39-0] ← TEMP2[31-0] « 8
TEMP5[39-0] ← TEMP3[39-0] XOR TEMP4[39-0]
```

```
TEMP6[31-0] ← TEMP5[39-0] MOD2 11EDC6F41H
DEST[31-0] ← BIT_REFLECT (TEMP6[31-0])
DEST[63-32] ← 00000000H
```

CRC32 instruction for 8-bit source operand and 32-bit destination operand:

```
TEMP1[7-0] ← BIT_REFLECT8(SRC[7-0])
TEMP2[31-0] ← BIT_REFLECT32 (DEST[31-0])
TEMP3[39-0] ← TEMP1[7-0] « 32
TEMP4[39-0] ← TEMP2[31-0] « 8
TEMP5[39-0] ← TEMP3[39-0] XOR TEMP4[39-0]
TEMP6[31-0] ← TEMP5[39-0] MOD2 11EDC6F41H
DEST[31-0] ← BIT_REFLECT (TEMP6[31-0])
```

Flags Affected

None

Intel C/C++ Compiler Intrinsic Equivalent

unsigned int _mm_crc32_u8(unsigned int crc, unsigned char data)
unsigned int _mm_crc32_u16(unsigned int crc, unsigned short data)
unsigned int _mm_crc32_u32(unsigned int crc, unsigned int data)
unsinged __int64 _mm_crc32_u64(unsinged __int64 crc, unsigned __int64 data)

SIMD Floating Point Exceptions

None

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS or GS segments.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF (fault-code)	For a page fault.
#UD	If CPUID.01H:ECX.SSE4_2 [Bit 20] = 0.
	If LOCK prefix is used.

Real Mode Exceptions

#GP(0)	If any part of the operand lies outside of the effective address space from 0 to 0FFFFH.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#UD	If CPUID.01H:ECX.SSE4_2 [Bit 20] = 0.
	If LOCK prefix is used.

Virtual 8086 Mode Exceptions

#GP(0)	If any part of the operand lies outside of the effective address space from 0 to 0FFFFH.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF (fault-code)	For a page fault.
#UD	If CPUID.01H:ECX.SSE4_2 [Bit 20] = 0.
	If LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in Protected Mode.

64-Bit Mode Exceptions

#GP(0)	If the memory address is in a non-canonical form.
#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#PF (fault-code)	For a page fault.
#UD	If CPUID.01H:ECX.SSE4_2 [Bit 20] = 0.
	If LOCK prefix is used.

CVTDQ2PD—Convert Packed Dword Integers to Packed Double-Precision FP Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F E6 CVTDQ2PD *xmm1, xmm2/m64*	A	V/V	SSE2	Convert two packed signed doubleword integers from *xmm2/m128* to two packed double-precision floating-point values in *xmm1*.
VEX.128.F3.0F.WIG E6 /r VCVTDQ2PD xmm1, xmm2/m64	A	V/V	AVX	Convert two packed signed doubleword integers from xmm2/mem to two packed double-precision floating-point values in xmm1.
VEX.256.F3.0F.WIG E6 /r VCVTDQ2PD ymm1, xmm2/m128	A	V/V	AVX	Convert four packed signed doubleword integers from xmm2/mem to four packed double-precision floating-point values in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts two packed signed doubleword integers in the source operand (second operand) to two packed double-precision floating-point values in the destination operand (first operand).

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The source operand is an XMM register or 64- bit memory location. The destination operation is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: The source operand is an XMM register or 64- bit memory location. The destination operation is a YMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The source operand is an XMM register or 128- bit memory location. The destination operation is a YMM register.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b, otherwise instructions will #UD.

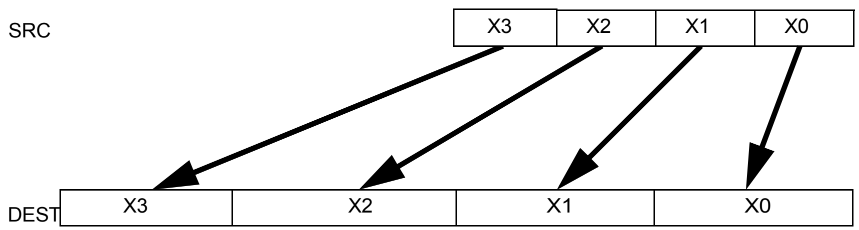


Figure 3-14. CVTDQ2PD (VEX.256 encoded version)

Operation

CVTDQ2PD (128-bit Legacy SSE version)

DEST[63:0] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC[31:0])
DEST[127:64] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC[63:32])
DEST[VLMAX-1:128] (unmodified)

VCVTDQ2PD (VEX.128 encoded version)

DEST[63:0] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC[31:0])
DEST[127:64] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC[63:32])
DEST[VLMAX-1:128] ← 0

VCVTDQ2PD (VEX.256 encoded version)

DEST[63:0] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC[31:0])
DEST[127:64] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC[63:32])
DEST[191:128] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC[95:64])
DEST[255:192] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC[127:96)

Intel C/C++ Compiler Intrinsic Equivalent

CVTDQ2PD __m128d _mm_cvtepi32_pd(__m128i a)

VCVTDQ2PD __m256d _mm256_cvtepi32_pd (__m128i src)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 5; additionally

#UD	If VEX.vvvv != 1111B.

CVTDQ2PS—Convert Packed Dword Integers to Packed Single-Precision FP Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 5B /*r* CVTDQ2PS *xmm1*, *xmm2/m128*	A	V/V	SSE2	Convert four packed signed doubleword integers from *xmm2/m128* to four packed single-precision floating-point values in *xmm1*.
VEX.128.0F.WIG 5B /r VCVTDQ2PS xmm1, xmm2/m128	A	V/V	AVX	Convert four packed signed doubleword integers from xmm2/mem to four packed single-precision floating-point values in xmm1.
VEX.256.0F.WIG 5B /r VCVTDQ2PS ymm1, ymm2/m256	A	V/V	AVX	Convert eight packed signed doubleword integers from ymm2/mem to eight packed single-precision floating-point values in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts four packed signed doubleword integers in the source operand (second operand) to four packed single-precision floating-point values in the destination operand (first operand).

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The source operand is an XMM register or 128- bit memory location. The destination operation is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: The source operand is an XMM register or 128- bit memory location. The destination operation is a YMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The source operand is a YMM register or 256- bit memory location. The destination operation is a YMM register.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b, otherwise instructions will #UD.

Operation

CVTDQ2PS (128-bit Legacy SSE version)

DEST[31:0] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[31:0])
DEST[63:32] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[63:32])
DEST[95:64] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[95:64])
DEST[127:96] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[127z:96)
DEST[VLMAX-1:128] (unmodified)

VCVTDQ2PS (VEX.128 encoded version)

DEST[31:0] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[31:0])
DEST[63:32] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[63:32])
DEST[95:64] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[95:64])
DEST[127:96] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[127z:96)
DEST[VLMAX-1:128] ← 0

VCVTDQ2PS (VEX.256 encoded version)

DEST[31:0] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[31:0])
DEST[63:32] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[63:32])
DEST[95:64] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[95:64])
DEST[127:96] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[127z:96)
DEST[159:128] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[159:128])
DEST[191:160] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[191:160])
DEST[223:192] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[223:192])
DEST[255:224] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[255:224)

Intel C/C++ Compiler Intrinsic Equivalent

CVTDQ2PS __m128 _mm_cvtepi32_ps(__m128i a)

VCVTDQ2PS __m256 _mm256_cvtepi32_ps (__m256i src)

SIMD Floating-Point Exceptions

Precision.

Other Exceptions

See Exceptions Type 2; additionally

#UD	If VEX.vvvv != 1111B.

CVTPD2DQ—Convert Packed Double-Precision FP Values to Packed Dword Integers

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F E6 CVTPD2DQ *xmm1*, *xmm2/m128*	A	V/V	SSE2	Convert two packed double-precision floating-point values from *xmm2/m128* to two packed signed doubleword integers in *xmm1*.
VEX.128.F2.0F.WIG E6 /r VCVTPD2DQ xmm1, xmm2/m128	A	V/V	AVX	Convert two packed double-precision floating-point values in xmm2/mem to two signed doubleword integers in xmm1.
VEX.256.F2.0F.WIG E6 /r VCVTPD2DQ xmm1, ymm2/m256	A	V/V	AVX	Convert four packed double-precision floating-point values in ymm2/mem to four signed doubleword integers in xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts two packed double-precision floating-point values in the source operand (second operand) to two packed signed doubleword integers in the destination operand (first operand).

The source operand can be an XMM register or a 128-bit memory location. The destination operand is an XMM register. The result is stored in the low quadword of the destination operand and the high quadword is cleared to all 0s.

When a conversion is inexact, the value returned is rounded according to the rounding control bits in the MXCSR register. If a converted result is larger than the maximum signed doubleword integer, the floating-point invalid exception is raised, and if this exception is masked, the indefinite integer value (80000000H) is returned.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The source operand is an XMM register or 128- bit memory location. The destination operation is an XMM register. Bits[127:64] of the destination XMM register are zeroed. However, the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: The source operand is an XMM register or 128- bit memory location. The destination operation is a YMM register. The upper bits (VLMAX-1:64) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The source operand is a YMM register or 256- bit memory location. The destination operation is an XMM register. The upper bits (255:128) of the corresponding YMM register destination are zeroed.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b, otherwise instructions will #UD.

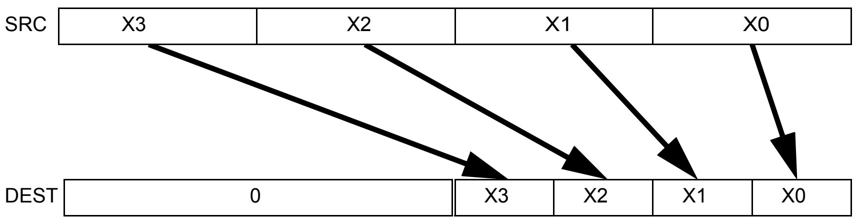


Figure 3-15. VCVTPD2DQ (VEX.256 encoded version)

Operation

CVTPD2DQ (128-bit Legacy SSE version)

DEST[31:0] ← Convert_Double_Precision_Floating_Point_To_Integer(SRC[63:0])
DEST[63:32] ← Convert_Double_Precision_Floating_Point_To_Integer(SRC[127:64])
DEST[127:64] ← 0
DEST[VLMAX-1:128] (unmodified)

VCVTPD2DQ (VEX.128 encoded version)

DEST[31:0] ← Convert_Double_Precision_Floating_Point_To_Integer(SRC[63:0])
DEST[63:32] ← Convert_Double_Precision_Floating_Point_To_Integer(SRC[127:64])
DEST[VLMAX-1:64] ← 0

VCVTPD2DQ (VEX.256 encoded version)

DEST[31:0] ← Convert_Double_Precision_Floating_Point_To_Integer(SRC[63:0])
DEST[63:32] ← Convert_Double_Precision_Floating_Point_To_Integer(SRC[127:64])

```
DEST[95:64] ← Convert_Double_Precision_Floating_Point_To_Integer(SRC[191:128])
DEST[127:96] ← Convert_Double_Precision_Floating_Point_To_Integer(SRC[255:192)
DEST[255:128]← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

CVTPD2DQ __m128i _mm_cvtpd_epi32 (__m128d src)

CVTPD2DQ __m128i _mm256_cvtpd_epi32 (__m256d src)

SIMD Floating-Point Exceptions

Invalid, Precision.

Other Exceptions

See Exceptions Type 2; additionally

#UD	If VEX.vvvv != 1111B.

CVTPD2PI—Convert Packed Double-Precision FP Values to Packed Dword Integers

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
66 0F 2D /*r*	CVTPD2PI *mm, xmm/m128*	A	Valid	Valid	Convert two packed double-precision floating-point values from *xmm/m128* to two packed signed doubleword integers in *mm*.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts two packed double-precision floating-point values in the source operand (second operand) to two packed signed doubleword integers in the destination operand (first operand).

The source operand can be an XMM register or a 128-bit memory location. The destination operand is an MMX technology register.

When a conversion is inexact, the value returned is rounded according to the rounding control bits in the MXCSR register. If a converted result is larger than the maximum signed doubleword integer, the floating-point invalid exception is raised, and if this exception is masked, the indefinite integer value (80000000H) is returned.

This instruction causes a transition from x87 FPU to MMX technology operation (that is, the x87 FPU top-of-stack pointer is set to 0 and the x87 FPU tag word is set to all 0s [valid]). If this instruction is executed while an x87 FPU floating-point exception is pending, the exception is handled before the CVTPD2PI instruction is executed.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Operation

```
DEST[31:0] ← Convert_Double_Precision_Floating_Point_To_Integer32(SRC[63:0]);
DEST[63:32] ← Convert_Double_Precision_Floating_Point_To_Integer32(SRC[127:64]);
```

Intel C/C++ Compiler Intrinsic Equivalent

```
CVTPD1PI __m64 _mm_cvtpd_pi32(__m128d a)
```

SIMD Floating-Point Exceptions

Invalid, Precision.

Protected Mode Exceptions

#GP(0)	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
	If a memory operand is not aligned on a 16-byte boundary, regardless of segment.
#SS(0)	For an illegal address in the SS segment.
#PF(fault-code)	For a page fault.
#MF	If there is a pending x87 FPU exception.
#NM	If CR0.TS[bit 3] = 1.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand is not aligned on a 16-byte boundary, regardless of segment.
	If any part of the operand lies outside the effective address space from 0 to FFFFH.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	For a page fault.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
	If memory operand is not aligned on a 16-byte boundary, regardless of segment.
#PF(fault-code)	For a page fault.
#MF	If there is a pending x87 FPU exception.
#NM	If CR0.TS[bit 3] = 1.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.

CVTPD2PS—Convert Packed Double-Precision FP Values to Packed Single-Precision FP Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 5A /*r* CVTPD2PS *xmm1*, *xmm2/m128*	A	V/V	SSE2	Convert two packed double-precision floating-point values in *xmm2/m128* to two packed single-precision floating-point values in *xmm1*.
VEX.128.66.0F.WIG 5A /r VCVTPD2PS xmm1, xmm2/m128	A	V/V	AVX	Convert two packed double-precision floating-point values in xmm2/mem to two single-precision floating-point values in xmm1.
VEX.256.66.0F.WIG 5A /r VCVTPD2PS xmm1, ymm2/m256	A	V/V	AVX	Convert four packed double-precision floating-point values in ymm2/mem to four single-precision floating-point values in xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts two packed double-precision floating-point values in the source operand (second operand) to two packed single-precision floating-point values in the destination operand (first operand).

When a conversion is inexact, the value returned is rounded according to the rounding control bits in the MXCSR register.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The source operand is an XMM register or 128- bit memory location. The destination operation is an XMM register. Bits[127:64] of the destination XMM register are zeroed. However, the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: The source operand is an XMM register or 128- bit memory location. The destination operation is a YMM register. The upper bits (VLMAX-1:64) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The source operand is a YMM register or 256- bit memory location. The destination operation is an XMM register. The upper bits (255:128) of the corresponding YMM register destination are zeroed.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

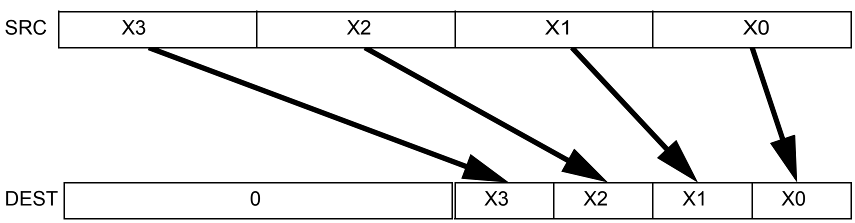


Figure 3-16. VCVTPD2PS (VEX.256 encoded version)

Operation

CVTPD2PS (128-bit Legacy SSE version)

DEST[31:0] ← Convert_Double_Precision_To_Single_Precision_Floating_Point(SRC[63:0])
DEST[63:32] ← Convert_Double_Precision_To_Single_Precision_Floating_Point(SRC[127:64])
DEST[127:64] ← 0
DEST[VLMAX-1:128] (unmodified)

VCVTPD2PS (VEX.128 encoded version)

DEST[31:0] ← Convert_Double_Precision_To_Single_Precision_Floating_Point(SRC[63:0])
DEST[63:32] ← Convert_Double_Precision_To_Single_Precision_Floating_Point(SRC[127:64])
DEST[VLMAX-1:64] ← 0

VCVTPD2PS (VEX.256 encoded version)

DEST[31:0] ← Convert_Double_Precision_To_Single_Precision_Floating_Point(SRC[63:0])
DEST[63:32] ← Convert_Double_Precision_To_Single_Precision_Floating_Point(SRC[127:64])
DEST[95:64] ← Convert_Double_Precision_To_Single_Precision_Floating_Point(SRC[191:128])
DEST[127:96] ← Convert_Double_Precision_To_Single_Precision_Floating_Point(SRC[255:192)

DEST[255:128]← 0

Intel C/C++ Compiler Intrinsic Equivalent

CVTPD2PS __m128 _mm_cvtpd_ps(__m128d a)

CVTPD2PS __m256 _mm256_cvtpd_ps (__m256d a)

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 2; additionally

#UD	If VEX.vvvv != 1111B.

CVTPI2PD—Convert Packed Dword Integers to Packed Double-Precision FP Values

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
66 0F 2A /*r*	CVTPI2PD *xmm, mm/m64**	A	Valid	Valid	Convert two packed signed doubleword integers from *mm/mem64* to two packed double-precision floating-point values in *xmm*.

NOTES:

*Operation is different for different operand sets; see the Description section.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts two packed signed doubleword integers in the source operand (second operand) to two packed double-precision floating-point values in the destination operand (first operand).

The source operand can be an MMX technology register or a 64-bit memory location. The destination operand is an XMM register. In addition, depending on the operand configuration:

- **For operands *xmm, mm*:** the instruction causes a transition from x87 FPU to MMX technology operation (that is, the x87 FPU top-of-stack pointer is set to 0 and the x87 FPU tag word is set to all 0s [valid]). If this instruction is executed while an x87 FPU floating-point exception is pending, the exception is handled before the CVTPI2PD instruction is executed.
- **For operands *xmm, m64*:** the instruction does not cause a transition to MMX technology and does not take x87 FPU exceptions.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Operation

```
DEST[63:0] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC[31:0]);
DEST[127:64] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC[63:32]);
```

Intel C/C++ Compiler Intrinsic Equivalent

```
CVTPI2PD __m128d _mm_cvtpi32_pd(__m64 a)
```

SIMD Floating-Point Exceptions

Precision.

Protected Mode Exceptions

#GP(0)	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
#SS(0)	For an illegal address in the SS segment.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

Real-Address Mode Exceptions

GP	If any part of the operand lies outside the effective address space from 0 to FFFFH.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	For a page fault.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.

#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

CVTPI2PS—Convert Packed Dword Integers to Packed Single-Precision FP Values

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 2A /*r*	CVTPI2PS *xmm*, *mm*/*m64*	A	Valid	Valid	Convert two signed doubleword integers from *mm*/*m64* to two single-precision floating-point values in *xmm*.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts two packed signed doubleword integers in the source operand (second operand) to two packed single-precision floating-point values in the destination operand (first operand).

The source operand can be an MMX technology register or a 64-bit memory location. The destination operand is an XMM register. The results are stored in the low quadword of the destination operand, and the high quadword remains unchanged. When a conversion is inexact, the value returned is rounded according to the rounding control bits in the MXCSR register.

This instruction causes a transition from x87 FPU to MMX technology operation (that is, the x87 FPU top-of-stack pointer is set to 0 and the x87 FPU tag word is set to all 0s [valid]). If this instruction is executed while an x87 FPU floating-point exception is pending, the exception is handled before the CVTPI2PS instruction is executed.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Operation

```
DEST[31:0] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[31:0]);
DEST[63:32] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[63:32]);
(* High quadword of destination unchanged *)
```

Intel C/C++ Compiler Intrinsic Equivalent

```
CVTPI2PS __m128 _mm_cvtpi32_ps(__m128 a, __m64 b)
```

SIMD Floating-Point Exceptions

Precision.

Protected Mode Exceptions

#GP(0)	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
#SS(0)	For an illegal address in the SS segment.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

Real-Address Mode Exceptions

GP	If any part of the operand lies outside the effective address space from 0 to FFFFH.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	For a page fault.

#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

CVTPS2DQ—Convert Packed Single-Precision FP Values to Packed Dword Integers

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 5B /*r* CVTPS2DQ *xmm1*, *xmm2/m128*	A	V/V	SSE2	Convert four packed single-precision floating-point values from *xmm2/m128* to four packed signed doubleword integers in *xmm1*.
VEX.128.66.0F.WIG 5B /r VCVTPS2DQ xmm1, xmm2/m128	A	V/V	AVX	Convert four packed single precision floating-point values from xmm2/mem to four packed signed doubleword values in xmm1.
VEX.256.66.0F.WIG 5B /r VCVTPS2DQ ymm1, ymm2/m256	A	V/V	AVX	Convert eight packed single precision floating-point values from ymm2/mem to eight packed signed doubleword values in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts four or eight packed single-precision floating-point values in the source operand to four or eight signed doubleword integers in the destination operand.

When a conversion is inexact, the value returned is rounded according to the rounding control bits in the MXCSR register. If a converted result is larger than the maximum signed doubleword integer, the floating-point invalid exception is raised, and if this exception is masked, the indefinite integer value (80000000H) is returned.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The source operand is an XMM register or 128- bit memory location. The destination operation is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: The source operand is an XMM register or 128- bit memory location. The destination operation is a YMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The source operand is a YMM register or 256- bit memory location. The destination operation is a YMM register.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

Operation

CVTPS2DQ (128-bit Legacy SSE version)

```
DEST[31:0] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[31:0])
DEST[63:32] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[63:32])
DEST[95:64] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[95:64])
DEST[127:96] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[127:96])
DEST[VLMAX-1:128] (unmodified)
```

VCVTPS2DQ (VEX.128 encoded version)

```
DEST[31:0] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[31:0])
DEST[63:32] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[63:32])
DEST[95:64] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[95:64])
DEST[127:96] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[127:96])
DEST[VLMAX-1:128] ← 0
```

VCVTPS2DQ (VEX.256 encoded version)

```
DEST[31:0] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[31:0])
DEST[63:32] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[63:32])
DEST[95:64] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[95:64])
DEST[127:96] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[127:96)
DEST[159:128] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[159:128])
DEST[191:160] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[191:160])
DEST[223:192] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[223:192])
DEST[255:224] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[255:224])
```

Intel C/C++ Compiler Intrinsic Equivalent

CVTPS2DQ __m128i _mm_cvtps_epi32(__m128 a)

VCVTPS2DQ __ m256i _mm256_cvtps_epi32 (__m256 a)

SIMD Floating-Point Exceptions

Invalid, Precision.

Other Exceptions

See Exceptions Type 2; additionally

#UD	If VEX.vvvv != 1111B.

CVTPS2PD—Convert Packed Single-Precision FP Values to Packed Double-Precision FP Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 5A /*r* CVTPS2PD *xmm1*, *xmm2/m64*	A	V/V	SSE2	Convert two packed single-precision floating-point values in *xmm2/m64* to two packed double-precision floating-point values in *xmm1*.
VEX.128.0F.WIG 5A /r VCVTPS2PD xmm1, xmm2/m64	A	V/V	AVX	Convert two packed single-precision floating-point values in xmm2/mem to two packed double-precision floating-point values in xmm1.
VEX.256.0F.WIG 5A /r VCVTPS2PD ymm1, xmm2/m128	A	V/V	AVX	Convert four packed single-precision floating-point values in xmm2/mem to four packed double-precision floating-point values in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts two or four packed single-precision floating-point values in the source operand (second operand) to two or four packed double-precision floating-point values in the destination operand (first operand).

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The source operand is an XMM register or 64- bit memory location. The destination operation is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: The source operand is an XMM register or 64- bit memory location. The destination operation is a YMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The source operand is an XMM register or 128- bit memory location. The destination operation is a YMM register.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

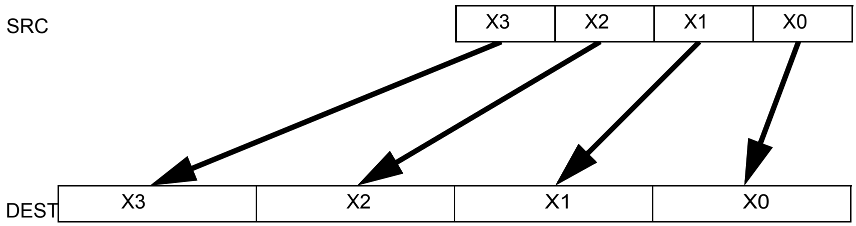


Figure 3-17. CVTPS2PD (VEX.256 encoded version)

Operation

CVTPS2PD (128-bit Legacy SSE version)

DEST[63:0] ← Convert_Single_Precision_To_Double_Precision_Floating_Point(SRC[31:0])
DEST[127:64] ← Convert_Single_Precision_To_Double_Precision_Floating_Point(SRC[63:32])
DEST[VLMAX-1:128] (unmodified)

VCVTPS2PD (VEX.128 encoded version)

DEST[63:0] ← Convert_Single_Precision_To_Double_Precision_Floating_Point(SRC[31:0])
DEST[127:64] ← Convert_Single_Precision_To_Double_Precision_Floating_Point(SRC[63:32])
DEST[VLMAX-1:128] ← 0

VCVTPS2PD (VEX.256 encoded version)

DEST[63:0] ← Convert_Single_Precision_To_Double_Precision_Floating_Point(SRC[31:0])
DEST[127:64] ← Convert_Single_Precision_To_Double_Precision_Floating_Point(SRC[63:32])
DEST[191:128] ← Convert_Single_Precision_To_Double_Precision_Floating_Point(SRC[95:64])
DEST[255:192] ← Convert_Single_Precision_To_Double_Precision_Floating_Point(SRC[127:96)

Intel C/C++ Compiler Intrinsic Equivalent

CVTPS2PD __m128d _mm_cvtps_pd(__m128 a)

VCVTPS2PD __m256d _mm256_cvtps_pd (__m128 a)

SIMD Floating-Point Exceptions

Invalid, Denormal.

Other Exceptions

See Exceptions Type 3; additionally

#UDIf VEX.vvvv != 1111B.

CVTPS2PI—Convert Packed Single-Precision FP Values to Packed Dword Integers

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 2D /*r*	CVTPS2PI *mm, xmm/m64*	A	Valid	Valid	Convert two packed single-precision floating-point values from *xmm/m64* to two packed signed doubleword integers in *mm*.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts two packed single-precision floating-point values in the source operand (second operand) to two packed signed doubleword integers in the destination operand (first operand).

The source operand can be an XMM register or a 128-bit memory location. The destination operand is an MMX technology register. When the source operand is an XMM register, the two single-precision floating-point values are contained in the low quadword of the register. When a conversion is inexact, the value returned is rounded according to the rounding control bits in the MXCSR register. If a converted result is larger than the maximum signed doubleword integer, the floating-point invalid exception is raised, and if this exception is masked, the indefinite integer value (80000000H) is returned.

CVTPS2PI causes a transition from x87 FPU to MMX technology operation (that is, the x87 FPU top-of-stack pointer is set to 0 and the x87 FPU tag word is set to all 0s [valid]). If this instruction is executed while an x87 FPU floating-point exception is pending, the exception is handled before the CVTPS2PI instruction is executed.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Operation

```
DEST[31:0] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[31:0]);
DEST[63:32] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[63:32]);
```

Intel C/C++ Compiler Intrinsic Equivalent

```
CVTPS2PI   __m64 _mm_cvtps_pi32(__m128 a)
```

SIMD Floating-Point Exceptions

Invalid, Precision.

Protected Mode Exceptions

#GP(0)	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
#SS(0)	For an illegal address in the SS segment.
#PF(fault-code)	For a page fault.
#MF	If there is a pending x87 FPU exception.
#NM	If CR0.TS[bit 3] = 1.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

Real-Address Mode Exceptions

GP	If any part of the operand lies outside the effective address space from 0 to FFFFH.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	For a page fault.

#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

CVTSD2SI—Convert Scalar Double-Precision FP Value to Integer

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 2D /*r* CVTSD2SI *r32, xmm/m64*	A	V/V	SSE2	Convert one double-precision floating-point value from *xmm/m64* to one signed doubleword integer *r32*.
F2 REX.W 0F 2D /*r* CVTSD2SI *r64, xmm/m64*	A	V/N.E.	SSE2	Convert one double-precision floating-point value from *xmm/m64* to one signed quadword integer sign-extended into *r64*.
VEX.LIG.F2.0F.W0 2D /r VCVTSD2SI r32, xmm1/m64	A	V/V	AVX	Convert one double precision floating-point value from xmm1/m64 to one signed doubleword integer r32.
VEX.LIG.F2.0F.W1 2D /r VCVTSD2SI r64, xmm1/m64	A	V/N.E.	AVX	Convert one double precision floating-point value from xmm1/m64 to one signed quadword integer sign-extended into r64.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts a double-precision floating-point value in the source operand (second operand) to a signed doubleword integer in the destination operand (first operand). The source operand can be an XMM register or a 64-bit memory location. The destination operand is a general-purpose register. When the source operand is an XMM register, the double-precision floating-point value is contained in the low quadword of the register.

When a conversion is inexact, the value returned is rounded according to the rounding control bits in the MXCSR register. If a converted result is larger than the maximum signed doubleword integer, the floating-point invalid exception is raised, and if this exception is masked, the indefinite integer value (80000000H) is returned.

In 64-bit mode, the instruction can access additional registers (XMM8-XMM15, R8-R15) when used with a REX.R prefix. Use of the REX.W prefix promotes the instruction to 64-bit operation. See the summary chart at the beginning of this section for encoding data and limits.

Legacy SSE instructions: Use of the REX.W prefix promotes the instruction to 64-bit operation. See the summary chart at the beginning of this section for encoding data and limits.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b, otherwise instructions will #UD.

Operation

```
IF 64-Bit Mode and OperandSize = 64
    THEN
        DEST[63:0] ← Convert_Double_Precision_Floating_Point_To_Integer64(SRC[63:0]);
    ELSE
        DEST[31:0] ← Convert_Double_Precision_Floating_Point_To_Integer32(SRC[63:0]);
FI;
```

Intel C/C++ Compiler Intrinsic Equivalent

int _mm_cvtsd_si32(__m128d a)

__int64 _mm_cvtsd_si64(__m128d a)

SIMD Floating-Point Exceptions

Invalid, Precision.

Other Exceptions

See Exceptions Type 3; additionally

#UD	If VEX.vvvv != 1111B.

CVTSD2SS—Convert Scalar Double-Precision FP Value to Scalar Single-Precision FP Value

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 5A /*r* CVTSD2SS *xmm1*, *xmm2/m64*	A	V/V	SSE2	Convert one double-precision floating-point value in *xmm2/m64* to one single-precision floating-point value in *xmm1*.
VEX.NDS.LIG.F2.0F.WIG 5A /r VCVTSD2SS xmm1,xmm2, xmm3/m64	B	V/V	AVX	Convert one double-precision floating-point value in xmm3/m64 to one single-precision floating-point value and merge with high bits in xmm2.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Converts a double-precision floating-point value in the source operand (second operand) to a single-precision floating-point value in the destination operand (first operand).

The source operand can be an XMM register or a 64-bit memory location. The destination operand is an XMM register. When the source operand is an XMM register, the double-precision floating-point value is contained in the low quadword of the register. The result is stored in the low doubleword of the destination operand, and the upper 3 doublewords are left unchanged. When the conversion is inexact, the value returned is rounded according to the rounding control bits in the MXCSR register.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The destination and first source operand are the same. Bits (VLMAX-1:32) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (127:64) of the XMM register destination are copied from corresponding bits in the first source operand. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

CVTSD2SS (128-bit Legacy SSE version)

```
DEST[31:0] ← Convert_Double_Precision_To_Single_Precision_Floating_Point(SRC[63:0]);
(* DEST[VLMAX-1:32] Unmodified *)
```

VCVTSD2SS (VEX.128 encoded version)

```
DEST[31:0] ← Convert_Double_Precision_To_Single_Precision_Floating_Point(SRC2[63:0]);
DEST[127:32] ← SRC1[127:32]
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

CVTSD2SS __m128 _mm_cvtsd_ss(__m128 a, __m128d b)

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 3.

CVTSI2SD—Convert Dword Integer to Scalar Double-Precision FP Value

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 2A /*r* CVTSI2SD *xmm*, *r/m32*	A	V/V	SSE2	Convert one signed doubleword integer from *r/m32* to one double-precision floating-point value in *xmm*.
F2 REX.W 0F 2A /*r* CVTSI2SD *xmm*, *r/m64*	A	V/N.E.	SSE2	Convert one signed quadword integer from *r/m64* to one double-precision floating-point value in *xmm*.
VEX.NDS.LIG.F2.0F.W0 2A /r VCVTSI2SD xmm1, xmm2, r/m32	B	V/V	AVX	Convert one signed doubleword integer from r/m32 to one double-precision floating-point value in xmm1.
VEX.NDS.LIG.F2.0F.W1 2A /r VCVTSI2SD xmm1, xmm2, r/m64	B	V/N.E.	AVX	Convert one signed quadword integer from r/m64 to one double-precision floating-point value in xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Converts a signed doubleword integer (or signed quadword integer if operand size is 64 bits) in the second source operand to a double-precision floating-point value in the destination operand. The result is stored in the low quadword of the destination operand, and the high quadword left unchanged. When conversion is inexact, the value returned is rounded according to the rounding control bits in the MXCSR register.

Legacy SSE instructions: Use of the REX.W prefix promotes the instruction to 64-bit operands. See the summary chart at the beginning of this section for encoding data and limits.

The second source operand can be a general-purpose register or a 32/64-bit memory location. The first source and destination operands are XMM registers.

128-bit Legacy SSE version: The destination and first source operand are the same. Bits (VLMAX-1:64) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (127:64) of the XMM register destination are copied from corresponding bits in the first source operand. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

CVTSI2SD

```
IF 64-Bit Mode And OperandSize = 64
THEN
    DEST[63:0] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC[63:0]);
ELSE
    DEST[63:0] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC[31:0]);
FI;
DEST[VLMAX-1:64] (Unmodified)
```

VCVTSI2SD

```
IF 64-Bit Mode And OperandSize = 64
THEN
    DEST[63:0] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC2[63:0]);
ELSE
    DEST[63:0] ← Convert_Integer_To_Double_Precision_Floating_Point(SRC2[31:0]);
FI;
DEST[127:64] ← SRC1[127:64]
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

CVTSI2SD __m128d _mm_cvtsi32_sd(__m128d a, int b)

CVTSI2SD __m128d _mm_cvtsi64_sd(__m128d a, __int64 b)

SIMD Floating-Point Exceptions

Precision.

Other Exceptions

See Exceptions Type 3.

CVTSI2SS—Convert Dword Integer to Scalar Single-Precision FP Value

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 2A /r CVTSI2SS *xmm*, *r/m32*	A	V/V	SSE	Convert one signed doubleword integer from *r/m32* to one single-precision floating-point value in *xmm*.
F3 REX.W 0F 2A /r CVTSI2SS *xmm*, *r/m64*	A	V/N.E.	SSE	Convert one signed quadword integer from *r/m64* to one single-precision floating-point value in *xmm*.
VEX.NDS.LIG.F3.0F.W0 2A /r VCVTSI2SS xmm1, xmm2, r/m32	B	V/V	AVX	Convert one signed doubleword integer from r/m32 to one single-precision floating-point value in xmm1.
VEX.NDS.LIG.F3.0F.W1 2A /r VCVTSI2SS xmm1, xmm2, r/m64	B	V/N.E.	AVX	Convert one signed quadword integer from r/m64 to one single-precision floating-point value in xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Converts a signed doubleword integer (or signed quadword integer if operand size is 64 bits) in the source operand (second operand) to a single-precision floating-point value in the destination operand (first operand). The source operand can be a general-purpose register or a memory location. The destination operand is an XMM register. The result is stored in the low doubleword of the destination operand, and the upper three doublewords are left unchanged. When a conversion is inexact, the value returned is rounded according to the rounding control bits in the MXCSR register.

Legacy SSE instructions: In 64-bit mode, the instruction can access additional registers (XMM8-XMM15, R8-R15) when used with a REX.R prefix. Use of the REX.W prefix promotes the instruction to 64-bit operands. See the summary chart at the beginning of this section for encoding data and limits.

128-bit Legacy SSE version: The destination and first source operand are the same. Bits (VLMAX-1:32) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (127:32) of the XMM register destination are copied from corresponding bits in the first source operand. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

CVTSI2SS (128-bit Legacy SSE version)

```
IF 64-Bit Mode And OperandSize = 64
THEN
    DEST[31:0] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[63:0]);
ELSE
    DEST[31:0] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[31:0]);
FI;
DEST[VLMAX-1:32] (Unmodified)
```

VCVTSI2SS (VEX.128 encoded version)

```
IF 64-Bit Mode And OperandSize = 64
THEN
    DEST[31:0] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[63:0]);
ELSE
    DEST[31:0] ← Convert_Integer_To_Single_Precision_Floating_Point(SRC[31:0]);
FI;
DEST[127:32] ← SRC1[127:32]
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

CVTSI2SS __m128 _mm_cvtsi32_ss(__m128 a, int b)

CVTSI2SS __m128 _mm_cvtsi64_ss(__m128 a, __int64 b)

SIMD Floating-Point Exceptions

Precision.

Other Exceptions

See Exceptions Type 3.

CVTSS2SD—Convert Scalar Single-Precision FP Value to Scalar Double-Precision FP Value

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 5A /*r* CVTSS2SD *xmm1*, *xmm2/m32*	A	V/V	SSE2	Convert one single-precision floating-point value in *xmm2/m32* to one double-precision floating-point value in *xmm1*.
VEX.NDS.LIG.F3.0F.WIG 5A /r VCVTSS2SD xmm1, xmm2, xmm3/m32	B	V/V	AVX	Convert one single-precision floating-point value in xmm3/m32 to one double-precision floating-point value and merge with high bits of xmm2.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Converts a single-precision floating-point value in the source operand (second operand) to a double-precision floating-point value in the destination operand (first operand). The source operand can be an XMM register or a 32-bit memory location. The destination operand is an XMM register. When the source operand is an XMM register, the single-precision floating-point value is contained in the low doubleword of the register. The result is stored in the low quadword of the destination operand, and the high quadword is left unchanged.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The destination and first source operand are the same. Bits (VLMAX-1:64) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (127:64) of the XMM register destination are copied from corresponding bits in the first source operand. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

CVTSS2SD (128-bit Legacy SSE version)

DEST[63:0] ← Convert_Single_Precision_To_Double_Precision_Floating_Point(SRC[31:0]);
DEST[VLMAX-1:64] (Unmodified)

VCVTSS2SD (VEX.128 encoded version)
DEST[63:0] ← Convert_Single_Precision_To_Double_Precision_Floating_Point(SRC2[31:0])
DEST[127:64] ← SRC1[127:64]
DEST[VLMAX-1:128] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

CVTSS2SD __m128d _mm_cvtss_sd(__m128d a, __m128 b)

SIMD Floating-Point Exceptions

Invalid, Denormal.

Other Exceptions

See Exceptions Type 3.

CVTSS2SI—Convert Scalar Single-Precision FP Value to Dword Integer

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 2D /*r* CVTSS2SI *r32, xmm/m32*	A	V/V	SSE	Convert one single-precision floating-point value from *xmm/m32* to one signed doubleword integer in *r32*.
F3 REX.W 0F 2D /*r* CVTSS2SI *r64, xmm/m32*	A	V/N.E.	SSE	Convert one single-precision floating-point value from *xmm/m32* to one signed quadword integer in *r64*.
VEX.LIG.F3.0F.W0 2D /r VCVTSS2SI r32, xmm1/m32	A	V/V	AVX	Convert one single-precision floating-point value from xmm1/m32 to one signed doubleword integer in r32.
VEX.LIG.F3.0F.W1 2D /r VCVTSS2SI r64, xmm1/m32	A	V/N.E.	AVX	Convert one single-precision floating-point value from xmm1/m32 to one signed quadword integer in r64.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts a single-precision floating-point value in the source operand (second operand) to a signed doubleword integer (or signed quadword integer if operand size is 64 bits) in the destination operand (first operand). The source operand can be an XMM register or a memory location. The destination operand is a general-purpose register. When the source operand is an XMM register, the single-precision floating-point value is contained in the low doubleword of the register.

When a conversion is inexact, the value returned is rounded according to the rounding control bits in the MXCSR register. If a converted result is larger than the maximum signed doubleword integer, the floating-point invalid exception is raised, and if this exception is masked, the indefinite integer value (80000000H) is returned.

In 64-bit mode, the instruction can access additional registers (XMM8-XMM15, R8-R15) when used with a REX.R prefix. Use of the REX.W prefix promotes the instruction to 64-bit operands. See the summary chart at the beginning of this section for encoding data and limits.

Legacy SSE instructions: In 64-bit mode, Use of the REX.W prefix promotes the instruction to 64-bit operands. See the summary chart at the beginning of this section for encoding data and limits.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b, otherwise instructions will #UD.

Operation

```
IF 64-bit Mode and OperandSize = 64
    THEN
        DEST[64:0] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[31:0]);
    ELSE
        DEST[31:0] ← Convert_Single_Precision_Floating_Point_To_Integer(SRC[31:0]);
FI;
```

Intel C/C++ Compiler Intrinsic Equivalent

int _mm_cvtss_si32(__m128d a)

__int64 _mm_cvtss_si64(__m128d a)

SIMD Floating-Point Exceptions

Invalid, Precision.

Other Exceptions

See Exceptions Type 3; additionally

#UFD	If VEX.vvvv != 1111B.

CVTTPD2DQ—Convert with Truncation Packed Double-Precision FP Values to Packed Dword Integers

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F E6 CVTTPD2DQ *xmm1*, *xmm2/m128*	A	V/V	SSE2	Convert two packed double-precision floating-point values from *xmm2/m128* to two packed signed doubleword integers in *xmm1* using truncation.
VEX.128.66.0F.WIG E6 /r VCVTTPD2DQ xmm1, xmm2/m128	A	V/V	AVX	Convert two packed double-precision floating-point values in xmm2/mem to two signed doubleword integers in xmm1 using truncation.
VEX.256.66.0F.WIG E6 /r VCVTTPD2DQ xmm1, ymm2/m256	A	V/V	AVX	Convert four packed double-precision floating-point values in ymm2/mem to four signed doubleword integers in xmm1 using truncation.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts two or four packed double-precision floating-point values in the source operand (second operand) to two or four packed signed doubleword integers in the destination operand (first operand).

When a conversion is inexact, a truncated (round toward zero) value is returned.If a converted result is larger than the maximum signed doubleword integer, the floating-point invalid exception is raised, and if this exception is masked, the indefinite integer value (80000000H) is returned.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The source operand is an XMM register or 128- bit memory location. The destination operation is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: The source operand is an XMM register or 128- bit memory location. The destination operation is a YMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The source operand is a YMM register or 256- bit memory location. The destination operation is an XMM register. The upper bits (255:128) of the corresponding YMM register destination are zeroed.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b, otherwise instructions will #UD.

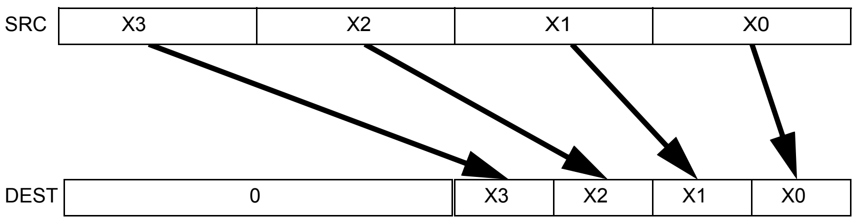


Figure 3-18. VCVTTPD2DQ (VEX.256 encoded version)

Operation

CVTTPD2DQ (128-bit Legacy SSE version)

DEST[31:0] ← Convert_Double_Precision_Floating_Point_To_Integer_Truncate(SRC[63:0])
DEST[63:32] ← Convert_Double_Precision_Floating_Point_To_Integer_Truncate(SRC[127:64])
DEST[127:64] ← 0
DEST[VLMAX-1:128] (unmodified)

VCVTTPD2DQ (VEX.128 encoded version)

DEST[31:0] ← Convert_Double_Precision_Floating_Point_To_Integer_Truncate(SRC[63:0])
DEST[63:32] ← Convert_Double_Precision_Floating_Point_To_Integer_Truncate(SRC[127:64])
DEST[VLMAX-1:64] ← 0

VCVTTPD2DQ (VEX.256 encoded version)

DEST[31:0] ← Convert_Double_Precision_Floating_Point_To_Integer_Truncate(SRC[63:0])
DEST[63:32] ← Convert_Double_Precision_Floating_Point_To_Integer_Truncate(SRC[127:64])
DEST[95:64] ← Convert_Double_Precision_Floating_Point_To_Integer_Truncate(SRC[191:128])
DEST[127:96] ← Convert_Double_Precision_Floating_Point_To_Integer_Truncate(SRC[255:192)

DEST[255:128]← 0

Intel C/C++ Compiler Intrinsic Equivalent

CVTTPD2DQ __m128i _mm_cvttpd_epi32(__m128d a)

VCVTTPD2DQ __m128i _mm256_cvttpd_epi32 (__m256d src)

SIMD Floating-Point Exceptions

Invalid, Precision.

Other Exceptions

See Exceptions Type 2; additionally

#UFD	If VEX.vvvv != 1111B.

CVTTPD2PI—Convert with Truncation Packed Double-Precision FP Values to Packed Dword Integers

Opcode/ Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
66 0F 2C /*r* CVTTPD2PI *mm, xmm/m128*	A	Valid	Valid	Convert two packer double-precision floating-point values from *xmm/m128* to two packed signed doubleword integers in *mm* using truncation.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts two packed double-precision floating-point values in the source operand (second operand) to two packed signed doubleword integers in the destination operand (first operand). The source operand can be an XMM register or a 128-bit memory location. The destination operand is an MMX technology register.

When a conversion is inexact, a truncated (round toward zero) result is returned. If a converted result is larger than the maximum signed doubleword integer, the floating-point invalid exception is raised, and if this exception is masked, the indefinite integer value (80000000H) is returned.

This instruction causes a transition from x87 FPU to MMX technology operation (that is, the x87 FPU top-of-stack pointer is set to 0 and the x87 FPU tag word is set to all 0s [valid]). If this instruction is executed while an x87 FPU floating-point exception is pending, the exception is handled before the CVTTPD2PI instruction is executed.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Operation

```
DEST[31:0] ← Convert_Double_Precision_Floating_Point_To_Integer32_Truncate(SRC[63:0]);
DEST[63:32] ← Convert_Double_Precision_Floating_Point_To_Integer32_
                    Truncate(SRC[127:64]);
```

Intel C/C++ Compiler Intrinsic Equivalent

CVTTPD1PI __m64 _mm_cvttpd_pi32(__m128d a)

SIMD Floating-Point Exceptions

Invalid, Precision.

Protected Mode Exceptions

#GP(0)	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
	If a memory operand is not aligned on a 16-byte boundary, regardless of segment.
#SS(0)	For an illegal address in the SS segment.
#PF(fault-code)	For a page fault.
#MF	If there is a pending x87 FPU exception.
#NM	If CR0.TS[bit 3] = 1.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand is not aligned on a 16-byte boundary, regardless of segment.
	If any part of the operand lies outside the effective address space from 0 to FFFFH.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	For a page fault.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
	If memory operand is not aligned on a 16-byte boundary, regardless of segment.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.

CVTTPS2DQ—Convert with Truncation Packed Single-Precision FP Values to Packed Dword Integers

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 5B /*r* CVTTPS2DQ *xmm1*, *xmm2/m128*	A	V/V	SSE2	Convert four single-precision floating-point values from *xmm2/m128* to four signed doubleword integers in *xmm1* using truncation.
VEX.128.F3.0F.WIG 5B /r VCVTTPS2DQ xmm1, xmm2/m128	A	V/V	AVX	Convert four packed single precision floating-point values from xmm2/mem to four packed signed doubleword values in xmm1 using truncation.
VEX.256.F3.0F.WIG 5B /r VCVTTPS2DQ ymm1, ymm2/m256	A	V/V	AVX	Convert eight packed single precision floating-point values from ymm2/mem to eight packed signed doubleword values in ymm1 using truncation.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts four or eight packed single-precision floating-point values in the source operand to four or eight signed doubleword integers in the destination operand.

When a conversion is inexact, a truncated (round toward zero) value is returned.If a converted result is larger than the maximum signed doubleword integer, the floating-point invalid exception is raised, and if this exception is masked, the indefinite integer value (80000000H) is returned.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The source operand is an XMM register or 128- bit memory location. The destination operation is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: The source operand is an XMM register or 128- bit memory location. The destination operation is a YMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The source operand is a YMM register or 256- bit memory location. The destination operation is a YMM register.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

Operation

CVTTPS2DQ (128-bit Legacy SSE version)

```
DEST[31:0] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[31:0])
DEST[63:32] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[63:32])
DEST[95:64] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[95:64])
DEST[127:96] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[127:96])
DEST[VLMAX-1:128] (unmodified)
```

VCVTTPS2DQ (VEX.128 encoded version)

```
DEST[31:0] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[31:0])
DEST[63:32] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[63:32])
DEST[95:64] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[95:64])
DEST[127:96] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[127:96])
DEST[VLMAX-1:128] ← 0
```

VCVTTPS2DQ (VEX.256 encoded version)

```
DEST[31:0] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[31:0])
DEST[63:32] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[63:32])
DEST[95:64] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[95:64])
DEST[127:96] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[127:96)
DEST[159:128] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[159:128])
DEST[191:160] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[191:160])
DEST[223:192] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[223:192])
DEST[255:224] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[255:224])
```

Intel C/C++ Compiler Intrinsic Equivalent

CVTTPS2DQ __m128i _mm_cvttps_epi32(__m128 a)

VCVTTPS2DQ __m256i _mm256_cvttps_epi32 (__m256 a)

SIMD Floating-Point Exceptions

Invalid, Precision.

Other Exceptions

See Exceptions Type 2; additionally

#UD	If VEX.vvvv != 1111B.

CVTTPS2PI—Convert with Truncation Packed Single-Precision FP Values to Packed Dword Integers

Opcode/ Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 2C /r CVTTPS2PI *mm, xmm/m64*	A	Valid	Valid	Convert two single-precision floating-point values from *xmm/m64* to two signed doubleword signed integers in *mm* using truncation.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts two packed single-precision floating-point values in the source operand (second operand) to two packed signed doubleword integers in the destination operand (first operand). The source operand can be an XMM register or a 64-bit memory location. The destination operand is an MMX technology register. When the source operand is an XMM register, the two single-precision floating-point values are contained in the low quadword of the register.

When a conversion is inexact, a truncated (round toward zero) result is returned. If a converted result is larger than the maximum signed doubleword integer, the floating-point invalid exception is raised, and if this exception is masked, the indefinite integer value (80000000H) is returned.

This instruction causes a transition from x87 FPU to MMX technology operation (that is, the x87 FPU top-of-stack pointer is set to 0 and the x87 FPU tag word is set to all 0s [valid]). If this instruction is executed while an x87 FPU floating-point exception is pending, the exception is handled before the CVTTPS2PI instruction is executed.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Operation

```
DEST[31:0] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[31:0]);
DEST[63:32] ← Convert_Single_Precision_Floating_Point_To_Integer_Truncate(SRC[63:32]);
```

Intel C/C++ Compiler Intrinsic Equivalent

CVTTPS2PI __m64 _mm_cvttps_pi32(__m128 a)

SIMD Floating-Point Exceptions

Invalid, Precision.

Protected Mode Exceptions

#GP(0)	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
#SS(0)	For an illegal address in the SS segment.
#PF(fault-code)	For a page fault.
#MF	If there is a pending x87 FPU exception.
#NM	If CR0.TS[bit 3] = 1.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

Real-Address Mode Exceptions

GP	If any part of the operand lies outside the effective address space from 0 to FFFFH.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	For a page fault.

#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#XM	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 1.
#UD	If an unmasked SIMD floating-point exception and CR4.OSXMMEXCPT[bit 10] = 0.
	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

CVTTSD2SI—Convert with Truncation Scalar Double-Precision FP Value to Signed Integer

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 2C /*r* CVTTSD2SI *r32, xmm/m64*	A	V/V	SSE2	Convert one double-precision floating-point value from *xmm/m64* to one signed doubleword integer in *r32* using truncation.
F2 REX.W 0F 2C /*r* CVTTSD2SI *r64, xmm/m64*	A	V/N.E.	SSE2	Convert one double precision floating-point value from *xmm/m64* to one signedquadword integer in *r64* using truncation.
VEX.LIG.F2.0F.W0 2C /r VCVTTSD2SI r32, xmm1/m64	A	V/V	AVX	Convert one double-precision floating-point value from xmm1/m64 to one signed doubleword integer in r32 using truncation.
VEX.LIG.F2.0F.W1 2C /r VCVTTSD2SI r64, xmm1/m64	A	V/N.E.	AVX	Convert one double precision floating-point value from xmm1/m64 to one signed quadword integer in r64 using truncation.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts a double-precision floating-point value in the source operand (second operand) to a signed doubleword integer (or signed quadword integer if operand size is 64 bits) in the destination operand (first operand). The source operand can be an XMM register or a 64-bit memory location. The destination operand is a general purpose register. When the source operand is an XMM register, the double-precision floating-point value is contained in the low quadword of the register.

When a conversion is inexact, a truncated (round toward zero) result is returned. If a converted result is larger than the maximum signed doubleword integer, the floating point invalid exception is raised. If this exception is masked, the indefinite integer value (80000000H) is returned.

Legacy SSE instructions: In 64-bit mode, Use of the REX.W prefix promotes the instruction to 64-bit operation. See the summary chart at the beginning of this section for encoding data and limits.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b, otherwise instructions will #UD.

Operation

```
IF 64-Bit Mode and OperandSize = 64
    THEN
        DEST[63:0] ← Convert_Double_Precision_Floating_Point_To_
                            Integer64_Truncate(SRC[63:0]);
    ELSE
        DEST[31:0] ← Convert_Double_Precision_Floating_Point_To_
                            Integer32_Truncate(SRC[63:0]);
FI;
```

Intel C/C++ Compiler Intrinsic Equivalent

int _mm_cvttsd_si32(__m128d a)

__int64 _mm_cvttsd_si64(__m128d a)

SIMD Floating-Point Exceptions

Invalid, Precision.

Other Exceptions

See Exceptions Type 3; additionally

#UD	If VEX.vvvv != 1111B.

CVTTSS2SI—Convert with Truncation Scalar Single-Precision FP Value to Dword Integer

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 2C /r CVTTSS2SI *r32, xmm/m32*	A	V/V	SSE	Convert one single-precision floating-point value from *xmm/m32* to one signed doubleword integer in *r32* using truncation.
F3 REX.W 0F 2C /r CVTTSS2SI *r64, xmm/m32*	A	V/N.E.	SSE	Convert one single-precision floating-point value from *xmm/m32* to one signed quadword integer in *r64* using truncation.
VEX.LIG.F3.0F.W0 2C /r VCVTTSS2SI r32, xmm1/m32	A	V/V	AVX	Convert one single-precision floating-point value from xmm1/m32 to one signed doubleword integer in r32 using truncation.
VEX.LIG.F3.0F.W1 2C /r VCVTTSS2SI r64, xmm1/m32	A	V/N.E.	AVX	Convert one single-precision floating-point value from xmm1/m32 to one signed quadword integer in r64 using truncation.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Converts a single-precision floating-point value in the source operand (second operand) to a signed doubleword integer (or signed quadword integer if operand size is 64 bits) in the destination operand (first operand). The source operand can be an XMM register or a 32-bit memory location. The destination operand is a general-purpose register. When the source operand is an XMM register, the single-precision floating-point value is contained in the low doubleword of the register.

When a conversion is inexact, a truncated (round toward zero) result is returned. If a converted result is larger than the maximum signed doubleword integer, the floating-point invalid exception is raised. If this exception is masked, the indefinite integer value (80000000H) is returned.

Legacy SSE instructions: In 64-bit mode, the instruction can access additional registers (XMM8-XMM15, R8-R15) when used with a REX.R prefix. Use of the REX.W prefix promotes the instruction to 64-bit operation. See the summary chart at the beginning of this section for encoding data and limits.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b, otherwise instructions will #UD.

Operation

```
IF 64-Bit Mode and OperandSize = 64
    THEN
        DEST[63:0] ← Convert_Single_Precision_Floating_Point_To_
                            Integer_Truncate(SRC[31:0]);
    ELSE
        DEST[31:0] ← Convert_Single_Precision_Floating_Point_To_
                            Integer_Truncate(SRC[31:0]);
FI;
```

Intel C/C++ Compiler Intrinsic Equivalent

int _mm_cvttss_si32(__m128d a)

__int64 _mm_cvttss_si64(__m128d a)

SIMD Floating-Point Exceptions

Invalid, Precision.

Other Exceptions

See Exceptions Type 3; additionally

#UD	If VEX.vvvv != 1111B.

CWD/CDQ/CQO—Convert Word to Doubleword/Convert Doubleword to Quadword

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
99	CWD	A	Valid	Valid	DX:AX ← sign-extend of AX.
99	CDQ	A	Valid	Valid	EDX:EAX ← sign-extend of EAX.
REX.W + 99	CQO	A	Valid	N.E.	RDX:RAX← sign-extend of RAX.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Doubles the size of the operand in register AX, EAX, or RAX (depending on the operand size) by means of sign extension and stores the result in registers DX:AX, EDX:EAX, or RDX:RAX, respectively. The CWD instruction copies the sign (bit 15) of the value in the AX register into every bit position in the DX register. The CDQ instruction copies the sign (bit 31) of the value in the EAX register into every bit position in the EDX register. The CQO instruction (available in 64-bit mode only) copies the sign (bit 63) of the value in the RAX register into every bit position in the RDX register.

The CWD instruction can be used to produce a doubleword dividend from a word before word division. The CDQ instruction can be used to produce a quadword dividend from a doubleword before doubleword division. The CQO instruction can be used to produce a double quadword dividend from a quadword before a quadword division.

The CWD and CDQ mnemonics reference the same opcode. The CWD instruction is intended for use when the operand-size attribute is 16 and the CDQ instruction for when the operand-size attribute is 32. Some assemblers may force the operand size to 16 when CWD is used and to 32 when CDQ is used. Others may treat these mnemonics as synonyms (CWD/CDQ) and use the current setting of the operand-size attribute to determine the size of values to be converted, regardless of the mnemonic used.

In 64-bit mode, use of the REX.W prefix promotes operation to 64 bits. The CQO mnemonics reference the same opcode as CWD/CDQ. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
IF OperandSize = 16 (* CWD instruction *)
    THEN
        DX ← SignExtend(AX);
    ELSE IF OperandSize = 32 (* CDQ instruction *)
        EDX ← SignExtend(EAX); FI;
    ELSE IF 64-Bit Mode and OperandSize = 64 (* CQO instruction*)
        RDX ← SignExtend(RAX); FI;
FI;
```

Flags Affected

None.

Exceptions (All Operating Modes)

#UD	If the LOCK prefix is used.

DAA—Decimal Adjust AL after Addition

Opcode	Instruction	Op/En	64-Bit Mode	Compat/Leg Mode	Description
27	DAA	A	Invalid	Valid	Decimal adjust AL after addition.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Adjusts the sum of two packed BCD values to create a packed BCD result. The AL register is the implied source and destination operand. The DAA instruction is only useful when it follows an ADD instruction that adds (binary addition) two 2-digit, packed BCD values and stores a byte result in the AL register. The DAA instruction then adjusts the contents of the AL register to contain the correct 2-digit, packed BCD result. If a decimal carry is detected, the CF and AF flags are set accordingly.

This instruction executes as described above in compatibility mode and legacy mode. It is not valid in 64-bit mode.

Operation

```
IF 64-Bit Mode
    THEN
        #UD;
    ELSE
        old_AL ← AL;
        old_CF ← CF;
        CF ← 0;
        IF (((AL AND 0FH) > 9) or AF = 1)
            THEN
                AL ← AL + 6;
                CF ← old_CF or (Carry from AL ← AL + 6);
                AF ← 1;
            ELSE
                AF ← 0;
        FI;
        IF ((old_AL > 99H) or (old_CF = 1))
            THEN
                AL ← AL + 60H;
                CF ← 1;
```

```
            ELSE
                CF ← 0;
        FI;
FI;
```

Example

```
ADD    AL, BL    Before: AL=79H BL=35H EFLAGS(OSZAPC)=XXXXXX
                 After: AL=AEH BL=35H EFLAGS(0SZAPC)=110000
DAA              Before: AL=AEH BL=35H EFLAGS(OSZAPC)=110000
                 After: AL=14H BL=35H EFLAGS(0SZAPC)=X00111
DAA              Before: AL=2EH BL=35H EFLAGS(OSZAPC)=110000
                 After: AL=34H BL=35H EFLAGS(0SZAPC)=X00101
```

Flags Affected

The CF and AF flags are set if the adjustment of the value results in a decimal carry in either digit of the result (see the "Operation" section above). The SF, ZF, and PF flags are set according to the result. The OF flag is undefined.

Protected Mode Exceptions

#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

#UD	If the LOCK prefix is used.

64-Bit Mode Exceptions

#UD	If in 64-bit mode.

DAS—Decimal Adjust AL after Subtraction

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
2F	DAS	A	Invalid	Valid	Decimal adjust AL after subtraction.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Adjusts the result of the subtraction of two packed BCD values to create a packed BCD result. The AL register is the implied source and destination operand. The DAS instruction is only useful when it follows a SUB instruction that subtracts (binary subtraction) one 2-digit, packed BCD value from another and stores a byte result in the AL register. The DAS instruction then adjusts the contents of the AL register to contain the correct 2-digit, packed BCD result. If a decimal borrow is detected, the CF and AF flags are set accordingly.

This instruction executes as described above in compatibility mode and legacy mode. It is not valid in 64-bit mode.

Operation

```
IF 64-Bit Mode
    THEN
        #UD;
    ELSE
        old_AL ← AL;
        old_CF ← CF;
        CF ← 0;
        IF (((AL AND 0FH) > 9) or AF = 1)
            THEN
                AL ← AL - 6;
                CF ← old_CF or (Borrow from AL ← AL − 6);
                AF ← 1;
            ELSE
                AF ← 0;
        FI;
        IF ((old_AL > 99H) or (old_CF = 1))
            THEN
                AL ← AL − 60H;
```

```
            CF ← 1;
    FI;
FI;
```

Example

SUB	AL, BL	Before: AL = 35H, BL = 47H, EFLAGS(OSZAPC) = XXXXXX
		After: AL = EEH, BL = 47H, EFLAGS(OSZAPC) = 010111
DAA		Before: AL = EEH, BL = 47H, EFLAGS(OSZAPC) = 010111
		After: AL = 88H, BL = 47H, EFLAGS(OSZAPC) = X10111

Flags Affected

The CF and AF flags are set if the adjustment of the value results in a decimal borrow in either digit of the result (see the "Operation" section above). The SF, ZF, and PF flags are set according to the result. The OF flag is undefined.

Protected Mode Exceptions

#UD If the LOCK prefix is used.

Real-Address Mode Exceptions

#UD If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#UD If the LOCK prefix is used.

Compatibility Mode Exceptions

#UD If the LOCK prefix is used.

64-Bit Mode Exceptions

#UD If in 64-bit mode.

DEC—Decrement by 1

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
FE /1	DEC *r/m8*	A	Valid	Valid	Decrement *r/m8* by 1.
REX + FE /1	DEC *r/m8**	A	Valid	N.E.	Decrement *r/m8* by 1.
FF /1	DEC *r/m16*	A	Valid	Valid	Decrement *r/m16* by 1.
FF /1	DEC *r/m32*	A	Valid	Valid	Decrement *r/m32* by 1.
REX.W + FF /1	DEC *r/m64*	A	Valid	N.E.	Decrement *r/m64* by 1.
48+rw	DEC *r16*	B	N.E.	Valid	Decrement *r16* by 1.
48+rd	DEC *r32*	B	N.E.	Valid	Decrement *r32* by 1.

NOTES:

* In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r, w)	NA	NA	NA
B	reg (r, w)	NA	NA	NA

Description

Subtracts 1 from the destination operand, while preserving the state of the CF flag. The destination operand can be a register or a memory location. This instruction allows a loop counter to be updated without disturbing the CF flag. (To perform a decrement operation that updates the CF flag, use a SUB instruction with an immediate operand of 1.)

This instruction can be used with a LOCK prefix to allow the instruction to be executed atomically.

In 64-bit mode, DEC r16 and DEC r32 are not encodable (because opcodes 48H through 4FH are REX prefixes). Otherwise, the instruction's 64-bit mode default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits.

See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
DEST ← DEST – 1;
```

Flags Affected

The CF flag is not affected. The OF, SF, ZF, AF, and PF flags are set according to the result.

Protected Mode Exceptions

#GP(0)	If the destination operand is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

DIV—Unsigned Divide

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
F6 /6	DIV *r/m8*	A	Valid	Valid	Unsigned divide AX by *r/m8*, with result stored in AL ← Quotient, AH ← Remainder.
REX + F6 /6	DIV *r/m8**	A	Valid	N.E.	Unsigned divide AX by *r/m8*, with result stored in AL ← Quotient, AH ← Remainder.
F7 /6	DIV *r/m16*	A	Valid	Valid	Unsigned divide DX:AX by *r/m16*, with result stored in AX ← Quotient, DX ← Remainder.
F7 /6	DIV *r/m32*	A	Valid	Valid	Unsigned divide EDX:EAX by *r/m32*, with result stored in EAX ← Quotient, EDX ← Remainder.
REX.W + F7 /6	DIV *r/m64*	A	Valid	N.E.	Unsigned divide RDX:RAX by *r/m64*, with result stored in RAX ← Quotient, RDX ← Remainder.

NOTES:

* In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (w)	NA	NA	NA

Description

Divides unsigned the value in the AX, DX:AX, EDX:EAX, or RDX:RAX registers (dividend) by the source operand (divisor) and stores the result in the AX (AH:AL), DX:AX, EDX:EAX, or RDX:RAX registers. The source operand can be a general-purpose register or a memory location. The action of this instruction depends on the operand size (dividend/divisor). Division using 64-bit operand is available only in 64-bit mode.

Non-integral results are truncated (chopped) towards 0. The remainder is always less than the divisor in magnitude. Overflow is indicated with the #DE (divide error) exception rather than with the CF flag.

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits. In 64-bit mode when REX.W is applied, the instruction divides the unsigned value in RDX:RAX by the source operand and stores the quotient in RAX, the remainder in RDX.

See the summary chart at the beginning of this section for encoding data and limits. See Table 3-25.

Table 3-25. DIV Action

Operand Size	Dividend	Divisor	Quotient	Remainder	Maximum Quotient
Word/byte	AX	r/m8	AL	AH	255
Doubleword/word	DX:AX	r/m16	AX	DX	65,535
Quadword/doubleword	EDX:EAX	r/m32	EAX	EDX	$2^{32} - 1$
Doublequadword/ quadword	RDX:RAX	r/m64	RAX	RDX	$2^{64} - 1$

Operation

```
IF SRC = 0
    THEN #DE; FI; (* Divide Error *)
IF OperandSize = 8 (* Word/Byte Operation *)
    THEN
        temp ← AX / SRC;
        IF temp > FFH
            THEN #DE; (* Divide error *)
            ELSE
                AL ← temp;
                AH ← AX MOD SRC;
        FI;
    ELSE IF OperandSize = 16 (* Doubleword/word operation *)
        THEN
            temp ← DX:AX / SRC;
            IF temp > FFFFH
                THEN #DE; (* Divide error *)
            ELSE
                AX ← temp;
                DX ← DX:AX MOD SRC;
            FI;
        FI;
    ELSE IF Operandsize = 32 (* Quadword/doubleword operation *)
        THEN
            temp ← EDX:EAX / SRC;
```

```
            IF temp > FFFFFFFFH
                THEN #DE; (* Divide error *)
            ELSE
                EAX ← temp;
                EDX ← EDX:EAX MOD SRC;
            FI;
        FI;
    ELSE IF 64-Bit Mode and Operandsize = 64 (* Doublequadword/quadword operation *)
        THEN
            temp ← RDX:RAX / SRC;
            IF temp > FFFFFFFFFFFFFFFFH
                THEN #DE; (* Divide error *)
            ELSE
                RAX ← temp;
                RDX ← RDX:RAX MOD SRC;
            FI;
        FI;
FI;
```

Flags Affected

The CF, OF, SF, ZF, AF, and PF flags are undefined.

Protected Mode Exceptions

#DE	If the source operand (divisor) is 0
	If the quotient is too large for the designated register.
#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#DE	If the source operand (divisor) is 0.
	If the quotient is too large for the designated register.
#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.

	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#DE	If the source operand (divisor) is 0.
	If the quotient is too large for the designated register.
#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#DE	If the source operand (divisor) is 0
	If the quotient is too large for the designated register.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

DIVPD—Divide Packed Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 5E /*r* DIVPD *xmm1*, *xmm2/m128*	A	V/V	SSE2	Divide packed double-precision floating-point values in *xmm1* by packed double-precision floating-point values *xmm2/m128.*
VEX.NDS.128.66.0F.WIG 5E /r VDIVPD xmm1, xmm2, xmm3/m128	B	V/V	AVX	Divide packed double-precision floating-point values in xmm2 by packed double-precision floating-point values in xmm3/mem.
VEX.NDS.256.66.0F.WIG 5E /r VDIVPD ymm1, ymm2, ymm3/m256	B	V/V	AVX	Divide packed double-precision floating-point values in ymm2 by packed double-precision floating-point values in ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs an SIMD divide of the two or four packed double-precision floating-point values in the first source operand by the two or four packed double-precision floating-point values in the second source operand. See Chapter 11 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for an overview of a SIMD double-precision floating-point operation.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

DIVPD (128-bit Legacy SSE version)

DEST[63:0] ← SRC1[63:0] / SRC2[63:0]
DEST[127:64] ← SRC1[127:64] / SRC2[127:64]
DEST[VLMAX-1:128] (Unmodified)

VDIVPD (VEX.128 encoded version)

DEST[63:0] ← SRC1[63:0] / SRC2[63:0]
DEST[127:64] ← SRC1[127:64] / SRC2[127:64]
DEST[VLMAX-1:128] ← 0

VDIVPD (VEX.256 encoded version)

DEST[63:0] ← SRC1[63:0] / SRC2[63:0]
DEST[127:64] ← SRC1[127:64] / SRC2[127:64]
DEST[191:128] ← SRC1[191:128] / SRC2[191:128]
DEST[255:192] ← SRC1[255:192] / SRC2[255:192]

Intel C/C++ Compiler Intrinsic Equivalent

DIVPD __m128d _mm_div_pd(__m128d a, __m128d b)

VDIVPD __m256d _mm256_div_pd (__m256d a, __m256d b);

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Divide-by-Zero, Precision, Denormal.

Other Exceptions

See Exceptions Type 2.

DIVPS—Divide Packed Single-Precision Floating-Point Values

Opcode	Instruction	Op/En	64/32-bit Mode	CPUID Feature Flag	Description
0F 5E /r	DIVPS *xmm1, xmm2/m128*	A	V/V	SSE	Divide packed single-precision floating-point values in *xmm1* by packed single-precision floating-point values *xmm2/m128*.
VEX.NDS.128.0F.WIG 5E /r VDIVPS xmm1, xmm2, xmm3/m128		B	V/V	AVX	Divide packed single-precision floating-point values in xmm2 by packed double-precision floating-point values in xmm3/mem.
VEX.NDS.256.0F.WIG 5E /r VDIVPS ymm1, ymm2, ymm3/m256		B	V/V	AVX	Divide packed single-precision floating-point values in ymm2 by packed double-precision floating-point values in ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs an SIMD divide of the four or eight packed single-precision floating-point values in the first source operand by the four or eight packed single-precision floating-point values in the second source operand. See Chapter 10 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for an overview of a SIMD single-precision floating-point operation.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

DIVPS (128-bit Legacy SSE version)

DEST[31:0] ← SRC1[31:0] / SRC2[31:0]
DEST[63:32] ← SRC1[63:32] / SRC2[63:32]
DEST[95:64] ← SRC1[95:64] / SRC2[95:64]
DEST[127:96] ← SRC1[127:96] / SRC2[127:96]
DEST[VLMAX-1:128] (Unmodified)

VDIVPS (VEX.128 encoded version)

DEST[31:0] ← SRC1[31:0] / SRC2[31:0]
DEST[63:32] ← SRC1[63:32] / SRC2[63:32]
DEST[95:64] ← SRC1[95:64] / SRC2[95:64]
DEST[127:96] ← SRC1[127:96] / SRC2[127:96]
DEST[VLMAX-1:128] ← 0

VDIVPS (VEX.256 encoded version)

DEST[31:0] ← SRC1[31:0] / SRC2[31:0]
DEST[63:32] ← SRC1[63:32] / SRC2[63:32]
DEST[95:64] ← SRC1[95:64] / SRC2[95:64]
DEST[127:96] ← SRC1[127:96] / SRC2[127:96]
DEST[159:128] ← SRC1[159:128] / SRC2[159:128]
DEST[191:160]← SRC1[191:160] / SRC2[191:160]
DEST[223:192] ← SRC1[223:192] / SRC2[223:192]
DEST[255:224] ← SRC1[255:224] / SRC2[255:224].

Intel C/C++ Compiler Intrinsic Equivalent

DIVPS __m128 _mm_div_ps(__m128 a, __m128 b)

VDIVPS __m256 _mm256_div_ps (__m256 a, __m256 b);

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Divide-by-Zero, Precision, Denormal.

Other Exceptions

See Exceptions Type 2.

DIVSD—Divide Scalar Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 5E /*r* DIVSD *xmm1*, *xmm2/m64*	A	V/V	SSE2	Divide low double-precision floating-point value n *xmm1* by low double-precision floating-point value in *xmm2/mem64.*
VEX.NDS.LIG.F2.0F.WIG 5E /r VDIVSD xmm1, xmm2, xmm3/m64	A	V/V	AVX	Divide low double-precision floating point values in xmm2 by low double precision floating-point value in xmm3/mem64.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Divides the low double-precision floating-point value in the first source operand by the low double-precision floating-point value in the second source operand, and stores the double-precision floating-point result in the destination operand. The second source operand can be an XMM register or a 64-bit memory location. The first source and destination hyperons are XMM registers. The high quadword of the destination operand is copied from the high quadword of the first source operand. See Chapter 11 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for an overview of a scalar double-precision floating-point operation.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The first source operand and the destination operand are the same. Bits (VLMAX-1:64) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

DIVSD (128-bit Legacy SSE version)

DEST[63:0] ← DEST[63:0] / SRC[63:0]

DEST[VLMAX-1:64] (Unmodified)

VDIVSD (VEX.128 encoded version)
DEST[63:0] ← SRC1[63:0] / SRC2[63:0]
DEST[127:64] ← SRC1[127:64]
DEST[VLMAX-1:128] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

DIVSD __m128d _mm_div_sd (m128d a, m128d b)

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Divide-by-Zero, Precision, Denormal.

Other Exceptions

See Exceptions Type 3.

DIVSS—Divide Scalar Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 5E /*r* DIVSS *xmm1*, *xmm2/m32*	A	V/V	SSE	Divide low single-precision floating-point value in *xmm1* by low single-precision floating-point value in *xmm2/m32.*
VEX.NDS.LIG.F3.0F.WIG 5E /r VDIVSS xmm1, xmm2, xmm3/m32	B	V/V	AVX	Divide low single-precision floating point value in xmm2 by low single precision floating-point value in xmm3/m32.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Divides the low single-precision floating-point value in the first source operand by the low single-precision floating-point value in the second source operand, and stores the single-precision floating-point result in the destination operand. The second source operand can be an XMM register or a 32-bit memory location. The first source and destination operands are XMM registers. The three high-order doublewords of the destination are copied from the same dwords of the first source operand. See Chapter 10 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for an overview of a scalar single-precision floating-point operation.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The first source operand and the destination operand are the same. Bits (VLMAX-1:32) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

DIVSS (128-bit Legacy SSE version)

DEST[31:0] ← DEST[31:0] / SRC[31:0]

DEST[VLMAX-1:32] (Unmodified)

VDIVSS (VEX.128 encoded version)
DEST[31:0] ← SRC1[31:0] / SRC2[31:0]
DEST[127:32] ← SRC1[127:32]
DEST[VLMAX-1:128] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

DIVSS __m128 _mm_div_ss(__m128 a, __m128 b)

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Divide-by-Zero, Precision, Denormal.

Other Exceptions

See Exceptions Type 3.

DPPD — Dot Product of Packed Double Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 3A 41 /r ib DPPD *xmm1, xmm2/m128, imm8*	A	V/V	SSE4_1	Selectively multiply packed DP floating-point values from *xmm1* with packed DP floating-point values from *xmm2*, add and selectively store the packed DP floating-point values to *xmm1*.
VEX.NDS.128.66.0F3A.WIG 41 /r ib VDPPD xmm1,xmm2, xmm3/m128, imm8	B	V/V	AVX	Selectively multiply packed DP floating-point values from xmm2 with packed DP floating-point values from xmm3, add and selectively store the packed DP floating-point values to xmm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	imm8	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Conditionally multiplies the packed double-precision floating-point values in the destination operand (first operand) with the packed double-precision floating-point values in the source (second operand) depending on a mask extracted from bits [5:4] of the immediate operand (third operand). If a condition mask bit is zero, the corresponding multiplication is replaced by a value of 0.0.

The two resulting double-precision values are summed into an intermediate result. The intermediate result is conditionally broadcasted to the destination using a broadcast mask specified by bits [1:0] of the immediate byte.

If a broadcast mask bit is "1", the intermediate result is copied to the corresponding qword element in the destination operand. If a broadcast mask bit is zero, the corresponding element in the destination is set to zero.

DPPS follows the NaN forwarding rules stated in the Software Developer's Manual, vol. 1, table 4.7. These rules do not cover horizontal prioritization of NaNs. Horizontal propagation of NaNs to the destination and the positioning of those NaNs in the destination is implementation dependent. NaNs on the input sources or computationally generated NaNs will have at least one NaN propagated to the destination.

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

If VDPPD is encoded with VEX.L= 1, an attempt to execute the instruction encoded with VEX.L= 1 will cause an #UD exception.

Operation

DP_primitive (SRC1, SRC2)

```
IF (imm8[4] = 1)
    THEN Temp1[63:0] ← DEST[63:0] * SRC[63:0];
    ELSE Temp1[63:0] ← +0.0; FI;
IF (imm8[5] = 1)
    THEN Temp1[127:64] ← DEST[127:64] * SRC[127:64];
    ELSE Temp1[127:64] ← +0.0; FI;

Temp2[63:0] ← Temp1[63:0] + Temp1[127:64];

IF (imm8[0] = 1)
    THEN DEST[63:0] ← Temp2[63:0];
    ELSE DEST[63:0] ← +0.0; FI;
IF (imm8[1] = 1)
    THEN DEST[127:64] ← Temp2[63:0];
    ELSE DEST[127:64] ← +0.0; FI;
```

DPPD (128-bit Legacy SSE version)

```
DEST[127:0]←DP_Primitive(SRC1[127:0], SRC2[127:0]);
DEST[VLMAX-1:128] (Unmodified)
```

VDPPD (VEX.128 encoded version)

```
DEST[127:0]←DP_Primitive(SRC1[127:0], SRC2[127:0]);
DEST[VLMAX-1:128] ← 0
```

Flags Affected

None

Intel C/C++ Compiler Intrinsic Equivalent

DPPD __m128d _mm_dp_pd (__m128d a, __m128d b, const int mask);

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal

Exceptions are determined separately for each add and multiply operation. Unmasked exceptions will leave the destination untouched.

Other Exceptions

See Exceptions Type 2; additionally

#UD	If VEX.L= 1.

DPPS — Dot Product of Packed Single Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 3A 40 /r ib DPPS *xmm1, xmm2/m128, imm8*	A	V/V	SSE4_1	Selectively multiply packed SP floating-point values from *xmm1* with packed SP floating-point values from *xmm2*, add and selectively store the packed SP floating-point values or zero values to *xmm1*.
VEX.NDS.128.66.0F3A.WIG 40 /r ib VDPPS xmm1,xmm2, xmm3/m128, imm8	B	V/V	AVX	Multiply packed SP floating point values from xmm1 with packed SP floating point values from xmm2/mem selectively add and store to xmm1.
VEX.NDS.256.66.0F3A.WIG 40 /r ib VDPPS ymm1, ymm2, ymm3/m256, imm8	B	V/V	AVX	Multiply packed single-precision floating-point values from ymm2 with packed SP floating point values from ymm3/mem, selectively add pairs of elements and store to ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	imm8	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Conditionally multiplies the packed single precision floating-point values in the destination operand (first operand) with the packed single-precision floats in the source (second operand) depending on a mask extracted from the high 4 bits of the immediate byte (third operand). If a condition mask bit in Imm8[7:4] is zero, the corresponding multiplication is replaced by a value of 0.0.

The four resulting single-precision values are summed into an intermediate result. The intermediate result is conditionally broadcasted to the destination using a broadcast mask specified by bits [3:0] of the immediate byte..

If a broadcast mask bit is "1", the intermediate result is copied to the corresponding dword element in the destination operand. If a broadcast mask bit is zero, the corresponding element in the destination is set to zero.

DPPS follows the NaN forwarding rules stated in the Software Developer's Manual, vol. 1, table 4.7. These rules do not cover horizontal prioritization of NaNs. Horizontal propagation of NaNs to the destination and the positioning of those NaNs in the destination is implementation dependent. NaNs on the input sources or computationally generated NaNs will have at least one NaN propagated to the destination.

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

```
DP_primitive (SRC1, SRC2)
IF (imm8[4] = 1)
    THEN Temp1[31:0] ← DEST[31:0] * SRC[31:0];
    ELSE Temp1[31:0] ← +0.0; FI;
IF (imm8[5] = 1)
    THEN Temp1[63:32] ← DEST[63:32] * SRC[63:32];
    ELSE Temp1[63:32] ← +0.0; FI;
IF (imm8[6] = 1)
    THEN Temp1[95:64] ← DEST[95:64] * SRC[95:64];
    ELSE Temp1[95:64] ← +0.0; FI;
IF (imm8[7] = 1)
    THEN Temp1[127:96] ← DEST[127:96] * SRC[127:96];
    ELSE Temp1[127:96] ← +0.0; FI;

Temp2[31:0] ← Temp1[31:0] + Temp1[63:32];
Temp3[31:0] ← Temp1[95:64] + Temp1[127:96];
Temp4[31:0] ← Temp2[31:0] + Temp3[31:0];

IF (imm8[0] = 1)
    THEN DEST[31:0] ← Temp4[31:0];
    ELSE DEST[31:0] ← +0.0; FI;
IF (imm8[1] = 1)
```

```
    THEN DEST[63:32] ← Temp4[31:0];
    ELSE DEST[63:32] ← +0.0; FI;
IF (imm8[2] = 1)
    THEN DEST[95:64] ← Temp4[31:0];
    ELSE DEST[95:64] ← +0.0; FI;
IF (imm8[3] = 1)
    THEN DEST[127:96] ← Temp4[31:0];
    ELSE DEST[127:96] ← +0.0; FI;
```

DPP (128-bit Legacy SSE version)

```
DEST[127:0]←DP_Primitive(SRC1[127:0], SRC2[127:0]);
DEST[VLMAX-1:128] (Unmodified)
```

VDPPS (VEX.128 encoded version)

```
DEST[127:0]←DP_Primitive(SRC1[127:0], SRC2[127:0]);
DEST[VLMAX-1:128] ← 0
```

VDPPS (VEX.256 encoded version)

```
DEST[127:0]←DP_Primitive(SRC1[127:0], SRC2[127:0]);
DEST[255:128]←DP_Primitive(SRC1[255:128], SRC2[255:128]);
```

Intel C/C++ Compiler Intrinsic Equivalent

```
(V)DPPS __m128 _mm_dp_ps ( __m128 a, __m128 b, const int mask);

VDPPS __m256 _mm256_dp_ps ( __m256 a, __m256 b, const int mask);
```

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal

Exceptions are determined separately for each add and multiply operation, in the order of their execution. Unmasked exceptions will leave the destination operands unchanged.

Other Exceptions

See Exceptions Type 2.

EMMS—Empty MMX Technology State

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 77	EMMS	A	Valid	Valid	Set the x87 FPU tag word to empty.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Sets the values of all the tags in the x87 FPU tag word to empty (all 1s). This operation marks the x87 FPU data registers (which are aliased to the MMX technology registers) as available for use by x87 FPU floating-point instructions. (See Figure 8-7 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for the format of the x87 FPU tag word.) All other MMX instructions (other than the EMMS instruction) set all the tags in x87 FPU tag word to valid (all 0s).

The EMMS instruction must be used to clear the MMX technology state at the end of all MMX technology procedures or subroutines and before calling other procedures or subroutines that may execute x87 floating-point instructions. If a floating-point instruction loads one of the registers in the x87 FPU data register stack before the x87 FPU tag word has been reset by the EMMS instruction, an x87 floating-point register stack overflow can occur that will result in an x87 floating-point exception or incorrect result.

EMMS operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
x87FPUTagWord ← FFFFH;
```

Intel C/C++ Compiler Intrinsic Equivalent

```
void _mm_empty()
```

Flags Affected

None.

Protected Mode Exceptions

#UD	If CR0.EM[bit 2] = 1.
#NM	If CR0.TS[bit 3] = 1.

#MF	If there is a pending FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

ENTER—Make Stack Frame for Procedure Parameters

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
C8 *iw* 00	ENTER *imm16*, 0	A	Valid	Valid	Create a stack frame for a procedure.
C8 *iw* 01	ENTER *imm16*,1	A	Valid	Valid	Create a nested stack frame for a procedure.
C8 *iw* ib	ENTER *imm16*, *imm8*	A	Valid	Valid	Create a nested stack frame for a procedure.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	iw	imm8	NA	NA

Description

Creates a stack frame for a procedure. The first operand (size operand) specifies the size of the stack frame (that is, the number of bytes of dynamic storage allocated on the stack for the procedure). The second operand (nesting level operand) gives the lexical nesting level (0 to 31) of the procedure. The nesting level determines the number of stack frame pointers that are copied into the "display area" of the new stack frame from the preceding frame. Both of these operands are immediate values.

The stack-size attribute determines whether the BP (16 bits), EBP (32 bits), or RBP (64 bits) register specifies the current frame pointer and whether SP (16 bits), ESP (32 bits), or RSP (64 bits) specifies the stack pointer. In 64-bit mode, stack-size attribute is always 64-bits.

The ENTER and companion LEAVE instructions are provided to support block structured languages. The ENTER instruction (when used) is typically the first instruction in a procedure and is used to set up a new stack frame for a procedure. The LEAVE instruction is then used at the end of the procedure (just before the RET instruction) to release the stack frame.

If the nesting level is 0, the processor pushes the frame pointer from the BP/EBP/RBP register onto the stack, copies the current stack pointer from the SP/ESP/RSP register into the BP/EBP/RBP register, and loads the SP/ESP/RSP register with the current stack-pointer value minus the value in the size operand. For nesting levels of 1 or greater, the processor pushes additional frame pointers on the stack before adjusting the stack pointer. These additional frame pointers provide the called procedure with access points to other nested frames on the stack. See "Procedure Calls for Block-Structured Languages" in Chapter 6 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for more information about the actions of the ENTER instruction.

The ENTER instruction causes a page fault whenever a write using the final value of the stack pointer (within the current stack segment) would do so.

In 64-bit mode, default operation size is 64 bits; 32-bit operation size cannot be encoded.

Operation

```
NestingLevel ← NestingLevel MOD 32
IF 64-Bit Mode (StackSize = 64)
    THEN
        Push(RBP);
        FrameTemp ← RSP;
    ELSE IF StackSize = 32
        THEN
            Push(EBP);
            FrameTemp ← ESP; FI;
    ELSE (* StackSize = 16 *)
            Push(BP);
            FrameTemp ← SP;
FI;
IF NestingLevel = 0
    THEN GOTO CONTINUE;
FI;

IF (NestingLevel > 1)
    THEN FOR i ← 1 to (NestingLevel - 1)
        DO
            IF 64-Bit Mode (StackSize = 64)
                THEN
                    RBP ← RBP - 8;
                    Push([RBP]); (* Quadword push *)
                ELSE IF OperandSize = 32
                    THEN
                        IF StackSize = 32
                            EBP ← EBP - 4;
                            Push([EBP]); (* Doubleword push *)
                        ELSE (* StackSize = 16 *)
                            BP ← BP - 4;
                            Push([BP]); (* Doubleword push *)
                        FI;
                    FI;
                ELSE (* OperandSize = 16 *)
                    IF StackSize = 32
                        THEN
```

```
                        EBP ← EBP - 2;
                        Push([EBP]); (* Word push *)
                    ELSE (* StackSize = 16 *)
                        BP ← BP - 2;
                        Push([BP]); (* Word push *)
                FI;
            FI;
    OD;
FI;

IF 64-Bit Mode (StackSize = 64)
    THEN
        Push(FrameTemp); (* Quadword push *)
    ELSE IF OperandSize = 32
        THEN
            Push(FrameTemp); FI; (* Doubleword push *)
    ELSE (* OperandSize = 16 *)
            Push(FrameTemp); (* Word push *)
FI;

CONTINUE:
IF 64-Bit Mode (StackSize = 64)
    THEN
            RBP ← FrameTemp;
            RSP ← RSP − Size;
    ELSE IF StackSize = 32
        THEN
            EBP ← FrameTemp;
            ESP ← ESP − Size; FI;
    ELSE (* StackSize = 16 *)
            BP ← FrameTemp;
            SP ← SP − Size;
FI;

END;
```

Flags Affected

None.

Protected Mode Exceptions

#SS(0)	If the new value of the SP or ESP register is outside the stack segment limit.

#PF(fault-code)	If a page fault occurs or if a write using the final value of the stack pointer (within the current stack segment) would cause a page fault.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#SS	If the new value of the SP or ESP register is outside the stack segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#SS(0)	If the new value of the SP or ESP register is outside the stack segment limit.
#PF(fault-code)	If a page fault occurs or if a write using the final value of the stack pointer (within the current stack segment) would cause a page fault.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If the stack address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs or if a write using the final value of the stack pointer (within the current stack segment) would cause a page fault.
#UD	If the LOCK prefix is used.

VEXTRACTF128 — Extract Packed Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
VEX.256.66.0F3A.W0 19 /r ib VEXTRACTF128 xmm1/m128, ymm2, imm8	A	V/V	AVX	Extract 128 bits of packed floating-point values from ymm2 and store results in xmm1/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

Extracts 128-bits of packed floating-point values from the source operand (second operand) at an 128-bit offset from imm8[0] into the destination operand (first operand). The destination may be either an XMM register or an 128-bit memory location.

VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

The high 7 bits of the immediate are ignored.

If VEXTRACTF128 is encoded with VEX.L= 0, an attempt to execute the instruction encoded with VEX.L= 0 will cause an #UD exception.

Operation

VEXTRACTF128 (memory destination form)

```
CASE (imm8[0]) OF
    0: DEST[127:0] ← SRC1[127:0]
    1: DEST[127:0] ← SRC1[255:128]
ESAC.
```

VEXTRACTF128 (register destination form)

```
CASE (imm8[0]) OF
    0: DEST[127:0] ← SRC1[127:0]
    1: DEST[127:0] ← SRC1[255:128]
ESAC.
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

VEXTRACTF128 __m128 _mm256_extractf128_ps (__m256 a, int offset);

VEXTRACTF128 __m128d _mm256_extractf128_pd (__m256d a, int offset);

VEXTRACTF128 __m128i_mm256_extractf128_si256(__m256i a, int offset);

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 6; additionally

#UD	If VEX.L= 0
	If VEX.W=1.

EXTRACTPS — Extract Packed Single Precision Floating-Point Value

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 3A 17 /r ib EXTRACTPS *reg/m32, xmm2, imm8*	A	V/V	SSE4_1	Extract a single-precision floating-point value from *xmm2* at the source offset specified by *imm8* and store the result to *reg or m32*. The upper 32 bits of r64 is zeroed if reg is r64.
VEX.128.66.0F3A.WIG 17 /r ib VEXTRACTPS r/m32, xmm1, imm8	A	V/V	AVX	Extract one single-precision floating-point value from xmm1 at the offset specified by imm8 and store the result in reg or m32. Zero extend the results in 64-bit register if applicable.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (w)	ModRM:reg (r)	imm8	NA

Description

Extracts a single-precision floating-point value from the source operand (second operand) at the 32-bit offset specified from imm8. Immediate bits higher than the most significant offset for the vector length are ignored.

The extracted single-precision floating-point value is stored in the low 32-bits of the destination operand

In 64-bit mode, destination register operand has default operand size of 64 bits. The upper 32-bits of the register are filled with zero. REX.W is ignored.

128-bit Legacy SSE version: When a REX.W prefix is used in 64-bit mode with a general purpose register (GPR) as a destination operand, the packed single quantity is zero extended to 64 bits.

VEX.128 encoded version: When VEX.128.66.0F3A.W1 17 form is used in 64-bit mode with a general purpose register (GPR) as a destination operand, the packed single quantity is zero extended to 64 bits. VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

The source register is an XMM register. Imm8[1:0] determine the starting DWORD offset from which to extract the 32-bit floating-point value.

If VEXTRACTPS is encoded with VEX.L= 1, an attempt to execute the instruction encoded with VEX.L= 1 will cause an #UD exception.

Operation

EXTRACTPS (128-bit Legacy SSE version)

```
SRC_OFFSET ← IMM8[1:0]
IF ( 64-Bit Mode and DEST is register)
    DEST[31:0] ← (SRC[127:0] » (SRC_OFFET*32)) AND 0FFFFFFFFh
    DEST[63:32] ← 0
ELSE
    DEST[31:0] ← (SRC[127:0] » (SRC_OFFET*32)) AND 0FFFFFFFFh
FI
```

VEXTRACTPS (VEX.128 encoded version)

```
SRC_OFFSET ← IMM8[1:0]
IF ( 64-Bit Mode and DEST is register)
    DEST[31:0] ← (SRC[127:0] » (SRC_OFFET*32)) AND 0FFFFFFFFh
    DEST[63:32] ← 0
ELSE
    DEST[31:0] ← (SRC[127:0] » (SRC_OFFET*32)) AND 0FFFFFFFFh
FI
```

Intel C/C++ Compiler Intrinsic Equivalent

EXTRACTPS _mm_extractmem_ps (float *dest, __m128 a, const int nidx);

EXTRACTPS __m128 _mm_extract_ps (__m128 a, const int nidx);

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 5; additionally

#UD	If VEX.L= 1.

F2XM1—Compute 2^x–1

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 F0	F2XM1	Valid	Valid	Replace ST(0) with $(2^{ST(0)} - 1)$.

Description

Computes the exponential value of 2 to the power of the source operand minus 1. The source operand is located in register ST(0) and the result is also stored in ST(0). The value of the source operand must lie in the range –1.0 to +1.0. If the source value is outside this range, the result is undefined.

The following table shows the results obtained when computing the exponential value of various classes of numbers, assuming that neither overflow nor underflow occurs.

Table 3-26. Results Obtained from F2XM1

ST(0) SRC	ST(0) DEST
– 1.0 to –0	– 0.5 to – 0
– 0	– 0
+ 0	+ 0
+ 0 to +1.0	+ 0 to 1.0

Values other than 2 can be exponentiated using the following formula:

$$x^y \leftarrow 2^{(y * \log_2 x)}$$

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

$ST(0) \leftarrow (2^{ST(0)} - 1);$

FPU Flags Affected

C1	Set to 0 if stack underflow occurred. Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Source operand is an SNaN value or unsupported format.
#D	Source is a denormal value.

#U	Result is too small for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FABS—Absolute Value

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 E1	FABS	Valid	Valid	Replace ST with its absolute value.

Description

Clears the sign bit of ST(0) to create the absolute value of the operand. The following table shows the results obtained when creating the absolute value of various classes of numbers.

Table 3-27. Results Obtained from FABS

ST(0) SRC	ST(0) DEST
– •	+ •
– F	+ F
– 0	+ 0
+ 0	+ 0
+ F	+ F
+ •	+ •
NaN	NaN

NOTES:

F Means finite floating-point value.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

ST(0) ← |ST(0)|;

FPU Flags Affected

C1	Set to 0 if stack underflow occurred; otherwise, set to 0.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FADD/FADDP/FIADD—Add

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D8 /0	FADD *m32fp*	Valid	Valid	Add *m32fp* to ST(0) and store result in ST(0).
DC /0	FADD *m64fp*	Valid	Valid	Add *m64fp* to ST(0) and store result in ST(0).
D8 C0+i	FADD ST(0), ST(i)	Valid	Valid	Add ST(0) to ST(i) and store result in ST(0).
DC C0+i	FADD ST(i), ST(0)	Valid	Valid	Add ST(i) to ST(0) and store result in ST(i).
DE C0+i	FADDP ST(i), ST(0)	Valid	Valid	Add ST(0) to ST(i), store result in ST(i), and pop the register stack.
DE C1	FADDP	Valid	Valid	Add ST(0) to ST(1), store result in ST(1), and pop the register stack.
DA /0	FIADD *m32int*	Valid	Valid	Add *m32int* to ST(0) and store result in ST(0).
DE /0	FIADD *m16int*	Valid	Valid	Add *m16int* to ST(0) and store result in ST(0).

Description

Adds the destination and source operands and stores the sum in the destination location. The destination operand is always an FPU register; the source operand can be a register or a memory location. Source operands in memory can be in single-precision or double-precision floating-point format or in word or doubleword integer format.

The no-operand version of the instruction adds the contents of the ST(0) register to the ST(1) register. The one-operand version adds the contents of a memory location (either a floating-point or an integer value) to the contents of the ST(0) register. The two-operand version, adds the contents of the ST(0) register to the ST(i) register or vice versa. The value in ST(0) can be doubled by coding:

```
FADD ST(0), ST(0);
```

The FADDP instructions perform the additional operation of popping the FPU register stack after storing the result. To pop the register stack, the processor marks the ST(0) register as empty and increments the stack pointer (TOP) by 1. (The no-operand version of the floating-point add instructions always results in the register stack being popped. In some assemblers, the mnemonic for this instruction is FADD rather than FADDP.)

The FIADD instructions convert an integer source operand to double extended-precision floating-point format before performing the addition.

The table on the following page shows the results obtained when adding various classes of numbers, assuming that neither overflow nor underflow occurs.

When the sum of two operands with opposite signs is 0, the result is +0, except for the round toward –∞ mode, in which case the result is –0. When the source operand is an integer 0, it is treated as a +0.

When both operand are infinities of the same sign, the result is ∞ of the expected sign. If both operands are infinities of opposite signs, an invalid-operation exception is generated. See Table 3-28.

Table 3-28. FADD/FADDP/FIADD Results

		DEST						
		- ∞	– F	– 0	+ 0	+ F	+ ∞	NaN
SRC	- ∞	- ∞	- ∞	- ∞	- ∞	- ∞	*	NaN
	– F or – I	- ∞	– F	SRC	SRC	± F or ± 0	+ ∞	NaN
	–0	- ∞	DEST	– 0	± 0	DEST	+ ∞	NaN
	+ 0	- ∞	DEST	± 0	+ 0	DEST	+ ∞	NaN
	+ F or + I	- ∞	± F or ± 0	SRC	SRC	+ F	+ ∞	NaN
	+ ∞	*	+ ∞	+ ∞	+ ∞	+ ∞	+ ∞	NaN
	NaN	NaN	NaN	NaN	NaN	NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

I Means integer.

* Indicates floating-point invalid-arithmetic-operand (#IA) exception.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF Instruction = FIADD
    THEN
        DEST ← DEST + ConvertToDoubleExtendedPrecisionFP(SRC);
    ELSE (* Source operand is floating-point value *)
        DEST ← DEST + SRC;
FI;

IF Instruction = FADDP
    THEN
        PopRegisterStack;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred. Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Operand is an SNaN value or unsupported format. Operands are infinities of unlike sign.
#D	Source operand is a denormal value.
#U	Result is too small for destination format.
#O	Result is too large for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit. If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FBLD—Load Binary Coded Decimal

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
DF /4	FBLD *m80 dec*	Valid	Valid	Convert BCD value to floating-point and push onto the FPU stack.

Description

Converts the BCD source operand into double extended-precision floating-point format and pushes the value onto the FPU stack. The source operand is loaded without rounding errors. The sign of the source operand is preserved, including that of –0.

The packed BCD digits are assumed to be in the range 0 through 9; the instruction does not check for invalid digits (AH through FH). Attempting to load an invalid encoding produces an undefined result.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
TOP ← TOP – 1;
ST(0) ← ConvertToDoubleExtendedPrecisionFP(SRC);
```

FPU Flags Affected

C1	Set to 1 if stack overflow occurred; otherwise, set to 0.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack overflow occurred.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FBSTP—Store BCD Integer and Pop

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
DF /6	FBSTP m80bcd	Valid	Valid	Store ST(0) in m80bcd and pop ST(0).

Description

Converts the value in the ST(0) register to an 18-digit packed BCD integer, stores the result in the destination operand, and pops the register stack. If the source value is a non-integral value, it is rounded to an integer value, according to rounding mode specified by the RC field of the FPU control word. To pop the register stack, the processor marks the ST(0) register as empty and increments the stack pointer (TOP) by 1.

The destination operand specifies the address where the first byte destination value is to be stored. The BCD value (including its sign bit) requires 10 bytes of space in memory.

The following table shows the results obtained when storing various classes of numbers in packed BCD format.

Table 3-29. FBSTP Results

ST(0)	DEST
– • or Value Too Large for DEST Format	*
$F \leq -1$	$-D$
$-1 < F < -0$	**
-0	-0
$+0$	$+0$
$+0 < F < +1$	**
$F \geq +1$	$+D$
+ • or Value Too Large for DEST Format	*
NaN	*

NOTES:

F Means finite floating-point value.

D Means packed-BCD number.

* Indicates floating-point invalid-operation (#IA) exception.

** ± 0 or ± 1, depending on the rounding mode.

If the converted value is too large for the destination format, or if the source operand is an ∞, SNaN, QNAN, or is in an unsupported format, an invalid-arithmetic-operand condition is signaled. If the invalid-operation exception is not masked, an invalid-arithmetic-operand exception (#IA) is generated and no value is stored in the destination operand. If the invalid-operation exception is masked, the packed BCD indefinite value is stored in memory.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
DEST ← BCD(ST(0));
PopRegisterStack;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred.
	Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Converted value that exceeds 18 BCD digits in length.
	Source operand is an SNaN, QNaN, ±∞, or in an unsupported format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#GP(0)	If a segment register is being loaded with a segment selector that points to a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FCHS—Change Sign

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 E0	FCHS	Valid	Valid	Complements sign of ST(0).

Description

Complements the sign bit of ST(0). This operation changes a positive value into a negative value of equal magnitude or vice versa. The following table shows the results obtained when changing the sign of various classes of numbers.

Table 3-30. FCHS Results

ST(0) SRC	ST(0) DEST
– •	+ •
– F	+ F
– 0	+ 0
+ 0	– 0
+ F	– F
+ •	– •
NaN	NaN

NOTES:

* F means finite floating-point value.

This instruction’s operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
SignBit(ST(0)) ← NOT (SignBit(ST(0)));
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred; otherwise, set to 0.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FCLEX/FNCLEX—Clear Exceptions

Opcode*	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
9B DB E2	FCLEX	Valid	Valid	Clear floating-point exception flags after checking for pending unmasked floating-point exceptions.
DB E2	FNCLEX*	Valid	Valid	Clear floating-point exception flags without checking for pending unmasked floating-point exceptions.

NOTES:

* See IA-32 Architecture Compatibility section below.

Description

Clears the floating-point exception flags (PE, UE, OE, ZE, DE, and IE), the exception summary status flag (ES), the stack fault flag (SF), and the busy flag (B) in the FPU status word. The FCLEX instruction checks for and handles any pending unmasked floating-point exceptions before clearing the exception flags; the FNCLEX instruction does not.

The assembler issues two instructions for the FCLEX instruction (an FWAIT instruction followed by an FNCLEX instruction), and the processor executes each of these instructions separately. If an exception is generated for either of these instructions, the save EIP points to the instruction that caused the exception.

IA-32 Architecture Compatibility

When operating a Pentium or Intel486 processor in MS-DOS* compatibility mode, it is possible (under unusual circumstances) for an FNCLEX instruction to be interrupted prior to being executed to handle a pending FPU exception. See the section titled "No-Wait FPU Instructions Can Get FPU Interrupt in Window" in Appendix D of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for a description of these circumstances. An FNCLEX instruction cannot be interrupted in this way on a Pentium 4, Intel Xeon, or P6 family processor.

This instruction affects only the x87 FPU floating-point exception flags. It does not affect the SIMD floating-point exception flags in the MXCRS register.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
FPUStatusWord[0:7] ← 0;
FPUStatusWord[15] ← 0;
```

FPU Flags Affected

The PE, UE, OE, ZE, DE, IE, ES, SF, and B flags in the FPU status word are cleared. The C0, C1, C2, and C3 flags are undefined.

Floating-Point Exceptions

None.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FCMOV*cc*—Floating-Point Conditional Move

Opcode*	Instruction	64-Bit Mode	Compat/ Leg Mode*	Description
DA C0+i	FCMOVB ST(0), ST(i)	Valid	Valid	Move if below (CF=1).
DA C8+i	FCMOVE ST(0), ST(i)	Valid	Valid	Move if equal (ZF=1).
DA D0+i	FCMOVBE ST(0), ST(i)	Valid	Valid	Move if below or equal (CF=1 or ZF=1).
DA D8+i	FCMOVU ST(0), ST(i)	Valid	Valid	Move if unordered (PF=1).
DB C0+i	FCMOVNB ST(0), ST(i)	Valid	Valid	Move if not below (CF=0).
DB C8+i	FCMOVNE ST(0), ST(i)	Valid	Valid	Move if not equal (ZF=0).
DB D0+i	FCMOVNBE ST(0), ST(i)	Valid	Valid	Move if not below or equal (CF=0 and ZF=0).
DB D8+i	FCMOVNU ST(0), ST(i)	Valid	Valid	Move if not unordered (PF=0).

NOTES:

* See IA-32 Architecture Compatibility section below.

Description

Tests the status flags in the EFLAGS register and moves the source operand (second operand) to the destination operand (first operand) if the given test condition is true. The condition for each mnemonic os given in the Description column above and in Chapter 8 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*. The source operand is always in the ST(i) register and the destination operand is always ST(0).

The FCMOV*cc* instructions are useful for optimizing small IF constructions. They also help eliminate branching overhead for IF operations and the possibility of branch mispredictions by the processor.

A processor may not support the FCMOV*cc* instructions. Software can check if the FCMOV*cc* instructions are supported by checking the processor's feature information with the CPUID instruction (see "COMISS—Compare Scalar Ordered Single-Precision Floating-Point Values and Set EFLAGS" in this chapter). If both the CMOV and FPU feature bits are set, the FCMOV*cc* instructions are supported.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

IA-32 Architecture Compatibility

The FCMOVcc instructions were introduced to the IA-32 Architecture in the P6 family processors and are not available in earlier IA-32 processors.

Operation

```
IF condition TRUE
    THEN ST(0) ← ST(i);
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.

Integer Flags Affected

None.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FCOM/FCOMP/FCOMPP—Compare Floating Point Values

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D8 /2	FCOM *m32fp*	Valid	Valid	Compare ST(0) with *m32fp*.
DC /2	FCOM *m64fp*	Valid	Valid	Compare ST(0) with *m64fp*.
D8 D0+i	FCOM ST(i)	Valid	Valid	Compare ST(0) with ST(i).
D8 D1	FCOM	Valid	Valid	Compare ST(0) with ST(1).
D8 /3	FCOMP *m32fp*	Valid	Valid	Compare ST(0) with *m32fp* and pop register stack.
DC /3	FCOMP *m64fp*	Valid	Valid	Compare ST(0) with m64fp and pop register stack.
D8 D8+i	FCOMP ST(i)	Valid	Valid	Compare ST(0) with ST(i) and pop register stack.
D8 D9	FCOMP	Valid	Valid	Compare ST(0) with ST(1) and pop register stack.
DE D9	FCOMPP	Valid	Valid	Compare ST(0) with ST(1) and pop register stack twice.

Description

Compares the contents of register ST(0) and source value and sets condition code flags C0, C2, and C3 in the FPU status word according to the results (see the table below). The source operand can be a data register or a memory location. If no source operand is given, the value in ST(0) is compared with the value in ST(1). The sign of zero is ignored, so that −0.0 is equal to +0.0.

Table 3-31. FCOM/FCOMP/FCOMPP Results

Condition	C3	C2	C0
ST(0) > SRC	0	0	0
ST(0) < SRC	0	0	1
ST(0) = SRC	1	0	0
Unordered*	1	1	1

NOTES:

* Flags not set if unmasked invalid-arithmetic-operand (#IA) exception is generated.

This instruction checks the class of the numbers being compared (see "FXAM—Examine ModR/M" in this chapter). If either operand is a NaN or is in an unsupported format, an invalid-arithmetic-operand exception (#IA) is raised and, if the exception is masked, the condition flags are set to "unordered." If the invalid-arithmetic-operand exception is unmasked, the condition code flags are not set.

The FCOMP instruction pops the register stack following the comparison operation and the FCOMPP instruction pops the register stack twice following the comparison operation. To pop the register stack, the processor marks the ST(0) register as empty and increments the stack pointer (TOP) by 1.

The FCOM instructions perform the same operation as the FUCOM instructions. The only difference is how they handle QNaN operands. The FCOM instructions raise an invalid-arithmetic-operand exception (#IA) when either or both of the operands is a NaN value or is in an unsupported format. The FUCOM instructions perform the same operation as the FCOM instructions, except that they do not generate an invalid-arithmetic-operand exception for QNaNs.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
CASE (relation of operands) OF
    ST > SRC:          C3, C2, C0 ← 000;
    ST < SRC:          C3, C2, C0 ← 001;
    ST = SRC:          C3, C2, C0 ← 100;
ESAC;

IF ST(0) or SRC = NaN or unsupported format
    THEN
        #IA
        IF FPUControlWord.IM = 1
            THEN
                C3, C2, C0 ← 111;
        FI;
FI;

IF Instruction = FCOMP
    THEN
        PopRegisterStack;
FI;

IF Instruction = FCOMPP
    THEN
        PopRegisterStack;
        PopRegisterStack;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred; otherwise, set to 0.
C0, C2, C3	See table on previous page.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	One or both operands are NaN values or have unsupported formats.
	Register is marked empty.
#D	One or both operands are denormal values.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FCOMI/FCOMIP/FUCOMI/FUCOMIP—Compare Floating Point Values and Set EFLAGS

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
DB F0+i	FCOMI ST, ST(i)	Valid	Valid	Compare ST(0) with ST(i) and set status flags accordingly.
DF F0+i	FCOMIP ST, ST(i)	Valid	Valid	Compare ST(0) with ST(i), set status flags accordingly, and pop register stack.
DB E8+i	FUCOMI ST, ST(i)	Valid	Valid	Compare ST(0) with ST(i), check for ordered values, and set status flags accordingly.
DF E8+i	FUCOMIP ST, ST(i)	Valid	Valid	Compare ST(0) with ST(i), check for ordered values, set status flags accordingly, and pop register stack.

Description

Performs an unordered comparison of the contents of registers ST(0) and ST(i) and sets the status flags ZF, PF, and CF in the EFLAGS register according to the results (see the table below). The sign of zero is ignored for comparisons, so that –0.0 is equal to +0.0.

Table 3-32. FCOMI/FCOMIP/ FUCOMI/FUCOMIP Results

Comparison Results*	ZF	PF	CF
ST0 > ST(i)	0	0	0
ST0 < ST(i)	0	0	1
ST0 = ST(i)	1	0	0
Unordered**	1	1	1

NOTES:

* See the IA-32 Architecture Compatibility section below.

** Flags not set if unmasked invalid-arithmetic-operand (#IA) exception is generated.

An unordered comparison checks the class of the numbers being compared (see "FXAM—Examine ModR/M" in this chapter). The FUCOMI/FUCOMIP instructions perform the same operations as the FCOMI/FCOMIP instructions. The only difference is that the FUCOMI/FUCOMIP instructions raise the invalid-arithmetic-operand exception (#IA) only when either or both operands are an SNaN or are in an unsupported format; QNaNs cause the condition code flags to be set to unordered, but do not cause an exception to be generated. The FCOMI/FCOMIP instructions raise an invalid-operation exception when either or both of the operands are a NaN value of any kind or are in an unsupported format.

If the operation results in an invalid-arithmetic-operand exception being raised, the status flags in the EFLAGS register are set only if the exception is masked.

The FCOMI/FCOMIP and FUCOMI/FUCOMIP instructions clear the OF flag in the EFLAGS register (regardless of whether an invalid-operation exception is detected).

The FCOMIP and FUCOMIP instructions also pop the register stack following the comparison operation. To pop the register stack, the processor marks the ST(0) register as empty and increments the stack pointer (TOP) by 1.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

IA-32 Architecture Compatibility

The FCOMI/FCOMIP/FUCOMI/FUCOMIP instructions were introduced to the IA-32 Architecture in the P6 family processors and are not available in earlier IA-32 processors.

Operation

```
CASE (relation of operands) OF
    ST(0) > ST(i):      ZF, PF, CF ← 000;
    ST(0) < ST(i):      ZF, PF, CF ← 001;
    ST(0) = ST(i):      ZF, PF, CF ← 100;
ESAC;

IF Instruction is FCOMI or FCOMIP
    THEN
        IF ST(0) or ST(i) = NaN or unsupported format
            THEN
                #IA
                IF FPUControlWord.IM = 1
                    THEN
                        ZF, PF, CF ← 111;
                FI;
        FI;
FI;

IF Instruction is FUCOMI or FUCOMIP
    THEN
        IF ST(0) or ST(i) = QNaN, but not SNaN or unsupported format
            THEN
                ZF, PF, CF ← 111;
            ELSE (* ST(0) or ST(i) is SNaN or unsupported format *)
                 #IA;
                IF FPUControlWord.IM = 1
                    THEN
```

```
                ZF, PF, CF ← 111;
            FI;
        FI;
FI;

IF Instruction is FCOMIP or FUCOMIP
    THEN
        PopRegisterStack;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred; otherwise, set to 0.
C0, C2, C3	Not affected.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	(FCOMI or FCOMIP instruction) One or both operands are NaN values or have unsupported formats.
	(FUCOMI or FUCOMIP instruction) One or both operands are SNaN values (but not QNaNs) or have undefined formats. Detection of a QNaN value does not raise an invalid-operand exception.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FCOS—Cosine

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 FF	FCOS	Valid	Valid	Replace ST(0) with its cosine.

Description

Computes the cosine of the source operand in register ST(0) and stores the result in ST(0). The source operand must be given in radians and must be within the range -2^{63} to $+2^{63}$. The following table shows the results obtained when taking the cosine of various classes of numbers.

Table 3-33. FCOS Results

ST(0) SRC	ST(0) DEST
– •	*
– F	–1 to +1
– 0	+ 1
+ 0	+ 1
+ F	– 1 to + 1
+ •	*
NaN	NaN

NOTES:

F Means finite floating-point value.

* Indicates floating-point invalid-arithmetic-operand (#IA) exception.

If the source operand is outside the acceptable range, the C2 flag in the FPU status word is set, and the value in register ST(0) remains unchanged. The instruction does not raise an exception when the source operand is out of range. It is up to the program to check the C2 flag for out-of-range conditions. Source values outside the range -2^{63} to $+2^{63}$ can be reduced to the range of the instruction by subtracting an appropriate integer multiple of 2π or by using the FPREM instruction with a divisor of 2π. See the section titled "Pi" in Chapter 8 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for a discussion of the proper value to use for π in performing such reductions.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

IF $|ST(0)| < 2^{63}$
THEN
    C2 ← 0;

```
    ST(0) ← cosine(ST(0));
ELSE (* Source operand is out-of-range *)
    C2 ← 1;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred.
	Set if result was rounded up; cleared otherwise.
	Undefined if C2 is 1.
C2	Set to 1 if outside range (-2^{63} < source operand < $+2^{63}$); otherwise, set to 0.
C0, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Source operand is an SNaN value, ∞, or unsupported format.
#D	Source is a denormal value.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FDECSTP—Decrement Stack-Top Pointer

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 F6	FDECSTP	Valid	Valid	Decrement TOP field in FPU status word.

Description

Subtracts one from the TOP field of the FPU status word (decrements the top-of-stack pointer). If the TOP field contains a 0, it is set to 7. The effect of this instruction is to rotate the stack by one position. The contents of the FPU data registers and tag register are not affected.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF TOP = 0
    THEN TOP ← 7;
    ELSE TOP ← TOP – 1;
FI;
```

FPU Flags Affected

The C1 flag is set to 0. The C0, C2, and C3 flags are undefined.

Floating-Point Exceptions

None.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FDIV/FDIVP/FIDIV—Divide

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D8 /6	FDIV *m32fp*	Valid	Valid	Divide ST(0) by *m32fp* and store result in ST(0).
DC /6	FDIV *m64fp*	Valid	Valid	Divide ST(0) by *m64fp* and store result in ST(0).
D8 F0+i	FDIV ST(0), ST(i)	Valid	Valid	Divide ST(0) by ST(i) and store result in ST(0).
DC F8+i	FDIV ST(i), ST(0)	Valid	Valid	Divide ST(i) by ST(0) and store result in ST(i).
DE F8+i	FDIVP ST(i), ST(0)	Valid	Valid	Divide ST(i) by ST(0), store result in ST(i), and pop the register stack.
DE F9	FDIVP	Valid	Valid	Divide ST(1) by ST(0), store result in ST(1), and pop the register stack.
DA /6	FIDIV *m32int*	Valid	Valid	Divide ST(0) by *m32int* and store result in ST(0).
DE /6	FIDIV *m16int*	Valid	Valid	Divide ST(0) by *m64int* and store result in ST(0).

Description

Divides the destination operand by the source operand and stores the result in the destination location. The destination operand (dividend) is always in an FPU register; the source operand (divisor) can be a register or a memory location. Source operands in memory can be in single-precision or double-precision floating-point format, word or doubleword integer format.

The no-operand version of the instruction divides the contents of the ST(1) register by the contents of the ST(0) register. The one-operand version divides the contents of the ST(0) register by the contents of a memory location (either a floating-point or an integer value). The two-operand version, divides the contents of the ST(0) register by the contents of the ST(i) register or vice versa.

The FDIVP instructions perform the additional operation of popping the FPU register stack after storing the result. To pop the register stack, the processor marks the ST(0) register as empty and increments the stack pointer (TOP) by 1. The no-operand version of the floating-point divide instructions always results in the register stack being popped. In some assemblers, the mnemonic for this instruction is FDIV rather than FDIVP.

The FIDIV instructions convert an integer source operand to double extended-precision floating-point format before performing the division. When the source operand is an integer 0, it is treated as a +0.

If an unmasked divide-by-zero exception (#Z) is generated, no result is stored; if the exception is masked, an ∞ of the appropriate sign is stored in the destination operand.

The following table shows the results obtained when dividing various classes of numbers, assuming that neither overflow nor underflow occurs.

Table 3-34. FDIV/FDIVP/FIDIV Results

	DEST							
		– •	– F	– 0	+ 0	+F	+ •	NaN
	– •	*	+ 0	+ 0	– 0	– 0	*	NaN
	–F	+ •	+ F	+ 0	– 0	– F	– •	NaN
	–I	+ •	+ F	+ 0	– 0	– F	– •	NaN
SRC	–0	+ •	**	*	*	**	– •	NaN
	+0	– •	**	*	*	**	+ •	NaN
	+I	– •	–F	–0	+0	+ F	+ •	NaN
	+F	– •	–F	–0	+0	+ F	+ •	NaN
	+ •	*	–0	–0	+0	+ 0	*	NaN
	NaN	NaN	NaN	NaN	NaN	NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

I Means integer.

* Indicates floating-point invalid-arithmetic-operand (#IA) exception.

** Indicates floating-point zero-divide (#Z) exception.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF SRC = 0
    THEN
        #Z;
    ELSE
        IF Instruction is FIDIV
            THEN
                DEST ← DEST / ConvertToDoubleExtendedPrecisionFP(SRC);
            ELSE (* Source operand is floating-point value *)
                DEST ← DEST / SRC;
        FI;
FI;
```

```
IF Instruction = FDIVP
    THEN
        PopRegisterStack;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred. Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Operand is an SNaN value or unsupported format. ±∞ / ±∞; ±0 / ±0
#D	Source is a denormal value.
#Z	DEST / ±0, where DEST is not equal to ±0.
#U	Result is too small for destination format.
#O	Result is too large for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit. If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FDIVR/FDIVRP/FIDIVR—Reverse Divide

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D8 /7	FDIVR *m32fp*	Valid	Valid	Divide *m32fp* by ST(0) and store result in ST(0).
DC /7	FDIVR *m64fp*	Valid	Valid	Divide *m64fp* by ST(0) and store result in ST(0).
D8 F8+i	FDIVR ST(0), ST(i)	Valid	Valid	Divide ST(i) by ST(0) and store result in ST(0).
DC F0+i	FDIVR ST(i), ST(0)	Valid	Valid	Divide ST(0) by ST(i) and store result in ST(i).
DE F0+i	FDIVRP ST(i), ST(0)	Valid	Valid	Divide ST(0) by ST(i), store result in ST(i), and pop the register stack.
DE F1	FDIVRP	Valid	Valid	Divide ST(0) by ST(1), store result in ST(1), and pop the register stack.
DA /7	FIDIVR *m32int*	Valid	Valid	Divide *m32int* by ST(0) and store result in ST(0).
DE /7	FIDIVR *m16int*	Valid	Valid	Divide *m16int* by ST(0) and store result in ST(0).

Description

Divides the source operand by the destination operand and stores the result in the destination location. The destination operand (divisor) is always in an FPU register; the source operand (dividend) can be a register or a memory location. Source operands in memory can be in single-precision or double-precision floating-point format, word or doubleword integer format.

These instructions perform the reverse operations of the FDIV, FDIVP, and FIDIV instructions. They are provided to support more efficient coding.

The no-operand version of the instruction divides the contents of the ST(0) register by the contents of the ST(1) register. The one-operand version divides the contents of a memory location (either a floating-point or an integer value) by the contents of the ST(0) register. The two-operand version, divides the contents of the ST(i) register by the contents of the ST(0) register or vice versa.

The FDIVRP instructions perform the additional operation of popping the FPU register stack after storing the result. To pop the register stack, the processor marks the ST(0) register as empty and increments the stack pointer (TOP) by 1. The no-operand version of the floating-point divide instructions always results in the register stack being popped. In some assemblers, the mnemonic for this instruction is FDIVR rather than FDIVRP.

The FIDIVR instructions convert an integer source operand to double extended-precision floating-point format before performing the division.

If an unmasked divide-by-zero exception (#Z) is generated, no result is stored; if the exception is masked, an ∞ of the appropriate sign is stored in the destination operand.

The following table shows the results obtained when dividing various classes of numbers, assuming that neither overflow nor underflow occurs.

Table 3-35. FDIVR/FDIVRP/FIDIVR Results

		DEST						
		– •	– F	– 0	+ 0	+ F	+ •	NaN
SRC	– •	*	+ •	+ •	– •	– •	*	NaN
	– F	+ 0	+ F	**	**	– F	– 0	NaN
	– I	+ 0	+ F	**	**	– F	– 0	NaN
	– 0	+ 0	+ 0	*	*	– 0	– 0	NaN
	+ 0	– 0	– 0	*	*	+ 0	+ 0	NaN
	+ I	– 0	– F	**	**	+ F	+ 0	NaN
	+ F	– 0	– F	**	**	+ F	+ 0	NaN
	+ •	*	– •	– •	+ •	+ •	*	NaN
	NaN	NaN	NaN	NaN	NaN	NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

I Means integer.

* Indicates floating-point invalid-arithmetic-operand (#IA) exception.

** Indicates floating-point zero-divide (#Z) exception.

When the source operand is an integer 0, it is treated as a +0. This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF DEST = 0
    THEN
        #Z;
    ELSE
        IF Instruction = FIDIVR
            THEN
                DEST ← ConvertToDoubleExtendedPrecisionFP(SRC) / DEST;
            ELSE (* Source operand is floating-point value *)
```

```
                DEST ← SRC / DEST;
        FI;
FI;

IF Instruction = FDIVRP
    THEN
        PopRegisterStack;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred.
	Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Operand is an SNaN value or unsupported format.
	$\pm\infty$ / $\pm\infty$; ± 0 / ± 0
#D	Source is a denormal value.
#Z	SRC / ± 0, where SRC is not equal to ± 0.
#U	Result is too small for destination format.
#O	Result is too large for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.

#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FFREE—Free Floating-Point Register

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
DD C0+i	FFREE ST(i)	Valid	Valid	Sets tag for ST(i) to empty.

Description

Sets the tag in the FPU tag register associated with register ST(i) to empty (11B). The contents of ST(i) and the FPU stack-top pointer (TOP) are not affected.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
TAG(i) ← 11B;
```

FPU Flags Affected

C0, C1, C2, C3 undefined.

Floating-Point Exceptions

None.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FICOM/FICOMP—Compare Integer

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
DE /2	FICOM *m16int*	Valid	Valid	Compare ST(0) with *m16int.*
DA /2	FICOM *m32int*	Valid	Valid	Compare ST(0) with *m32int.*
DE /3	FICOMP *m16int*	Valid	Valid	Compare ST(0) with *m16int* and pop stack register.
DA /3	FICOMP *m32int*	Valid	Valid	Compare ST(0) with *m32int* and pop stack register.

Description

Compares the value in ST(0) with an integer source operand and sets the condition code flags C0, C2, and C3 in the FPU status word according to the results (see table below). The integer value is converted to double extended-precision floating-point format before the comparison is made.

Table 3-36. FICOM/FICOMP Results

Condition	C3	C2	C0
ST(0) > SRC	0	0	0
ST(0) < SRC	0	0	1
ST(0) = SRC	1	0	0
Unordered	1	1	1

These instructions perform an “unordered comparison.” An unordered comparison also checks the class of the numbers being compared (see “FXAM—Examine ModR/M” in this chapter). If either operand is a NaN or is in an undefined format, the condition flags are set to “unordered.”

The sign of zero is ignored, so that –0.0 ← +0.0.

The FICOMP instructions pop the register stack following the comparison. To pop the register stack, the processor marks the ST(0) register empty and increments the stack pointer (TOP) by 1.

This instruction’s operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
CASE (relation of operands) OF
    ST(0) > SRC:        C3, C2, C0 ← 000;
    ST(0) < SRC:        C3, C2, C0 ← 001;
    ST(0) = SRC:        C3, C2, C0 ← 100;
    Unordered:          C3, C2, C0 ← 111;
```

```
ESAC;

IF Instruction = FICOMP
    THEN
        PopRegisterStack;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred; otherwise, set to 0.
C0, C2, C3	See table on previous page.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	One or both operands are NaN values or have unsupported formats.
#D	One or both operands are denormal values.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.

#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FILD—Load Integer

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
DF /0	FILD *m16int*	Valid	Valid	Push *m16int* onto the FPU register stack.
DB /0	FILD *m32int*	Valid	Valid	Push *m32int* onto the FPU register stack.
DF /5	FILD *m64int*	Valid	Valid	Push *m64int* onto the FPU register stack.

Description

Converts the signed-integer source operand into double extended-precision floating-point format and pushes the value onto the FPU register stack. The source operand can be a word, doubleword, or quadword integer. It is loaded without rounding errors. The sign of the source operand is preserved.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
TOP ← TOP − 1;
ST(0) ← ConvertToDoubleExtendedPrecisionFP(SRC);
```

FPU Flags Affected

C1	Set to 1 if stack overflow occurred; set to 0 otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack overflow occurred.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FINCSTP—Increment Stack-Top Pointer

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 F7	FINCSTP	Valid	Valid	Increment the TOP field in the FPU status register.

Description

Adds one to the TOP field of the FPU status word (increments the top-of-stack pointer). If the TOP field contains a 7, it is set to 0. The effect of this instruction is to rotate the stack by one position. The contents of the FPU data registers and tag register are not affected. This operation is not equivalent to popping the stack, because the tag for the previous top-of-stack register is not marked empty.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF TOP = 7
    THEN TOP ← 0;
    ELSE TOP ← TOP + 1;
FI;
```

FPU Flags Affected

The C1 flag is set to 0. The C0, C2, and C3 flags are undefined.

Floating-Point Exceptions

None.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FINIT/FNINIT—Initialize Floating-Point Unit

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
9B DB E3	FINIT	Valid	Valid	Initialize FPU after checking for pending unmasked floating-point exceptions.
DB E3	FNINIT*	Valid	Valid	Initialize FPU without checking for pending unmasked floating-point exceptions.

NOTES:

* See IA-32 Architecture Compatibility section below.

Description

Sets the FPU control, status, tag, instruction pointer, and data pointer registers to their default states. The FPU control word is set to 037FH (round to nearest, all exceptions masked, 64-bit precision). The status word is cleared (no exception flags set, TOP is set to 0). The data registers in the register stack are left unchanged, but they are all tagged as empty (11B). Both the instruction and data pointers are cleared.

The FINIT instruction checks for and handles any pending unmasked floating-point exceptions before performing the initialization; the FNINIT instruction does not.

The assembler issues two instructions for the FINIT instruction (an FWAIT instruction followed by an FNINIT instruction), and the processor executes each of these instructions in separately. If an exception is generated for either of these instructions, the save EIP points to the instruction that caused the exception.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

IA-32 Architecture Compatibility

When operating a Pentium or Intel486 processor in MS-DOS compatibility mode, it is possible (under unusual circumstances) for an FNINIT instruction to be interrupted prior to being executed to handle a pending FPU exception. See the section titled "No-Wait FPU Instructions Can Get FPU Interrupt in Window" in Appendix D of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for a description of these circumstances. An FNINIT instruction cannot be interrupted in this way on a Pentium 4, Intel Xeon, or P6 family processor.

In the Intel387 math coprocessor, the FINIT/FNINIT instruction does not clear the instruction and data pointers.

This instruction affects only the x87 FPU. It does not affect the XMM and MXCSR registers.

Operation

```
FPUControlWord ← 037FH;
FPUStatusWord ← 0;
FPUTagWord ← FFFFH;
FPUDataPointer ← 0;
FPUInstructionPointer ← 0;
FPULastInstructionOpcode ← 0;
```

FPU Flags Affected

C0, C1, C2, C3 set to 0.

Floating-Point Exceptions

None.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FIST/FISTP—Store Integer

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
DF /2	FIST *m16int*	Valid	Valid	Store ST(0) in *m16int.*
DB /2	FIST *m32int*	Valid	Valid	Store ST(0) in *m32int.*
DF /3	FISTP *m16int*	Valid	Valid	Store ST(0) in *m16int* and pop register stack.
DB /3	FISTP *m32int*	Valid	Valid	Store ST(0) in *m32int* and pop register stack.
DF /7	FISTP *m64int*	Valid	Valid	Store ST(0) in *m64int* and pop register stack.

Description

The FIST instruction converts the value in the ST(0) register to a signed integer and stores the result in the destination operand. Values can be stored in word or doubleword integer format. The destination operand specifies the address where the first byte of the destination value is to be stored.

The FISTP instruction performs the same operation as the FIST instruction and then pops the register stack. To pop the register stack, the processor marks the ST(0) register as empty and increments the stack pointer (TOP) by 1. The FISTP instruction also stores values in quadword integer format.

The following table shows the results obtained when storing various classes of numbers in integer format.

Table 3-37. FIST/FISTP Results

ST(0)	DEST
–• or Value Too Large for DEST Format	*
F ≤ –1	– I
–1 < F < –0	**
– 0	0
+ 0	0
+ 0 < F < + 1	**
F ≥ + 1	+ I
+• or Value Too Large for DEST Format	*

Table 3-37. FIST/FISTP Results (Contd.)

ST(0)	DEST
NaN	*

NOTES:

F Means finite floating-point value.

I Means integer.

* Indicates floating-point invalid-operation (#IA) exception.

** 0 or ±1, depending on the rounding mode.

If the source value is a non-integral value, it is rounded to an integer value, according to the rounding mode specified by the RC field of the FPU control word.

If the converted value is too large for the destination format, or if the source operand is an ∞, SNaN, QNAN, or is in an unsupported format, an invalid-arithmetic-operand condition is signaled. If the invalid-operation exception is not masked, an invalid-arithmetic-operand exception (#IA) is generated and no value is stored in the destination operand. If the invalid-operation exception is masked, the integer indefinite value is stored in memory.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
DEST ← Integer(ST(0));

IF Instruction = FISTP
    THEN
        PopRegisterStack;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred.
	Indicates rounding direction of if the inexact exception (#P) is generated: 0 ← not roundup; 1 ← roundup.
	Set to 0 otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Converted value is too large for the destination format.
	Source operand is an SNaN, QNaN, ±∞, or unsupported format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#GP(0)	If the destination is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FISTTP—Store Integer with Truncation

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
DF /1	FISTTP *m16int*	Valid	Valid	Store ST(0) in *m16int with truncation.*
DB /1	FISTTP *m32int*	Valid	Valid	Store ST(0) in *m32int with truncation.*
DD /1	FISTTP *m64int*	Valid	Valid	Store ST(0) in *m64int* with truncation.

Description

FISTTP converts the value in ST into a signed integer using truncation (chop) as rounding mode, transfers the result to the destination, and pop ST. FISTTP accepts word, short integer, and long integer destinations.

The following table shows the results obtained when storing various classes of numbers in integer format.

Table 3-38. FISTTP Results

ST(0)	DEST
–• or Value Too Large for DEST Format	*
F ≤ – 1	– I
– 1 < F < + 1	0
F Š + 1	+ I
+• or Value Too Large for DEST Format	*
NaN	*

NOTES:

F Means finite floating-point value.

I Means integer.

∗ Indicates floating-point invalid-operation (#IA) exception.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
DEST ← ST;
pop ST;
```

Flags Affected

C1 is cleared; C0, C2, C3 undefined.

Numeric Exceptions

Invalid, Stack Invalid (stack underflow), Precision.

Protected Mode Exceptions

#GP(0)	If the destination is in a nonwritable segment.
	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
#SS(0)	For an illegal address in the SS segment.
#PF(fault-code)	For a page fault.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#NM	If CR0.EM[bit 2] = 1.
	If CR0.TS[bit 3] = 1.
#UD	If CPUID.01H:ECX.SSE3[bit 0] = 0.
	If the LOCK prefix is used.

Real Address Mode Exceptions

GP(0)	If any part of the operand would lie outside of the effective address space from 0 to 0FFFFH.
#NM	If CR0.EM[bit 2] = 1.
	If CR0.TS[bit 3] = 1.
#UD	If CPUID.01H:ECX.SSE3[bit 0] = 0.
	If the LOCK prefix is used.

Virtual 8086 Mode Exceptions

GP(0)	If any part of the operand would lie outside of the effective address space from 0 to 0FFFFH.
#NM	If CR0.EM[bit 2] = 1.
	If CR0.TS[bit 3] = 1.
#UD	If CPUID.01H:ECX.SSE3[bit 0] = 0.
	If the LOCK prefix is used.
#PF(fault-code)	For a page fault.
#AC(0)	For unaligned memory reference if the current privilege is 3.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
	If the LOCK prefix is used.

FLD—Load Floating Point Value

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 /0	FLD *m32fp*	Valid	Valid	Push *m32fp* onto the FPU register stack.
DD /0	FLD *m64fp*	Valid	Valid	Push *m64fp* onto the FPU register stack.
DB /5	FLD *m80fp*	Valid	Valid	Push *m80fp* onto the FPU register stack.
D9 C0+i	FLD ST(i)	Valid	Valid	Push ST(i) onto the FPU register stack.

Description

Pushes the source operand onto the FPU register stack. The source operand can be in single-precision, double-precision, or double extended-precision floating-point format. If the source operand is in single-precision or double-precision floating-point format, it is automatically converted to the double extended-precision floating-point format before being pushed on the stack.

The FLD instruction can also push the value in a selected FPU register [ST(i)] onto the stack. Here, pushing register ST(0) duplicates the stack top.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF SRC is ST(i)
    THEN
        temp ← ST(i);
FI;

TOP ← TOP − 1;

IF SRC is memory-operand
    THEN
        ST(0) ← ConvertToDoubleExtendedPrecisionFP(SRC);
    ELSE (* SRC is ST(i) *)
        ST(0) ← temp;
FI;
```

FPU Flags Affected

C1	Set to 1 if stack overflow occurred; otherwise, set to 0.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow or overflow occurred.

#IA	Source operand is an SNaN. Does not occur if the source operand is in double extended-precision floating-point format (FLD m80fp or FLD ST(i)).
#D	Source operand is a denormal value. Does not occur if the source operand is in double extended-precision floating-point format.

Protected Mode Exceptions

#GP(0)	If destination is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FLD1/FLDL2T/FLDL2E/FLDPI/FLDLG2/FLDLN2/FLDZ—Load Constant

Opcode*	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 E8	FLD1	Valid	Valid	Push +1.0 onto the FPU register stack.
D9 E9	FLDL2T	Valid	Valid	Push $\log_2 10$ onto the FPU register stack.
D9 EA	FLDL2E	Valid	Valid	Push $\log_2 e$ onto the FPU register stack.
D9 EB	FLDPI	Valid	Valid	Push π onto the FPU register stack.
D9 EC	FLDLG2	Valid	Valid	Push $\log_{10} 2$ onto the FPU register stack.
D9 ED	FLDLN2	Valid	Valid	Push $\log_e 2$ onto the FPU register stack.
D9 EE	FLDZ	Valid	Valid	Push +0.0 onto the FPU register stack.

NOTES:

* See IA-32 Architecture Compatibility section below.

Description

Push one of seven commonly used constants (in double extended-precision floating-point format) onto the FPU register stack. The constants that can be loaded with these instructions include +1.0, +0.0, $\log_2 10$, $\log_2 e$, π, $\log_{10} 2$, and $\log_e 2$. For each constant, an internal 66-bit constant is rounded (as specified by the RC field in the FPU control word) to double extended-precision floating-point format. The inexact-result exception (#P) is not generated as a result of the rounding, nor is the C1 flag set in the x87 FPU status word if the value is rounded up.

See the section titled “Pi” in Chapter 8 of the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*, for a description of the π constant.

This instruction’s operation is the same in non-64-bit modes and 64-bit mode.

IA-32 Architecture Compatibility

When the RC field is set to round-to-nearest, the FPU produces the same constants that is produced by the Intel 8087 and Intel 287 math coprocessors.

Operation

```
TOP ← TOP − 1;
ST(0) ← CONSTANT;
```

FPU Flags Affected

C1	Set to 1 if stack overflow occurred; otherwise, set to 0.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack overflow occurred.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FLDCW—Load x87 FPU Control Word

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 /5	FLDCW m2byte	Valid	Valid	Load FPU control word from *m2byte*.

Description

Loads the 16-bit source operand into the FPU control word. The source operand is a memory location. This instruction is typically used to establish or change the FPU's mode of operation.

If one or more exception flags are set in the FPU status word prior to loading a new FPU control word and the new control word unmasks one or more of those exceptions, a floating-point exception will be generated upon execution of the next floating-point instruction (except for the no-wait floating-point instructions, see the section titled "Software Exception Handling" in Chapter 8 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*). To avoid raising exceptions when changing FPU operating modes, clear any pending exceptions (using the FCLEX or FNCLEX instruction) before loading the new control word.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
FPUControlWord ← SRC;
```

FPU Flags Affected

C0, C1, C2, C3 undefined.

Floating-Point Exceptions

None; however, this operation might unmask a pending exception in the FPU status word. That exception is then generated upon execution of the next "waiting" floating-point instruction.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.

#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FLDENV—Load x87 FPU Environment

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 /4	FLDENV *m14/28byte*	Valid	Valid	Load FPU environment from *m14byte* or *m28byte.*

Description

Loads the complete x87 FPU operating environment from memory into the FPU registers. The source operand specifies the first byte of the operating-environment data in memory. This data is typically written to the specified memory location by a FSTENV or FNSTENV instruction.

The FPU operating environment consists of the FPU control word, status word, tag word, instruction pointer, data pointer, and last opcode. Figures 8-9 through 8-12 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, show the layout in memory of the loaded environment, depending on the operating mode of the processor (protected or real) and the current operand-size attribute (16-bit or 32-bit). In virtual-8086 mode, the real mode layouts are used.

The FLDENV instruction should be executed in the same operating mode as the corresponding FSTENV/FNSTENV instruction.

If one or more unmasked exception flags are set in the new FPU status word, a floating-point exception will be generated upon execution of the next floating-point instruction (except for the no-wait floating-point instructions, see the section titled "Software Exception Handling" in Chapter 8 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*). To avoid generating exceptions when loading a new environment, clear all the exception flags in the FPU status word that is being loaded.

If a page or limit fault occurs during the execution of this instruction, the state of the x87 FPU registers as seen by the fault handler may be different than the state being loaded from memory. In such situations, the fault handler should ignore the status of the x87 FPU registers, handle the fault, and return. The FLDENV instruction will then complete the loading of the x87 FPU registers with no resulting context inconsistency.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
FPUControlWord ← SRC[FPUControlWord];
FPUStatusWord ← SRC[FPUStatusWord];
FPUTagWord ← SRC[FPUTagWord];
FPUDataPointer ← SRC[FPUDataPointer];
FPUInstructionPointer ← SRC[FPUInstructionPointer];
FPULastInstructionOpcode ← SRC[FPULastInstructionOpcode];
```

FPU Flags Affected

The C0, C1, C2, C3 flags are loaded.

Floating-Point Exceptions

None; however, if an unmasked exception is loaded in the status word, it is generated upon execution of the next “waiting” floating-point instruction.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FMUL/FMULP/FIMUL—Multiply

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D8 /1	FMUL *m32fp*	Valid	Valid	Multiply ST(0) by *m32fp* and store result in ST(0).
DC /1	FMUL *m64fp*	Valid	Valid	Multiply ST(0) by *m64fp* and store result in ST(0).
D8 C8+i	FMUL ST(0), ST(i)	Valid	Valid	Multiply ST(0) by ST(i) and store result in ST(0).
DC C8+i	FMUL ST(i), ST(0)	Valid	Valid	Multiply ST(i) by ST(0) and store result in ST(i).
DE C8+i	FMULP ST(i), ST(0)	Valid	Valid	Multiply ST(i) by ST(0), store result in ST(i), and pop the register stack.
DE C9	FMULP	Valid	Valid	Multiply ST(1) by ST(0), store result in ST(1), and pop the register stack.
DA /1	FIMUL *m32int*	Valid	Valid	Multiply ST(0) by *m32int* and store result in ST(0).
DE /1	FIMUL *m16int*	Valid	Valid	Multiply ST(0) by *m16int* and store result in ST(0).

Description

Multiplies the destination and source operands and stores the product in the destination location. The destination operand is always an FPU data register; the source operand can be an FPU data register or a memory location. Source operands in memory can be in single-precision or double-precision floating-point format or in word or doubleword integer format.

The no-operand version of the instruction multiplies the contents of the ST(1) register by the contents of the ST(0) register and stores the product in the ST(1) register. The one-operand version multiplies the contents of the ST(0) register by the contents of a memory location (either a floating point or an integer value) and stores the product in the ST(0) register. The two-operand version, multiplies the contents of the ST(0) register by the contents of the ST(i) register, or vice versa, with the result being stored in the register specified with the first operand (the destination operand).

The FMULP instructions perform the additional operation of popping the FPU register stack after storing the product. To pop the register stack, the processor marks the ST(0) register as empty and increments the stack pointer (TOP) by 1. The no-operand version of the floating-point multiply instructions always results in the register stack being popped. In some assemblers, the mnemonic for this instruction is FMUL rather than FMULP.

The FIMUL instructions convert an integer source operand to double extended-precision floating-point format before performing the multiplication.

The sign of the result is always the exclusive-OR of the source signs, even if one or more of the values being multiplied is 0 or ∞. When the source operand is an integer 0, it is treated as a +0.

The following table shows the results obtained when multiplying various classes of numbers, assuming that neither overflow nor underflow occurs.

Table 3-39. FMUL/FMULP/FIMUL Results

		DEST						
		–•	–F	–0	+0	+F	+•	NaN
	–•	+•	+•	*	*	–•	–•	NaN
	–F	+•	+F	+0	–0	– F	–•	NaN
	–I	+•	+F	+0	–0	–F	–•	NaN
SRC	–0	*	+0	+0	–0	–0	*	NaN
	+0	*	–0	–0	+0	+0	*	NaN
	+I	–•	–F	–0	+0	+F	+•	NaN
	+F	–•	–F	–0	+0	+F	+•	NaN
	+•	–•	–•	*	*	+•	+•	NaN
	NaN	NaN	NaN	NaN	NaN	NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

I Means Integer.

* Indicates invalid-arithmetic-operand (#IA) exception.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF Instruction = FIMUL
    THEN
        DEST ← DEST ∗ ConvertToDoubleExtendedPrecisionFP(SRC);
    ELSE (* Source operand is floating-point value *)
        DEST ← DEST ∗ SRC;
FI;

IF Instruction = FMULP
    THEN
        PopRegisterStack;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred. Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Operand is an SNaN value or unsupported format. One operand is ±0 and the other is ±∞.
#D	Source operand is a denormal value.
#U	Result is too small for destination format.
#O	Result is too large for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit. If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FNOP—No Operation

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 D0	FNOP	Valid	Valid	No operation is performed.

Description

Performs no FPU operation. This instruction takes up space in the instruction stream but does not affect the FPU or machine context, except the EIP register.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

FPU Flags Affected

C0, C1, C2, C3 undefined.

Floating-Point Exceptions

None.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FPATAN—Partial Arctangent

Opcode*	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 F3	FPATAN	Valid	Valid	Replace ST(1) with arctan(ST(1)/ST(0)) and pop the register stack.

NOTES:

* See IA-32 Architecture Compatibility section below.

Description

Computes the arctangent of the source operand in register ST(1) divided by the source operand in register ST(0), stores the result in ST(1), and pops the FPU register stack. The result in register ST(0) has the same sign as the source operand ST(1) and a magnitude less than $+\pi$.

The FPATAN instruction returns the angle between the X axis and the line from the origin to the point (X,Y), where Y (the ordinate) is ST(1) and X (the abscissa) is ST(0). The angle depends on the sign of X and Y independently, not just on the sign of the ratio Y/X. This is because a point (–X,Y) is in the second quadrant, resulting in an angle between $\pi/2$ and π, while a point (X,–Y) is in the fourth quadrant, resulting in an angle between 0 and $-\pi/2$. A point (–X,–Y) is in the third quadrant, giving an angle between $-\pi/2$ and $-\pi$.

The following table shows the results obtained when computing the arctangent of various classes of numbers, assuming that underflow does not occur.

Table 3-40. FPATAN Results

		ST(0)						
		-•	– F	– 0	+ 0	+ F	+ •	NaN
ST(1)	-•	– 3π/4*	– π/2	– π/2	– π/2	– π/2	– π/4*	NaN
	–F	-p	–π to –π/2	–π/2	–π/2	–π/2 to – 0	- 0	NaN
	–0	-p	-p	-p*	– 0*	– 0	– 0	NaN
	+0	+p	+ p	+ π*	+ 0*	+ 0	+ 0	NaN
	+F	+p	+π to +π/2	+ π/2	+π/2	+π/2 to +0	+ 0	NaN
	+•	+3π/4*	+π/2	+π/2	+π/2	+ π/2	+ π/4*	NaN
	NaN	NaN	NaN	NaN	NaN	NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

* Table 8-10 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, specifies that the ratios 0/0 and •/• generate the floating-point invalid arithmetic-operation exception and, if this exception is masked, the floating-point QNaN indefinite value is returned. With the FPATAN instruction, the 0/0 or •/• value is actually not calculated using division. Instead, the arctangent of the two variables is derived from a standard mathematical formulation that is generalized to allow complex numbers as arguments. In this complex variable formulation, arctangent(0,0) etc. has well defined values. These values are needed to develop a library to compute transcendental functions with complex arguments, based on the FPU functions that only allow floating-point values as arguments.

There is no restriction on the range of source operands that FPATAN can accept.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

IA-32 Architecture Compatibility

The source operands for this instruction are restricted for the 80287 math coprocessor to the following range:

$0 \le |ST(1)| < |ST(0)| < +\infty$

Operation

```
ST(1) ← arctan(ST(1) / ST(0));
PopRegisterStack;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred. Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Source operand is an SNaN value or unsupported format.
#D	Source operand is a denormal value.
#U	Result is too small for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FPREM—Partial Remainder

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 F8	FPREM	Valid	Valid	Replace ST(0) with the remainder obtained from dividing ST(0) by ST(1).

Description

Computes the remainder obtained from dividing the value in the ST(0) register (the dividend) by the value in the ST(1) register (the divisor or **modulus**), and stores the result in ST(0). The remainder represents the following value:

Remainder ← ST(0) – (Q ∗ ST(1))

Here, Q is an integer value that is obtained by truncating the floating-point number quotient of [ST(0) / ST(1)] toward zero. The sign of the remainder is the same as the sign of the dividend. The magnitude of the remainder is less than that of the modulus, unless a partial remainder was computed (as described below).

This instruction produces an exact result; the inexact-result exception does not occur and the rounding control has no effect. The following table shows the results obtained when computing the remainder of various classes of numbers, assuming that underflow does not occur.

Table 3-41. FPREM Results

		ST(1)						
		-•	-F	-0	+0	+F	+•	NaN
ST(0)	-•	*	*	*	*	*	*	NaN
	-F	ST(0)	-F or -0	**	**	-F or -0	ST(0)	NaN
	-0	-0	-0	*	*	-0	-0	NaN
	+0	+0	+0	*	*	+0	+0	NaN
	+F	ST(0)	+F or +0	**	**	+F or +0	ST(0)	NaN
	+•	*	*	*	*	*	*	NaN
	NaN	NaN	NaN	NaN	NaN	NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

* Indicates floating-point invalid-arithmetic-operand (#IA) exception.

** Indicates floating-point zero-divide (#Z) exception.

When the result is 0, its sign is the same as that of the dividend. When the modulus is ∞, the result is equal to the value in ST(0).

The FPREM instruction does not compute the remainder specified in IEEE Std 754. The IEEE specified remainder can be computed with the FPREM1 instruction. The FPREM instruction is provided for compatibility with the Intel 8087 and Intel287 math coprocessors.

The FPREM instruction gets its name “partial remainder” because of the way it computes the remainder. This instruction arrives at a remainder through iterative subtraction. It can, however, reduce the exponent of ST(0) by no more than 63 in one execution of the instruction. If the instruction succeeds in producing a remainder that is less than the modulus, the operation is complete and the C2 flag in the FPU status word is cleared. Otherwise, C2 is set, and the result in ST(0) is called the **partial remainder**. The exponent of the partial remainder will be less than the exponent of the original dividend by at least 32. Software can re-execute the instruction (using the partial remainder in ST(0) as the dividend) until C2 is cleared. (Note that while executing such a remainder-computation loop, a higher-priority interrupting routine that needs the FPU can force a context switch in-between the instructions in the loop.)

An important use of the FPREM instruction is to reduce the arguments of periodic functions. When reduction is complete, the instruction stores the three least-significant bits of the quotient in the C3, C1, and C0 flags of the FPU status word. This information is important in argument reduction for the tangent function (using a modulus of $\pi/4$), because it locates the original angle in the correct one of eight sectors of the unit circle.

This instruction’s operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
D ← exponent(ST(0)) – exponent(ST(1));

IF D < 64
    THEN
        Q ← Integer(TruncateTowardZero(ST(0) / ST(1)));
        ST(0) ← ST(0) – (ST(1) ∗ Q);
        C2 ← 0;
        C0, C3, C1 ← LeastSignificantBits(Q); (* Q2, Q1, Q0 *)
    ELSE
        C2 ← 1;
        N ← An implementation-dependent number between 32 and 63;
        QQ ← Integer(TruncateTowardZero((ST(0)  / ST(1)) / 2^(D – N)));
        ST(0) ← ST(0) – (ST(1) ∗ QQ ∗ 2^(D – N));
FI;
```

FPU Flags Affected

C0	Set to bit 2 (Q2) of the quotient.
C1	Set to 0 if stack underflow occurred; otherwise, set to least significant bit of quotient (Q0).
C2	Set to 0 if reduction complete; set to 1 if incomplete.
C3	Set to bit 1 (Q1) of the quotient.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Source operand is an SNaN value, modulus is 0, dividend is ∞, or unsupported format.
#D	Source operand is a denormal value.
#U	Result is too small for destination format.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FPREM1—Partial Remainder

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 F5	FPREM1	Valid	Valid	Replace ST(0) with the IEEE remainder obtained from dividing ST(0) by ST(1).

Description

Computes the IEEE remainder obtained from dividing the value in the ST(0) register (the dividend) by the value in the ST(1) register (the divisor or **modulus**), and stores the result in ST(0). The remainder represents the following value:

Remainder ← ST(0) – (Q ∗ ST(1))

Here, Q is an integer value that is obtained by rounding the floating-point number quotient of [ST(0) / ST(1)] toward the nearest integer value. The magnitude of the remainder is less than or equal to half the magnitude of the modulus, unless a partial remainder was computed (as described below).

This instruction produces an exact result; the precision (inexact) exception does not occur and the rounding control has no effect. The following table shows the results obtained when computing the remainder of various classes of numbers, assuming that underflow does not occur.

Table 3-42. FPREM1 Results

		ST(1)						
		- •	– F	– 0	+ 0	+ F	+•	NaN
	- •	*	*	*	*	*	*	NaN
ST(0)	– F	ST(0)	±F or –0	**	**	± F or – 0	ST(0)	NaN
	– 0	– 0	– 0	*	*	– 0	-0	NaN
	+ 0	+ 0	+ 0	*	*	+ 0	+0	NaN
	+ F	ST(0)	± F or + 0	**	**	± F or + 0	ST(0)	NaN
	+•	*	*	*	*	*	*	NaN
	NaN	NaN	NaN	NaN	NaN	NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

* Indicates floating-point invalid-arithmetic-operand (#IA) exception.

** Indicates floating-point zero-divide (#Z) exception.

When the result is 0, its sign is the same as that of the dividend. When the modulus is ∞, the result is equal to the value in ST(0).

The FPREM1 instruction computes the remainder specified in IEEE Standard 754. This instruction operates differently from the FPREM instruction in the way that it rounds the quotient of ST(0) divided by ST(1) to an integer (see the “Operation” section below).

Like the FPREM instruction, FPREM1 computes the remainder through iterative subtraction, but can reduce the exponent of ST(0) by no more than 63 in one execution of the instruction. If the instruction succeeds in producing a remainder that is less than one half the modulus, the operation is complete and the C2 flag in the FPU status word is cleared. Otherwise, C2 is set, and the result in ST(0) is called the **partial remainder**. The exponent of the partial remainder will be less than the exponent of the original dividend by at least 32. Software can re-execute the instruction (using the partial remainder in ST(0) as the dividend) until C2 is cleared. (Note that while executing such a remainder-computation loop, a higher-priority interrupting routine that needs the FPU can force a context switch in-between the instructions in the loop.)

An important use of the FPREM1 instruction is to reduce the arguments of periodic functions. When reduction is complete, the instruction stores the three least-significant bits of the quotient in the C3, C1, and C0 flags of the FPU status word. This information is important in argument reduction for the tangent function (using a modulus of $\pi/4$), because it locates the original angle in the correct one of eight sectors of the unit circle.

This instruction’s operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
D ← exponent(ST(0)) – exponent(ST(1));

IF D < 64
    THEN
        Q ← Integer(RoundTowardNearestInteger(ST(0) / ST(1)));
        ST(0) ← ST(0) – (ST(1) ∗ Q);
        C2 ← 0;
        C0, C3, C1 ← LeastSignificantBits(Q); (* Q2, Q1, Q0 *)
    ELSE
        C2 ← 1;
        N ← An implementation-dependent number between 32 and 63;
        QQ ← Integer(TruncateTowardZero((ST(0)  / ST(1)) / 2^(D − N)));
        ST(0) ← ST(0) – (ST(1) ∗ QQ ∗ 2^(D − N));
FI;
```

FPU Flags Affected

C0	Set to bit 2 (Q2) of the quotient.
C1	Set to 0 if stack underflow occurred; otherwise, set to least significant bit of quotient (Q0).

C2	Set to 0 if reduction complete; set to 1 if incomplete.
C3	Set to bit 1 (Q1) of the quotient.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Source operand is an SNaN value, modulus (divisor) is 0, dividend is ∞, or unsupported format.
#D	Source operand is a denormal value.
#U	Result is too small for destination format.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FPTAN—Partial Tangent

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 F2	FPTAN	Valid	Valid	Replace ST(0) with its tangent and push 1 onto the FPU stack.

Description

Computes the tangent of the source operand in register ST(0), stores the result in ST(0), and pushes a 1.0 onto the FPU register stack. The source operand must be given in radians and must be less than $\pm 2^{63}$. The following table shows the unmasked results obtained when computing the partial tangent of various classes of numbers, assuming that underflow does not occur.

Table 3-43. FPTAN Results

ST(0) SRC	ST(0) DEST
-•	*
- F	– F to + F
– 0	- 0
+ 0	+ 0
+ F	– F to + F
+•	*
NaN	NaN

NOTES:

F Means finite floating-point value.

* Indicates floating-point invalid-arithmetic-operand (#IA) exception.

If the source operand is outside the acceptable range, the C2 flag in the FPU status word is set, and the value in register ST(0) remains unchanged. The instruction does not raise an exception when the source operand is out of range. It is up to the program to check the C2 flag for out-of-range conditions. Source values outside the range -2^{63} to $+2^{63}$ can be reduced to the range of the instruction by subtracting an appropriate integer multiple of 2π or by using the FPREM instruction with a divisor of 2π. See the section titled “Pi” in Chapter 8 of the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*, for a discussion of the proper value to use for π in performing such reductions.

The value 1.0 is pushed onto the register stack after the tangent has been computed to maintain compatibility with the Intel 8087 and Intel287 math coprocessors. This operation also simplifies the calculation of other trigonometric functions. For instance, the cotangent (which is the reciprocal of the tangent) can be computed by executing a FDIVR instruction after the FPTAN instruction.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF ST(0) < 2^63
    THEN
        C2 ← 0;
        ST(0) ← tan(ST(0));
        TOP ← TOP − 1;
        ST(0) ← 1.0;
    ELSE (* Source operand is out-of-range *)
        C2 ← 1;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred; set to 1 if stack overflow occurred. Set if result was rounded up; cleared otherwise.
C2	Set to 1 if outside range (-2^{63} < source operand < $+2^{63}$); otherwise, set to 0.
C0, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow or overflow occurred.
#IA	Source operand is an SNaN value, ∞, or unsupported format.
#D	Source operand is a denormal value.
#U	Result is too small for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FRNDINT—Round to Integer

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 FC	FRNDINT	Valid	Valid	Round ST(0) to an integer.

Description

Rounds the source value in the ST(0) register to the nearest integral value, depending on the current rounding mode (setting of the RC field of the FPU control word), and stores the result in ST(0).

If the source value is ∞, the value is not changed. If the source value is not an integral value, the floating-point inexact-result exception (#P) is generated.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
ST(0) ← RoundToIntegralValue(ST(0));
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred.
	Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Source operand is an SNaN value or unsupported format.
#D	Source operand is a denormal value.
#P	Source operand is not an integral value.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FRSTOR—Restore x87 FPU State

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
DD /4	FRSTOR *m94/108byte*	Valid	Valid	Load FPU state from *m94byte* or *m108byte*.

Description

Loads the FPU state (operating environment and register stack) from the memory area specified with the source operand. This state data is typically written to the specified memory location by a previous FSAVE/FNSAVE instruction.

The FPU operating environment consists of the FPU control word, status word, tag word, instruction pointer, data pointer, and last opcode. Figures 8-9 through 8-12 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, show the layout in memory of the stored environment, depending on the operating mode of the processor (protected or real) and the current operand-size attribute (16-bit or 32-bit). In virtual-8086 mode, the real mode layouts are used. The contents of the FPU register stack are stored in the 80 bytes immediately following the operating environment image.

The FRSTOR instruction should be executed in the same operating mode as the corresponding FSAVE/FNSAVE instruction.

If one or more unmasked exception bits are set in the new FPU status word, a floating-point exception will be generated. To avoid raising exceptions when loading a new operating environment, clear all the exception flags in the FPU status word that is being loaded.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
FPUControlWord ← SRC[FPUControlWord];
FPUStatusWord ← SRC[FPUStatusWord];
FPUTagWord ← SRC[FPUTagWord];
FPUDataPointer ← SRC[FPUDataPointer];
FPUInstructionPointer ← SRC[FPUInstructionPointer];
FPULastInstructionOpcode ← SRC[FPULastInstructionOpcode];

ST(0) ← SRC[ST(0)];
ST(1) ← SRC[ST(1)];
ST(2) ← SRC[ST(2)];
ST(3) ← SRC[ST(3)];
ST(4) ← SRC[ST(4)];
ST(5) ← SRC[ST(5)];
ST(6) ← SRC[ST(6)];
```

ST(7) ← SRC[ST(7)];

FPU Flags Affected

The C0, C1, C2, C3 flags are loaded.

Floating-Point Exceptions

None; however, this operation might unmask an existing exception that has been detected but not generated, because it was masked. Here, the exception is generated at the completion of the instruction.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit. If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.

#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FSAVE/FNSAVE—Store x87 FPU State

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
9B DD /6	FSAVE *m94/108byte*	Valid	Valid	Store FPU state to *m94byte* or *m108byte* after checking for pending unmasked floating-point exceptions. Then re-initialize the FPU.
DD /6	FNSAVE* *m94/108byte*	Valid	Valid	Store FPU environment to *m94byte* or *m108byte* without checking for pending unmasked floating-point exceptions. Then re-initialize the FPU.

NOTES:

* See IA-32 Architecture Compatibility section below.

Description

Stores the current FPU state (operating environment and register stack) at the specified destination in memory, and then re-initializes the FPU. The FSAVE instruction checks for and handles pending unmasked floating-point exceptions before storing the FPU state; the FNSAVE instruction does not.

The FPU operating environment consists of the FPU control word, status word, tag word, instruction pointer, data pointer, and last opcode. Figures 8-9 through 8-12 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, show the layout in memory of the stored environment, depending on the operating mode of the processor (protected or real) and the current operand-size attribute (16-bit or 32-bit). In virtual-8086 mode, the real mode layouts are used. The contents of the FPU register stack are stored in the 80 bytes immediately follow the operating environment image.

The saved image reflects the state of the FPU after all floating-point instructions preceding the FSAVE/FNSAVE instruction in the instruction stream have been executed.

After the FPU state has been saved, the FPU is reset to the same default values it is set to with the FINIT/FNINIT instructions (see "FINIT/FNINIT—Initialize Floating-Point Unit" in this chapter).

The FSAVE/FNSAVE instructions are typically used when the operating system needs to perform a context switch, an exception handler needs to use the FPU, or an application program needs to pass a "clean" FPU to a procedure.

The assembler issues two instructions for the FSAVE instruction (an FWAIT instruction followed by an FNSAVE instruction), and the processor executes each of these instructions separately. If an exception is generated for either of these instructions, the save EIP points to the instruction that caused the exception.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

IA-32 Architecture Compatibility

For Intel math coprocessors and FPUs prior to the Intel Pentium processor, an FWAIT instruction should be executed before attempting to read from the memory image stored with a prior FSAVE/FNSAVE instruction. This FWAIT instruction helps ensure that the storage operation has been completed.

When operating a Pentium or Intel486 processor in MS-DOS compatibility mode, it is possible (under unusual circumstances) for an FNSAVE instruction to be interrupted prior to being executed to handle a pending FPU exception. See the section titled "No-Wait FPU Instructions Can Get FPU Interrupt in Window" in Appendix D of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for a description of these circumstances. An FNSAVE instruction cannot be interrupted in this way on a Pentium 4, Intel Xeon, or P6 family processor.

Operation

```
(* Save FPU State and Registers *)

DEST[FPUControlWord] ← FPUControlWord;
DEST[FPUStatusWord] ← FPUStatusWord;
DEST[FPUTagWord] ← FPUTagWord;
DEST[FPUDataPointer] ← FPUDataPointer;
DEST[FPUInstructionPointer] ← FPUInstructionPointer;
DEST[FPULastInstructionOpcode] ← FPULastInstructionOpcode;

DEST[ST(0)] ← ST(0);
DEST[ST(1)] ← ST(1);
DEST[ST(2)] ← ST(2);
DEST[ST(3)] ← ST(3);
DEST[ST(4)]← ST(4);
DEST[ST(5)] ← ST(5);
DEST[ST(6)] ← ST(6);
DEST[ST(7)] ← ST(7);

(* Initialize FPU *)

FPUControlWord ← 037FH;
FPUStatusWord ← 0;
FPUTagWord ← FFFFH;
FPUDataPointer ← 0;
FPUInstructionPointer ← 0;
FPULastInstructionOpcode ← 0;
```

FPU Flags Affected

The C0, C1, C2, and C3 flags are saved and then cleared.

Floating-Point Exceptions

None.

Protected Mode Exceptions

#GP(0)	If destination is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

FSCALE—Scale

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 FD	FSCALE	Valid	Valid	Scale ST(0) by ST(1).

Description

Truncates the value in the source operand (toward 0) to an integral value and adds that value to the exponent of the destination operand. The destination and source operands are floating-point values located in registers ST(0) and ST(1), respectively. This instruction provides rapid multiplication or division by integral powers of 2. The following table shows the results obtained when scaling various classes of numbers, assuming that neither overflow nor underflow occurs.

Table 3-44. FSCALE Results

	ST(1)							
		-•	– F	– 0	+ 0	+ F	+•	NaN
ST(0)	- •	NaN	-•	-•	-•	-•	-•	NaN
	– F	– 0	– F	– F	– F	– F	-•	NaN
	– 0	-0	-0	-0	-0	-0	NaN	NaN
	+ 0	+ 0	+ 0	+ 0	+ 0	+ 0	NaN	NaN
	+ F	+ 0	+ F	+ F	+ F	+ F	+•	NaN
	+•	NaN	+•	+•	+•	+•	+•	NaN
	NaN	NaN	NaN	NaN	NaN	NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

In most cases, only the exponent is changed and the mantissa (significand) remains unchanged. However, when the value being scaled in ST(0) is a denormal value, the mantissa is also changed and the result may turn out to be a normalized number. Similarly, if overflow or underflow results from a scale operation, the resulting mantissa will differ from the source's mantissa.

The FSCALE instruction can also be used to reverse the action of the FXTRACT instruction, as shown in the following example:

```
FXTRACT;
FSCALE;
FSTP ST(1);
```

In this example, the FXTRACT instruction extracts the significand and exponent from the value in ST(0) and stores them in ST(0) and ST(1) respectively. The FSCALE then scales the significand in ST(0) by the exponent in ST(1), recreating the original value before the FXTRACT operation was performed. The FSTP ST(1) instruction overwrites the exponent (extracted by the FXTRACT instruction) with the recreated value, which returns the stack to its original state with only one register [ST(0)] occupied.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

$ST(0) \leftarrow ST(0) * 2^{RoundTowardZero(ST(1))}$;

FPU Flags Affected

C1	Set to 0 if stack underflow occurred.
	Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Source operand is an SNaN value or unsupported format.
#D	Source operand is a denormal value.
#U	Result is too small for destination format.
#O	Result is too large for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FSIN—Sine

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 FE	FSIN	Valid	Valid	Replace ST(0) with its sine.

Description

Computes the sine of the source operand in register ST(0) and stores the result in ST(0). The source operand must be given in radians and must be within the range -2^{63} to $+2^{63}$. The following table shows the results obtained when taking the sine of various classes of numbers, assuming that underflow does not occur.

Table 3-45. FSIN Results

SRC (ST(0))	DEST (ST(0))
- •	*
− F	− 1 to + 1
− 0	−0
+ 0	+ 0
+ F	- 1 to +1
+ •	*
NaN	NaN

NOTES:

F Means finite floating-point value.

* Indicates floating-point invalid-arithmetic-operand (#IA) exception.

If the source operand is outside the acceptable range, the C2 flag in the FPU status word is set, and the value in register ST(0) remains unchanged. The instruction does not raise an exception when the source operand is out of range. It is up to the program to check the C2 flag for out-of-range conditions. Source values outside the range -2^{63} to $+2^{63}$ can be reduced to the range of the instruction by subtracting an appropriate integer multiple of 2π or by using the FPREM instruction with a divisor of 2π. See the section titled "Pi" in Chapter 8 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for a discussion of the proper value to use for π in performing such reductions.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF ST(0) < 2^63
    THEN
        C2 ← 0;
```

```
        ST(0) ← sin(ST(0));
    ELSE (* Source operand out of range *)
        C2 ← 1;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred.
	Set if result was rounded up; cleared otherwise.
C2	Set to 1 if outside range (-2^{63} < source operand < $+2^{63}$); otherwise, set to 0.
C0, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Source operand is an SNaN value, ∞, or unsupported format.
#D	Source operand is a denormal value.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FSINCOS—Sine and Cosine

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 FB	FSINCOS	Valid	Valid	Compute the sine and cosine of ST(0); replace ST(0) with the sine, and push the cosine onto the register stack.

Description

Computes both the sine and the cosine of the source operand in register ST(0), stores the sine in ST(0), and pushes the cosine onto the top of the FPU register stack. (This instruction is faster than executing the FSIN and FCOS instructions in succession.)

The source operand must be given in radians and must be within the range -2^{63} to $+2^{63}$. The following table shows the results obtained when taking the sine and cosine of various classes of numbers, assuming that underflow does not occur.

Table 3-46. FSINCOS Results

SRC	DEST	
ST(0)	ST(1) Cosine	ST(0) Sine
-•	*	*
– F	– 1 to + 1	– 1 to + 1
– 0	+ 1	– 0
+ 0	+ 1	+ 0
+ F	– 1 to + 1	– 1 to + 1
+•	*	*
NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

* Indicates floating-point invalid-arithmetic-operand (#IA) exception.

If the source operand is outside the acceptable range, the C2 flag in the FPU status word is set, and the value in register ST(0) remains unchanged. The instruction does not raise an exception when the source operand is out of range. It is up to the program to check the C2 flag for out-of-range conditions. Source values outside the range -2^{63} to $+2^{63}$ can be reduced to the range of the instruction by subtracting an appropriate integer multiple of 2π or by using the FPREM instruction with a divisor of 2π. See the section titled “Pi” in Chapter 8 of the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*, for a discussion of the proper value to use for π in performing such reductions.

This instruction’s operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF ST(0) < 2^63
    THEN
        C2 ← 0;
        TEMP ← cosine(ST(0));
        ST(0) ← sine(ST(0));
        TOP ← TOP − 1;
        ST(0) ← TEMP;
    ELSE (* Source operand out of range *)
        C2 ← 1;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred; set to 1 of stack overflow occurs.
	Set if result was rounded up; cleared otherwise.
C2	Set to 1 if outside range (-2^{63} < source operand < $+2^{63}$); otherwise, set to 0.
C0, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow or overflow occurred.
#IA	Source operand is an SNaN value, ∞, or unsupported format.
#D	Source operand is a denormal value.
#U	Result is too small for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FSQRT—Square Root

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 FA	FSQRT	Valid	Valid	Computes square root of ST(0) and stores the result in ST(0).

Description

Computes the square root of the source value in the ST(0) register and stores the result in ST(0).

The following table shows the results obtained when taking the square root of various classes of numbers, assuming that neither overflow nor underflow occurs.

Table 3-47. FSQRT Results

SRC (ST(0))	DEST (ST(0))
-•	*
– F	*
– 0	– 0
+ 0	+ 0
+ F	+ F
+•	+•
NaN	NaN

NOTES:

F Means finite floating-point value.

* Indicates floating-point invalid-arithmetic-operand (#IA) exception.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

ST(0) ← SquareRoot(ST(0));

FPU Flags Affected

C1	Set to 0 if stack underflow occurred.
	Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Source operand is an SNaN value or unsupported format.

	Source operand is a negative value (except for –0).
#D	Source operand is a denormal value.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FST/FSTP—Store Floating Point Value

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 /2	FST *m32fp*	Valid	Valid	Copy ST(0) to *m32fp*.
DD /2	FST *m64fp*	Valid	Valid	Copy ST(0) to *m64fp*.
DD D0+i	FST ST(i)	Valid	Valid	Copy ST(0) to ST(i).
D9 /3	FSTP *m32fp*	Valid	Valid	Copy ST(0) to *m32fp* and pop register stack.
DD /3	FSTP *m64fp*	Valid	Valid	Copy ST(0) to *m64fp* and pop register stack.
DB /7	FSTP *m80fp*	Valid	Valid	Copy ST(0) to *m80fp* and pop register stack.
DD D8+i	FSTP ST(i)	Valid	Valid	Copy ST(0) to ST(i) and pop register stack.

Description

The FST instruction copies the value in the ST(0) register to the destination operand, which can be a memory location or another register in the FPU register stack. When storing the value in memory, the value is converted to single-precision or double-precision floating-point format.

The FSTP instruction performs the same operation as the FST instruction and then pops the register stack. To pop the register stack, the processor marks the ST(0) register as empty and increments the stack pointer (TOP) by 1. The FSTP instruction can also store values in memory in double extended-precision floating-point format.

If the destination operand is a memory location, the operand specifies the address where the first byte of the destination value is to be stored. If the destination operand is a register, the operand specifies a register in the register stack relative to the top of the stack.

If the destination size is single-precision or double-precision, the significand of the value being stored is rounded to the width of the destination (according to the rounding mode specified by the RC field of the FPU control word), and the exponent is converted to the width and bias of the destination format. If the value being stored is too large for the destination format, a numeric overflow exception (#O) is generated and, if the exception is unmasked, no value is stored in the destination operand. If the value being stored is a denormal value, the denormal exception (#D) is not generated. This condition is simply signaled as a numeric underflow exception (#U) condition.

If the value being stored is ±0, ±∞, or a NaN, the least-significant bits of the significand and the exponent are truncated to fit the destination format. This operation preserves the value's identity as a 0, ∞, or NaN.

If the destination operand is a non-empty register, the invalid-operation exception is not generated.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
DEST ← ST(0);

IF Instruction = FSTP
    THEN
        PopRegisterStack;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred.
	Indicates rounding direction of if the floating-point inexact exception (#P) is generated: 0 ← not roundup; 1 ← roundup.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Source operand is an SNaN value or unsupported format. Does not occur if the source operand is in double extended-precision floating-point format.
#U	Result is too small for the destination format.
#O	Result is too large for the destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#GP(0)	If the destination is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FSTCW/FNSTCW—Store x87 FPU Control Word

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
9B D9 /7	FSTCW *m2byte*	Valid	Valid	Store FPU control word to *m2byte* after checking for pending unmasked floating-point exceptions.
D9 /7	FNSTCW* *m2byte*	Valid	Valid	Store FPU control word to *m2byte* without checking for pending unmasked floating-point exceptions.

NOTES:

* See IA-32 Architecture Compatibility section below.

Description

Stores the current value of the FPU control word at the specified destination in memory. The FSTCW instruction checks for and handles pending unmasked floating-point exceptions before storing the control word; the FNSTCW instruction does not.

The assembler issues two instructions for the FSTCW instruction (an FWAIT instruction followed by an FNSTCW instruction), and the processor executes each of these instructions in separately. If an exception is generated for either of these instructions, the save EIP points to the instruction that caused the exception.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

IA-32 Architecture Compatibility

When operating a Pentium or Intel486 processor in MS-DOS compatibility mode, it is possible (under unusual circumstances) for an FNSTCW instruction to be interrupted prior to being executed to handle a pending FPU exception. See the section titled "No-Wait FPU Instructions Can Get FPU Interrupt in Window" in Appendix D of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for a description of these circumstances. An FNSTCW instruction cannot be interrupted in this way on a Pentium 4, Intel Xeon, or P6 family processor.

Operation

```
DEST ← FPUControlWord;
```

FPU Flags Affected

The C0, C1, C2, and C3 flags are undefined.

Floating-Point Exceptions

None.

Protected Mode Exceptions

#GP(0)	If the destination is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FSTENV/FNSTENV—Store x87 FPU Environment

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
9B D9 /6	FSTENV *m14/28byte*	Valid	Valid	Store FPU environment to *m14byte* or *m28byte* after checking for pending unmasked floating-point exceptions. Then mask all floating-point exceptions.
D9 /6	FNSTENV* *m14/28byte*	Valid	Valid	Store FPU environment to *m14byte* or *m28byte* without checking for pending unmasked floating-point exceptions. Then mask all floating-point exceptions.

NOTES:

* See IA-32 Architecture Compatibility section below.

Description

Saves the current FPU operating environment at the memory location specified with the destination operand, and then masks all floating-point exceptions. The FPU operating environment consists of the FPU control word, status word, tag word, instruction pointer, data pointer, and last opcode. Figures 8-9 through 8-12 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, show the layout in memory of the stored environment, depending on the operating mode of the processor (protected or real) and the current operand-size attribute (16-bit or 32-bit). In virtual-8086 mode, the real mode layouts are used.

The FSTENV instruction checks for and handles any pending unmasked floating-point exceptions before storing the FPU environment; the FNSTENV instruction does not. The saved image reflects the state of the FPU after all floating-point instructions preceding the FSTENV/FNSTENV instruction in the instruction stream have been executed.

These instructions are often used by exception handlers because they provide access to the FPU instruction and data pointers. The environment is typically saved in the stack. Masking all exceptions after saving the environment prevents floating-point exceptions from interrupting the exception handler.

The assembler issues two instructions for the FSTENV instruction (an FWAIT instruction followed by an FNSTENV instruction), and the processor executes each of these instructions separately. If an exception is generated for either of these instructions, the save EIP points to the instruction that caused the exception.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

IA-32 Architecture Compatibility

When operating a Pentium or Intel486 processor in MS-DOS compatibility mode, it is possible (under unusual circumstances) for an FNSTENV instruction to be interrupted prior to being executed to handle a pending FPU exception. See the section titled "No-Wait FPU Instructions Can Get FPU Interrupt in Window" in Appendix D of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for a description of these circumstances. An FNSTENV instruction cannot be interrupted in this way on a Pentium 4, Intel Xeon, or P6 family processor.

Operation

```
DEST[FPUControlWord] ← FPUControlWord;
DEST[FPUStatusWord] ← FPUStatusWord;
DEST[FPUTagWord] ← FPUTagWord;
DEST[FPUDataPointer] ← FPUDataPointer;
DEST[FPUInstructionPointer] ← FPUInstructionPointer;
DEST[FPULastInstructionOpcode] ← FPULastInstructionOpcode;
```

FPU Flags Affected

The C0, C1, C2, and C3 are undefined.

Floating-Point Exceptions

None.

Protected Mode Exceptions

#GP(0)	If the destination is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.

#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FSTSW/FNSTSW—Store x87 FPU Status Word

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
9B DD /7	FSTSW *m2byte*	Valid	Valid	Store FPU status word at *m2byte* after checking for pending unmasked floating-point exceptions.
9B DF E0	FSTSW AX	Valid	Valid	Store FPU status word in AX register after checking for pending unmasked floating-point exceptions.
DD /7	FNSTSW* *m2byte*	Valid	Valid	Store FPU status word at *m2byte* without checking for pending unmasked floating-point exceptions.
DF E0	FNSTSW* AX	Valid	Valid	Store FPU status word in AX register without checking for pending unmasked floating-point exceptions.

NOTES:

* See IA-32 Architecture Compatibility section below.

Description

Stores the current value of the x87 FPU status word in the destination location. The destination operand can be either a two-byte memory location or the AX register. The FSTSW instruction checks for and handles pending unmasked floating-point exceptions before storing the status word; the FNSTSW instruction does not.

The FNSTSW AX form of the instruction is used primarily in conditional branching (for instance, after an FPU comparison instruction or an FPREM, FPREM1, or FXAM instruction), where the direction of the branch depends on the state of the FPU condition code flags. (See the section titled “Branching and Conditional Moves on FPU Condition Codes” in Chapter 8 of the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*.) This instruction can also be used to invoke exception handlers (by examining the exception flags) in environments that do not use interrupts. When the FNSTSW AX instruction is executed, the AX register is updated before the processor executes any further instructions. The status stored in the AX register is thus guaranteed to be from the completion of the prior FPU instruction.

The assembler issues two instructions for the FSTSW instruction (an FWAIT instruction followed by an FNSTSW instruction), and the processor executes each of these instructions separately. If an exception is generated for either of these instructions, the save EIP points to the instruction that caused the exception.

This instruction’s operation is the same in non-64-bit modes and 64-bit mode.

IA-32 Architecture Compatibility

When operating a Pentium or Intel486 processor in MS-DOS compatibility mode, it is possible (under unusual circumstances) for an FNSTSW instruction to be interrupted prior to being executed to handle a pending FPU exception. See the section titled "No-Wait FPU Instructions Can Get FPU Interrupt in Window" in Appendix D of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for a description of these circumstances. An FNSTSW instruction cannot be interrupted in this way on a Pentium 4, Intel Xeon, or P6 family processor.

Operation

```
DEST ← FPUStatusWord;
```

FPU Flags Affected

The C0, C1, C2, and C3 are undefined.

Floating-Point Exceptions

None.

Protected Mode Exceptions

#GP(0)	If the destination is located in a non-writable segment. If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit. If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FSUB/FSUBP/FISUB—Subtract

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D8 /4	FSUB *m32fp*	Valid	Valid	Subtract *m32fp* from ST(0) and store result in ST(0).
DC /4	FSUB *m64fp*	Valid	Valid	Subtract *m64fp* from ST(0) and store result in ST(0).
D8 E0+i	FSUB ST(0), ST(i)	Valid	Valid	Subtract ST(i) from ST(0) and store result in ST(0).
DC E8+i	FSUB ST(i), ST(0)	Valid	Valid	Subtract ST(0) from ST(i) and store result in ST(i).
DE E8+i	FSUBP ST(i), ST(0)	Valid	Valid	Subtract ST(0) from ST(i), store result in ST(i), and pop register stack.
DE E9	FSUBP	Valid	Valid	Subtract ST(0) from ST(1), store result in ST(1), and pop register stack.
DA /4	FISUB *m32int*	Valid	Valid	Subtract *m32int* from ST(0) and store result in ST(0).
DE /4	FISUB *m16int*	Valid	Valid	Subtract *m16int* from ST(0) and store result in ST(0).

Description

Subtracts the source operand from the destination operand and stores the difference in the destination location. The destination operand is always an FPU data register; the source operand can be a register or a memory location. Source operands in memory can be in single-precision or double-precision floating-point format or in word or doubleword integer format.

The no-operand version of the instruction subtracts the contents of the ST(0) register from the ST(1) register and stores the result in ST(1). The one-operand version subtracts the contents of a memory location (either a floating-point or an integer value) from the contents of the ST(0) register and stores the result in ST(0). The two-operand version, subtracts the contents of the ST(0) register from the ST(i) register or vice versa.

The FSUBP instructions perform the additional operation of popping the FPU register stack following the subtraction. To pop the register stack, the processor marks the ST(0) register as empty and increments the stack pointer (TOP) by 1. The no-operand version of the floating-point subtract instructions always results in the register stack being popped. In some assemblers, the mnemonic for this instruction is FSUB rather than FSUBP.

The FISUB instructions convert an integer source operand to double extended-precision floating-point format before performing the subtraction.

Table 3-48 shows the results obtained when subtracting various classes of numbers from one another, assuming that neither overflow nor underflow occurs. Here, the SRC value is subtracted from the DEST value (DEST – SRC = result).

When the difference between two operands of like sign is 0, the result is +0, except for the round toward –∞ mode, in which case the result is –0. This instruction also guarantees that +0 – (–0) = +0, and that –0 – (+0) = –0. When the source operand is an integer 0, it is treated as a +0.

When one operand is ∞, the result is ∞ of the expected sign. If both operands are ∞ of the same sign, an invalid-operation exception is generated.

Table 3-48. FSUB/FSUBP/FISUB Results

		SRC						
		- •	– F or – I	– 0	+ 0	+ F or + I	+•	NaN
DEST	- •	*	-•	- •	- •	- •	-•	NaN
	– F	+•	±F or ±0	DEST	DEST	– F	-•	NaN
	– 0	+•	–SRC	±0	– 0	– SRC	-•	NaN
	+ 0	+•	–SRC	+ 0	±0	– SRC	-•	NaN
	+ F	+•	+ F	DEST	DEST	±F or ±0	-•	NaN
	+•	+•	+•	+•	+•	+•	*	NaN
	NaN	NaN	NaN	NaN	NaN	NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

I Means integer.

* Indicates floating-point invalid-arithmetic-operand (#IA) exception.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF Instruction = FISUB
    THEN
        DEST ← DEST − ConvertToDoubleExtendedPrecisionFP(SRC);
    ELSE (* Source operand is floating-point value *)
        DEST ← DEST − SRC;
FI;
```

```
IF Instruction = FSUBP
    THEN
        PopRegisterStack;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred.
	Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Operand is an SNaN value or unsupported format.
	Operands are infinities of like sign.
#D	Source operand is a denormal value.
#U	Result is too small for destination format.
#O	Result is too large for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FSUBR/FSUBRP/FISUBR—Reverse Subtract

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D8 /5	FSUBR *m32fp*	Valid	Valid	Subtract ST(0) from *m32fp* and store result in ST(0).
DC /5	FSUBR *m64fp*	Valid	Valid	Subtract ST(0) from *m64fp* and store result in ST(0).
D8 E8+i	FSUBR ST(0), ST(i)	Valid	Valid	Subtract ST(0) from ST(i) and store result in ST(0).
DC E0+i	FSUBR ST(i), ST(0)	Valid	Valid	Subtract ST(i) from ST(0) and store result in ST(i).
DE E0+i	FSUBRP ST(i), ST(0)	Valid	Valid	Subtract ST(i) from ST(0), store result in ST(i), and pop register stack.
DE E1	FSUBRP	Valid	Valid	Subtract ST(1) from ST(0), store result in ST(1), and pop register stack.
DA /5	FISUBR *m32int*	Valid	Valid	Subtract ST(0) from *m32int* and store result in ST(0).
DE /5	FISUBR *m16int*	Valid	Valid	Subtract ST(0) from *m16int* and store result in ST(0).

Description

Subtracts the destination operand from the source operand and stores the difference in the destination location. The destination operand is always an FPU register; the source operand can be a register or a memory location. Source operands in memory can be in single-precision or double-precision floating-point format or in word or doubleword integer format.

These instructions perform the reverse operations of the FSUB, FSUBP, and FISUB instructions. They are provided to support more efficient coding.

The no-operand version of the instruction subtracts the contents of the ST(1) register from the ST(0) register and stores the result in ST(1). The one-operand version subtracts the contents of the ST(0) register from the contents of a memory location (either a floating-point or an integer value) and stores the result in ST(0). The two-operand version, subtracts the contents of the ST(i) register from the ST(0) register or vice versa.

The FSUBRP instructions perform the additional operation of popping the FPU register stack following the subtraction. To pop the register stack, the processor marks the ST(0) register as empty and increments the stack pointer (TOP) by 1. The no-operand version of the floating-point reverse subtract instructions always results in the register stack being popped. In some assemblers, the mnemonic for this instruction is FSUBR rather than FSUBRP.

The FISUBR instructions convert an integer source operand to double extended-precision floating-point format before performing the subtraction.

The following table shows the results obtained when subtracting various classes of numbers from one another, assuming that neither overflow nor underflow occurs. Here, the DEST value is subtracted from the SRC value (SRC – DEST = result).

When the difference between two operands of like sign is 0, the result is +0, except for the round toward –∞ mode, in which case the result is –0. This instruction also guarantees that +0 – (–0) = +0, and that –0 – (+0) = –0. When the source operand is an integer 0, it is treated as a +0.

When one operand is ∞, the result is ∞ of the expected sign. If both operands are ∞ of the same sign, an invalid-operation exception is generated.

Table 3-49. FSUBR/FSUBRP/FISUBR Results

		SRC						
		-•	–F or –I	–0	+0	+F or +I	+•	NaN
	- •	*	+•	+•	+•	+•	+•	NaN
	– F	-•	±F or ±0	–DEST	–DEST	+ F	+•	NaN
DEST	– 0	-•	SRC	±0	+ 0	SRC	+•	NaN
	+ 0	-•	SRC	– 0	±0	SRC	+•	NaN
	+ F	-•	– F	–DEST	–DEST	±F or ±0	+•	NaN
	+•	-•	-•	-•	-•	-•	*	NaN
	NaN	NaN	NaN	NaN	NaN	NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

I Means integer.

* Indicates floating-point invalid-arithmetic-operand (#IA) exception.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF Instruction = FISUBR
    THEN
        DEST ← ConvertToDoubleExtendedPrecisionFP(SRC) − DEST;
    ELSE (* Source operand is floating-point value *)
        DEST ← SRC − DEST; FI;
```

```
IF Instruction = FSUBRP
    THEN
        PopRegisterStack; FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred. Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Operand is an SNaN value or unsupported format. Operands are infinities of like sign.
#D	Source operand is a denormal value.
#U	Result is too small for destination format.
#O	Result is too large for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit. If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

FTST—TEST

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 E4	FTST	Valid	Valid	Compare ST(0) with 0.0.

Description

Compares the value in the ST(0) register with 0.0 and sets the condition code flags C0, C2, and C3 in the FPU status word according to the results (see table below).

Table 3-50. FTST Results

Condition	C3	C2	C0
ST(0) > 0.0	0	0	0
ST(0) < 0.0	0	0	1
ST(0) = 0.0	1	0	0
Unordered	1	1	1

This instruction performs an "unordered comparison." An unordered comparison also checks the class of the numbers being compared (see "FXAM—Examine ModR/M" in this chapter). If the value in register ST(0) is a NaN or is in an undefined format, the condition flags are set to "unordered" and the invalid operation exception is generated.

The sign of zero is ignored, so that (– 0.0 ← +0.0).

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
CASE (relation of operands) OF
    Not comparable:    C3, C2, C0 ← 111;
    ST(0) > 0.0:       C3, C2, C0 ← 000;
    ST(0) < 0.0:       C3, C2, C0 ← 001;
    ST(0) = 0.0:       C3, C2, C0 ← 100;
ESAC;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred; otherwise, set to 0.
C0, C2, C3	See Table 3-50.

Floating-Point Exceptions

#IS	Stack underflow occurred.

#IA	The source operand is a NaN value or is in an unsupported format.
#D	The source operand is a denormal value.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FUCOM/FUCOMP/FUCOMPP—Unordered Compare Floating Point Values

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
DD E0+i	FUCOM ST(i)	Valid	Valid	Compare ST(0) with ST(i).
DD E1	FUCOM	Valid	Valid	Compare ST(0) with ST(1).
DD E8+i	FUCOMP ST(i)	Valid	Valid	Compare ST(0) with ST(i) and pop register stack.
DD E9	FUCOMP	Valid	Valid	Compare ST(0) with ST(1) and pop register stack.
DA E9	FUCOMPP	Valid	Valid	Compare ST(0) with ST(1) and pop register stack twice.

Description

Performs an unordered comparison of the contents of register ST(0) and ST(i) and sets condition code flags C0, C2, and C3 in the FPU status word according to the results (see the table below). If no operand is specified, the contents of registers ST(0) and ST(1) are compared. The sign of zero is ignored, so that –0.0 is equal to +0.0.

Table 3-51. FUCOM/FUCOMP/FUCOMPP Results

Comparison Results*	C3	C2	C0
ST0 > ST(i)	0	0	0
ST0 < ST(i)	0	0	1
ST0 = ST(i)	1	0	0
Unordered	1	1	1

NOTES:

* Flags not set if unmasked invalid-arithmetic-operand (#IA) exception is generated.

An unordered comparison checks the class of the numbers being compared (see "FXAM—Examine ModR/M" in this chapter). The FUCOM/FUCOMP/FUCOMPP instructions perform the same operations as the FCOM/FCOMP/FCOMPP instructions. The only difference is that the FUCOM/FUCOMP/FUCOMPP instructions raise the invalid-arithmetic-operand exception (#IA) only when either or both operands are an SNaN or are in an unsupported format; QNaNs cause the condition code flags to be set to unordered, but do not cause an exception to be generated. The FCOM/FCOMP/FCOMPP instructions raise an invalid-operation exception when either or both of the operands are a NaN value of any kind or are in an unsupported format.

As with the FCOM/FCOMP/FCOMPP instructions, if the operation results in an invalid-arithmetic-operand exception being raised, the condition code flags are set only if the exception is masked.

The FUCOMP instruction pops the register stack following the comparison operation and the FUCOMPP instruction pops the register stack twice following the comparison operation. To pop the register stack, the processor marks the ST(0) register as empty and increments the stack pointer (TOP) by 1.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
CASE (relation of operands) OF
    ST > SRC:          C3, C2, C0 ← 000;
    ST < SRC:          C3, C2, C0 ← 001;
    ST = SRC:          C3, C2, C0 ← 100;
ESAC;

IF ST(0) or SRC = QNaN, but not SNaN or unsupported format
    THEN
        C3, C2, C0 ← 111;
    ELSE (* ST(0) or SRC is SNaN or unsupported format *)
         #IA;
        IF FPUControlWord.IM = 1
            THEN
                C3, C2, C0 ← 111;
        FI;
FI;

IF Instruction = FUCOMP
    THEN
        PopRegisterStack;
FI;

IF Instruction = FUCOMPP
    THEN
        PopRegisterStack;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred.
C0, C2, C3	See Table 3-51.

Floating-Point Exceptions

#IS	Stack underflow occurred.

#IA	One or both operands are SNaN values or have unsupported formats. Detection of a QNaN value in and of itself does not raise an invalid-operand exception.
#D	One or both operands are denormal values.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FXAM—Examine ModR/M

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 E5	FXAM	Valid	Valid	Classify value or number in ST(0).

Description

Examines the contents of the ST(0) register and sets the condition code flags C0, C2, and C3 in the FPU status word to indicate the class of value or number in the register (see the table below).

Table 3-52. FXAM Results

Class	C3	C2	C0
Unsupported	0	0	0
NaN	0	0	1
Normal finite number	0	1	0
Infinity	0	1	1
Zero	1	0	0
Empty	1	0	1
Denormal number	1	1	0

The C1 flag is set to the sign of the value in ST(0), regardless of whether the register is empty or full.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
C1 ← sign bit of ST; (* 0 for positive, 1 for negative *)

CASE (class of value or number in ST(0)) OF
    Unsupported:C3, C2, C0 ← 000;
    NaN:         C3, C2, C0 ← 001;
    Normal:      C3, C2, C0 ← 010;
    Infinity:    C3, C2, C0 ← 011;
    Zero:        C3, C2, C0 ← 100;
    Empty:       C3, C2, C0 ← 101;
    Denormal:    C3, C2, C0 ← 110;
ESAC;
```

FPU Flags Affected

C1	Sign of value in ST(0).
C0, C2, C3	See Table 3-52.

Floating-Point Exceptions

None.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FXCH—Exchange Register Contents

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 C8+i	FXCH ST(i)	Valid	Valid	Exchange the contents of ST(0) and ST(i).
D9 C9	FXCH	Valid	Valid	Exchange the contents of ST(0) and ST(1).

Description

Exchanges the contents of registers ST(0) and ST(i). If no source operand is specified, the contents of ST(0) and ST(1) are exchanged.

This instruction provides a simple means of moving values in the FPU register stack to the top of the stack [ST(0)], so that they can be operated on by those floating-point instructions that can only operate on values in ST(0). For example, the following instruction sequence takes the square root of the third register from the top of the register stack:

```
FXCH ST(3);
FSQRT;
FXCH ST(3);
```

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF (Number-of-operands) is 1
    THEN
        temp ← ST(0);
        ST(0) ← SRC;
        SRC ← temp;
    ELSE
        temp ← ST(0);
        ST(0) ← ST(1);
        ST(1) ← temp;
FI;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred; otherwise, set to 1.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FXRSTOR—Restore x87 FPU, MMX , XMM, and MXCSR State

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F AE /1	FXRSTOR *m512byte*	A	Valid	Valid	Restore the x87 FPU, MMX, XMM, and MXCSR register state from *m512byte*.
REX.W+ 0F AE /1	FXRSTOR64 *m512byte*	A	Valid	N.E.	Restore the x87 FPU, MMX, XMM, and MXCSR register state from *m512byte*.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r)	NA	NA	NA

Description

Reloads the x87 FPU, MMX technology, XMM, and MXCSR registers from the 512-byte memory image specified in the source operand. This data should have been written to memory previously using the FXSAVE instruction, and in the same format as required by the operating modes. The first byte of the data should be located on a 16-byte boundary. There are three distinct layouts of the FXSAVE state map: one for legacy and compatibility mode, a second format for 64-bit mode FXSAVE/FXRSTOR with REX.W=0, and the third format is for 64-bit mode with FXSAVE64/FXRSTOR64. Table 3-53 shows the layout of the legacy/compatibility mode state information in memory and describes the fields in the memory image for the FXRSTOR and FXSAVE instructions. Table 3-56 shows the layout of the 64-bit mode state information when REX.W is set (FXSAVE64/FXRSTOR64). Table 3-57 shows the layout of the 64-bit mode state information when REX.W is clear (FXSAVE/FXRSTOR).

The state image referenced with an FXRSTOR instruction must have been saved using an FXSAVE instruction or be in the same format as required by Table 3-53, Table 3-56, or Table 3-57. Referencing a state image saved with an FSAVE, FNSAVE instruction or incompatible field layout will result in an incorrect state restoration.

The FXRSTOR instruction does not flush pending x87 FPU exceptions. To check and raise exceptions when loading x87 FPU state information with the FXRSTOR instruction, use an FWAIT instruction after the FXRSTOR instruction.

If the OSFXSR bit in control register CR4 is not set, the FXRSTOR instruction may not restore the states of the XMM and MXCSR registers. This behavior is implementation dependent.

If the MXCSR state contains an unmasked exception with a corresponding status flag also set, loading the register with the FXRSTOR instruction will not result in a SIMD floating-point error condition being generated. Only the next occurrence of this unmasked exception will result in the exception being generated.

Bits 16 through 32 of the MXCSR register are defined as reserved and should be set to 0. Attempting to write a 1 in any of these bits from the saved state image will result in a general protection exception (#GP) being generated.

Bytes 464:511 of an FXSAVE image are available for software use. FXRSTOR ignores the content of bytes 464:511 in an FXSAVE state image.

Operation

```
(x87 FPU, MMX, XMM7-XMM0, MXCSR) ← Load(SRC);
```

x87 FPU and SIMD Floating-Point Exceptions

None.

Protected Mode Exceptions

#GP(0)	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
	If a memory operand is not aligned on a 16-byte boundary, regardless of segment. (See alignment check exception [#AC] below.)
	For an attempt to set reserved bits in MXCSR.
#SS(0)	For an illegal address in the SS segment.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
	If CR0.EM[bit 2] = 1.
#UD	If CPUID.01H:EDX.FXSR[bit 24] = 0.
	If instruction is preceded by a LOCK prefix.
#AC	If this exception is disabled a general protection exception (#GP) is signaled if the memory operand is not aligned on a 16-byte boundary, as described above. If the alignment check exception (#AC) is enabled (and the CPL is 3), signaling of #AC is not guaranteed and may vary with implementation, as follows. In all implementations where #AC is not signaled, a general protection exception is signaled in its place. In addition, the width of the alignment check may also vary with implementation. For instance, for a given implementation, an alignment check exception might be signaled for a 2-byte misalignment, whereas a general protection exception might be signaled for all other misalignments (4-, 8-, or 16-byte misalignments).
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand is not aligned on a 16-byte boundary, regardless of segment.
	If any part of the operand lies outside the effective address space from 0 to FFFFH.
	For an attempt to set reserved bits in MXCSR.
#NM	If CR0.TS[bit 3] = 1.
	If CR0.EM[bit 2] = 1.
#UD	If CPUID.01H:EDX.FXSR[bit 24] = 0.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	For a page fault.
#AC	For unaligned memory reference.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
	If memory operand is not aligned on a 16-byte boundary, regardless of segment.
	For an attempt to set reserved bits in MXCSR.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
	If CR0.EM[bit 2] = 1.
#UD	If CPUID.01H:EDX.FXSR[bit 24] = 0.
	If instruction is preceded by a LOCK prefix.
#AC	If this exception is disabled a general protection exception (#GP) is signaled if the memory operand is not aligned on a 16-byte boundary, as described above. If the alignment check exception (#AC) is enabled (and the CPL is 3), signaling of #AC is not guaranteed and may vary with implementation, as follows. In all implementations where #AC is not signaled, a general protection exception is signaled in its place. In addition, the width of the alignment check may also vary with implemen-

tation. For instance, for a given implementation, an alignment check exception might be signaled for a 2-byte misalignment, whereas a general protection exception might be signaled for all other misalignments (4-, 8-, or 16-byte misalignments).

FXSAVE—Save x87 FPU, MMX Technology, and SSE State

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F AE /0	FXSAVE *m512byte*	A	Valid	Valid	Save the x87 FPU, MMX, XMM, and MXCSR register state to *m512byte*.
REX.W+ 0F AE /0	FXSAVE64 *m512byte*	A	Valid	N.E.	Save the x87 FPU, MMX, XMM, and MXCSR register state to *m512byte*.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (w)	NA	NA	NA

Description

Saves the current state of the x87 FPU, MMX technology, XMM, and MXCSR registers to a 512-byte memory location specified in the destination operand. The content layout of the 512 byte region depends on whether the processor is operating in non-64-bit operating modes or 64-bit sub-mode of IA-32e mode.

Bytes 464:511 are available to software use. The processor does not write to bytes 464:511 of an FXSAVE area.

The operation of FXSAVE in non-64-bit modes is described first.

Non-64-Bit Mode Operation

Table 3-53 shows the layout of the state information in memory when the processor is operating in legacy modes.

Table 3-53. Non-64-bit-Mode Layout of FXSAVE and FXRSTOR Memory Region

<table>
<tr><th>15 14</th><th>13 12</th><th>11 10</th><th>9 8</th><th>7 6</th><th>5</th><th>4</th><th>3 2</th><th>1 0</th><th></th></tr>
<tr><td>Rsrvd</td><td>CS</td><td colspan="2">FPU IP</td><td>FOP</td><td>Rsrvd</td><td>FTW</td><td>FSW</td><td>FCW</td><td>0</td></tr>
<tr><td colspan="2">MXCSR_MASK</td><td colspan="2">MXCSR</td><td>Rsrvd</td><td colspan="2">DS</td><td colspan="2">FPU DP</td><td>16</td></tr>
<tr><td colspan="3">Reserved</td><td colspan="6">ST0/MM0</td><td>32</td></tr>
<tr><td colspan="3">Reserved</td><td colspan="6">ST1/MM1</td><td>48</td></tr>
<tr><td colspan="3">Reserved</td><td colspan="6">ST2/MM2</td><td>64</td></tr>
<tr><td colspan="3">Reserved</td><td colspan="6">ST3/MM3</td><td>80</td></tr>
<tr><td colspan="3">Reserved</td><td colspan="6">ST4/MM4</td><td>96</td></tr>
</table>

Table 3-53. Non-64-bit-Mode Layout of FXSAVE and FXRSTOR Memory Region (Contd.)

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
Reserved						ST5/MM5										112
Reserved						ST6/MM6										128
Reserved						ST7/MM7										144
XMM0																160
XMM1																176
XMM2																192
XMM3																208
XMM4																224
XMM5																240
XMM6																256
XMM7																272
Reserved																288
Reserved																304
Reserved																320
Reserved																336
Reserved																352
Reserved																368
Reserved																384
Reserved																400
Reserved																416
Reserved																432
Reserved																448
Available																464
Available																480
Available																496

The destination operand contains the first byte of the memory image, and it must be aligned on a 16-byte boundary. A misaligned destination operand will result in a general-protection (#GP) exception being generated (or in some cases, an alignment check exception [#AC]).

The FXSAVE instruction is used when an operating system needs to perform a context switch or when an exception handler needs to save and examine the current state of the x87 FPU, MMX technology, and/or XMM and MXCSR registers.

The fields in Table 3-53 are defined in Table 3-54.

Table 3-54. Field Definitions

Field	Definition
FCW	x87 FPU Control Word (16 bits). See Figure 8-6 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for the layout of the x87 FPU control word.
FSW	x87 FPU Status Word (16 bits). See Figure 8-4 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for the layout of the x87 FPU status word.
Abridged FTW	x87 FPU Tag Word (8 bits). The tag information saved here is abridged, as described in the following paragraphs.
FOP	x87 FPU Opcode (16 bits). The lower 11 bits of this field contain the opcode, upper 5 bits are reserved. See Figure 8-8 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for the layout of the x87 FPU opcode field.
FPU IP	x87 FPU Instruction Pointer Offset (32 bits). The contents of this field differ depending on the current addressing mode (32-bit or 16-bit) of the processor when the FXSAVE instruction was executed: 32-bit mode — 32-bit IP offset. 16-bit mode — low 16 bits are IP offset; high 16 bits are reserved. See "x87 FPU Instruction and Operand (Data) Pointers" in Chapter 8 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for a description of the x87 FPU instruction pointer.
CS	x87 FPU Instruction Pointer Selector (16 bits).
FPU DP	x87 FPU Instruction Operand (Data) Pointer Offset (32 bits). The contents of this field differ depending on the current addressing mode (32-bit or 16-bit) of the processor when the FXSAVE instruction was executed: 32-bit mode — 32-bit DP offset. 16-bit mode — low 16 bits are DP offset; high 16 bits are reserved. See "x87 FPU Instruction and Operand (Data) Pointers" in Chapter 8 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for a description of the x87 FPU operand pointer.
DS	x87 FPU Instruction Operand (Data) Pointer Selector (16 bits).

Table 3-54. Field Definitions (Contd.)

Field	Definition
MXCSR	MXCSR Register State (32 bits). See Figure 10-3 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for the layout of the MXCSR register. If the OSFXSR bit in control register CR4 is not set, the FXSAVE instruction may not save this register. This behavior is implementation dependent.
MXCSR_ MASK	MXCSR_MASK (32 bits). This mask can be used to adjust values written to the MXCSR register, ensuring that reserved bits are set to 0. Set the mask bits and flags in MXCSR to the mode of operation desired for SSE and SSE2 SIMD floating-point instructions. See "Guidelines for Writing to the MXCSR Register" in Chapter 11 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for instructions for how to determine and use the MXCSR_MASK value.
ST0/MM0 through ST7/MM7	x87 FPU or MMX technology registers. These 80-bit fields contain the x87 FPU data registers or the MMX technology registers, depending on the state of the processor prior to the execution of the FXSAVE instruction. If the processor had been executing x87 FPU instruction prior to the FXSAVE instruction, the x87 FPU data registers are saved; if it had been executing MMX instructions (or SSE or SSE2 instructions that operated on the MMX technology registers), the MMX technology registers are saved. When the MMX technology registers are saved, the high 16 bits of the field are reserved.
XMM0 through XMM7	XMM registers (128 bits per field). If the OSFXSR bit in control register CR4 is not set, the FXSAVE instruction may not save these registers. This behavior is implementation dependent.

The FXSAVE instruction saves an abridged version of the x87 FPU tag word in the FTW field (unlike the FSAVE instruction, which saves the complete tag word). The tag information is saved in physical register order (R0 through R7), rather than in top-of-stack (TOS) order. With the FXSAVE instruction, however, only a single bit (1 for valid or 0 for empty) is saved for each tag. For example, assume that the tag word is currently set as follows:

```
R7  R6  R5  R4  R3  R2  R1  R0
11  xx  xx  xx  11  11  11  11
```

Here, 11B indicates empty stack elements and "xx" indicates valid (00B), zero (01B), or special (10B).

For this example, the FXSAVE instruction saves only the following 8 bits of information:

```
R7  R6  R5  R4  R3  R2  R1  R0
0   1   1   1   0   0   0   0
```

Here, a 1 is saved for any valid, zero, or special tag, and a 0 is saved for any empty tag.

The operation of the FXSAVE instruction differs from that of the FSAVE instruction, the as follows:

- FXSAVE instruction does not check for pending unmasked floating-point exceptions. (The FXSAVE operation in this regard is similar to the operation of the FNSAVE instruction).
- After the FXSAVE instruction has saved the state of the x87 FPU, MMX technology, XMM, and MXCSR registers, the processor retains the contents of the registers. Because of this behavior, the FXSAVE instruction cannot be used by an application program to pass a "clean" x87 FPU state to a procedure, since it retains the current state. To clean the x87 FPU state, an application must explicitly execute an FINIT instruction after an FXSAVE instruction to reinitialize the x87 FPU state.
- The format of the memory image saved with the FXSAVE instruction is the same regardless of the current addressing mode (32-bit or 16-bit) and operating mode (protected, real address, or system management). This behavior differs from the FSAVE instructions, where the memory image format is different depending on the addressing mode and operating mode. Because of the different image formats, the memory image saved with the FXSAVE instruction cannot be restored correctly with the FRSTOR instruction, and likewise the state saved with the FSAVE instruction cannot be restored correctly with the FXRSTOR instruction.

The FSAVE format for FTW can be recreated from the FTW valid bits and the stored 80-bit FP data (assuming the stored data was not the contents of MMX technology registers) using Table 3-55.

Table 3-55. Recreating FSAVE Format

Exponent all 1's	Exponent all 0's	Fraction all 0's	J and M bits	FTW valid bit	x87 FTW	
0	0	0	0x	1	Special	10
0	0	0	1x	1	Valid	00
0	0	1	00	1	Special	10
0	0	1	10	1	Valid	00
0	1	0	0x	1	Special	10
0	1	0	1x	1	Special	10
0	1	1	00	1	Zero	01
0	1	1	10	1	Special	10
1	0	0	1x	1	Special	10
1	0	0	1x	1	Special	10

Table 3-55. Recreating FSAVE Format (Contd.)

Exponent all 1's	Exponent all 0's	Fraction all 0's	J and M bits	FTW valid bit	x87 FTW	
1	0	1	00	1	Special	10
1	0	1	10	1	Special	10
For all legal combinations above.				0	Empty	11

The J-bit is defined to be the 1-bit binary integer to the left of the decimal place in the significand. The M-bit is defined to be the most significant bit of the fractional portion of the significand (i.e., the bit immediately to the right of the decimal place).

When the M-bit is the most significant bit of the fractional portion of the significand, it must be 0 if the fraction is all 0's.

IA-32e Mode Operation

In compatibility sub-mode of IA-32e mode, legacy SSE registers, XMM0 through XMM7, are saved according to the legacy FXSAVE map. In 64-bit mode, all of the SSE registers, XMM0 through XMM15, are saved. Additionally, there are two different layouts of the FXSAVE map in 64-bit mode, corresponding to FXSAVE64 (which requires REX.W=1) and FXSAVE (REX.W=0). In the FXSAVE64 map (Table 3-56), the FPU IP and FPU DP pointers are 64-bit wide. In the FXSAVE map for 64-bit mode (Table 3-57), the FPU IP and FPU DP pointers are 32-bits.

Table 3-56. Layout of the 64-bit-mode FXSAVE64 Map (requires REX.W = 1)

15 14	13 12	11 10	9 8	7 6	5	4	3 2	1 0	
FPU IP				FOP	Re-served	FTW	FSW	FCW	**0**
MXCSR_MASK		MXCSR		FPU DP					**16**
Reserved			ST0/MM0						**32**
Reserved			ST1/MM1						**48**
Reserved			ST2/MM2						**64**
Reserved			ST3/MM3						**80**
Reserved			ST4/MM4						**96**
Reserved			ST5/MM5						**112**
Reserved			ST6/MM6						**128**
Reserved			ST7/MM7						**144**
XMM0									**160**
XMM1									**176**

Table 3-56. Layout of the 64-bit-mode FXSAVE64 Map (requires REX.W = 1) (Contd.)

15 14	13 12	11 10	9 8	7 6	5	4	3 2	1 0	
XMM2									**192**
XMM3									**208**
XMM4									**224**
XMM5									**240**
XMM6									**256**
XMM7									**272**
XMM8									**288**
XMM9									**304**
XMM10									**320**
XMM11									**336**
XMM12									**352**
XMM13									**368**
XMM14									**384**
XMM15									**400**
Reserved									**416**
Reserved									**432**
Reserved									**448**
Available									**464**
Available									**480**
Available									**496**

Table 3-57. Layout of the 64-bit-mode FXSAVE Map (REX.W = 0)

15 14	13 12	11 10	9 8	7 6	5	4	3 2	1 0	
Reserved	CS	FPU IP		FOP	Re-served	FTW	FSW	FCW	**0**
MXCSR_MASK		MXCSR		Re-served			FPU DP		**16**
Reserved			ST0/MM0						**32**
Reserved			ST1/MM1						**48**
Reserved			ST2/MM2						**64**
Reserved			ST3/MM3						**80**

Table 3-57. Layout of the 64-bit-mode FXSAVE Map (REX.W = 0) (Contd.) (Contd.)

15 14	13 12	11 10	9 8	7 6	5	4	3 2	1 0	
Reserved			ST4/MM4						**96**
Reserved			ST5/MM5						**112**
Reserved			ST6/MM6						**128**
Reserved			ST7/MM7						**144**
XMM0									**160**
XMM1									**176**
XMM2									**192**
XMM3									**208**
XMM4									**224**
XMM5									**240**
XMM6									**256**
XMM7									**272**
XMM8									**288**
XMM9									**304**
XMM10									**320**
XMM11									**336**
XMM12									**352**
XMM13									**368**
XMM14									**384**
XMM15									**400**
Reserved									**416**
Reserved									**432**
Reserved									**448**
Available									**464**
Available									**480**
Available									**496**

Operation

```
IF 64-Bit Mode
    THEN
        IF REX.W = 1
            THEN
```

```
                DEST ← Save64BitPromotedFxsave(x87 FPU, MMX, XMM7-XMM0,
                MXCSR);
            ELSE
                DEST ← Save64BitDefaultFxsave(x87 FPU, MMX, XMM7-XMM0, MXCSR);
        FI;
    ELSE
        DEST ← SaveLegacyFxsave(x87 FPU, MMX, XMM7-XMM0, MXCSR);
FI;
```

Protected Mode Exceptions

#GP(0)	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
	If a memory operand is not aligned on a 16-byte boundary, regardless of segment. (See the description of the alignment check exception [#AC] below.)
#SS(0)	For an illegal address in the SS segment.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
	If CR0.EM[bit 2] = 1.
#UD	If CPUID.01H:EDX.FXSR[bit 24] = 0.
#UD	If the LOCK prefix is used.
#AC	If this exception is disabled a general protection exception (#GP) is signaled if the memory operand is not aligned on a 16-byte boundary, as described above. If the alignment check exception (#AC) is enabled (and the CPL is 3), signaling of #AC is not guaranteed and may vary with implementation, as follows. In all implementations where #AC is not signaled, a general protection exception is signaled in its place. In addition, the width of the alignment check may also vary with implementation. For instance, for a given implementation, an alignment check exception might be signaled for a 2-byte misalignment, whereas a general protection exception might be signaled for all other misalignments (4-, 8-, or 16-byte misalignments).

Real-Address Mode Exceptions

#GP	If a memory operand is not aligned on a 16-byte boundary, regardless of segment.
	If any part of the operand lies outside the effective address space from 0 to FFFFH.
#NM	If CR0.TS[bit 3] = 1.
	If CR0.EM[bit 2] = 1.

#UD	If CPUID.01H:EDX.FXSR[bit 24] = 0.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	For a page fault.
#AC	For unaligned memory reference.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
	If memory operand is not aligned on a 16-byte boundary, regardless of segment.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
	If CR0.EM[bit 2] = 1.
#UD	If CPUID.01H:EDX.FXSR[bit 24] = 0.
	If the LOCK prefix is used.
#AC	If this exception is disabled a general protection exception (#GP) is signaled if the memory operand is not aligned on a 16-byte boundary, as described above. If the alignment check exception (#AC) is enabled (and the CPL is 3), signaling of #AC is not guaranteed and may vary with implementation, as follows. In all implementations where #AC is not signaled, a general protection exception is signaled in its place. In addition, the width of the alignment check may also vary with implementation. For instance, for a given implementation, an alignment check exception might be signaled for a 2-byte misalignment, whereas a general protection exception might be signaled for all other misalignments (4-, 8-, or 16-byte misalignments).

Implementation Note

The order in which the processor signals general-protection (#GP) and page-fault (#PF) exceptions when they both occur on an instruction boundary is given in Table 5-2 in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 3B*. This order vary for FXSAVE for different processor implementations.

FXTRACT—Extract Exponent and Significand

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 F4	FXTRACT	Valid	Valid	Separate value in ST(0) into exponent and significand, store exponent in ST(0), and push the significand onto the register stack.

Description

Separates the source value in the ST(0) register into its exponent and significand, stores the exponent in ST(0), and pushes the significand onto the register stack. Following this operation, the new top-of-stack register ST(0) contains the value of the original significand expressed as a floating-point value. The sign and significand of this value are the same as those found in the source operand, and the exponent is 3FFFH (biased value for a true exponent of zero). The ST(1) register contains the value of the original operand's true (unbiased) exponent expressed as a floating-point value. (The operation performed by this instruction is a superset of the IEEE-recommended logb(*x*) function.)

This instruction and the F2XM1 instruction are useful for performing power and range scaling operations. The FXTRACT instruction is also useful for converting numbers in double extended-precision floating-point format to decimal representations (e.g., for printing or displaying).

If the floating-point zero-divide exception (#Z) is masked and the source operand is zero, an exponent value of –∞ is stored in register ST(1) and 0 with the sign of the source operand is stored in register ST(0).

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
TEMP ← Significand(ST(0));
ST(0) ← Exponent(ST(0));
TOP← TOP − 1;
ST(0) ← TEMP;
```

FPU Flags Affected

C1	Set to 0 if stack underflow occurred; set to 1 if stack overflow occurred.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow or overflow occurred.

#IA	Source operand is an SNaN value or unsupported format.
#Z	ST(0) operand is ±0.
#D	Source operand is a denormal value.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FYL2X—Compute y ∗ $\log_2 x$

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 F1	FYL2X	Valid	Valid	Replace ST(1) with (ST(1) ∗ $\log_2$ST(0)) and pop the register stack.

Description

Computes (ST(1) ∗ $\log_2$ (ST(0))), stores the result in resister ST(1), and pops the FPU register stack. The source operand in ST(0) must be a non-zero positive number.

The following table shows the results obtained when taking the log of various classes of numbers, assuming that neither overflow nor underflow occurs.

Table 3-58. FYL2X Results

		ST(0)							
		-•	–F	±0	+0<+F<+1	+ 1	+ F > + 1	+•	NaN
ST(1)	-•	*	*	+•	+•	*	- •	-•	NaN
	–F	*	*	**	+ F	– 0	– F	-•	NaN
	–0	*	*	*	+ 0	– 0	– 0	*	NaN
	+0	*	*	*	– 0	+ 0	+ 0	*	NaN
	+F	*	*	**	– F	+ 0	+ F	+•	NaN
	+•	*	*	-•	-•	*	+•	+•	NaN
	NaN	NaN	NaN	NaN	NaN	NaN	NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

* Indicates floating-point invalid-operation (#IA) exception.

** Indicates floating-point zero-divide (#Z) exception.

If the divide-by-zero exception is masked and register ST(0) contains ±0, the instruction returns ∞ with a sign that is the opposite of the sign of the source operand in register ST(1).

The FYL2X instruction is designed with a built-in multiplication to optimize the calculation of logarithms with an arbitrary positive base (b):

$$\log_b x \leftarrow (\log_2 b)^{-1} * \log_2 x$$

This instruction’s operation is the same in non-64-bit modes and 64-bit mode.

Operation

$ST(1) \leftarrow ST(1) * \log_2 ST(0);$
PopRegisterStack;

FPU Flags Affected

C1	Set to 0 if stack underflow occurred. Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Either operand is an SNaN or unsupported format. Source operand in register ST(0) is a negative finite value (not -0).
#Z	Source operand in register ST(0) is ±0.
#D	Source operand is a denormal value.
#U	Result is too small for destination format.
#O	Result is too large for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

FYL2XP1—Compute y $*$ $\log_2(x + 1)$

Opcode	Instruction	64-Bit Mode	Compat/ Leg Mode	Description
D9 F9	FYL2XP1	Valid	Valid	Replace ST(1) with ST(1) $*$ $\log_2$(ST(0) + 1.0) and pop the register stack.

Description

Computes (ST(1) $*$ $\log_2$(ST(0) + 1.0)), stores the result in register ST(1), and pops the FPU register stack. The source operand in ST(0) must be in the range:

$$-(1-\sqrt{2}/2))\text{to}(1-\sqrt{2}/2)$$

The source operand in ST(1) can range from $-\infty$ to $+\infty$. If the ST(0) operand is outside of its acceptable range, the result is undefined and software should not rely on an exception being generated. Under some circumstances exceptions may be generated when ST(0) is out of range, but this behavior is implementation specific and not guaranteed.

The following table shows the results obtained when taking the log epsilon of various classes of numbers, assuming that underflow does not occur.

Table 3-59. FYL2XP1 Results

		ST(0)				
		$-(1-(\sqrt{2}/2))$ to -0	-0	+0	+0 to +(1 - ($\sqrt{2}/2$))	NaN
ST(1)	-•	+•	*	*	-•	NaN
	-F	+F	+0	-0	-F	NaN
	-0	+0	+0	-0	-0	NaN
	+0	-0	-0	+0	+0	NaN
	+F	-F	-0	+0	+F	NaN
	+•	-•	*	*	+•	NaN
	NaN	NaN	NaN	NaN	NaN	NaN

NOTES:

F Means finite floating-point value.

* Indicates floating-point invalid-operation (#IA) exception.

This instruction provides optimal accuracy for values of epsilon [the value in register ST(0)] that are close to 0. For small epsilon (ε) values, more significant digits can be retained by using the FYL2XP1 instruction than by using (ε+1) as an argument to the FYL2X instruction. The (ε+1) expression is commonly found in compound interest and annuity calculations. The result can be simply converted into a value in another logarithm base by including a scale factor in the ST(1) source operand. The following equation is used to calculate the scale factor for a particular logarithm base, where n is the logarithm base desired for the result of the FYL2XP1 instruction:

scale factor $\leftarrow \log_n 2$

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

$ST(1) \leftarrow ST(1) * \log_2(ST(0) + 1.0);$
PopRegisterStack;

FPU Flags Affected

C1	Set to 0 if stack underflow occurred. Set if result was rounded up; cleared otherwise.
C0, C2, C3	Undefined.

Floating-Point Exceptions

#IS	Stack underflow occurred.
#IA	Either operand is an SNaN value or unsupported format.
#D	Source operand is a denormal value.
#U	Result is too small for destination format.
#O	Result is too large for destination format.
#P	Value cannot be represented exactly in destination format.

Protected Mode Exceptions

#NM	CR0.EM[bit 2] or CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

HADDPD—Packed Double-FP Horizontal Add

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 7C /*r* HADDPD *xmm1, xmm2/m128*	A	V/V	SSE3	Horizontal add packed double-precision floating-point values from *xmm2/m128* to *xmm1*.
VEX.NDS.128.66.0F.WIG 7C /r VHADDPD xmm1,xmm2, xmm3/m128	B	V/V	AVX	Horizontal add packed double-precision floating-point values from xmm2 and xmm3/mem.
VEX.NDS.256.66.0F.WIG 7C /r VHADDPD ymm1, ymm2, ymm3/m256	B	V/V	AVX	Horizontal add packed double-precision floating-point values from ymm2 and ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Adds the double-precision floating-point values in the high and low quadwords of the destination operand and stores the result in the low quadword of the destination operand.

Adds the double-precision floating-point values in the high and low quadwords of the source operand and stores the result in the high quadword of the destination operand.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

See Figure 3-19 for HADDPD; see Figure 3-20 for VHADDPD.

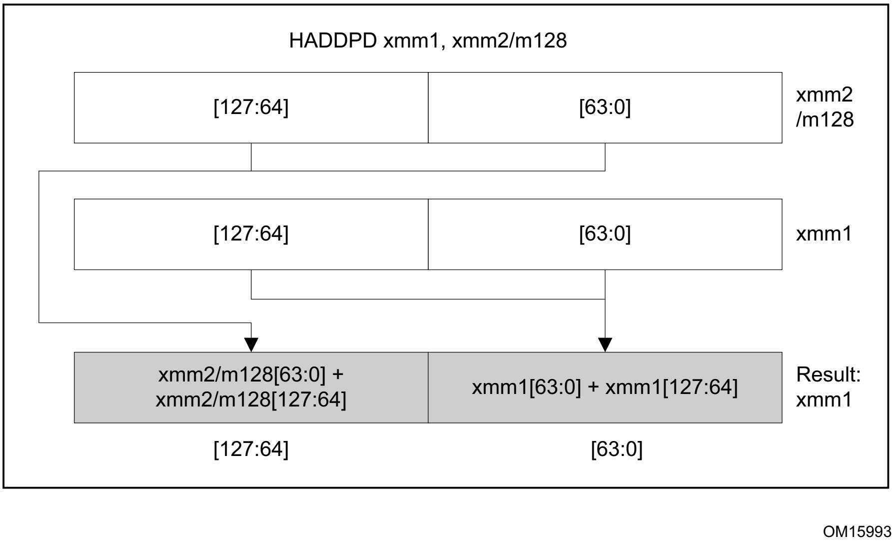


Figure 3-19. HADDPD—Packed Double-FP Horizontal Add

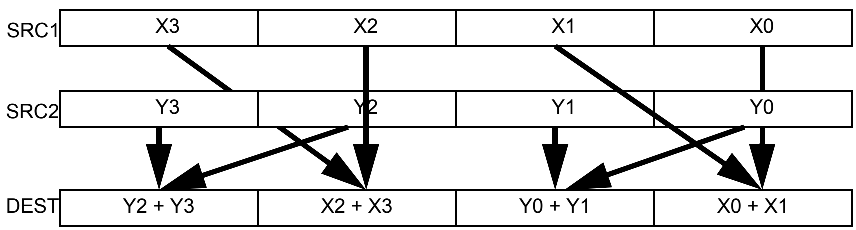


Figure 3-20. VHADDPD operation

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

HADDPD (128-bit Legacy SSE version)

DEST[63:0] ← SRC1[127:64] + SRC1[63:0]
DEST[127:64] ← SRC2[127:64] + SRC2[63:0]
DEST[VLMAX-1:128] (Unmodified)

VHADDPD (VEX.128 encoded version)

DEST[63:0] ← SRC1[127:64] + SRC1[63:0]
DEST[127:64] ← SRC2[127:64] + SRC2[63:0]
DEST[VLMAX-1:128] ← 0

VHADDPD (VEX.256 encoded version)

DEST[63:0] ← SRC1[127:64] + SRC1[63:0]
DEST[127:64] ← SRC2[127:64] + SRC2[63:0]
DEST[191:128] ← SRC1[255:192] + SRC1[191:128]
DEST[255:192] ← SRC2[255:192] + SRC2[191:128]

Intel C/C++ Compiler Intrinsic Equivalent

VHADDPD __m256d _mm256_hadd_pd (__m256d a, __m256d b);

HADDPD __m128d _mm_hadd_pd (__m128d a, __m128d b);

Exceptions

When the source operand is a memory operand, the operand must be aligned on a 16-byte boundary or a general-protection exception (#GP) will be generated.

Numeric Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 2.

HADDPS—Packed Single-FP Horizontal Add

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 7C /*r* HADDPS *xmm1, xmm2/m128*	A	V/V	SSE3	Horizontal add packed single-precision floating-point values from *xmm2/m128* to *xmm1*.
VEX.NDS.128.F2.0F.WIG 7C /r VHADDPS xmm1, xmm2, xmm3/m128	B	V/V	AVX	Horizontal add packed single-precision floating-point values from xmm2 and xmm3/mem.
VEX.NDS.256.F2.0F.WIG 7C /r VHADDPS ymm1, ymm2, ymm3/m256	B	V/V	AVX	Horizontal add packed single-precision floating-point values from ymm2 and ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Adds the single-precision floating-point values in the first and second dwords of the destination operand and stores the result in the first dword of the destination operand.

Adds single-precision floating-point values in the third and fourth dword of the destination operand and stores the result in the second dword of the destination operand.

Adds single-precision floating-point values in the first and second dword of the source operand and stores the result in the third dword of the destination operand.

Adds single-precision floating-point values in the third and fourth dword of the source operand and stores the result in the fourth dword of the destination operand.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

See Figure 3-21 for HADDPS; see Figure 3-22 for VHADDPS.

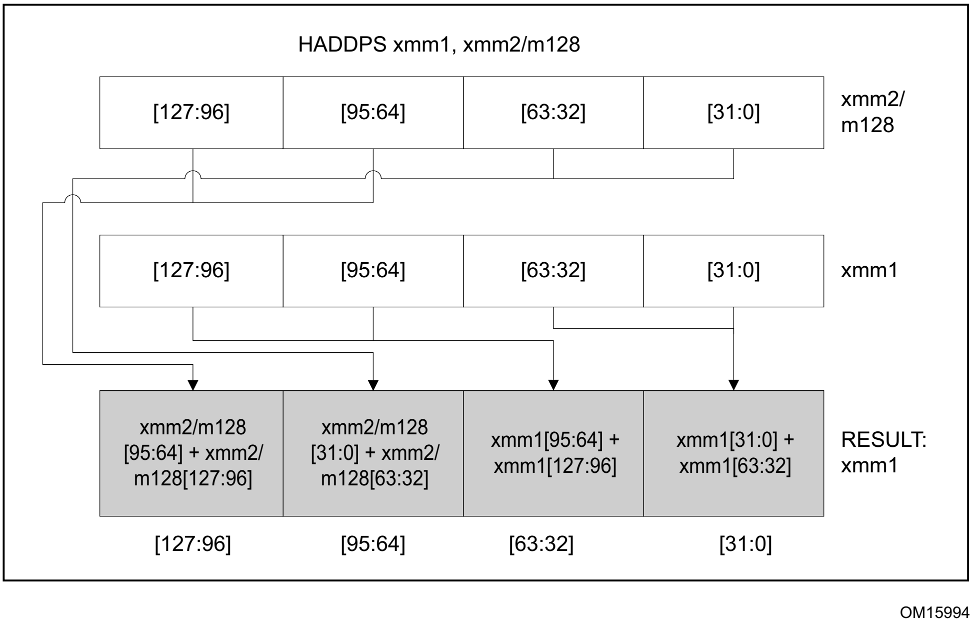


Figure 3-21. HADDPS—Packed Single-FP Horizontal Add

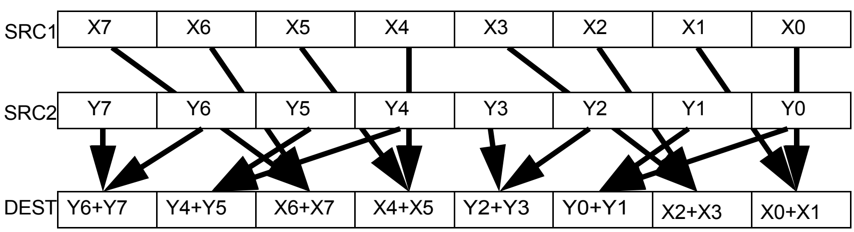


Figure 3-22. VHADDPS operation

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

HADDPS (128-bit Legacy SSE version)

DEST[31:0] ← SRC1[63:32] + SRC1[31:0]
DEST[63:32] ← SRC1[127:96] + SRC1[95:64]
DEST[95:64] ← SRC2[63:32] + SRC2[31:0]
DEST[127:96] ← SRC2[127:96] + SRC2[95:64]
DEST[VLMAX-1:128] (Unmodified)

VHADDPS (VEX.128 encoded version)

DEST[31:0] ← SRC1[63:32] + SRC1[31:0]
DEST[63:32] ← SRC1[127:96] + SRC1[95:64]
DEST[95:64] ← SRC2[63:32] + SRC2[31:0]
DEST[127:96] ← SRC2[127:96] + SRC2[95:64]
DEST[VLMAX-1:128] ← 0

VHADDPS (VEX.256 encoded version)

DEST[31:0] ← SRC1[63:32] + SRC1[31:0]
DEST[63:32] ← SRC1[127:96] + SRC1[95:64]
DEST[95:64] ← SRC2[63:32] + SRC2[31:0]
DEST[127:96] ← SRC2[127:96] + SRC2[95:64]
DEST[159:128] ← SRC1[191:160] + SRC1[159:128]
DEST[191:160] ← SRC1[255:224] + SRC1[223:192]
DEST[223:192] ← SRC2[191:160] + SRC2[159:128]
DEST[255:224] ← SRC2[255:224] + SRC2[223:192]

Intel C/C++ Compiler Intrinsic Equivalent

HADDPS __m128 _mm_hadd_ps (__m128 a, __m128 b);

VHADDPS __m256 _mm256_hadd_ps (__m256 a, __m256 b);

Exceptions

When the source operand is a memory operand, the operand must be aligned on a 16-byte boundary or a general-protection exception (#GP) will be generated.

Numeric Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 2.

HLT—Halt

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
F4	HLT	A	Valid	Valid	Halt

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Stops instruction execution and places the processor in a HALT state. An enabled interrupt (including NMI and SMI), a debug exception, the BINIT# signal, the INIT# signal, or the RESET# signal will resume execution. If an interrupt (including NMI) is used to resume execution after a HLT instruction, the saved instruction pointer (CS:EIP) points to the instruction following the HLT instruction.

When a HLT instruction is executed on an Intel 64 or IA-32 processor supporting Intel Hyper-Threading Technology, only the logical processor that executes the instruction is halted. The other logical processors in the physical processor remain active, unless they are each individually halted by executing a HLT instruction.

The HLT instruction is a privileged instruction. When the processor is running in protected or virtual-8086 mode, the privilege level of a program or procedure must be 0 to execute the HLT instruction.

This instruction’s operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
Enter Halt state;
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If the current privilege level is not 0.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

None.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

HSUBPD—Packed Double-FP Horizontal Subtract

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 7D /*r* HSUBPD *xmm1, xmm2/m128*	A	V/V	SSE3	Horizontal subtract packed double-precision floating-point values from *xmm2/m128* to *xmm1*.
VEX.NDS.128.66.0F.WIG 7D /r VHSUBPD xmm1,xmm2, xmm3/m128	B	V/V	AVX	Horizontal subtract packed double-precision floating-point values from xmm2 and xmm3/mem.
VEX.NDS.256.66.0F.WIG 7D /r VHSUBPD ymm1, ymm2, ymm3/m256	B	V/V	AVX	Horizontal subtract packed double-precision floating-point values from ymm2 and ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

The HSUBPD instruction subtracts horizontally the packed DP FP numbers of both operands.

Subtracts the double-precision floating-point value in the high quadword of the destination operand from the low quadword of the destination operand and stores the result in the low quadword of the destination operand.

Subtracts the double-precision floating-point value in the high quadword of the source operand from the low quadword of the source operand and stores the result in the high quadword of the destination operand.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

See Figure 3-23 for HSUBPD; see Figure 3-24 for VHSUBPD.

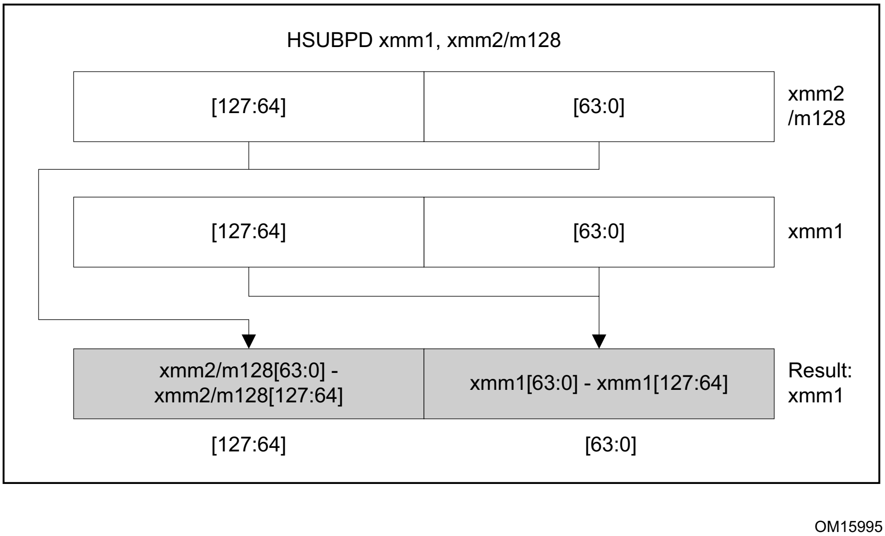


Figure 3-23. HSUBPD—Packed Double-FP Horizontal Subtract

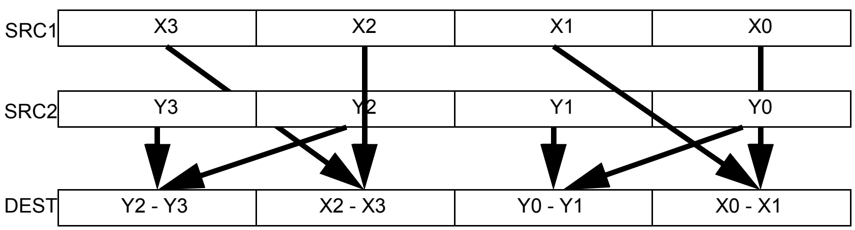


Figure 3-24. VHSUBPD operation

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

HSUBPD (128-bit Legacy SSE version)

```
DEST[63:0] ← SRC1[63:0] - SRC1[127:64]
DEST[127:64] ← SRC2[63:0] - SRC2[127:64]
DEST[VLMAX-1:128] (Unmodified)
```

VHSUBPD (VEX.128 encoded version)

```
DEST[63:0] ← SRC1[63:0] - SRC1[127:64]
DEST[127:64] ← SRC2[63:0] - SRC2[127:64]
DEST[VLMAX-1:128] ← 0
```

VHSUBPD (VEX.256 encoded version)

```
DEST[63:0] ← SRC1[63:0] - SRC1[127:64]
DEST[127:64] ← SRC2[63:0] - SRC2[127:64]
DEST[191:128] ← SRC1[191:128] - SRC1[255:192]
DEST[255:192] ← SRC2[191:128] - SRC2[255:192]
```

Intel C/C++ Compiler Intrinsic Equivalent

```
HSUBPD __m128d _mm_hsub_pd(__m128d a, __m128d b)

VHSUBPD __m256d _mm256_hsub_pd (__m256d a, __m256d b);
```

Exceptions

When the source operand is a memory operand, the operand must be aligned on a 16-byte boundary or a general-protection exception (#GP) will be generated.

Numeric Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 2.

HSUBPS—Packed Single-FP Horizontal Subtract

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 7D /*r* HSUBPS *xmm1, xmm2/m128*	A	V/V	SSE3	Horizontal subtract packed single-precision floating-point values from *xmm2/m128* to *xmm1*.
VEX.NDS.128.F2.0F.WIG 7D /r VHSUBPS xmm1, xmm2, xmm3/m128	B	V/V	AVX	Horizontal subtract packed single-precision floating-point values from xmm2 and xmm3/mem.
VEX.NDS.256.F2.0F.WIG 7D /r VHSUBPS ymm1, ymm2, ymm3/m256	B	V/V	AVX	Horizontal subtract packed single-precision floating-point values from ymm2 and ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Subtracts the single-precision floating-point value in the second dword of the destination operand from the first dword of the destination operand and stores the result in the first dword of the destination operand.

Subtracts the single-precision floating-point value in the fourth dword of the destination operand from the third dword of the destination operand and stores the result in the second dword of the destination operand.

Subtracts the single-precision floating-point value in the second dword of the source operand from the first dword of the source operand and stores the result in the third dword of the destination operand.

Subtracts the single-precision floating-point value in the fourth dword of the source operand from the third dword of the source operand and stores the result in the fourth dword of the destination operand.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

See Figure 3-25 for HSUBPS; see Figure 3-26 for VHSUBPS.

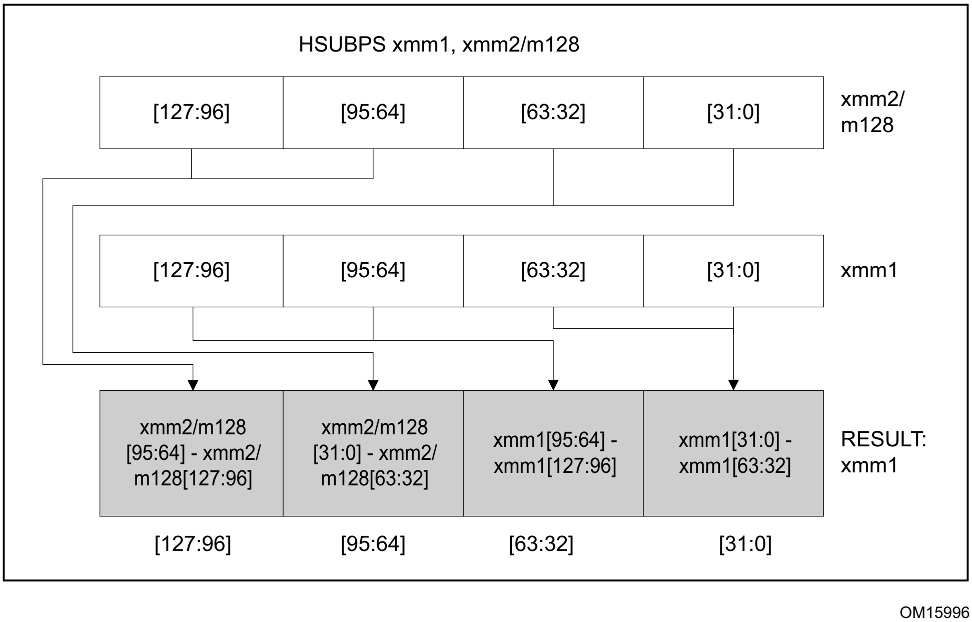


Figure 3-25. HSUBPS—Packed Single-FP Horizontal Subtract

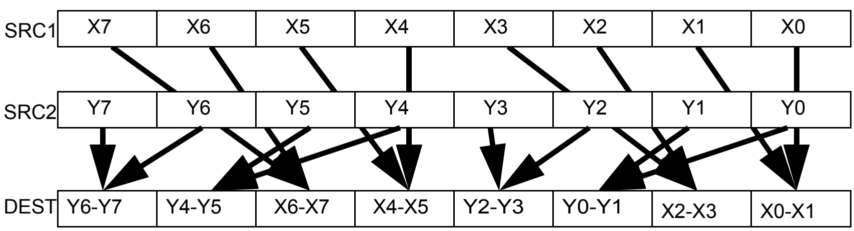


Figure 3-26. VHSUBPS operation

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

HSUBPS (128-bit Legacy SSE version)

```
DEST[31:0] ← SRC1[31:0] - SRC1[63:32]
DEST[63:32] ← SRC1[95:64] - SRC1[127:96]
DEST[95:64] ← SRC2[31:0] - SRC2[63:32]
DEST[127:96] ← SRC2[95:64] - SRC2[127:96]
DEST[VLMAX-1:128] (Unmodified)
```

VHSUBPS (VEX.128 encoded version)

```
DEST[31:0] ← SRC1[31:0] - SRC1[63:32]
DEST[63:32] ← SRC1[95:64] - SRC1[127:96]
DEST[95:64] ← SRC2[31:0] - SRC2[63:32]
DEST[127:96] ← SRC2[95:64] - SRC2[127:96]
DEST[VLMAX-1:128] ← 0
```

VHSUBPS (VEX.256 encoded version)

```
DEST[31:0] ← SRC1[31:0] - SRC1[63:32]
DEST[63:32] ← SRC1[95:64] - SRC1[127:96]
DEST[95:64] ← SRC2[31:0] - SRC2[63:32]
DEST[127:96] ← SRC2[95:64] - SRC2[127:96]
DEST[159:128] ← SRC1[159:128] - SRC1[191:160]
DEST[191:160] ← SRC1[223:192] - SRC1[255:224]
DEST[223:192] ← SRC2[159:128] - SRC2[191:160]
DEST[255:224] ← SRC2[223:192] - SRC2[255:224]
```

Intel C/C++ Compiler Intrinsic Equivalent

HSUBPS __m128 _mm_hsub_ps(__m128 a, __m128 b);

VHSUBPS __m256 _mm256_hsub_ps (__m256 a, __m256 b);

Exceptions

When the source operand is a memory operand, the operand must be aligned on a 16-byte boundary or a general-protection exception (#GP) will be generated.

Numeric Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 2.

IDIV—Signed Divide

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
F6 /7	IDIV *r/m8*	A	Valid	Valid	Signed divide AX by *r/m8*, with result stored in: AL ← Quotient, AH ← Remainder.
REX + F6 /7	IDIV *r/m8**	A	Valid	N.E.	Signed divide AX by *r/m8*, with result stored in AL ← Quotient, AH ← Remainder.
F7 /7	IDIV *r/m16*	A	Valid	Valid	Signed divide DX:AX by *r/m16*, with result stored in AX ← Quotient, DX ← Remainder.
F7 /7	IDIV *r/m32*	A	Valid	Valid	Signed divide EDX:EAX by *r/m32*, with result stored in EAX ← Quotient, EDX ← Remainder.
REX.W + F7 /7	IDIV *r/m64*	A	Valid	N.E.	Signed divide RDX:RAX by *r/m64*, with result stored in RAX ← Quotient, RDX ← Remainder.

NOTES:

* In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r)	NA	NA	NA

Description

Divides the (signed) value in the AX, DX:AX, or EDX:EAX (dividend) by the source operand (divisor) and stores the result in the AX (AH:AL), DX:AX, or EDX:EAX registers. The source operand can be a general-purpose register or a memory location. The action of this instruction depends on the operand size (dividend/divisor).

Non-integral results are truncated (chopped) towards 0. The remainder is always less than the divisor in magnitude. Overflow is indicated with the #DE (divide error) exception rather than with the CF flag.

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits. In 64-bit mode when REX.W is applied, the instruction divides the signed value in RDX:RAX by the source operand. RAX contains a 64-bit quotient; RDX contains a 64-bit remainder.

See the summary chart at the beginning of this section for encoding data and limits. See Table 3-60.

Table 3-60. IDIV Results

Operand Size	Dividend	Divisor	Quotient	Remainder	Quotient Range
Word/byte	AX	r/m8	AL	AH	−128 to +127
Doubleword/word	DX:AX	r/m16	AX	DX	−32,768 to +32,767
Quadword/doubleword	EDX:EAX	r/m32	EAX	EDX	-2^{31} to $2^{32} - 1$
Doublequadword/ quadword	RDX:RAX	r/m64	RAX	RDX	-2^{63} to $2^{64} - 1$

Operation

```
IF SRC = 0
    THEN #DE; (* Divide error *)
FI;

IF OperandSize = 8 (* Word/byte operation *)
    THEN
        temp ← AX / SRC; (* Signed division *)
        IF (temp > 7FH) or (temp < 80H)
        (* If a positive result is greater than 7FH or a negative result is less than 80H *)
            THEN #DE; (* Divide error *)
            ELSE
                AL ← temp;
                AH ← AX SignedModulus SRC;
        FI;
    ELSE IF OperandSize = 16 (* Doubleword/word operation *)
        THEN
            temp ← DX:AX / SRC; (* Signed division *)
            IF (temp > 7FFFH) or (temp < 8000H)
            (* If a positive result is greater than 7FFFH
            or a negative result is less than 8000H *)
                THEN
                    #DE; (* Divide error *)
                ELSE
                    AX ← temp;
                    DX ← DX:AX SignedModulus SRC;
            FI;
        FI;
```

```
    ELSE IF OperandSize = 32 (* Quadword/doubleword operation *)
            temp ← EDX:EAX / SRC; (* Signed division *)
            IF (temp > 7FFFFFFFH) or (temp < 80000000H)
            (* If a positive result is greater than 7FFFFFFFH
            or a negative result is less than 80000000H *)
                THEN
                    #DE; (* Divide error *)
                ELSE
                    EAX ← temp;
                    EDX ← EDXE:AX SignedModulus SRC;
            FI;
        FI;
    ELSE IF OperandSize = 64 (* Doublequadword/quadword operation *)
            temp ← RDX:RAX / SRC; (* Signed division *)
            IF (temp > 7FFFFFFFFFFFH) or (temp < 8000000000000000H)
            (* If a positive result is greater than 7FFFFFFFFFFFH
            or a negative result is less than 8000000000000000H *)
                THEN
                    #DE; (* Divide error *)
                ELSE
                    RAX ← temp;
                    RDX ← RDE:RAX SignedModulus SRC;
            FI;
        FI;
FI;
```

Flags Affected

The CF, OF, SF, ZF, AF, and PF flags are undefined.

Protected Mode Exceptions

#DE	If the source operand (divisor) is 0.
	The signed result (quotient) is too large for the destination.
#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#DE	If the source operand (divisor) is 0.
	The signed result (quotient) is too large for the destination.
#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#DE	If the source operand (divisor) is 0.
	The signed result (quotient) is too large for the destination.
#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#DE	If the source operand (divisor) is 0
	If the quotient is too large for the designated register.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

IMUL—Signed Multiply

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
F6 /5	IMUL *r/m8**	A	Valid	Valid	AX← AL ∗ *r/m* byte.
F7 /5	IMUL *r/m16*	A	Valid	Valid	DX:AX ← AX ∗ *r/m* word.
F7 /5	IMUL *r/m32*	A	Valid	Valid	EDX:EAX ← EAX ∗ *r/m*32.
REX.W + F7 /5	IMUL *r/m64*	A	Valid	N.E.	RDX:RAX ← RAX ∗ *r/m*64.
0F AF /*r*	IMUL *r16, r/m16*	B	Valid	Valid	word register ← word register ∗ *r/m*16.
0F AF /*r*	IMUL *r32, r/m32*	B	Valid	Valid	doubleword register ← doubleword register ∗ *r/m32.*
REX.W + 0F AF /*r*	IMUL *r64, r/m64*	B	Valid	N.E.	Quadword register ← Quadword register ∗ *r/m64*.
6B /*r ib*	IMUL *r16, r/m16, imm8*	C	Valid	Valid	word register ← *r/m16* ∗ sign-extended immediate byte.
6B /*r ib*	IMUL *r32, r/m32, imm8*	C	Valid	Valid	doubleword register ← *r/m32* ∗ sign-extended immediate byte.
REX.W + 6B /*r ib*	IMUL *r64, r/m64, imm8*	C	Valid	N.E.	Quadword register ← *r/m64* ∗ sign-extended immediate byte.
69 /*r iw*	IMUL *r16, r/m16, imm16*	C	Valid	Valid	word register ← *r/m16* ∗ immediate word.
69 /*r id*	IMUL *r32, r/m32, imm32*	C	Valid	Valid	doubleword register ← *r/m32* ∗ immediate doubleword.
REX.W + 69 /*r id*	IMUL *r64, r/m64, imm32*	C	Valid	N.E.	Quadword register ← *r/m64* ∗ immediate doubleword.

NOTES:

* In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r, w)	NA	NA	NA
B	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
C	ModRM:reg (r, w)	ModRM:r/m (r)	imm8/16/32	NA

Description

Performs a signed multiplication of two operands. This instruction has three forms, depending on the number of operands.

- **One-operand form** — This form is identical to that used by the MUL instruction. Here, the source operand (in a general-purpose register or memory location) is multiplied by the value in the AL, AX, EAX, or RAX register (depending on the operand size) and the product is stored in the AX, DX:AX, EDX:EAX, or RDX:RAX registers, respectively.
- **Two-operand form** — With this form the destination operand (the first operand) is multiplied by the source operand (second operand). The destination operand is a general-purpose register and the source operand is an immediate value, a general-purpose register, or a memory location. The product is then stored in the destination operand location.
- **Three-operand form** — This form requires a destination operand (the first operand) and two source operands (the second and the third operands). Here, the first source operand (which can be a general-purpose register or a memory location) is multiplied by the second source operand (an immediate value). The product is then stored in the destination operand (a general-purpose register).

When an immediate value is used as an operand, it is sign-extended to the length of the destination operand format.

The CF and OF flags are set when significant bit (including the sign bit) are carried into the upper half of the result. The CF and OF flags are cleared when the result (including the sign bit) fits exactly in the lower half of the result.

The three forms of the IMUL instruction are similar in that the length of the product is calculated to twice the length of the operands. With the one-operand form, the product is stored exactly in the destination. With the two- and three- operand forms, however, the result is truncated to the length of the destination before it is stored in the destination register. Because of this truncation, the CF or OF flag should be tested to ensure that no significant bits are lost.

The two- and three-operand forms may also be used with unsigned operands because the lower half of the product is the same regardless if the operands are signed or unsigned. The CF and OF flags, however, cannot be used to determine if the upper half of the result is non-zero.

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits. Use of REX.W modifies the three forms of the instruction as follows.

- One-operand form —The source operand (in a 64-bit general-purpose register or memory location) is multiplied by the value in the RAX register and the product is stored in the RDX:RAX registers.
- **Two-operand form** — The source operand is promoted to 64 bits if it is a register or a memory location. If the source operand is an immediate, it is sign extended to 64 bits. The destination operand is promoted to 64 bits.
- **Three-operand form** — The first source operand (either a register or a memory location) and destination operand are promoted to 64 bits.

Operation

```
IF (NumberOfOperands = 1)
   THEN IF (OperandSize = 8)
       THEN
           AX ← AL ∗ SRC (* Signed multiplication *)
           IF AL = AX
               THEN CF ← 0; OF ← 0;
               ELSE CF ← 1; OF ← 1; FI;
       ELSE IF OperandSize = 16
           THEN
               DX:AX ← AX ∗ SRC (* Signed multiplication *)
               IF sign_extend_to_32 (AX) = DX:AX
                   THEN CF ← 0; OF ← 0;
                   ELSE CF ← 1; OF ← 1; FI;
           ELSE IF OperandSize = 32
               THEN
                   EDX:EAX ← EAX ∗ SRC (* Signed multiplication *)
                   IF EAX = EDX:EAX
                       THEN CF ← 0; OF ← 0;
                       ELSE CF ← 1; OF ← 1; FI;
               ELSE (* OperandSize = 64 *)
                   RDX:RAX ← RAX ∗ SRC (* Signed multiplication *)
                   IF RAX = RDX:RAX
                       THEN CF ← 0; OF ← 0;
                       ELSE CF ← 1; OF ← 1; FI;
               FI;
       FI;
   ELSE IF (NumberOfOperands = 2)
       THEN
           temp ← DEST ∗ SRC (* Signed multiplication; temp is double DEST size *)
           DEST ← DEST ∗ SRC (* Signed multiplication *)
           IF temp ≠ DEST
               THEN CF ← 1; OF ← 1;
               ELSE CF ← 0; OF ← 0; FI;
```

```
        ELSE (* NumberOfOperands = 3 *)
            DEST ← SRC1 ∗ SRC2 (* Signed multiplication *)
            temp ← SRC1 ∗ SRC2 (* Signed multiplication; temp is double SRC1 size *)
            IF temp ≠ DEST
                THEN CF ← 1; OF ← 1;
                ELSE CF ← 0; OF ← 0; FI;
    FI;
FI;
```

Flags Affected

For the one operand form of the instruction, the CF and OF flags are set when significant bits are carried into the upper half of the result and cleared when the result fits exactly in the lower half of the result. For the two- and three-operand forms of the instruction, the CF and OF flags are set when the result must be truncated to fit in the destination operand size and cleared when the result fits exactly in the destination operand size. The SF, ZF, AF, and PF flags are undefined.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.

#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

IN—Input from Port

Opcode	Instruction	Op/En	64-Bit Mode	Compat/Leg Mode	Description
E4 *ib*	IN AL, *imm8*	A	Valid	Valid	Input byte from *imm8* I/O port address into AL.
E5 *ib*	IN AX, *imm8*	A	Valid	Valid	Input word from *imm8* I/O port address into AX.
E5 *ib*	IN EAX, *imm8*	A	Valid	Valid	Input dword from *imm8* I/O port address into EAX.
EC	IN AL,DX	B	Valid	Valid	Input byte from I/O port in DX into AL.
ED	IN AX,DX	B	Valid	Valid	Input word from I/O port in DX into AX.
ED	IN EAX,DX	B	Valid	Valid	Input doubleword from I/O port in DX into EAX.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	imm8	NA	NA	NA
B	NA	NA	NA	NA

Description

Copies the value from the I/O port specified with the second operand (source operand) to the destination operand (first operand). The source operand can be a byte-immediate or the DX register; the destination operand can be register AL, AX, or EAX, depending on the size of the port being accessed (8, 16, or 32 bits, respectively). Using the DX register as a source operand allows I/O port addresses from 0 to 65,535 to be accessed; using a byte immediate allows I/O port addresses 0 to 255 to be accessed.

When accessing an 8-bit I/O port, the opcode determines the port size; when accessing a 16- and 32-bit I/O port, the operand-size attribute determines the port size. At the machine code level, I/O instructions are shorter when accessing 8-bit I/O ports. Here, the upper eight bits of the port address will be 0.

This instruction is only useful for accessing I/O ports located in the processor's I/O address space. See Chapter 13, "Input/Output," in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for more information on accessing I/O ports in the I/O address space.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
IF ((PE = 1) and ((CPL > IOPL) or (VM = 1)))
    THEN (* Protected mode with CPL > IOPL or virtual-8086 mode *)
        IF (Any I/O Permission Bit for I/O port being accessed = 1)
            THEN (* I/O operation is not allowed *)
                #GP(0);
            ELSE ( * I/O operation is allowed *)
                DEST ← SRC; (* Read from selected I/O port *)
        FI;
    ELSE (Real Mode or Protected Mode with CPL ≤ IOPL *)
        DEST ← SRC; (* Read from selected I/O port *)
FI;
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If the CPL is greater than (has less privilege) the I/O privilege level (IOPL) and any of the corresponding I/O permission bits in TSS for the I/O port being accessed is 1.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If any of the I/O permission bits in the TSS for the I/O port being accessed is 1.
#PF(fault-code)	If a page fault occurs.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(0)	If the CPL is greater than (has less privilege) the I/O privilege level (IOPL) and any of the corresponding I/O permission bits in TSS for the I/O port being accessed is 1.
#UD	If the LOCK prefix is used.

INC—Increment by 1

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
FE /0	INC *r/m8*	A	Valid	Valid	Increment *r/m* byte by 1.
REX + FE /0	INC *r/m8**	A	Valid	N.E.	Increment *r/m* byte by 1.
FF /0	INC *r/m16*	A	Valid	Valid	Increment *r/m* word by 1.
FF /0	INC *r/m32*	A	Valid	Valid	Increment *r/m* doubleword by 1.
REX.W + FF /0	INC *r/m64*	A	Valid	N.E.	Increment *r/m* quadword by 1.
40+ *rw***	INC *r16*	B	N.E.	Valid	Increment word register by 1.
40+ *rd*	INC *r32*	B	N.E.	Valid	Increment doubleword register by 1.

NOTES:

* In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

** 40H through 47H are REX prefixes in 64-bit mode.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r, w)	NA	NA	NA
B	reg (r, w)	NA	NA	NA

Description

Adds 1 to the destination operand, while preserving the state of the CF flag. The destination operand can be a register or a memory location. This instruction allows a loop counter to be updated without disturbing the CF flag. (Use a ADD instruction with an immediate operand of 1 to perform an increment operation that does updates the CF flag.)

This instruction can be used with a LOCK prefix to allow the instruction to be executed atomically.

In 64-bit mode, INC r16 and INC r32 are not encodable (because opcodes 40H through 47H are REX prefixes). Otherwise, the instruction’s 64-bit mode default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits.

Operation

DEST ← DEST + 1;

AFlags Affected

The CF flag is not affected. The OF, SF, ZF, AF, and PF flags are set according to the result.

Protected Mode Exceptions

#GP(0)	If the destination operand is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULLsegment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used but the destination is not a memory operand.

INS/INSB/INSW/INSD—Input from Port to String

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
6C	INS *m8*, DX	A	Valid	Valid	Input byte from I/O port specified in DX into memory location specified in ES:(E)DI or RDI.*
6D	INS *m16*, DX	A	Valid	Valid	Input word from I/O port specified in DX into memory location specified in ES:(E)DI or RDI.[1]
6D	INS *m32*, DX	A	Valid	Valid	Input doubleword from I/O port specified in DX into memory location specified in ES:(E)DI or RDI.[1]
6C	INSB	A	Valid	Valid	Input byte from I/O port specified in DX into memory location specified with ES:(E)DI or RDI.[1]
6D	INSW	A	Valid	Valid	Input word from I/O port specified in DX into memory location specified in ES:(E)DI or RDI.[1]
6D	INSD	A	Valid	Valid	Input doubleword from I/O port specified in DX into memory location specified in ES:(E)DI or RDI.[1]

NOTES:

* In 64-bit mode, only 64-bit (RDI) and 32-bit (EDI) address sizes are supported. In non-64-bit mode, only 32-bit (EDI) and 16-bit (DI) address sizes are supported.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Copies the data from the I/O port specified with the source operand (second operand) to the destination operand (first operand). The source operand is an I/O port address (from 0 to 65,535) that is read from the DX register. The destination operand is a memory location, the address of which is read from either the ES:DI, ES:EDI or the RDI registers (depending on the address-size attribute of the instruction, 16, 32 or 64, respectively). (The ES segment cannot be overridden with a segment override prefix.) The size of the I/O port being accessed (that is, the size of the source and destination operands) is determined by the opcode for an 8-bit I/O port or by the operand-size attribute of the instruction for a 16- or 32-bit I/O port.

At the assembly-code level, two forms of this instruction are allowed: the “explicit-operands” form and the “no-operands” form. The explicit-operands form (specified with the INS mnemonic) allows the source and destination operands to be specified explicitly. Here, the source operand must be “DX,” and the destination operand should be a symbol that indicates the size of the I/O port and the destination address. This explicit-operands form is provided to allow documentation; however, note that the documentation provided by this form can be misleading. That is, the destination operand symbol must specify the correct **type** (size) of the operand (byte, word, or doubleword), but it does not have to specify the correct **location**. The location is always specified by the ES:(E)DI registers, which must be loaded correctly before the INS instruction is executed.

The no-operands form provides “short forms” of the byte, word, and doubleword versions of the INS instructions. Here also DX is assumed by the processor to be the source operand and ES:(E)DI is assumed to be the destination operand. The size of the I/O port is specified with the choice of mnemonic: INSB (byte), INSW (word), or INSD (doubleword).

After the byte, word, or doubleword is transfer from the I/O port to the memory location, the DI/EDI/RDI register is incremented or decremented automatically according to the setting of the DF flag in the EFLAGS register. (If the DF flag is 0, the (E)DI register is incremented; if the DF flag is 1, the (E)DI register is decremented.) The (E)DI register is incremented or decremented by 1 for byte operations, by 2 for word operations, or by 4 for doubleword operations.

The INS, INSB, INSW, and INSD instructions can be preceded by the REP prefix for block input of ECX bytes, words, or doublewords. See “REP/REPE/REPZ /REPNE/REPNZ—Repeat String Operation Prefix” in Chapter 4 of the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 2B*, for a description of the REP prefix.

These instructions are only useful for accessing I/O ports located in the processor’s I/O address space. See Chapter 13, “Input/Output,” in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*, for more information on accessing I/O ports in the I/O address space.

In 64-bit mode, default address size is 64 bits, 32 bit address size is supported using the prefix 67H. The address of the memory destination is specified by RDI or EDI. 16-bit address size is not supported in 64-bit mode. The operand size is not promoted.

Operation

```
IF ((PE = 1) and ((CPL > IOPL) or (VM = 1)))
    THEN (* Protected mode with CPL > IOPL or virtual-8086 mode *)
```

```
        IF (Any I/O Permission Bit for I/O port being accessed = 1)
            THEN (* I/O operation is not allowed *)
                #GP(0);
            ELSE (* I/O operation is allowed *)
                DEST ← SRC; (* Read from I/O port *)
        FI;
    ELSE (Real Mode or Protected Mode with CPL IOPL *)
        DEST ← SRC; (* Read from I/O port *)
FI;
```

Non-64-bit Mode:

```
IF (Byte transfer)
    THEN IF DF = 0
        THEN (E)DI ← (E)DI + 1;
        ELSE (E)DI ← (E)DI – 1; FI;
    ELSE IF (Word transfer)
        THEN IF DF = 0
            THEN (E)DI ← (E)DI + 2;
            ELSE (E)DI ← (E)DI – 2; FI;
        ELSE (* Doubleword transfer *)
            THEN IF DF = 0
                THEN (E)DI ← (E)DI + 4;
                ELSE (E)DI ← (E)DI – 4; FI;
        FI;
FI;
```

FI64-bit Mode:

```
IF (Byte transfer)
    THEN IF DF = 0
        THEN (E|R)DI ← (E|R)DI + 1;
        ELSE (E|R)DI ← (E|R)DI – 1; FI;
    ELSE IF (Word transfer)
        THEN IF DF = 0
            THEN (E)DI ← (E)DI + 2;
            ELSE (E)DI ← (E)DI – 2; FI;
        ELSE (* Doubleword transfer *)
            THEN IF DF = 0
                THEN (E|R)DI ← (E|R)DI + 4;
                ELSE (E|R)DI ← (E|R)DI – 4; FI;
        FI;
FI;
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If the CPL is greater than (has less privilege) the I/O privilege level (IOPL) and any of the corresponding I/O permission bits in TSS for the I/O port being accessed is 1.
	If the destination is located in a non-writable segment.
	If an illegal memory operand effective address in the ES segments is given.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If any of the I/O permission bits in the TSS for the I/O port being accessed is 1.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the CPL is greater than (has less privilege) the I/O privilege level (IOPL) and any of the corresponding I/O permission bits in TSS for the I/O port being accessed is 1.
	If the memory address is in a non-canonical form.

#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

VINSERTF128 — Insert Packed Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
VEX.NDS.256.66.0F3A.W0 18 /r ib VINSERTF128 ymm1, ymm2, xmm3/m128, imm8	A	V/V	AVX	Insert a single precision floating-point value selected by *imm8* from *xmm2/m32* into xmm1 at the specified destination element specified by *imm8* and zero out destination elements in *xmm1* as indicated in *imm8*.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs an insertion of 128-bits of packed floating-point values from the second source operand (third operand) into an the destination operand (first operand) at an 128-bit offset from imm8[0]. The remaining portions of the destination are written by the corresponding fields of the first source operand (second operand). The second source operand can be either an XMM register or a 128-bit memory location.

The high 7 bits of the immediate are ignored.

Operation

```
TEMP[255:0] ← SRC1[255:0]
CASE (imm8[0]) OF
    0: TEMP[127:0] ← SRC2[127:0]
    1: TEMP[255:128] ← SRC2[127:0]
ESAC
DEST ←TEMP
```

Intel C/C++ Compiler Intrinsic Equivalent

INSERTF128 __m256 _mm256_insertf128_ps (__m256 a, __m128 b, int offset);

INSERTF128 __m256d _mm256_insertf128_pd (__m256d a, __m128d b, int offset);

INSERTF128 __m256i _mm256_insertf128_si256 (__m256i a, __m128i b, int offset);

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 6; additionally

#UD	If VEX.W = 1.

INSERTPS — Insert Packed Single Precision Floating-Point Value

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 3A 21 /r ib INSERTPS *xmm1, xmm2/m32, imm8*	A	V/V	SSE4_1	Insert a single precision floating-point value selected by *imm8* from *xmm2/m32* into xmm1 at the specified destination element specified by *imm8* and zero out destination elements in *xmm1* as indicated in *imm8*.
VEX.NDS.128.66.0F3A.WIG 21 /r ib VINSERTPS xmm1, xmm2, xmm3/m32, imm8	B	V/V	AVX	Insert a single precision floating point value selected by imm8 from xmm3/m32 and merge into xmm2 at the specified destination element specified by imm8 and zero out destination elements in xmm1 as indicated in imm8.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	imm8	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

(register source form)

Select a single precision floating-point element from second source as indicated by Count_S bits of the immediate operand and insert it into the first source at the location indicated by the Count_D bits of the immediate operand. Store in the destination and zero out destination elements based on the ZMask bits of the immediate operand.

(memory source form)

Load a floating-point element from a 32-bit memory location and insert it into the first source at the location indicated by the Count_D bits of the immediate operand. Store in the destination and zero out destination elements based on the ZMask bits of the immediate operand.

128-bit Legacy SSE version: The first source register is an XMM register. The second source operand is either an XMM register or a 32-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version. The destination and first source register is an XMM register. The second source operand is either an XMM register or a 32-bit memory location. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

If VINSERTPS is encoded with VEX.L= 1, an attempt to execute the instruction encoded with VEX.L= 1 will cause an #UD exception.

Operation

INSERTPS (128-bit Legacy SSE version)

```
IF (SRC == REG) THEN COUNT_S ← imm8[7:6]
    ELSE COUNT_S ← 0
COUNT_D ← imm8[5:4]
ZMASK ← imm8[3:0]
CASE (COUNT_S) OF
    0: TMP ← SRC[31:0]
    1: TMP ← SRC[63:32]
    2: TMP ← SRC[95:64]
    3: TMP ← SRC[127:96]
ESAC;

CASE (COUNT_D) OF
    0: TMP2[31:0] ← TMP
         TMP2[127:32] ← DEST[127:32]
    1: TMP2[63:32] ← TMP
         TMP2[31:0] ← DEST[31:0]
         TMP2[127:64] ← DEST[127:64]
    2: TMP2[95:64] ← TMP
         TMP2[63:0] ← DEST[63:0]
         TMP2[127:96] ← DEST[127:96]
    3: TMP2[127:96] ← TMP
         TMP2[95:0] ← DEST[95:0]
ESAC;

IF (ZMASK[0] == 1) THEN DEST[31:0] ← 00000000H
    ELSE DEST[31:0] ← TMP2[31:0]
IF (ZMASK[1] == 1) THEN DEST[63:32] ← 00000000H
    ELSE DEST[63:32] ← TMP2[63:32]
```

```
IF (ZMASK[2] == 1) THEN DEST[95:64] ← 00000000H
    ELSE DEST[95:64] ← TMP2[95:64]
IF (ZMASK[3] == 1) THEN DEST[127:96] ← 00000000H
    ELSE DEST[127:96] ← TMP2[127:96]
DEST[VLMAX-1:128] (Unmodified)
```

VINSERTPS (VEX.128 encoded version)

```
IF (SRC == REG) THEN COUNT_S ← imm8[7:6]
    ELSE COUNT_S ← 0
COUNT_D ← imm8[5:4]
ZMASK ← imm8[3:0]
CASE (COUNT_S) OF
    0: TMP ← SRC2[31:0]
    1: TMP ← SRC2[63:32]
    2: TMP ← SRC2[95:64]
    3: TMP ← SRC2[127:96]
ESAC;
CASE (COUNT_D) OF
    0: TMP2[31:0] ← TMP
        TMP2[127:32] ← SRC1[127:32]
    1: TMP2[63:32] ← TMP
        TMP2[31:0] ← SRC1[31:0]
        TMP2[127:64] ← SRC1[127:64]
    2: TMP2[95:64] ← TMP
        TMP2[63:0] ← SRC1[63:0]
        TMP2[127:96] ← SRC1[127:96]
    3: TMP2[127:96] ← TMP
        TMP2[95:0] ← SRC1[95:0]
ESAC;

IF (ZMASK[0] == 1) THEN DEST[31:0] ← 00000000H
    ELSE DEST[31:0] ← TMP2[31:0]
IF (ZMASK[1] == 1) THEN DEST[63:32] ← 00000000H
    ELSE DEST[63:32] ← TMP2[63:32]
IF (ZMASK[2] == 1) THEN DEST[95:64] ← 00000000H
    ELSE DEST[95:64] ← TMP2[95:64]
IF (ZMASK[3] == 1) THEN DEST[127:96] ← 00000000H
    ELSE DEST[127:96] ← TMP2[127:96]
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

INSERTPS __m128 _mm_insert_ps(__m128 dst, __m128 src, const int ndx);

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 5.

INT *n*/INTO/INT 3—Call to Interrupt Procedure

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
CC	INT 3	A	Valid	Valid	Interrupt 3—trap to debugger.
CD *ib*	INT *imm8*	B	Valid	Valid	Interrupt vector number specified by immediate byte.
CE	INTO	A	Invalid	Valid	Interrupt 4—if overflow flag is 1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA
B	imm8	NA	NA	NA

Description

The INT *n* instruction generates a call to the interrupt or exception handler specified with the destination operand (see the section titled "Interrupts and Exceptions" in Chapter 6 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*). The destination operand specifies an interrupt vector number from 0 to 255, encoded as an 8-bit unsigned intermediate value. Each interrupt vector number provides an index to a gate descriptor in the IDT. The first 32 interrupt vector numbers are reserved by Intel for system use. Some of these interrupts are used for internally generated exceptions.

The INT *n* instruction is the general mnemonic for executing a software-generated call to an interrupt handler. The INTO instruction is a special mnemonic for calling overflow exception (#OF), interrupt vector number 4. The overflow interrupt checks the OF flag in the EFLAGS register and calls the overflow interrupt handler if the OF flag is set to 1. (The INTO instruction cannot be used in 64-bit mode.)

The INT 3 instruction generates a special one byte opcode (CC) that is intended for calling the debug exception handler. (This one byte form is valuable because it can be used to replace the first byte of any instruction with a breakpoint, including other one byte instructions, without over-writing other code). To further support its function as a debug breakpoint, the interrupt generated with the CC opcode also differs from the regular software interrupts as follows:

- Interrupt redirection does not happen when in VME mode; the interrupt is handled by a protected-mode handler.
- The virtual-8086 mode IOPL checks do not occur. The interrupt is taken without faulting at any IOPL level.

Note that the “normal” 2-byte opcode for INT 3 (CD03) does not have these special features. Intel and Microsoft assemblers will not generate the CD03 opcode from any mnemonic, but this opcode can be created by direct numeric code definition or by self-modifying code.

The action of the INT *n* instruction (including the INTO and INT 3 instructions) is similar to that of a far call made with the CALL instruction. The primary difference is that with the INT *n* instruction, the EFLAGS register is pushed onto the stack before the return address. (The return address is a far address consisting of the current values of the CS and EIP registers.) Returns from interrupt procedures are handled with the IRET instruction, which pops the EFLAGS information and return address from the stack.

The interrupt vector number specifies an interrupt descriptor in the interrupt descriptor table (IDT); that is, it provides index into the IDT. The selected interrupt descriptor in turn contains a pointer to an interrupt or exception handler procedure. In protected mode, the IDT contains an array of 8-byte descriptors, each of which is an interrupt gate, trap gate, or task gate. In real-address mode, the IDT is an array of 4-byte far pointers (2-byte code segment selector and a 2-byte instruction pointer), each of which point directly to a procedure in the selected segment. (Note that in real-address mode, the IDT is called the **interrupt vector table**, and its pointers are called interrupt vectors.)

The following decision table indicates which action in the lower portion of the table is taken given the conditions in the upper portion of the table. Each Y in the lower section of the decision table represents a procedure defined in the “Operation” section for this instruction (except #GP).

Table 3-61. Decision Table

PE	0	1	1	1	1	1	1	1
VM	–	–	–	–	–	0	1	1
IOPL	–	–	–	–	–	–	<3	=3
DPL/CPL RELATIONSHIP	–	DPL< CPL	–	DPL> CPL	DPL= CPL or C	DPL< CPL & NC	–	–
INTERRUPT TYPE	–	S/W	–	–	–	–	–	–
GATE TYPE	–	–	Task	Trap or Interrupt	Trap or Interrupt	Trap or Interrupt	Trap or Interrupt	Trap or Interrupt
REAL-ADDRESS-MODE	Y							
PROTECTED-MODE		Y	Y	Y	Y	Y	Y	Y
TRAP-OR-INTERRUPT-GATE				Y	Y	Y	Y	Y
INTER-PRIVILEGE-LEVEL-INTERRUPT						Y		
INTRA-PRIVILEGE-LEVEL-INTERRUPT					Y			

Table 3-61. Decision Table (Contd.)

INTERRUPT-FROM-VIRTUAL-8086-MODE								Y
TASK-GATE			Y					
#GP		Y		Y			Y	

NOTES:

– Don't Care.

Y Yes, action taken.

Blank Action not taken.

When the processor is executing in virtual-8086 mode, the IOPL determines the action of the INT n instruction. If the IOPL is less than 3, the processor generates a #GP(selector) exception; if the IOPL is 3, the processor executes a protected mode interrupt to privilege level 0. The interrupt gate's DPL must be set to 3 and the target CPL of the interrupt handler procedure must be 0 to execute the protected mode interrupt to privilege level 0.

The interrupt descriptor table register (IDTR) specifies the base linear address and limit of the IDT. The initial base address value of the IDTR after the processor is powered up or reset is 0.

Operation

The following operational description applies not only to the INT *n* and INTO instructions, but also to external interrupts, nonmaskable interrupts (NMIs), and exceptions. Some of these events push onto the stack an error code.

The operational description specifies numerous checks whose failure may result in delivery of a nested exception. In these cases, the original event is not delivered.

The operational description specifies the error code delivered by any nested exception. In some cases, the error code is specified with a pseudofunction error_code(num,idt,ext), where idt and ext are bit values. The pseudofunction produces an error code as follows: (1) if idt is 0, the error code is (num & FCH) | ext; (2) if idt is 1, the error code is (num « 3) | 2 | ext.

In many cases, the pseudofunction error_code is invoked with a pseudovariable EXT. The value of EXT depends on the nature of the event whose delivery encountered a nested exception: if that event is a software interrupt, EXT is 0; otherwise, EXT is 1.

```
IF PE = 0
    THEN
        GOTO REAL-ADDRESS-MODE;
    ELSE (* PE = 1 *)
        IF (VM = 1 and IOPL < 3 AND INT n)
            THEN
                #GP(0); (* Bit 0 of error code is 0 because INT n *)
```

```
            ELSE (* Protected mode, IA-32e mode, or virtual-8086 mode interrupt *)
                IF (IA32_EFER.LMA = 0)
                    THEN (* Protected mode, or virtual-8086 mode interrupt *)
                        GOTO PROTECTED-MODE;
                ELSE (* IA-32e mode interrupt *)
                    GOTO IA-32e-MODE;
                FI;
        FI;
FI;
REAL-ADDRESS-MODE:
    IF ((vector_number « 2) + 3) is not within IDT limit
        THEN #GP; FI;
    IF stack not large enough for a 6-byte return information
        THEN #SS; FI;
    Push (EFLAGS[15:0]);
    IF ← 0; (* Clear interrupt flag *)
    TF ← 0; (* Clear trap flag *)
    AC ← 0; (* Clear AC flag *)
    Push(CS);
    Push(IP);
    (* No error codes are pushed in real-address mode*)
    CS ← IDT(Descriptor (vector_number « 2), selector));
    EIP ← IDT(Descriptor (vector_number « 2), offset)); (* 16 bit offset AND 0000FFFFH *)
END;
PROTECTED-MODE:
    IF ((vector_number « 3) + 7) is not within IDT limits
    or selected IDT descriptor is not an interrupt-, trap-, or task-gate type
        THEN #GP(error_code(vector_number,1,EXT)); FI;
        (* idt operand to error_code set because vector is used *)
    IF software interrupt (* Generated by INT n, INT3, or INTO *)
        THEN
            IF gate DPL < CPL (* PE = 1, DPL < CPL, software interrupt *)
                THEN #GP(error_code(vector_number,1,0)); FI;
                (* idt operand to error_code set because vector is used *)
                (* ext operand to error_code is 0 because INT n, INT3, or INTO*)
    FI;
    IF gate not present
        THEN #NP(error_code(vector_number,1,EXT)); FI;
        (* idt operand to error_code set because vector is used *)
    IF task gate (* Specified in the selected interrupt table descriptor *)
        THEN GOTO TASK-GATE;
        ELSE GOTO TRAP-OR-INTERRUPT-GATE; (* PE = 1, trap/interrupt gate *)
    FI;
```

```
END;
IA-32e-MODE:
    IF INTO and CS.L = 1 (64-bit mode)
        THEN #UD;
    FI;
    IF ((vector_number « 4) + 15) is not in IDT limits
    or selected IDT descriptor is not an interrupt-, or trap-gate type
        THEN #GP(error_code(vector_number,1,EXT));
        (* idt operand to error_code set because vector is used *)
    FI;
    IF software interrupt (* Generated by INT n, INT 3, or INTO *)
        THEN
            IF gate DPL < CPL (* PE = 1, DPL < CPL, software interrupt *)
                THEN #GP(error_code(vector_number,1,0));
                (* idt operand to error_code set because vector is used *)
                (* ext operand to error_code is 0 because INT n, INT3, or INTO*)
            FI;
    FI;
    IF gate not present
        THEN #NP(error_code(vector_number,1,EXT));
        (* idt operand to error_code set because vector is used *)
    FI;
    GOTO TRAP-OR-INTERRUPT-GATE; (* Trap/interrupt gate *)
END;
TASK-GATE: (* PE = 1, task gate *)
    Read TSS selector in task gate (IDT descriptor);
        IF local/global bit is set to local or index not within GDT limits
            THEN #GP(error_code(TSS selector,0,EXT)); FI;
            (* idt operand to error_code is 0 because selector is used *)
        Access TSS descriptor in GDT;
        IF TSS descriptor specifies that the TSS is busy (low-order 5 bits set to 00001)
            THEN #GP(TSS selector,0,EXT)); FI;
            (* idt operand to error_code is 0 because selector is used *)
        IF TSS not present
            THEN #NP(TSS selector,0,EXT)); FI;
            (* idt operand to error_code is 0 because selector is used *)
    SWITCH-TASKS (with nesting) to TSS;
    IF interrupt caused by fault with error code
        THEN
            IF stack limit does not allow push of error code
                THEN #SS(EXT); FI;
            Push(error code);
    FI;
```

```
    IF EIP not within code segment limit
        THEN #GP(EXT); FI;
END;
TRAP-OR-INTERRUPT-GATE:
    Read new code-segment selector for trap or interrupt gate (IDT descriptor);
    IF new code-segment selector is NULL
        THEN #GP(EXT); FI; (* Error code contains NULL selector *)
    IF new code-segment selector is not within its descriptor table limits
        THEN #GP(error_code(new code-segment selector,0,EXT)); FI;
        (* idt operand to error_code is 0 because selector is used *)
    Read descriptor referenced by new code-segment selector;
    IF descriptor does not indicate a code segment or new code-segment DPL > CPL
        THEN #GP(error_code(new code-segment selector,0,EXT)); FI;
        (* idt operand to error_code is 0 because selector is used *)
    IF new code-segment descriptor is not present,
        THEN #NP(error_code(new code-segment selector,0,EXT)); FI;
        (* idt operand to error_code is 0 because selector is used *)
    IF new code segment is non-conforming with DPL < CPL
        THEN
            IF VM = 0
                THEN
                    GOTO INTER-PRIVILEGE-LEVEL-INTERRUPT;
                    (* PE = 1, VM = 0, interrupt or trap gate, nonconforming code segment,
                    DPL < CPL *)
                ELSE (* VM = 1 *)
                    IF new code-segment DPL ≠ 0
                        THEN #GP(error_code(new code-segment selector,0,EXT));
                        (* idt operand to error_code is 0 because selector is used *)
                    GOTO INTERRUPT-FROM-VIRTUAL-8086-MODE; FI;
                    (* PE = 1, interrupt or trap gate, DPL < CPL, VM = 1 *)
            FI;
        ELSE (* PE = 1, interrupt or trap gate, DPL ≥ CPL *)
            IF VM = 1
                THEN #GP(error_code(new code-segment selector,0,EXT));
                (* idt operand to error_code is 0 because selector is used *)
            IF new code segment is conforming or new code-segment DPL = CPL
                THEN
                    GOTO INTRA-PRIVILEGE-LEVEL-INTERRUPT;
                ELSE (* PE = 1, interrupt or trap gate, nonconforming code segment, DPL > CPL *)
                    #GP(error_code(new code-segment selector,0,EXT));
                    (* idt operand to error_code is 0 because selector is used *)
            FI;
    FI;
```

```
END;
INTER-PRIVILEGE-LEVEL-INTERRUPT:
   (* PE = 1, interrupt or trap gate, non-conforming code segment, DPL < CPL *)
   IF (IA32_EFER.LMA = 0) (* Not IA-32e mode *)
       THEN
       (* Identify stack-segment selector for new privilege level in current TSS *)
           IF current TSS is 32-bit
               THEN
                   TSSstackAddress ← (new code-segment DPL « 3) + 4;
                   IF (TSSstackAddress + 5) > current TSS limit
                       THEN #TS(error_code(current TSS selector,0,EXT)); FI;
                       (* idt operand to error_code is 0 because selector is used *)
                   NewSS ← 2 bytes loaded from (TSS base + TSSstackAddress + 4);
                   NewESP ← 4 bytes loaded from (TSS base + TSSstackAddress);
               ELSE    (* current TSS is 16-bit *)
                   TSSstackAddress ← (new code-segment DPL « 2) + 2
                   IF (TSSstackAddress + 3) > current TSS limit
                       THEN #TS(error_code(current TSS selector,0,EXT)); FI;
                       (* idt operand to error_code is 0 because selector is used *)
                   NewSS ← 2 bytes loaded from (TSS base + TSSstackAddress + 2);
                   NewESP ← 2 bytes loaded from (TSS base + TSSstackAddress);
           FI;
           IF NewSS is NULL
               THEN #TS(EXT); FI;
           IF NewSS index is not within its descriptor-table limits
           or NewSS RPL ≠ new code-segment DPL
               THEN #TS(error_code(NewSS,0,EXT)); FI;
               (* idt operand to error_code is 0 because selector is used *)
           Read new stack-segment descriptor for NewSS in GDT or LDT;
           IF new stack-segment DPL ≠ new code-segment DPL
           or new stack-segment Type does not indicate writable data segment
               THEN #TS(error_code(NewSS,0,EXT)); FI;
               (* idt operand to error_code is 0 because selector is used *)
           IF NewSS is not present
               THEN #SS(error_code(NewSS,0,EXT)); FI;
               (* idt operand to error_code is 0 because selector is used *)
       ELSE (* IA-32e mode *)
           IF IDT-gate IST = 0
               THEN TSSstackAddress ← (new code-segment DPL « 3) + 4;
               ELSE TSSstackAddress ← (IDT gate IST « 3) + 28;
           FI;
           IF (TSSstackAddress + 7) > current TSS limit
               THEN #TS(error_code(current TSS selector,0,EXT); FI;
```

```
                (* idt operand to error_code is 0 because selector is used *)
            NewRSP ← 8 bytes loaded from (current TSS base + TSSstackAddress);
            NewSS ← new code-segment DPL; (* NULL selector with RPL = new CPL *)
FI;
IF IDT gate is 32-bit
            THEN
                IF new stack does not have room for 24 bytes (error code pushed)
                or 20 bytes (no error code pushed)
                    THEN #SS(error_code(NewSS,0,EXT)); FI;
                    (* idt operand to error_code is 0 because selector is used *)
            FI
    ELSE
        IF IDT gate is 16-bit
            THEN
                IF new stack does not have room for 12 bytes (error code pushed)
                or 10 bytes (no error code pushed);
                    THEN #SS(error_code(NewSS,0,EXT)); FI;
                    (* idt operand to error_code is 0 because selector is used *)
        ELSE (* 64-bit IDT gate*)
            IF StackAddress is non-canonical
                THEN #SS(EXT); FI; (* Error code contains NULL selector *)
    FI;
FI;
IF (IA32_EFER.LMA = 0) (* Not IA-32e mode *)
    THEN
        IF instruction pointer from IDT gate is not within new code-segment limits
            THEN #GP(EXT); FI; (* Error code contains NULL selector *)
        ESP ← NewESP;
        SS ← NewSS; (* Segment descriptor information also loaded *)
    ELSE (* IA-32e mode *)
        IF instruction pointer from IDT gate contains a non-canonical address
            THEN #GP(EXT); FI; (* Error code contains NULL selector *)
        RSP ← NewRSP & FFFFFFFFFFFFFFF0H;
        SS ← NewSS;
FI;
IF IDT gate is 32-bit
    THEN
        CS:EIP ← Gate(CS:EIP); (* Segment descriptor information also loaded *)
    ELSE
        IF IDT gate 16-bit
            THEN
                CS:IP ← Gate(CS:IP);
                (* Segment descriptor information also loaded *)
```

```
                ELSE (* 64-bit IDT gate *)
                    CS:RIP ← Gate(CS:RIP);
                    (* Segment descriptor information also loaded *)
            FI;
    FI;
    IF IDT gate is 32-bit
            THEN
                Push(far pointer to old stack);
                (* Old SS and ESP, 3 words padded to 4 *)
                Push(EFLAGS);
                Push(far pointer to return instruction);
                (* Old CS and EIP, 3 words padded to 4 *)
                Push(ErrorCode); (* If needed, 4 bytes *)
            ELSE
                IF IDT gate 16-bit
                    THEN
                        Push(far pointer to old stack);
                        (* Old SS and SP, 2 words *)
                        Push(EFLAGS(15-0]);
                        Push(far pointer to return instruction);
                        (* Old CS and IP, 2 words *)
                        Push(ErrorCode); (* If needed, 2 bytes *)
                    ELSE (* 64-bit IDT gate *)
                        Push(far pointer to old stack);
                        (* Old SS and SP, each an 8-byte push *)
                        Push(RFLAGS); (* 8-byte push *)
                        Push(far pointer to return instruction);
                        (* Old CS and RIP, each an 8-byte push *)
                        Push(ErrorCode); (* If needed, 8-bytes *)
            FI;
    FI;
    CPL ← new code-segment DPL;
    CS(RPL) ← CPL;
    IF IDT gate is interrupt gate
        THEN IF ← 0 (* Interrupt flag set to 0, interrupts disabled *); FI;
    TF ← 0;
    VM ← 0;
    RF ← 0;
    NT ← 0;
END;
INTERRUPT-FROM-VIRTUAL-8086-MODE:
    (* Identify stack-segment selector for privilege level 0 in current TSS *)
    IF current TSS is 32-bit
```

```
        THEN
            IF TSS limit < 9
                THEN #TS(error_code(current TSS selector,0,EXT)); FI;
                (* idt operand to error_code is 0 because selector is used *)
            NewSS ← 2 bytes loaded from (current TSS base + 8);
            NewESP ← 4 bytes loaded from (current TSS base + 4);
        ELSE (* current TSS is 16-bit *)
            IF TSS limit < 5
                THEN #TS(error_code(current TSS selector,0,EXT)); FI;
                (* idt operand to error_code is 0 because selector is used *)
            NewSS ← 2 bytes loaded from (current TSS base + 4);
            NewESP ← 2 bytes loaded from (current TSS base + 2);
    FI;
    IF NewSS is NULL
        THEN #TS(EXT); FI; (* Error code contains NULL selector *)
    IF NewSS index is not within its descriptor table limits
    or NewSS RPL ≠ 0
        THEN #TS(error_code(NewSS,0,EXT)); FI;
        (* idt operand to error_code is 0 because selector is used *)
    Read new stack-segment descriptor for NewSS in GDT or LDT;
    IF new stack-segment DPL ≠ 0 or stack segment does not indicate writable data segment
        THEN #TS(error_code(NewSS,0,EXT)); FI;
        (* idt operand to error_code is 0 because selector is used *)
    IF new stack segment not present
        THEN #SS(error_code(NewSS,0,EXT)); FI;
        (* idt operand to error_code is 0 because selector is used *)
    IF IDT gate is 32-bit
        THEN
            IF new stack does not have room for 40 bytes (error code pushed)
            or 36 bytes (no error code pushed)
                THEN #SS(error_code(NewSS,0,EXT)); FI;
                (* idt operand to error_code is 0 because selector is used *)
        ELSE (* IDT gate is 16-bit)
            IF new stack does not have room for 20 bytes (error code pushed)
            or 18 bytes (no error code pushed)
                THEN #SS(error_code(NewSS,0,EXT)); FI;
                (* idt operand to error_code is 0 because selector is used *)
    FI;
    IF instruction pointer from IDT gate is not within new code-segment limits
        THEN #GP(EXT); FI; (* Error code contains NULL selector *)
    tempEFLAGS ← EFLAGS;
    VM ← 0;
    TF ← 0;
```

```
    RF ← 0;
    NT ← 0;
    IF service through interrupt gate
        THEN IF = 0; FI;
    TempSS ← SS;
    TempESP ← ESP;
    SS ← NewSS;
    ESP ← NewESP;
    (* Following pushes are 16 bits for 16-bit IDT gates and 32 bits for 32-bit IDT gates;
    Segment selector pushes in 32-bit mode are padded to two words *)
    Push(GS);
    Push(FS);
    Push(DS);
    Push(ES);
    Push(TempSS);
    Push(TempESP);
    Push(TempEFlags);
    Push(CS);
    Push(EIP);
    GS ← 0; (* Segment registers made NULL, invalid for use in protected mode *)
    FS ← 0;
    DS ← 0;
    ES ← 0;
    CS:IP ← Gate(CS); (* Segment descriptor information also loaded *)
    IF OperandSize = 32
        THEN
            EIP ← Gate(instruction pointer);
        ELSE (* OperandSize is 16 *)
            EIP ← Gate(instruction pointer) AND 0000FFFFH;
    FI;
    (* Start execution of new routine in Protected Mode *)
END;
INTRA-PRIVILEGE-LEVEL-INTERRUPT:
    (* PE = 1, DPL = CPL or conforming segment *)
    IF IA32_EFER.LMA = 1 (* IA-32e mode *)
        IF IDT-descriptor IST ≠ 0
            THEN
                TSSstackAddress ← (IDT-descriptor IST « 3) + 28;
                IF (TSSstackAddress + 7) > TSS limit
                    THEN #TS(error_code(current TSS selector,0,EXT)); FI;
                    (* idt operand to error_code is 0 because selector is used *)
                NewRSP ← 8 bytes loaded from (current TSS base + TSSstackAddress);
    FI;
```

```
IF 32-bit gate (* implies IA32_EFER.LMA = 0 *)
    THEN
        IF current stack does not have room for 16 bytes (error code pushed)
        or 12 bytes (no error code pushed)
            THEN #SS(EXT); FI; (* Error code contains NULL selector *)
    ELSE IF 16-bit gate (* implies IA32_EFER.LMA = 0 *)
        IF current stack does not have room for 8 bytes (error code pushed)
        or 6 bytes (no error code pushed)
            THEN #SS(EXT); FI; (* Error code contains NULL selector *)
    ELSE (* IA32_EFER.LMA = 1, 64-bit gate*)
            IF NewRSP contains a non-canonical address
                THEN #SS(EXT); (* Error code contains NULL selector *)
    FI;
FI;
IF (IA32_EFER.LMA = 0) (* Not IA-32e mode *)
    THEN
        IF instruction pointer from IDT gate is not within new code-segment limit
            THEN #GP(EXT); FI; (* Error code contains NULL selector *)
    ELSE
        IF instruction pointer from IDT gate contains a non-canonical address
            THEN #GP(EXT); FI; (* Error code contains NULL selector *)
        RSP ← NewRSP & FFFFFFFFFFFFFFF0H;
FI;
IF IDT gate is 32-bit (* implies IA32_EFER.LMA = 0 *)
    THEN
        Push (EFLAGS);
        Push (far pointer to return instruction); (* 3 words padded to 4 *)
        CS:EIP ← Gate(CS:EIP); (* Segment descriptor information also loaded *)
        Push (ErrorCode); (* If any *)
    ELSE
        IF IDT gate is 16-bit (* implies IA32_EFER.LMA = 0 *)
            THEN
                Push (FLAGS);
                Push (far pointer to return location); (* 2 words *)
                CS:IP ← Gate(CS:IP);
                (* Segment descriptor information also loaded *)
                Push (ErrorCode); (* If any *)
            ELSE (* IA32_EFER.LMA = 1, 64-bit gate*)
                Push(far pointer to old stack);
                (* Old SS and SP, each an 8-byte push *)
                Push(RFLAGS); (* 8-byte push *)
                Push(far pointer to return instruction);
                (* Old CS and RIP, each an 8-byte push *)
```

```
                        Push(ErrorCode); (* If needed, 8 bytes *)
                        CS:RIP ← GATE(CS:RIP);
                        (* Segment descriptor information also loaded *)
                FI;
        FI;
        CS(RPL) ← CPL;
        IF IDT gate is interrupt gate
            THEN IF ← 0; FI; (* Interrupt flag set to 0; interrupts disabled *)
        TF ← 0;
        NT ← 0;
        VM ← 0;
        RF ← 0;
END;
```

Flags Affected

The EFLAGS register is pushed onto the stack. The IF, TF, NT, AC, RF, and VM flags may be cleared, depending on the mode of operation of the processor when the INT instruction is executed (see the "Operation" section). If the interrupt uses a task gate, any flags may be set or cleared, controlled by the EFLAGS image in the new task's TSS.

Protected Mode Exceptions

#GP(error_code)	If the instruction pointer in the IDT or in the interrupt-, trap-, or task gate is beyond the code segment limits.
	If the segment selector in the interrupt-, trap-, or task gate is NULL.
	If an interrupt-, trap-, or task gate, code segment, or TSS segment selector index is outside its descriptor table limits.
	If the interrupt vector number is outside the IDT limits.
	If an IDT descriptor is not an interrupt-, trap-, or task-descriptor.
	If an interrupt is generated by the INT *n*, INT 3, or INTO instruction and the DPL of an interrupt-, trap-, or task-descriptor is less than the CPL.
	If the segment selector in an interrupt- or trap-gate does not point to a segment descriptor for a code segment.
	If the segment selector for a TSS has its local/global bit set for local.
	If a TSS segment descriptor specifies that the TSS is busy or not available.
#SS(error_code)	If pushing the return address, flags, or error code onto the stack exceeds the bounds of the stack segment and no stack switch occurs.

	If the SS register is being loaded and the segment pointed to is marked not present.
	If pushing the return address, flags, error code, or stack segment pointer exceeds the bounds of the new stack segment when a stack switch occurs.
#NP(error_code)	If code segment, interrupt-, trap-, or task gate, or TSS is not present.
#TS(error_code)	If the RPL of the stack segment selector in the TSS is not equal to the DPL of the code segment being accessed by the interrupt or trap gate.
	If DPL of the stack segment descriptor pointed to by the stack segment selector in the TSS is not equal to the DPL of the code segment descriptor for the interrupt or trap gate.
	If the stack segment selector in the TSS is NULL.
	If the stack segment for the TSS is not a writable data segment.
	If segment-selector index for stack segment is outside descriptor table limits.
#PF(fault-code)	If a page fault occurs.
#UD	If the LOCK prefix is used.
#AC(EXT)	If alignment checking is enabled, the gate DPL is 3, and a stack push is unaligned.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the interrupt vector number is outside the IDT limits.
#SS	If stack limit violation on push.
	If pushing the return address, flags, or error code onto the stack exceeds the bounds of the stack segment.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(error_code)	(For INT *n,* INTO, or BOUND instruction) If the IOPL is less than 3 or the DPL of the interrupt-, trap-, or task-gate descriptor is not equal to 3.
	If the instruction pointer in the IDT or in the interrupt-, trap-, or task gate is beyond the code segment limits.
	If the segment selector in the interrupt-, trap-, or task gate is NULL.
	If a interrupt-, trap-, or task gate, code segment, or TSS segment selector index is outside its descriptor table limits.

	If the interrupt vector number is outside the IDT limits.
	If an IDT descriptor is not an interrupt-, trap-, or task-descriptor.
	If an interrupt is generated by the INT *n* instruction and the DPL of an interrupt-, trap-, or task-descriptor is less than the CPL.
	If the segment selector in an interrupt- or trap-gate does not point to a segment descriptor for a code segment.
	If the segment selector for a TSS has its local/global bit set for local.
#SS(error_code)	If the SS register is being loaded and the segment pointed to is marked not present.
	If pushing the return address, flags, error code, stack segment pointer, or data segments exceeds the bounds of the stack segment.
#NP(error_code)	If code segment, interrupt-, trap-, or task gate, or TSS is not present.
#TS(error_code)	If the RPL of the stack segment selector in the TSS is not equal to the DPL of the code segment being accessed by the interrupt or trap gate.
	If DPL of the stack segment descriptor for the TSS's stack segment is not equal to the DPL of the code segment descriptor for the interrupt or trap gate.
	If the stack segment selector in the TSS is NULL.
	If the stack segment for the TSS is not a writable data segment.
	If segment-selector index for stack segment is outside descriptor table limits.
#PF(fault-code)	If a page fault occurs.
#BP	If the INT 3 instruction is executed.
#OF	If the INTO instruction is executed and the OF flag is set.
#UD	If the LOCK prefix is used.
#AC(EXT)	If alignment checking is enabled, the gate DPL is 3, and a stack push is unaligned.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(error_code)	If the instruction pointer in the 64-bit interrupt gate or 64-bit trap gate is non-canonical.
	If the segment selector in the 64-bit interrupt or trap gate is NULL.
	If the interrupt vector number is outside the IDT limits.

	If the interrupt vector number points to a gate which is in non-canonical space.
	If the interrupt vector number points to a descriptor which is not a 64-bit interrupt gate or 64-bit trap gate.
	If the descriptor pointed to by the gate selector is outside the descriptor table limit.
	If the descriptor pointed to by the gate selector is in non-canonical space.
	If the descriptor pointed to by the gate selector is not a code segment.
	If the descriptor pointed to by the gate selector doesn't have the L-bit set, or has both the L-bit and D-bit set.
	If the descriptor pointed to by the gate selector has DPL > CPL.
#SS(error_code)	If a push of the old EFLAGS, CS selector, EIP, or error code is in non-canonical space with no stack switch.
	If a push of the old SS selector, ESP, EFLAGS, CS selector, EIP, or error code is in non-canonical space on a stack switch (either CPL change or no-CPL with IST).
#NP(error_code)	If the 64-bit interrupt-gate, 64-bit trap-gate, or code segment is not present.
#TS(error_code)	If an attempt to load RSP from the TSS causes an access to non-canonical space.
	If the RSP from the TSS is outside descriptor table limits.
#PF(fault-code)	If a page fault occurs.
#UD	If the LOCK prefix is used.
#AC(EXT)	If alignment checking is enabled, the gate DPL is 3, and a stack push is unaligned.

INVD—Invalidate Internal Caches

Opcode	Instruction	Op/En	64-Bit Mode	Compat/Leg Mode	Description
0F 08	INVD	A	Valid	Valid	Flush internal caches; initiate flushing of external caches.

NOTES:

* See the IA-32 Architecture Compatibility section below.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Invalidates (flushes) the processor's internal caches and issues a special-function bus cycle that directs external caches to also flush themselves. Data held in internal caches is not written back to main memory.

After executing this instruction, the processor does not wait for the external caches to complete their flushing operation before proceeding with instruction execution. It is the responsibility of hardware to respond to the cache flush signal.

The INVD instruction is a privileged instruction. When the processor is running in protected mode, the CPL of a program or procedure must be 0 to execute this instruction.

Use this instruction with care. Data cached internally and not written back to main memory will be lost. Unless there is a specific requirement or benefit to flushing caches without writing back modified cache lines (for example, testing or fault recovery where cache coherency with main memory is not a concern), software should use the WBINVD instruction.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

IA-32 Architecture Compatibility

The INVD instruction is implementation dependent; it may be implemented differently on different families of Intel 64 or IA-32 processors. This instruction is not supported on IA-32 processors earlier than the Intel486 processor.

Operation

```
Flush(InternalCaches);
SignalFlush(ExternalCaches);
Continue (* Continue execution *)
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If the current privilege level is not 0.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	The INVD instruction cannot be executed in virtual-8086 mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

INVLPG—Invalidate TLB Entry

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 01/7	INVLPG *m*	A	Valid	Valid	Invalidate TLB Entry for page that contains *m.*

NOTES:

* See the IA-32 Architecture Compatibility section below.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r)	NA	NA	NA

Description

Invalidates (flushes) the translation lookaside buffer (TLB) entry specified with the source operand. The source operand is a memory address. The processor determines the page that contains that address and flushes the TLB entry for that page.

The INVLPG instruction is a privileged instruction. When the processor is running in protected mode, the CPL of a program or procedure must be 0 to execute this instruction.

The INVLPG instruction normally flushes the TLB entry only for the specified page; however, in some cases, it flushes the entire TLB. See "MOV—Move to/from Control Registers" in this chapter for further information on operations that flush the TLB.

This instruction's operation is the same in all non-64-bit modes. It also operates the same in 64-bit mode, except if the memory address is in non-canonical form. In this case, INVLPG is the same as a NOP.

IA-32 Architecture Compatibility

The INVLPG instruction is implementation dependent, and its function may be implemented differently on different families of Intel 64 or IA-32 processors. This instruction is not supported on IA-32 processors earlier than the Intel486 processor.

Operation

```
Flush(RelevantTLBEntries);
Continue; (* Continue execution *)
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If the current privilege level is not 0.
#UD	Operand is a register.
	If the LOCK prefix is used.

Real-Address Mode Exceptions

#UD	Operand is a register.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	The INVLPG instruction cannot be executed at the virtual-8086 mode.

64-Bit Mode Exceptions

#GP(0)	If the current privilege level is not 0.
#UD	Operand is a register.
	If the LOCK prefix is used.

IRET/IRETD—Interrupt Return

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
CF	IRET	A	Valid	Valid	Interrupt return (16-bit operand size).
CF	IRETD	A	Valid	Valid	Interrupt return (32-bit operand size).
REX.W + CF	IRETQ	A	Valid	N.E.	Interrupt return (64-bit operand size).

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Returns program control from an exception or interrupt handler to a program or procedure that was interrupted by an exception, an external interrupt, or a software-generated interrupt. These instructions are also used to perform a return from a nested task. (A nested task is created when a CALL instruction is used to initiate a task switch or when an interrupt or exception causes a task switch to an interrupt or exception handler.) See the section titled “Task Linking” in Chapter 7 of the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 3A*.

IRET and IRETD are mnemonics for the same opcode. The IRETD mnemonic (interrupt return double) is intended for use when returning from an interrupt when using the 32-bit operand size; however, most assemblers use the IRET mnemonic interchangeably for both operand sizes.

In Real-Address Mode, the IRET instruction preforms a far return to the interrupted program or procedure. During this operation, the processor pops the return instruction pointer, return code segment selector, and EFLAGS image from the stack to the EIP, CS, and EFLAGS registers, respectively, and then resumes execution of the interrupted program or procedure.

In Protected Mode, the action of the IRET instruction depends on the settings of the NT (nested task) and VM flags in the EFLAGS register and the VM flag in the EFLAGS image stored on the current stack. Depending on the setting of these flags, the processor performs the following types of interrupt returns:

- Return from virtual-8086 mode.
- Return to virtual-8086 mode.
- Intra-privilege level return.
- Inter-privilege level return.
- Return from nested task (task switch).

If the NT flag (EFLAGS register) is cleared, the IRET instruction performs a far return from the interrupt procedure, without a task switch. The code segment being returned to must be equally or less privileged than the interrupt handler routine (as indicated by the RPL field of the code segment selector popped from the stack).

As with a real-address mode interrupt return, the IRET instruction pops the return instruction pointer, return code segment selector, and EFLAGS image from the stack to the EIP, CS, and EFLAGS registers, respectively, and then resumes execution of the interrupted program or procedure. If the return is to another privilege level, the IRET instruction also pops the stack pointer and SS from the stack, before resuming program execution. If the return is to virtual-8086 mode, the processor also pops the data segment registers from the stack.

If the NT flag is set, the IRET instruction performs a task switch (return) from a nested task (a task called with a CALL instruction, an interrupt, or an exception) back to the calling or interrupted task. The updated state of the task executing the IRET instruction is saved in its TSS. If the task is re-entered later, the code that follows the IRET instruction is executed.

If the NT flag is set and the processor is in IA-32e mode, the IRET instruction causes a general protection exception.

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.W prefix promotes operation to 64 bits (IRETQ). See the summary chart at the beginning of this section for encoding data and limits.

See "Changes to Instruction Behavior in VMX Non-Root Operation" in Chapter 22 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3B*, for more information about the behavior of this instruction in VMX non-root operation.

Operation

```
IF PE = 0
    THEN
         GOTO REAL-ADDRESS-MODE;
    ELSE
         IF (IA32_EFER.LMA = 0)
                   THEN (* Protected mode *)
                        GOTO PROTECTED-MODE;
                   ELSE (* IA-32e mode *)
                        GOTO IA-32e-MODE;
         FI;
FI;
REAL-ADDRESS-MODE;
    IF OperandSize = 32
         THEN
              IF top 12 bytes of stack not within stack limits
                   THEN #SS; FI;
```

```
            tempEIP ← 4 bytes at end of stack
            IF tempEIP[31:16] is not zero THEN #GP(0); FI;
            EIP ← Pop();
            CS ← Pop(); (* 32-bit pop, high-order 16 bits discarded *)
            tempEFLAGS ← Pop();
            EFLAGS ← (tempEFLAGS AND 257FD5H) OR (EFLAGS AND 1A0000H);
        ELSE (* OperandSize = 16 *)
            IF top 6 bytes of stack are not within stack limits
                THEN #SS; FI;
            EIP ← Pop(); (* 16-bit pop; clear upper 16 bits *)
            CS ← Pop(); (* 16-bit pop *)
            EFLAGS[15:0] ← Pop();
    FI;
    END;
PROTECTED-MODE:
    IF VM = 1 (* Virtual-8086 mode: PE = 1, VM = 1 *)
        THEN
            GOTO RETURN-FROM-VIRTUAL-8086-MODE; (* PE = 1, VM = 1 *)
    FI;
    IF NT = 1
        THEN
            GOTO TASK-RETURN; (* PE = 1, VM = 0, NT = 1 *)
    FI;
    IF OperandSize = 32
        THEN
            IF top 12 bytes of stack not within stack limits
                THEN #SS(0); FI;
            tempEIP ← Pop();
            tempCS ← Pop();
            tempEFLAGS ← Pop();
        ELSE (* OperandSize = 16 *)
            IF top 6 bytes of stack are not within stack limits
                THEN #SS(0); FI;
            tempEIP ← Pop();
            tempCS ← Pop();
            tempEFLAGS ← Pop();
            tempEIP ← tempEIP AND FFFFH;
            tempEFLAGS ← tempEFLAGS AND FFFFH;
    FI;
    IF tempEFLAGS(VM) = 1 and CPL = 0
        THEN
            GOTO RETURN-TO-VIRTUAL-8086-MODE;
        ELSE
```

```
                GOTO PROTECTED-MODE-RETURN;
    FI;
IA-32e-MODE:
    IF NT = 1
        THEN #GP(0);
    ELSE IF OperandSize = 32
        THEN
            IF top 12 bytes of stack not within stack limits
                THEN #SS(0); FI;
            tempEIP ← Pop();
            tempCS ← Pop();
            tempEFLAGS ← Pop();
        ELSE IF OperandSize = 16
            THEN
                IF top 6 bytes of stack are not within stack limits
                    THEN #SS(0); FI;
                tempEIP ← Pop();
                tempCS ← Pop();
                tempEFLAGS ← Pop();
                tempEIP ← tempEIP AND FFFFH;
                tempEFLAGS ← tempEFLAGS AND FFFFH;
            FI;
        ELSE (* OperandSize = 64 *)
            THEN
                    tempRIP ← Pop();
                    tempCS ← Pop();
                    tempEFLAGS ← Pop();
                    tempRSP ← Pop();
                    tempSS ← Pop();
    FI;
    GOTO IA-32e-MODE-RETURN;

RETURN-FROM-VIRTUAL-8086-MODE:
(* Processor is in virtual-8086 mode when IRET is executed and stays in virtual-8086 mode *)
    IF IOPL = 3 (* Virtual mode: PE = 1, VM = 1, IOPL = 3 *)
        THEN IF OperandSize = 32
            THEN
                IF top 12 bytes of stack not within stack limits
                    THEN #SS(0); FI;
                IF instruction pointer not within code segment limits
                    THEN #GP(0); FI;
                EIP ← Pop();
                CS ← Pop(); (* 32-bit pop, high-order 16 bits discarded *)
                EFLAGS ← Pop();
```

```
                    (* VM, IOPL,VIP and VIF EFLAG bits not modified by pop *)
                ELSE (* OperandSize = 16 *)
                    IF top 6 bytes of stack are not within stack limits
                        THEN #SS(0); FI;
                    IF instruction pointer not within code segment limits
                        THEN #GP(0); FI;
                    EIP ← Pop();
                    EIP ← EIP AND 0000FFFFH;
                    CS ← Pop(); (* 16-bit pop *)
                    EFLAGS[15:0] ← Pop(); (* IOPL in EFLAGS not modified by pop *)
                FI;
            ELSE
                #GP(0); (* Trap to virtual-8086 monitor: PE = 1, VM = 1, IOPL < 3 *)
    FI;
END;

RETURN-TO-VIRTUAL-8086-MODE:
    (* Interrupted procedure was in virtual-8086 mode: PE = 1, CPL=0, VM = 1 in flag image *)
    IF top 24 bytes of stack are not within stack segment limits
        THEN #SS(0); FI;
    IF instruction pointer not within code segment limits
        THEN #GP(0); FI;
    CS ← tempCS;
    EIP ← tempEIP & FFFFH;
    EFLAGS ← tempEFLAGS;
    TempESP ← Pop();
    TempSS ← Pop();
    ES ← Pop(); (* Pop 2 words; throw away high-order word *)
    DS ← Pop(); (* Pop 2 words; throw away high-order word *)
    FS ← Pop(); (* Pop 2 words; throw away high-order word *)
    GS ← Pop(); (* Pop 2 words; throw away high-order word *)
    SS:ESP ← TempSS:TempESP;
    CPL ← 3;
    (* Resume execution in Virtual-8086 mode *)
END;

TASK-RETURN: (* PE = 1, VM = 0, NT = 1 *)
    Read segment selector in link field of current TSS;
    IF local/global bit is set to local
    or index not within GDT limits
        THEN #TS (TSS selector); FI;
    Access TSS for task specified in link field of current TSS;
    IF TSS descriptor type is not TSS or if the TSS is marked not busy
```

```
            THEN #TS (TSS selector); FI;
    IF TSS not present
            THEN #NP(TSS selector); FI;
    SWITCH-TASKS (without nesting) to TSS specified in link field of current TSS;
    Mark the task just abandoned as NOT BUSY;
    IF EIP is not within code segment limit
            THEN #GP(0); FI;
END;

PROTECTED-MODE-RETURN: (* PE = 1 *)
    IF return code segment selector is NULL
            THEN GP(0); FI;
    IF return code segment selector addresses descriptor beyond descriptor table limit
            THEN GP(selector); FI;
    Read segment descriptor pointed to by the return code segment selector;
    IF return code segment descriptor is not a code segment
            THEN #GP(selector); FI;
    IF return code segment selector RPL < CPL
            THEN #GP(selector); FI;
    IF return code segment descriptor is conforming
    and return code segment DPL > return code segment selector RPL
            THEN #GP(selector); FI;
    IF return code segment descriptor is not present
            THEN #NP(selector); FI;
    IF return code segment selector RPL > CPL
            THEN GOTO RETURN-OUTER-PRIVILEGE-LEVEL;
            ELSE GOTO RETURN-TO-SAME-PRIVILEGE-LEVEL; FI;
END;

RETURN-TO-SAME-PRIVILEGE-LEVEL: (* PE = 1, RPL = CPL *)
    IF new mode ≠ 64-Bit Mode
            THEN
                    IF tempEIP is not within code segment limits
                            THEN #GP(0); FI;
                    EIP ← tempEIP;
            ELSE (* new mode = 64-bit mode *)
                    IF tempRIP is non-canonical
                                    THEN #GP(0); FI;
                    RIP ← tempRIP;
    FI;
    CS ← tempCS; (* Segment descriptor information also loaded *)
    EFLAGS (CF, PF, AF, ZF, SF, TF, DF, OF, NT) ← tempEFLAGS;
    IF OperandSize = 32 or OperandSize = 64
```

```
            THEN EFLAGS(RF, AC, ID) ← tempEFLAGS; FI;
    IF CPL ≤ IOPL
            THEN EFLAGS(IF) ← tempEFLAGS; FI;
    IF CPL = 0
            THEN (* VM = 0 in flags image *)
            EFLAGS(IOPL) ← tempEFLAGS;
            IF OperandSize = 32 or OperandSize = 64
                THEN EFLAGS(VIF, VIP) ← tempEFLAGS; FI;
    FI;
END;

RETURN-TO-OUTER-PRIVILEGE-LEVEL:
    IF OperandSize = 32
        THEN
            IF top 8 bytes on stack are not within limits
                THEN #SS(0); FI;
        ELSE (* OperandSize = 16 *)
            IF top 4 bytes on stack are not within limits
                THEN #SS(0); FI;
    FI;
    Read return segment selector;
    IF stack segment selector is NULL
        THEN #GP(0); FI;
    IF return stack segment selector index is not within its descriptor table limits
        THEN #GP(SSselector); FI;
    Read segment descriptor pointed to by return segment selector;
    IF stack segment selector RPL ≠ RPL of the return code segment selector
    or the stack segment descriptor does not indicate a a writable data segment;
    or the stack segment DPL ≠ RPL of the return code segment selector
        THEN #GP(SS selector); FI;
    IF stack segment is not present
        THEN #SS(SS selector); FI;
    IF new mode ≠ 64-Bit Mode
        THEN
            IF tempEIP is not within code segment limits
                THEN #GP(0); FI;
            EIP ← tempEIP;
        ELSE (* new mode = 64-bit mode *)
            IF tempRIP is non-canonical
                    THEN #GP(0); FI;
            RIP ← tempRIP;
    FI;
    CS ← tempCS;
```

```
    EFLAGS (CF, PF, AF, ZF, SF, TF, DF, OF, NT) ← tempEFLAGS;
    IF OperandSize = 32
        THEN EFLAGS(RF, AC, ID) ← tempEFLAGS; FI;
    IF CPL ≤ IOPL
        THEN EFLAGS(IF) ← tempEFLAGS; FI;
    IF CPL = 0
        THEN
            EFLAGS(IOPL) ← tempEFLAGS;
            IF OperandSize = 32
                THEN EFLAGS(VM, VIF, VIP) ← tempEFLAGS; FI;
            IF OperandSize = 64
                THEN EFLAGS(VIF, VIP) ← tempEFLAGS; FI;
    FI;
    CPL ← RPL of the return code segment selector;
    FOR each of segment register (ES, FS, GS, and DS)
        DO
            IF segment register points to data or non-conforming code segment
            and CPL > segment descriptor DPL (* Stored in hidden part of segment register *)
                THEN (* Segment register invalid *)
                    SegmentSelector ← 0; (* NULL segment selector *)
            FI;
        OD;
END;

IA-32e-MODE-RETURN: (* IA32_EFER.LMA = 1, PE = 1 *)
    IF ( (return code segment selector is NULL) or (return RIP is non-canonical) or
            (SS selector is NULL going back to compatibility mode) or
            (SS selector is NULL going back to CPL3 64-bit mode) or
            (RPL <> CPL going back to non-CPL3 64-bit mode for a NULL SS selector) )
        THEN GP(0); FI;
    IF return code segment selector addresses descriptor beyond descriptor table limit
        THEN GP(selector); FI;
    Read segment descriptor pointed to by the return code segment selector;
    IF return code segment descriptor is not a code segment
        THEN #GP(selector); FI;
    IF return code segment selector RPL < CPL
        THEN #GP(selector); FI;
    IF return code segment descriptor is conforming
    and return code segment DPL > return code segment selector RPL
        THEN #GP(selector); FI;
    IF return code segment descriptor is not present
        THEN #NP(selector); FI;
    IF return code segment selector RPL > CPL
```

```
        THEN GOTO RETURN-OUTER-PRIVILEGE-LEVEL;
        ELSE GOTO RETURN-TO-SAME-PRIVILEGE-LEVEL; FI;
END;
```

Flags Affected

All the flags and fields in the EFLAGS register are potentially modified, depending on the mode of operation of the processor. If performing a return from a nested task to a previous task, the EFLAGS register will be modified according to the EFLAGS image stored in the previous task's TSS.

Protected Mode Exceptions

#GP(0)	If the return code or stack segment selector is NULL.
	If the return instruction pointer is not within the return code segment limit.
#GP(selector)	If a segment selector index is outside its descriptor table limits.
	If the return code segment selector RPL is greater than the CPL.
	If the DPL of a conforming-code segment is greater than the return code segment selector RPL.
	If the DPL for a nonconforming-code segment is not equal to the RPL of the code segment selector.
	If the stack segment descriptor DPL is not equal to the RPL of the return code segment selector.
	If the stack segment is not a writable data segment.
	If the stack segment selector RPL is not equal to the RPL of the return code segment selector.
	If the segment descriptor for a code segment does not indicate it is a code segment.
	If the segment selector for a TSS has its local/global bit set for local.
	If a TSS segment descriptor specifies that the TSS is not busy.
	If a TSS segment descriptor specifies that the TSS is not available.
#SS(0)	If the top bytes of stack are not within stack limits.
#NP(selector)	If the return code or stack segment is not present.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If an unaligned memory reference occurs when the CPL is 3 and alignment checking is enabled.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If the return instruction pointer is not within the return code segment limit.
#SS	If the top bytes of stack are not within stack limits.

Virtual-8086 Mode Exceptions

#GP(0)	If the return instruction pointer is not within the return code segment limit.
	IF IOPL not equal to 3.
#PF(fault-code)	If a page fault occurs.
#SS(0)	If the top bytes of stack are not within stack limits.
#AC(0)	If an unaligned memory reference occurs and alignment checking is enabled.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

#GP(0)	If EFLAGS.NT[bit 14] = 1.

Other exceptions same as in Protected Mode.

64-Bit Mode Exceptions

#GP(0)	If EFLAGS.NT[bit 14] = 1.
	If the return code segment selector is NULL.
	If the stack segment selector is NULL going back to compatibility mode.
	If the stack segment selector is NULL going back to CPL3 64-bit mode.
	If a NULL stack segment selector RPL is not equal to CPL going back to non-CPL3 64-bit mode.
	If the return instruction pointer is not within the return code segment limit.
	If the return instruction pointer is non-canonical.
#GP(Selector)	If a segment selector index is outside its descriptor table limits.
	If a segment descriptor memory address is non-canonical.
	If the segment descriptor for a code segment does not indicate it is a code segment.
	If the proposed new code segment descriptor has both the D-bit and L-bit set.
	If the DPL for a nonconforming-code segment is not equal to the RPL of the code segment selector.
	If CPL is greater than the RPL of the code segment selector.

	If the DPL of a conforming-code segment is greater than the return code segment selector RPL.
	If the stack segment is not a writable data segment.
	If the stack segment descriptor DPL is not equal to the RPL of the return code segment selector.
	If the stack segment selector RPL is not equal to the RPL of the return code segment selector.
#SS(0)	If an attempt to pop a value off the stack violates the SS limit.
	If an attempt to pop a value off the stack causes a non-canonical address to be referenced.
#NP(selector)	If the return code or stack segment is not present.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If an unaligned memory reference occurs when the CPL is 3 and alignment checking is enabled.
#UD	If the LOCK prefix is used.

Jcc—Jump if Condition Is Met

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
77 *cb*	JA *rel8*	A	Valid	Valid	Jump short if above (CF=0 and ZF=0).
73 *cb*	JAE *rel8*	A	Valid	Valid	Jump short if above or equal (CF=0).
72 *cb*	JB *rel8*	A	Valid	Valid	Jump short if below (CF=1).
76 *cb*	JBE *rel8*	A	Valid	Valid	Jump short if below or equal (CF=1 or ZF=1).
72 *cb*	JC *rel8*	A	Valid	Valid	Jump short if carry (CF=1).
E3 *cb*	JCXZ *rel8*	A	N.E.	Valid	Jump short if CX register is 0.
E3 *cb*	JECXZ *rel8*	A	Valid	Valid	Jump short if ECX register is 0.
E3 *cb*	JRCXZ *rel8*	A	Valid	N.E.	Jump short if RCX register is 0.
74 *cb*	JE *rel8*	A	Valid	Valid	Jump short if equal (ZF=1).
7F *cb*	JG *rel8*	A	Valid	Valid	Jump short if greater (ZF=0 and SF=OF).
7D *cb*	JGE *rel8*	A	Valid	Valid	Jump short if greater or equal (SF=OF).
7C *cb*	JL *rel8*	A	Valid	Valid	Jump short if less (SF≠ OF).
7E *cb*	JLE *rel8*	A	Valid	Valid	Jump short if less or equal (ZF=1 or SF≠ OF).
76 *cb*	JNA *rel8*	A	Valid	Valid	Jump short if not above (CF=1 or ZF=1).
72 *cb*	JNAE *rel8*	A	Valid	Valid	Jump short if not above or equal (CF=1).
73 *cb*	JNB *rel8*	A	Valid	Valid	Jump short if not below (CF=0).
77 *cb*	JNBE *rel8*	A	Valid	Valid	Jump short if not below or equal (CF=0 and ZF=0).
73 *cb*	JNC *rel8*	A	Valid	Valid	Jump short if not carry (CF=0).
75 *cb*	JNE *rel8*	A	Valid	Valid	Jump short if not equal (ZF=0).
7E *cb*	JNG *rel8*	A	Valid	Valid	Jump short if not greater (ZF=1 or SF≠ OF).

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
7C *cb*	JNGE *rel8*	A	Valid	Valid	Jump short if not greater or equal (SF≠ OF).
7D *cb*	JNL *rel8*	A	Valid	Valid	Jump short if not less (SF=OF).
7F *cb*	JNLE *rel8*	A	Valid	Valid	Jump short if not less or equal (ZF=0 and SF=OF).
71 *cb*	JNO *rel8*	A	Valid	Valid	Jump short if not overflow (OF=0).
7B *cb*	JNP *rel8*	A	Valid	Valid	Jump short if not parity (PF=0).
79 *cb*	JNS *rel8*	A	Valid	Valid	Jump short if not sign (SF=0).
75 *cb*	JNZ *rel8*	A	Valid	Valid	Jump short if not zero (ZF=0).
70 *cb*	JO *rel8*	A	Valid	Valid	Jump short if overflow (OF=1).
7A *cb*	JP *rel8*	A	Valid	Valid	Jump short if parity (PF=1).
7A *cb*	JPE *rel8*	A	Valid	Valid	Jump short if parity even (PF=1).
7B *cb*	JPO *rel8*	A	Valid	Valid	Jump short if parity odd (PF=0).
78 *cb*	JS *rel8*	A	Valid	Valid	Jump short if sign (SF=1).
74 *cb*	JZ *rel8*	A	Valid	Valid	Jump short if zero (ZF ← 1).
0F 87 *cw*	JA *rel16*	A	N.S.	Valid	Jump near if above (CF=0 and ZF=0). Not supported in 64-bit mode.
0F 87 *cd*	JA *rel32*	A	Valid	Valid	Jump near if above (CF=0 and ZF=0).
0F 83 *cw*	JAE *rel16*	A	N.S.	Valid	Jump near if above or equal (CF=0). Not supported in 64-bit mode.
0F 83 *cd*	JAE *rel32*	A	Valid	Valid	Jump near if above or equal (CF=0).
0F 82 *cw*	JB *rel16*	A	N.S.	Valid	Jump near if below (CF=1). Not supported in 64-bit mode.
0F 82 *cd*	JB *rel32*	A	Valid	Valid	Jump near if below (CF=1).

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 86 *cw*	JBE *rel16*	A	N.S.	Valid	Jump near if below or equal (CF=1 or ZF=1). Not supported in 64-bit mode.
0F 86 *cd*	JBE *rel32*	A	Valid	Valid	Jump near if below or equal (CF=1 or ZF=1).
0F 82 *cw*	JC *rel16*	A	N.S.	Valid	Jump near if carry (CF=1). Not supported in 64-bit mode.
0F 82 *cd*	JC *rel32*	A	Valid	Valid	Jump near if carry (CF=1).
0F 84 *cw*	JE *rel16*	A	N.S.	Valid	Jump near if equal (ZF=1). Not supported in 64-bit mode.
0F 84 *cd*	JE *rel32*	A	Valid	Valid	Jump near if equal (ZF=1).
0F 84 *cw*	JZ *rel16*	A	N.S.	Valid	Jump near if 0 (ZF=1). Not supported in 64-bit mode.
0F 84 *cd*	JZ *rel32*	A	Valid	Valid	Jump near if 0 (ZF=1).
0F 8F *cw*	JG *rel16*	A	N.S.	Valid	Jump near if greater (ZF=0 and SF=OF). Not supported in 64-bit mode.
0F 8F *cd*	JG *rel32*	A	Valid	Valid	Jump near if greater (ZF=0 and SF=OF).
0F 8D *cw*	JGE *rel16*	A	N.S.	Valid	Jump near if greater or equal (SF=OF). Not supported in 64-bit mode.
0F 8D *cd*	JGE *rel32*	A	Valid	Valid	Jump near if greater or equal (SF=OF).
0F 8C *cw*	JL *rel16*	A	N.S.	Valid	Jump near if less (SF≠ OF). Not supported in 64-bit mode.
0F 8C *cd*	JL *rel32*	A	Valid	Valid	Jump near if less (SF≠ OF).
0F 8E *cw*	JLE *rel16*	A	N.S.	Valid	Jump near if less or equal (ZF=1 or SF≠ OF). Not supported in 64-bit mode.
0F 8E *cd*	JLE *rel32*	A	Valid	Valid	Jump near if less or equal (ZF=1 or SF≠ OF).

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 86 *cw*	JNA *rel16*	A	N.S.	Valid	Jump near if not above (CF=1 or ZF=1). Not supported in 64-bit mode.
0F 86 *cd*	JNA *rel32*	A	Valid	Valid	Jump near if not above (CF=1 or ZF=1).
0F 82 *cw*	JNAE *rel16*	A	N.S.	Valid	Jump near if not above or equal (CF=1). Not supported in 64-bit mode.
0F 82 *cd*	JNAE *rel32*	A	Valid	Valid	Jump near if not above or equal (CF=1).
0F 83 *cw*	JNB *rel16*	A	N.S.	Valid	Jump near if not below (CF=0). Not supported in 64-bit mode.
0F 83 *cd*	JNB *rel32*	A	Valid	Valid	Jump near if not below (CF=0).
0F 87 *cw*	JNBE *rel16*	A	N.S.	Valid	Jump near if not below or equal (CF=0 and ZF=0). Not supported in 64-bit mode.
0F 87 *cd*	JNBE *rel32*	A	Valid	Valid	Jump near if not below or equal (CF=0 and ZF=0).
0F 83 *cw*	JNC *rel16*	A	N.S.	Valid	Jump near if not carry (CF=0). Not supported in 64-bit mode.
0F 83 *cd*	JNC *rel32*	A	Valid	Valid	Jump near if not carry (CF=0).
0F 85 *cw*	JNE *rel16*	A	N.S.	Valid	Jump near if not equal (ZF=0). Not supported in 64-bit mode.
0F 85 *cd*	JNE *rel32*	A	Valid	Valid	Jump near if not equal (ZF=0).
0F 8E *cw*	JNG *rel16*	A	N.S.	Valid	Jump near if not greater (ZF=1 or SF≠ OF). Not supported in 64-bit mode.
0F 8E *cd*	JNG *rel32*	A	Valid	Valid	Jump near if not greater (ZF=1 or SF≠ OF).
0F 8C *cw*	JNGE *rel16*	A	N.S.	Valid	Jump near if not greater or equal (SF≠ OF). Not supported in 64-bit mode.

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 8C *cd*	JNGE *rel32*	A	Valid	Valid	Jump near if not greater or equal (SF≠ OF).
0F 8D *cw*	JNL *rel16*	A	N.S.	Valid	Jump near if not less (SF=OF). Not supported in 64-bit mode.
0F 8D *cd*	JNL *rel32*	A	Valid	Valid	Jump near if not less (SF=OF).
0F 8F *cw*	JNLE *rel16*	A	N.S.	Valid	Jump near if not less or equal (ZF=0 and SF=OF). Not supported in 64-bit mode.
0F 8F *cd*	JNLE *rel32*	A	Valid	Valid	Jump near if not less or equal (ZF=0 and SF=OF).
0F 81 *cw*	JNO *rel16*	A	N.S.	Valid	Jump near if not overflow (OF=0). Not supported in 64-bit mode.
0F 81 *cd*	JNO *rel32*	A	Valid	Valid	Jump near if not overflow (OF=0).
0F 8B *cw*	JNP *rel16*	A	N.S.	Valid	Jump near if not parity (PF=0). Not supported in 64-bit mode.
0F 8B *cd*	JNP *rel32*	A	Valid	Valid	Jump near if not parity (PF=0).
0F 89 *cw*	JNS *rel16*	A	N.S.	Valid	Jump near if not sign (SF=0). Not supported in 64-bit mode.
0F 89 *cd*	JNS *rel32*	A	Valid	Valid	Jump near if not sign (SF=0).
0F 85 *cw*	JNZ *rel16*	A	N.S.	Valid	Jump near if not zero (ZF=0). Not supported in 64-bit mode.
0F 85 *cd*	JNZ *rel32*	A	Valid	Valid	Jump near if not zero (ZF=0).
0F 80 *cw*	JO *rel16*	A	N.S.	Valid	Jump near if overflow (OF=1). Not supported in 64-bit mode.
0F 80 *cd*	JO *rel32*	A	Valid	Valid	Jump near if overflow (OF=1).

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 8A *cw*	JP *rel16*	A	N.S.	Valid	Jump near if parity (PF=1). Not supported in 64-bit mode.
0F 8A *cd*	JP *rel32*	A	Valid	Valid	Jump near if parity (PF=1).
0F 8A *cw*	JPE *rel16*	A	N.S.	Valid	Jump near if parity even (PF=1). Not supported in 64-bit mode.
0F 8A *cd*	JPE *rel32*	A	Valid	Valid	Jump near if parity even (PF=1).
0F 8B *cw*	JPO *rel16*	A	N.S.	Valid	Jump near if parity odd (PF=0). Not supported in 64-bit mode.
0F 8B *cd*	JPO *rel32*	A	Valid	Valid	Jump near if parity odd (PF=0).
0F 88 *cw*	JS *rel16*	A	N.S.	Valid	Jump near if sign (SF=1). Not supported in 64-bit mode.
0F 88 *cd*	JS *rel32*	A	Valid	Valid	Jump near if sign (SF=1).
0F 84 *cw*	JZ *rel16*	A	N.S.	Valid	Jump near if 0 (ZF=1). Not supported in 64-bit mode.
0F 84 *cd*	JZ *rel32*	A	Valid	Valid	Jump near if 0 (ZF=1).

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	Offset	NA	NA	NA

Description

Checks the state of one or more of the status flags in the EFLAGS register (CF, OF, PF, SF, and ZF) and, if the flags are in the specified state (condition), performs a jump to the target instruction specified by the destination operand. A condition code (*cc*) is associated with each instruction to indicate the condition being tested for. If the condition is not satisfied, the jump is not performed and execution continues with the instruction following the J*cc* instruction.

The target instruction is specified with a relative offset (a signed offset relative to the current value of the instruction pointer in the EIP register). A relative offset (*rel8*, *rel16,* or *rel32*) is generally specified as a label in assembly code, but at the machine code level, it is encoded as a signed, 8-bit or 32-bit immediate value, which is added to the instruction pointer. Instruction coding is most efficient for offsets of –128 to

+127. If the operand-size attribute is 16, the upper two bytes of the EIP register are cleared, resulting in a maximum instruction pointer size of 16 bits.

The conditions for each J*cc* mnemonic are given in the "Description" column of the table on the preceding page. The terms "less" and "greater" are used for comparisons of signed integers and the terms "above" and "below" are used for unsigned integers.

Because a particular state of the status flags can sometimes be interpreted in two ways, two mnemonics are defined for some opcodes. For example, the JA (jump if above) instruction and the JNBE (jump if not below or equal) instruction are alternate mnemonics for the opcode 77H.

The J*cc* instruction does not support far jumps (jumps to other code segments). When the target for the conditional jump is in a different segment, use the opposite condition from the condition being tested for the J*cc* instruction, and then access the target with an unconditional far jump (JMP instruction) to the other segment. For example, the following conditional far jump is illegal:

```
JZ FARLABEL;
```

To accomplish this far jump, use the following two instructions:

```
JNZ BEYOND;
JMP FARLABEL;
BEYOND:
```

The JRCXZ, JECXZ and JCXZ instructions differ from other J*cc* instructions because they do not check status flags. Instead, they check RCX, ECX or CX for 0. The register checked is determined by the address-size attribute. These instructions are useful when used at the beginning of a loop that terminates with a conditional loop instruction (such as LOOPNE). They can be used to prevent an instruction sequence from entering a loop when RCX, ECX or CX is 0. This would cause the loop to execute 2^{64}, 2^{32} or 64K times (not zero times).

All conditional jumps are converted to code fetches of one or two cache lines, regardless of jump address or cacheability.

In 64-bit mode, operand size is fixed at 64 bits. JMP Short is RIP = RIP + 8-bit offset sign extended to 64 bits. JMP Near is RIP = RIP + 32-bit offset sign extended to 64-bits.

Operation

```
IF condition
    THEN
        tempEIP ← EIP + SignExtend(DEST);
        IF OperandSize = 16
            THEN tempEIP ← tempEIP AND 0000FFFFH;
        FI;
    IF tempEIP is not within code segment limit
        THEN #GP(0);
```

```
        ELSE EIP ← tempEIP
    FI;
FI;
```

Protected Mode Exceptions

#GP(0)	If the offset being jumped to is beyond the limits of the CS segment.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If the offset being jumped to is beyond the limits of the CS segment or is outside of the effective address space from 0 to FFFFH. This condition can occur if a 32-bit address size override prefix is used.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(0)	If the memory address is in a non-canonical form.
#UD	If the LOCK prefix is used.

JMP—Jump

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
EB *cb*	JMP *rel8*	A	Valid	Valid	Jump short, RIP = RIP + 8-bit displacement sign extended to 64-bits
E9 *cw*	JMP *rel16*	A	N.S.	Valid	Jump near, relative, displacement relative to next instruction. Not supported in 64-bit mode.
E9 *cd*	JMP *rel32*	A	Valid	Valid	Jump near, relative, RIP = RIP + 32-bit displacement sign extended to 64-bits
FF /4	JMP *r/m16*	B	N.S.	Valid	Jump near, absolute indirect, address = zero-extended *r/m16*. Not supported in 64-bit mode.
FF /4	JMP *r/m32*	B	N.S.	Valid	Jump near, absolute indirect, address given in *r/m32*. Not supported in 64-bit mode.
FF /4	JMP *r/m64*	B	Valid	N.E.	Jump near, absolute indirect, RIP = 64-Bit offset from register or memory
EA *cd*	JMP *ptr16:16*	A	Inv.	Valid	Jump far, absolute, address given in operand
EA *cp*	JMP *ptr16:32*	A	Inv.	Valid	Jump far, absolute, address given in operand
FF /5	JMP *m16:16*	A	Valid	Valid	Jump far, absolute indirect, address given in *m16:16*
FF /5	JMP *m16:32*	A	Valid	Valid	Jump far, absolute indirect, address given in *m16:32.*
REX.W + FF /5	JMP *m16:64*	A	Valid	N.E.	Jump far, absolute indirect, address given in *m16:64.*

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	Offset	NA	NA	NA
B	ModRM:r/m (r)	NA	NA	NA

Description

Transfers program control to a different point in the instruction stream without recording return information. The destination (target) operand specifies the address of the instruction being jumped to. This operand can be an immediate value, a general-purpose register, or a memory location.

This instruction can be used to execute four different types of jumps:

- Near jump—A jump to an instruction within the current code segment (the segment currently pointed to by the CS register), sometimes referred to as an intrasegment jump.
- Short jump—A near jump where the jump range is limited to –128 to +127 from the current EIP value.
- Far jump—A jump to an instruction located in a different segment than the current code segment but at the same privilege level, sometimes referred to as an intersegment jump.
- Task switch—A jump to an instruction located in a different task.

A task switch can only be executed in protected mode (see Chapter 7, in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*, for information on performing task switches with the JMP instruction).

Near and Short Jumps. When executing a near jump, the processor jumps to the address (within the current code segment) that is specified with the target operand. The target operand specifies either an absolute offset (that is an offset from the base of the code segment) or a relative offset (a signed displacement relative to the current value of the instruction pointer in the EIP register). A near jump to a relative offset of 8-bits (*rel8*) is referred to as a short jump. The CS register is not changed on near and short jumps.

An absolute offset is specified indirectly in a general-purpose register or a memory location (*r/m16* or *r/m32*). The operand-size attribute determines the size of the target operand (16 or 32 bits). Absolute offsets are loaded directly into the EIP register. If the operand-size attribute is 16, the upper two bytes of the EIP register are cleared, resulting in a maximum instruction pointer size of 16 bits.

A relative offset (*rel8, rel16*, or *rel32*) is generally specified as a label in assembly code, but at the machine code level, it is encoded as a signed 8-, 16-, or 32-bit immediate value. This value is added to the value in the EIP register. (Here, the EIP register contains the address of the instruction following the JMP instruction). When using relative offsets, the opcode (for short vs. near jumps) and the operand-size attribute (for near relative jumps) determines the size of the target operand (8, 16, or 32 bits).

Far Jumps in Real-Address or Virtual-8086 Mode. When executing a far jump in real-address or virtual-8086 mode, the processor jumps to the code segment and offset specified with the target operand. Here the target operand specifies an absolute far address either directly with a pointer (*ptr16:16* or *ptr16:32*) or indirectly with a memory location (*m16:16* or *m16:32*). With the pointer method, the segment and address of the called procedure is encoded in the instruction, using a 4-byte (16-bit operand size) or 6-byte (32-bit operand size) far address immediate. With the indirect method, the target operand specifies a memory location that contains a 4-byte (16-bit operand size) or 6-byte (32-bit operand size) far address. The far address is loaded directly into the CS and EIP registers. If the operand-size attribute is 16, the upper two bytes of the EIP register are cleared.

Far Jumps in Protected Mode. When the processor is operating in protected mode, the JMP instruction can be used to perform the following three types of far jumps:

- A far jump to a conforming or non-conforming code segment.
- A far jump through a call gate.
- A task switch.

(The JMP instruction cannot be used to perform inter-privilege-level far jumps.)

In protected mode, the processor always uses the segment selector part of the far address to access the corresponding descriptor in the GDT or LDT. The descriptor type (code segment, call gate, task gate, or TSS) and access rights determine the type of jump to be performed.

If the selected descriptor is for a code segment, a far jump to a code segment at the same privilege level is performed. (If the selected code segment is at a different privilege level and the code segment is non-conforming, a general-protection exception is generated.) A far jump to the same privilege level in protected mode is very similar to one carried out in real-address or virtual-8086 mode. The target operand specifies an absolute far address either directly with a pointer (*ptr16:16* or *ptr16:32*) or indirectly with a memory location (*m16:16* or *m16:32*). The operand-size attribute determines the size of the offset (16 or 32 bits) in the far address. The new code segment selector and its descriptor are loaded into CS register, and the offset from the instruction is loaded into the EIP register. Note that a call gate (described in the next paragraph) can also be used to perform far call to a code segment at the same privilege level. Using this mechanism provides an extra level of indirection and is the preferred method of making jumps between 16-bit and 32-bit code segments.

When executing a far jump through a call gate, the segment selector specified by the target operand identifies the call gate. (The offset part of the target operand is ignored.) The processor then jumps to the code segment specified in the call gate descriptor and begins executing the instruction at the offset specified in the call gate. No stack switch occurs. Here again, the target operand can specify the far address of the call gate either directly with a pointer (*ptr16:16* or *ptr16:32*) or indirectly with a memory location (*m16:16* or *m16:32*).

Executing a task switch with the JMP instruction is somewhat similar to executing a jump through a call gate. Here the target operand specifies the segment selector of the task gate for the task being switched to (and the offset part of the target operand is ignored). The task gate in turn points to the TSS for the task, which contains the segment selectors for the task's code and stack segments. The TSS also contains the EIP value for the next instruction that was to be executed before the task was suspended. This instruction pointer value is loaded into the EIP register so that the task begins executing again at this next instruction.

The JMP instruction can also specify the segment selector of the TSS directly, which eliminates the indirection of the task gate. See Chapter 7 in *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*, for detailed information on the mechanics of a task switch.

Note that when you execute at task switch with a JMP instruction, the nested task flag (NT) is not set in the EFLAGS register and the new TSS's previous task link field is not loaded with the old task's TSS selector. A return to the previous task can thus not be carried out by executing the IRET instruction. Switching tasks with the JMP instruction differs in this regard from the CALL instruction which does set the NT flag and save the previous task link information, allowing a return to the calling task with an IRET instruction.

In 64-Bit Mode — The instruction's operation size is fixed at 64 bits. If a selector points to a gate, then RIP equals the 64-bit displacement taken from gate; else RIP equals the zero-extended offset from the far pointer referenced in the instruction.

See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
IF near jump
    IF 64-bit Mode
        THEN
            IF near relative jump
             THEN
                tempRIP ← RIP + DEST; (* RIP is instruction following JMP instruction*)
             ELSE (* Near absolute jump *)
                tempRIP ← DEST;
            FI;
        ELSE
            IF near relative jump
             THEN
                tempEIP ← EIP + DEST; (* EIP is instruction following JMP instruction*)
             ELSE (* Near absolute jump *)
                tempEIP ← DEST;
            FI;
    FI;
    IF (IA32_EFER.LMA = 0 or target mode = Compatibility mode)
    and tempEIP outside code segment limit
        THEN #GP(0); FI
    IF 64-bit mode and tempRIP is not canonical
        THEN #GP(0);
    FI;
    IF OperandSize = 32
         THEN
            EIP ← tempEIP;
```

```
            ELSE
                IF OperandSize = 16
                    THEN (* OperandSize = 16 *)
                        EIP ← tempEIP AND 0000FFFFH;
                    ELSE (* OperandSize = 64)
                        RIP ← tempRIP;
                FI;
        FI;
FI;
IF far jump and (PE = 0 or (PE = 1 AND VM = 1)) (* Real-address or virtual-8086 mode *)
    THEN
        tempEIP ← DEST(Offset); (* DEST is ptr16:32 or [m16:32] *)
        IF tempEIP is beyond code segment limit
            THEN #GP(0); FI;
        CS ← DEST(segment selector); (* DEST is ptr16:32 or [m16:32] *)
        IF OperandSize = 32
            THEN
                EIP ← tempEIP; (* DEST is ptr16:32 or [m16:32] *)
            ELSE (* OperandSize = 16 *)
                EIP ← tempEIP AND 0000FFFFH; (* Clear upper 16 bits *)
        FI;
FI;
IF far jump and (PE = 1 and VM = 0)
(* IA-32e mode or protected mode, not virtual-8086 mode *)
    THEN
        IF effective address in the CS, DS, ES, FS, GS, or SS segment is illegal
        or segment selector in target operand NULL
                THEN #GP(0); FI;
        IF segment selector index not within descriptor table limits
            THEN #GP(new selector); FI;
        Read type and access rights of segment descriptor;
        IF (EFER.LMA = 0)
            THEN
                IF segment type is not a conforming or nonconforming code
                segment, call gate, task gate, or TSS
                    THEN #GP(segment selector); FI;
            ELSE
                IF segment type is not a conforming or nonconforming code segment
                call gate
                    THEN #GP(segment selector); FI;
        FI;
        Depending on type and access rights:
            GO TO CONFORMING-CODE-SEGMENT;
            GO TO NONCONFORMING-CODE-SEGMENT;
```

```
                GO TO CALL-GATE;
                GO TO TASK-GATE;
                GO TO TASK-STATE-SEGMENT;
     ELSE
            #GP(segment selector);
FI;
CONFORMING-CODE-SEGMENT:
    IF L-Bit = 1 and D-BIT = 1 and IA32_EFER.LMA = 1
          THEN GP(new code segment selector); FI;
     IF DPL > CPL
          THEN #GP(segment selector); FI;
     IF segment not present
          THEN #NP(segment selector); FI;
    tempEIP ← DEST(Offset);
    IF OperandSize = 16
           THEN tempEIP ← tempEIP AND 0000FFFFH;
    FI;
    IF (IA32_EFER.LMA = 0 or target mode = Compatibility mode) and
    tempEIP outside code segment limit
          THEN #GP(0); FI
    IF tempEIP is non-canonical
          THEN #GP(0); FI;
    CS ← DEST[segment selector]; (* Segment descriptor information also loaded *)
    CS(RPL) ← CPL
    EIP ← tempEIP;
END;
NONCONFORMING-CODE-SEGMENT:
    IF L-Bit = 1 and D-BIT = 1 and IA32_EFER.LMA = 1
          THEN GP(new code segment selector); FI;
    IF (RPL > CPL) OR (DPL ≠ CPL)
          THEN #GP(code segment selector); FI;
    IF segment not present
          THEN #NP(segment selector); FI;
    tempEIP ← DEST(Offset);
    IF OperandSize = 16
           THEN tempEIP ← tempEIP AND 0000FFFFH; FI;
    IF (IA32_EFER.LMA = 0 OR target mode = Compatibility mode)
    and tempEIP outside code segment limit
          THEN #GP(0); FI
    IF tempEIP is non-canonical THEN #GP(0); FI;
    CS ← DEST[segment selector]; (* Segment descriptor information also loaded *)
    CS(RPL) ← CPL;
    EIP ← tempEIP;
END;
```

```
CALL-GATE:
    IF call gate DPL < CPL
    or call gate DPL < call gate segment-selector RPL
                THEN #GP(call gate selector); FI;
    IF call gate not present
          THEN #NP(call gate selector); FI;
    IF call gate code-segment selector is NULL
          THEN #GP(0); FI;
    IF call gate code-segment selector index outside descriptor table limits
          THEN #GP(code segment selector); FI;
    Read code segment descriptor;
    IF code-segment segment descriptor does not indicate a code segment
    or code-segment segment descriptor is conforming and DPL > CPL
    or code-segment segment descriptor is non-conforming and DPL ≠ CPL
                THEN #GP(code segment selector); FI;
    IF IA32_EFER.LMA = 1 and (code-segment descriptor is not a 64-bit code segment
    or code-segment segment descriptor has both L-Bit and D-bit set)
                THEN #GP(code segment selector); FI;
    IF code segment is not present
          THEN #NP(code-segment selector); FI;
     IF instruction pointer is not within code-segment limit
          THEN #GP(0); FI;
     tempEIP ← DEST(Offset);
     IF GateSize = 16
           THEN tempEIP ← tempEIP AND 0000FFFFH; FI;
    IF (IA32_EFER.LMA = 0 OR target mode = Compatibility mode) AND tempEIP
    outside code segment limit
          THEN #GP(0); FI
    CS ← DEST[SegmentSelector]; (* Segment descriptor information also loaded *)
    CS(RPL) ← CPL;
    EIP ← tempEIP;
END;
TASK-GATE:
    IF task gate DPL < CPL
    or task gate DPL < task gate segment-selector RPL
          THEN #GP(task gate selector); FI;
    IF task gate not present
          THEN #NP(gate selector); FI;
    Read the TSS segment selector in the task-gate descriptor;
    IF TSS segment selector local/global bit is set to local
    or index not within GDT limits
    or TSS descriptor specifies that the TSS is busy
          THEN #GP(TSS selector); FI;
```

```
    IF TSS not present
        THEN #NP(TSS selector); FI;
    SWITCH-TASKS to TSS;
    IF EIP not within code segment limit
        THEN #GP(0); FI;
END;
TASK-STATE-SEGMENT:
    IF TSS DPL < CPL
    or TSS DPL < TSS segment-selector RPL
    or TSS descriptor indicates TSS not available
        THEN #GP(TSS selector); FI;
    IF TSS is not present
        THEN #NP(TSS selector); FI;
    SWITCH-TASKS to TSS;
    IF EIP not within code segment limit
        THEN #GP(0); FI;
END;
```

Flags Affected

All flags are affected if a task switch occurs; no flags are affected if a task switch does not occur.

Protected Mode Exceptions

#GP(0)	If offset in target operand, call gate, or TSS is beyond the code segment limits.
	If the segment selector in the destination operand, call gate, task gate, or TSS is NULL.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#GP(selector)	If the segment selector index is outside descriptor table limits.
	If the segment descriptor pointed to by the segment selector in the destination operand is not for a conforming-code segment, nonconforming-code segment, call gate, task gate, or task state segment.
	If the DPL for a nonconforming-code segment is not equal to the CPL
	(When not using a call gate.) If the RPL for the segment's segment selector is greater than the CPL.
	If the DPL for a conforming-code segment is greater than the CPL.

	If the DPL from a call-gate, task-gate, or TSS segment descriptor is less than the CPL or than the RPL of the call-gate, task-gate, or TSS's segment selector.
	If the segment descriptor for selector in a call gate does not indicate it is a code segment.
	If the segment descriptor for the segment selector in a task gate does not indicate an available TSS.
	If the segment selector for a TSS has its local/global bit set for local.
	If a TSS segment descriptor specifies that the TSS is busy or not available.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NP (selector)	If the code segment being accessed is not present.
	If call gate, task gate, or TSS not present.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3. (Only occurs when fetching target from memory.)
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If the target operand is beyond the code segment limits.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made. (Only occurs when fetching target from memory.)
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same as 64-bit mode exceptions.

64-Bit Mode Exceptions

#GP(0)	If a memory address is non-canonical.
	If target offset in destination operand is non-canonical.
	If target offset in destination operand is beyond the new code segment limit.
	If the segment selector in the destination operand is NULL.
	If the code segment selector in the 64-bit gate is NULL.
#GP(selector)	If the code segment or 64-bit call gate is outside descriptor table limits.
	If the code segment or 64-bit call gate overlaps non-canonical space.
	If the segment descriptor from a 64-bit call gate is in non-canonical space.
	If the segment descriptor pointed to by the segment selector in the destination operand is not for a conforming-code segment, nonconforming-code segment, 64-bit call gate.
	If the segment descriptor pointed to by the segment selector in the destination operand is a code segment, and has both the D-bit and the L-bit set.
	If the DPL for a nonconforming-code segment is not equal to the CPL, or the RPL for the segment's segment selector is greater than the CPL.
	If the DPL for a conforming-code segment is greater than the CPL.
	If the DPL from a 64-bit call-gate is less than the CPL or than the RPL of the 64-bit call-gate.
	If the upper type field of a 64-bit call gate is not 0x0.
	If the segment selector from a 64-bit call gate is beyond the descriptor table limits.
	If the code segment descriptor pointed to by the selector in the 64-bit gate doesn't have the L-bit set and the D-bit clear.
	If the segment descriptor for a segment selector from the 64-bit call gate does not indicate it is a code segment.
	If the code segment is non-confirming and CPL $\neq$ DPL.
	If the code segment is confirming and CPL < DPL.
#NP(selector)	If a code segment or 64-bit call gate is not present.
#UD	(64-bit mode only) If a far jump is direct to an absolute address in memory.

	If the LOCK prefix is used.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

LAHF—Load Status Flags into AH Register

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
9F	LAHF	A	Invalid*	Valid	Load: AH ← EFLAGS(SF:ZF:0:AF:0:PF:1:CF).

NOTES:

*Valid in specific steppings. See Description section.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

This instruction executes as described above in compatibility mode and legacy mode. It is valid in 64-bit mode only if CPUID.80000001H:ECX.LAHF-SAHF[bit 0] = 1.

Operation

```
IF 64-Bit Mode
    THEN
        IF CPUID.80000001H:ECX.LAHF-SAHF[bit 0] = 1;
            THEN AH ← RFLAGS(SF:ZF:0:AF:0:PF:1:CF);
            ELSE #UD;
        FI;
    ELSE
        AH ← EFLAGS(SF:ZF:0:AF:0:PF:1:CF);
FI;
```

Flags Affected

None. The state of the flags in the EFLAGS register is not affected.

Protected Mode Exceptions

#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#UD	If CPUID.80000001H:ECX.LAHF-SAHF[bit 0] = 0.
	If the LOCK prefix is used.

LAR—Load Access Rights Byte

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 02 /r	LAR *r16, r16/m16*	A	Valid	Valid	*r16* ← *r16/m16* masked by FF00H.
0F 02 /r	LAR *r32, r32/m16*[1]	A	Valid	Valid	*r32* ← *r32/m16* masked by 00FxFF00H
REX.W + 0F 02 /r	LAR *r64, r32/m16*[1]	A	Valid	N.E.	*r64* ← *r32/m16* masked by 00FxFF00H and zero extended

NOTES:

1. For all loads (regardless of source or destination sizing) only bits 16-0 are used. Other bits are ignored.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Loads the access rights from the segment descriptor specified by the second operand (source operand) into the first operand (destination operand) and sets the ZF flag in the flag register. The source operand (which can be a register or a memory location) contains the segment selector for the segment descriptor being accessed. If the source operand is a memory address, only 16 bits of data are accessed. The destination operand is a general-purpose register.

The processor performs access checks as part of the loading process. Once loaded in the destination register, software can perform additional checks on the access rights information.

When the operand size is 32 bits, the access rights for a segment descriptor include the type and DPL fields and the S, P, AVL, D/B, and G flags, all of which are located in the second doubleword (bytes 4 through 7) of the segment descriptor. The doubleword is masked by 00FXFF00H before it is loaded into the destination operand. When the operand size is 16 bits, the access rights include the type and DPL fields. Here, the two lower-order bytes of the doubleword are masked by FF00H before being loaded into the destination operand.

This instruction performs the following checks before it loads the access rights in the destination register:

- Checks that the segment selector is not NULL.
- Checks that the segment selector points to a descriptor that is within the limits of the GDT or LDT being accessed

- Checks that the descriptor type is valid for this instruction. All code and data segment descriptors are valid for (can be accessed with) the LAR instruction. The valid system segment and gate descriptor types are given in Table 3-62.
- If the segment is not a conforming code segment, it checks that the specified segment descriptor is visible at the CPL (that is, if the CPL and the RPL of the segment selector are less than or equal to the DPL of the segment selector).

If the segment descriptor cannot be accessed or is an invalid type for the instruction, the ZF flag is cleared and no access rights are loaded in the destination operand.

The LAR instruction can only be executed in protected mode and IA-32e mode.

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.W prefix permits access to 64-bit registers as destination.

When the destination operand size is 64 bits, the access rights are loaded from the second doubleword (bytes 4 through 7) of the segment descriptor. The doubleword is masked by 00FXFF00H and zero extended to 64 bits before it is loaded into the destination operand.

Table 3-62. Segment and Gate Types

Type	Protected Mode		IA-32e Mode	
	Name	Valid	Name	Valid
0	Reserved	No	Reserved	No
1	Available 16-bit TSS	Yes	Reserved	No
2	LDT	Yes	LDT	No
3	Busy 16-bit TSS	Yes	Reserved	No
4	16-bit call gate	Yes	Reserved	No
5	16-bit/32-bit task gate	Yes	Reserved	No
6	16-bit interrupt gate	No	Reserved	No
7	16-bit trap gate	No	Reserved	No
8	Reserved	No	Reserved	No
9	Available 32-bit TSS	Yes	Available 64-bit TSS	Yes
A	Reserved	No	Reserved	No
B	Busy 32-bit TSS	Yes	Busy 64-bit TSS	Yes
C	32-bit call gate	Yes	64-bit call gate	Yes
D	Reserved	No	Reserved	No
E	32-bit interrupt gate	No	64-bit interrupt gate	No
F	32-bit trap gate	No	64-bit trap gate	No

Operation

```
IF Offset(SRC) > descriptor table limit
    THEN
        ZF = 0;
    ELSE
        IF SegmentDescriptor(Type) ≠ conforming code segment
        and (CPL > DPL) or (RPL > DPL)
        or segment type is not valid for instruction
            THEN
                ZF ← 0
            ELSE
                TEMP ← Read segment descriptor ;
                IF OperandSize = 64
                    THEN
                        DEST ← (ACCESSRIGHTWORD(TEMP) AND 00000000_00FxFF00H);
                    ELSE (* OperandSize = 32*)
                        DEST ← (ACCESSRIGHTWORD(TEMP) AND 00FxFF00H);
                    ELSE (* OperandSize = 16 *)
                        DEST ← (ACCESSRIGHTWORD(TEMP) AND FF00H);
                FI;
        FI;
FI;
```

Flags Affected

The ZF flag is set to 1 if the access rights are loaded successfully; otherwise, it is set to 0.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and the memory operand effective address is unaligned while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#UD	The LAR instruction is not recognized in real-address mode.

Virtual-8086 Mode Exceptions

#UD	The LAR instruction cannot be executed in virtual-8086 mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If the memory operand effective address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory operand effective address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and the memory operand effective address is unaligned while the current privilege level is 3.
#UD	If the LOCK prefix is used.

LDDQU—Load Unaligned Integer 128 Bits

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F F0 /*r* LDDQU *xmm1*, *mem*	A	V/V	SSE3	Load unaligned data from *mem* and return double quadword in *xmm1*.
VEX.128.F2.0F.WIG F0 /r VLDDQU xmm1, m128	A	V/V	AVX	Load unaligned packed integer values from mem to xmm1.
VEX.256.F2.0F.WIG F0 /r VLDDQU ymm1, m256	A	V/V	AVX	Load unaligned packed integer values from mem to ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

The instruction is *functionally similar* to (V)MOVDQU ymm/xmm, m256/m128 for loading from memory. That is: 32/16 bytes of data starting at an address specified by the source memory operand (second operand) are fetched from memory and placed in a destination register (first operand). The source operand need not be aligned on a 32/16-byte boundary. Up to 64/32 bytes may be loaded from memory; this is implementation dependent.

This instruction may improve performance relative to (V)MOVDQU if the source operand crosses a cache line boundary. In situations that require the data loaded by (V)LDDQU be modified and stored to the same location, use (V)MOVDQU or (V)MOVDQA instead of (V)LDDQU. To move a double quadword to or from memory locations that are known to be aligned on 16-byte boundaries, use the (V)MOVDQA instruction.

Implementation Notes

- If the source is aligned to a 32/16-byte boundary, based on the implementation, the 32/16 bytes may be loaded more than once. For that reason, the usage of (V)LDDQU should be avoided when using uncached or write-combining (WC) memory regions. For uncached or WC memory regions, keep using (V)MOVDQU.
- This instruction is a replacement for (V)MOVDQU (load) in situations where cache line splits significantly affect performance. It should not be used in situations where store-load forwarding is performance critical. If performance of store-load forwarding is critical to the application, use (V)MOVDQA store-load pairs when

data is 256/128-bit aligned or (V)MOVDQU store-load pairs when data is 256/128-bit unaligned.

- If the memory address is not aligned on 32/16-byte boundary, some implementations may load up to 64/32 bytes and return 32/16 bytes in the destination. Some processor implementations may issue multiple loads to access the appropriate 32/16 bytes. Developers of multi-threaded or multi-processor software should be aware that on these processors the loads will be performed in a non-atomic way.
- If alignment checking is enabled (CR0.AM = 1, RFLAGS.AC = 1, and CPL = 3), an alignment-check exception (#AC) may or may not be generated (depending on processor implementation) when the memory address is not aligned on an 8-byte boundary.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

Operation

LDDQU (128-bit Legacy SSE version)

DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] (Unmodified)

VLDDQU (VEX.128 encoded version)

DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] ← 0

VLDDQU (VEX.256 encoded version)

DEST[255:0] ← SRC[255:0]

Intel C/C++ Compiler Intrinsic Equivalent

LDDQU __m128i _mm_lddqu_si128 (__m128i * p);

LDDQU __m256i _mm256_lddqu_si256 (__m256i * p);

Numeric Exceptions

None.

Other Exceptions

See Exceptions Type 4;

Note treatment of #AC varies.

LDMXCSR—Load MXCSR Register

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F,AE,/2 LDMXCSR *m32*	A	V/V	SSE	Load MXCSR register from *m32*.
VEX.LZ.0F.WIG AE /2 VLDMXCSR *m32*	A	V/V	AVX	Load MXCSR register from *m32.*

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r)	NA	NA	NA

Description

Loads the source operand into the MXCSR control/status register. The source operand is a 32-bit memory location. See “MXCSR Control and Status Register” in Chapter 10, of the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*, for a description of the MXCSR register and its contents.

The LDMXCSR instruction is typically used in conjunction with the (V)STMXCSR instruction, which stores the contents of the MXCSR register in memory.

The default MXCSR value at reset is 1F80H.

If a (V)LDMXCSR instruction clears a SIMD floating-point exception mask bit and sets the corresponding exception flag bit, a SIMD floating-point exception will not be immediately generated. The exception will be generated only upon the execution of the next instruction that meets both conditions below:

- the instruction must operate on an XMM or YMM register operand,
- the instruction causes that particular SIMD floating-point exception to be reported.

This instruction’s operation is the same in non-64-bit modes and 64-bit mode.

If VLDMXCSR is encoded with VEX.L= 1, an attempt to execute the instruction encoded with VEX.L= 1 will cause an #UD exception.

Operation

MXCSR ← m32;

C/C++ Compiler Intrinsic Equivalent

_mm_setcsr(unsigned int i)

Numeric Exceptions

None.

Other Exceptions

See Exceptions Type 9; additionally

#GP	For an attempt to set reserved bits in MXCSR.

LDS/LES/LFS/LGS/LSS—Load Far Pointer

Opcode	Instruction	Op/En	64-Bit Mode	Compat/Leg Mode	Description
C5 /*r*	LDS *r16,m16:16*	A	Invalid	Valid	Load DS:*r16* with far pointer from memory.
C5 /*r*	LDS *r32,m16:32*	A	Invalid	Valid	Load DS:*r32* with far pointer from memory.
0F B2 /*r*	LSS *r16,m16:16*	A	Valid	Valid	Load SS:*r16* with far pointer from memory.
0F B2 /*r*	LSS *r32,m16:32*	A	Valid	Valid	Load SS:*r32* with far pointer from memory.
REX + 0F B2 /*r*	LSS *r64,m16:64*	A	Valid	N.E.	Load SS:*r64* with far pointer from memory.
C4 /*r*	LES *r16,m16:16*	A	Invalid	Valid	Load ES:*r16* with far pointer from memory.
C4 /*r*	LES *r32,m16:32*	A	Invalid	Valid	Load ES:*r32* with far pointer from memory.
0F B4 /*r*	LFS *r16,m16:16*	A	Valid	Valid	Load FS:*r16* with far pointer from memory.
0F B4 /*r*	LFS *r32,m16:32*	A	Valid	Valid	Load FS:*r32* with far pointer from memory.
REX + 0F B4 /*r*	LFS *r64,m16:64*	A	Valid	N.E.	Load FS:*r64* with far pointer from memory.
0F B5 /*r*	LGS *r16,m16:16*	A	Valid	Valid	Load GS:*r16* with far pointer from memory.
0F B5 /*r*	LGS *r32,m16:32*	A	Valid	Valid	Load GS:*r32* with far pointer from memory.
REX + 0F B5 /*r*	LGS *r64,m16:64*	A	Valid	N.E.	Load GS:*r64* with far pointer from memory.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Loads a far pointer (segment selector and offset) from the second operand (source operand) into a segment register and the first operand (destination operand). The source operand specifies a 48-bit or a 32-bit pointer in memory depending on the current setting of the operand-size attribute (32 bits or 16 bits, respectively). The instruction opcode and the destination operand specify a segment register/general-purpose register pair. The 16-bit segment selector from the source operand is loaded into the segment register specified with the opcode (DS, SS, ES, FS, or GS). The 32-bit or 16-bit offset is loaded into the register specified with the destination operand.

If one of these instructions is executed in protected mode, additional information from the segment descriptor pointed to by the segment selector in the source operand is loaded in the hidden part of the selected segment register.

Also in protected mode, a NULL selector (values 0000 through 0003) can be loaded into DS, ES, FS, or GS registers without causing a protection exception. (Any subsequent reference to a segment whose corresponding segment register is loaded with a NULL selector, causes a general-protection exception (#GP) and no memory reference to the segment occurs.)

In 64-bit mode, the instruction's default operation size is 32 bits. Using a REX prefix in the form of REX.W promotes operation to specify a source operand referencing an 80-bit pointer (16-bit selector, 64-bit offset) in memory. Using a REX prefix in the form of REX.R permits access to additional registers (R8-R15). See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
64-BIT_MODE
   IF SS is loaded
       THEN
           IF SegmentSelector = NULL and ( (RPL = 3) or
                   (RPL ≠ 3 and RPL ≠ CPL) )
               THEN #GP(0);
           ELSE IF descriptor is in non-canonical space
               THEN #GP(0); FI;
           ELSE IF Segment selector index is not within descriptor table limits
                   or segment selector RPL ≠ CPL
                   or access rights indicate nonwritable data segment
                   or DPL ≠ CPL
               THEN #GP(selector); FI;
           ELSE IF Segment marked not present
               THEN #SS(selector); FI;
           FI;
           SS ← SegmentSelector(SRC);
           SS ← SegmentDescriptor([SRC]);
   ELSE IF attempt to load DS, or ES
       THEN #UD;
   ELSE IF FS, or GS is loaded with non-NULL segment selector
       THEN IF Segment selector index is not within descriptor table limits
           or access rights indicate segment neither data nor readable code segment
```

```
                or segment is data or nonconforming-code segment
                and ( RPL > DPL or CPL > DPL)
                    THEN #GP(selector); FI;
                ELSE IF Segment marked not present
                    THEN #NP(selector); FI;
                FI;
                SegmentRegister ← SegmentSelector(SRC) ;
                SegmentRegister ← SegmentDescriptor([SRC]);
        FI;
    ELSE IF FS, or GS is loaded with a NULL selector:
        THEN
            SegmentRegister ← NULLSelector;
            SegmentRegister(DescriptorValidBit) ← 0; FI; (* Hidden flag;
                not accessible by software *)
    FI;
    DEST ← Offset(SRC);

PREOTECTED MODE OR COMPATIBILITY MODE;
    IF SS is loaded
        THEN
            IF SegementSelector = NULL
                THEN #GP(0);
            ELSE IF Segment selector index is not within descriptor table limits
                    or segment selector RPL ≠ CPL
                    or access rights indicate nonwritable data segment
                    or DPL ≠ CPL
                THEN #GP(selector); FI;
            ELSE IF Segment marked not present
                THEN #SS(selector); FI;
            FI;
            SS ← SegmentSelector(SRC);
            SS ← SegmentDescriptor([SRC]);
    ELSE IF DS, ES, FS, or GS is loaded with non-NULL segment selector
        THEN IF Segment selector index is not within descriptor table limits
            or access rights indicate segment neither data nor readable code segment
            or segment is data or nonconforming-code segment
            and (RPL > DPL or CPL > DPL)
                THEN #GP(selector); FI;
            ELSE IF Segment marked not present
                THEN #NP(selector); FI;
            FI;
            SegmentRegister ← SegmentSelector(SRC) AND RPL;
            SegmentRegister ← SegmentDescriptor([SRC]);
        FI;
```

```
    ELSE IF DS, ES, FS, or GS is loaded with a NULL selector:
        THEN
            SegmentRegister ← NULLSelector;
            SegmentRegister(DescriptorValidBit) ← 0; FI; (* Hidden flag;
                not accessible by software *)
    FI;
    DEST ← Offset(SRC);

Real-Address or Virtual-8086 Mode
    SegmentRegister ← SegmentSelector(SRC); FI;
    DEST ← Offset(SRC);
```

Flags Affected

None.

Protected Mode Exceptions

#UD	If source operand is not a memory location. If the LOCK prefix is used.
#GP(0)	If a NULL selector is loaded into the SS register. If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit. If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#GP(selector)	If the SS register is being loaded and any of the following is true: the segment selector index is not within the descriptor table limits, the segment selector RPL is not equal to CPL, the segment is a non-writable data segment, or DPL is not equal to CPL. If the DS, ES, FS, or GS register is being loaded with a non-NULL segment selector and any of the following is true: the segment selector index is not within descriptor table limits, the segment is neither a data nor a readable code segment, or the segment is a data or nonconforming-code segment and both RPL and CPL are greater than DPL.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#SS(selector)	If the SS register is being loaded and the segment is marked not present.
#NP(selector)	If DS, ES, FS, or GS register is being loaded with a non-NULL segment selector and the segment is marked not present.
#PF(fault-code)	If a page fault occurs.

#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If source operand is not a memory location.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#UD	If source operand is not a memory location.
	If the LOCK prefix is used.
#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(0)	If the memory address is in a non-canonical form.
	If a NULL selector is attempted to be loaded into the SS register in compatibility mode.
	If a NULL selector is attempted to be loaded into the SS register in CPL3 and 64-bit mode.
	If a NULL selector is attempted to be loaded into the SS register in non-CPL3 and 64-bit mode where its RPL is not equal to CPL.
#GP(Selector)	If the FS, or GS register is being loaded with a non-NULL segment selector and any of the following is true: the segment selector index is not within descriptor table limits, the memory address of the descriptor is non-canonical, the segment is neither a data nor a readable code segment, or the segment is a data or nonconforming-code segment and both RPL and CPL are greater than DPL.

	If the SS register is being loaded and any of the following is true: the segment selector index is not within the descriptor table limits, the memory address of the descriptor is non-canonical, the segment selector RPL is not equal to CPL, the segment is a nonwritable data segment, or DPL is not equal to CPL.
#SS(0)	If a memory operand effective address is non-canonical
#SS(Selector)	If the SS register is being loaded and the segment is marked not present.
#NP(selector)	If FS, or GS register is being loaded with a non-NULL segment selector and the segment is marked not present.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If source operand is not a memory location.
	If the LOCK prefix is used.

LEA—Load Effective Address

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
8D /*r*	LEA *r16,m*	A	Valid	Valid	Store effective address for *m* in register *r16.*
8D /*r*	LEA *r32,m*	A	Valid	Valid	Store effective address for *m* in register *r32.*
REX.W + 8D /*r*	LEA *r64,m*	A	Valid	N.E.	Store effective address for *m* in register *r64.*

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Computes the effective address of the second operand (the source operand) and stores it in the first operand (destination operand). The source operand is a memory address (offset part) specified with one of the processors addressing modes; the destination operand is a general-purpose register. The address-size and operand-size attributes affect the action performed by this instruction, as shown in the following table. The operand-size attribute of the instruction is determined by the chosen register; the address-size attribute is determined by the attribute of the code segment.

Table 3-63. Non-64-bit Mode LEA Operation with Address and Operand Size Attributes

Operand Size	Address Size	Action Performed
16	16	16-bit effective address is calculated and stored in requested 16-bit register destination.
16	32	32-bit effective address is calculated. The lower 16 bits of the address are stored in the requested 16-bit register destination.
32	16	16-bit effective address is calculated. The 16-bit address is zero-extended and stored in the requested 32-bit register destination.
32	32	32-bit effective address is calculated and stored in the requested 32-bit register destination.

Different assemblers may use different algorithms based on the size attribute and symbolic reference of the source operand.

In 64-bit mode, the instruction's destination operand is governed by operand size attribute, the default operand size is 32 bits. Address calculation is governed by address size attribute, the default address size is 64-bits. In 64-bit mode, address size of 16 bits is not encodable. See Table 3-64.

Table 3-64. 64-bit Mode LEA Operation with Address and Operand Size Attributes

Operand Size	Address Size	Action Performed
16	32	32-bit effective address is calculated (using 67H prefix). The lower 16 bits of the address are stored in the requested 16-bit register destination (using 66H prefix).
16	64	64-bit effective address is calculated (default address size). The lower 16 bits of the address are stored in the requested 16-bit register destination (using 66H prefix).
32	32	32-bit effective address is calculated (using 67H prefix) and stored in the requested 32-bit register destination.
32	64	64-bit effective address is calculated (default address size) and the lower 32 bits of the address are stored in the requested 32-bit register destination.
64	32	32-bit effective address is calculated (using 67H prefix), zero-extended to 64-bits, and stored in the requested 64-bit register destination (using REX.W).
64	64	64-bit effective address is calculated (default address size) and all 64-bits of the address are stored in the requested 64-bit register destination (using REX.W).

Operation

```
IF OperandSize = 16 and AddressSize = 16
    THEN
        DEST ← EffectiveAddress(SRC); (* 16-bit address *)
    ELSE IF OperandSize = 16 and AddressSize = 32
        THEN
            temp ← EffectiveAddress(SRC); (* 32-bit address *)
            DEST ← temp[0:15]; (* 16-bit address *)
        FI;
    ELSE IF OperandSize = 32 and AddressSize = 16
        THEN
            temp ← EffectiveAddress(SRC); (* 16-bit address *)
            DEST ← ZeroExtend(temp); (* 32-bit address *)
        FI;
    ELSE IF OperandSize = 32 and AddressSize = 32
        THEN
            DEST ← EffectiveAddress(SRC); (* 32-bit address *)
```

```
        FI;
    ELSE IF OperandSize = 16 and AddressSize = 64
        THEN
            temp ← EffectiveAddress(SRC); (* 64-bit address *)
            DEST ← temp[0:15]; (* 16-bit address *)
        FI;
    ELSE IF OperandSize = 32 and AddressSize = 64
        THEN
            temp ← EffectiveAddress(SRC); (* 64-bit address *)
            DEST ← temp[0:31]; (* 16-bit address *)
        FI;
    ELSE IF OperandSize = 64 and AddressSize = 64
        THEN
            DEST ← EffectiveAddress(SRC); (* 64-bit address *)
        FI;
FI;
```

Flags Affected

None.

Protected Mode Exceptions

#UD	If source operand is not a memory location.
	If the LOCK prefix is used.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

LEAVE—High Level Procedure Exit

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
C9	LEAVE	A	Valid	Valid	Set SP to BP, then pop BP.
C9	LEAVE	A	N.E.	Valid	Set ESP to EBP, then pop EBP.
C9	LEAVE	A	Valid	N.E.	Set RSP to RBP, then pop RBP.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Releases the stack frame set up by an earlier ENTER instruction. The LEAVE instruction copies the frame pointer (in the EBP register) into the stack pointer register (ESP), which releases the stack space allocated to the stack frame. The old frame pointer (the frame pointer for the calling procedure that was saved by the ENTER instruction) is then popped from the stack into the EBP register, restoring the calling procedure's stack frame.

A RET instruction is commonly executed following a LEAVE instruction to return program control to the calling procedure.

See "Procedure Calls for Block-Structured Languages" in Chapter 7 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for detailed information on the use of the ENTER and LEAVE instructions.

In 64-bit mode, the instruction's default operation size is 64 bits; 32-bit operation cannot be encoded. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
IF StackAddressSize = 32
    THEN
        ESP ← EBP;
    ELSE IF StackAddressSize = 64
        THEN RSP ← RBP; FI;
    ELSE IF StackAddressSize = 16
        THEN SP ← BP; FI;
FI;

IF OperandSize = 32
```

```
    THEN EBP ← Pop();
    ELSE IF OperandSize = 64
        THEN RBP ← Pop(); FI;
    ELSE IF OperandSize = 16
        THEN BP ← Pop(); FI;
FI;
```

Flags Affected

None.

Protected Mode Exceptions

#SS(0)	If the EBP register points to a location that is not within the limits of the current stack segment.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If the EBP register points to a location outside of the effective address space from 0 to FFFFH.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If the EBP register points to a location outside of the effective address space from 0 to FFFFH.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If the stack address is in a non-canonical form.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

LFENCE—Load Fence

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F AE /5	LFENCE	A	Valid	Valid	Serializes load operations.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Performs a serializing operation on all load-from-memory instructions that were issued prior the LFENCE instruction. Specifically, LFENCE does not execute until all prior instructions have completed locally, and no later instruction begins execution until LFENCE completes. In particular, an instruction that loads from memory and that precedes an LFENCE receives data from memory prior to completion of the LFENCE. (An LFENCE that follows an instruction that stores to memory might complete **before** the data being stored have become globally visible.) Instructions following an LFENCE may be fetched from memory before the LFENCE, but they will not execute until the LFENCE completes.

Weakly ordered memory types can be used to achieve higher processor performance through such techniques as out-of-order issue and speculative reads. The degree to which a consumer of data recognizes or knows that the data is weakly ordered varies among applications and may be unknown to the producer of this data. The LFENCE instruction provides a performance-efficient way of ensuring load ordering between routines that produce weakly-ordered results and routines that consume that data.

Processors are free to fetch and cache data speculatively from regions of system memory that use the WB, WC, and WT memory types. This speculative fetching can occur at any time and is not tied to instruction execution. Thus, it is not ordered with respect to executions of the LFENCE instruction; data can be brought into the caches speculatively just before, during, or after the execution of an LFENCE instruction.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
Wait_On_Following_Instructions_Until(preceding_instructions_complete);
```

Intel C/C++ Compiler Intrinsic Equivalent

```
void _mm_lfence(void)
```

Exceptions (All Modes of Operation)

#UD	If CPUID.01H:EDX.SSE2[bit 26] = 0.

If the LOCK prefix is used.

LGDT/LIDT—Load Global/Interrupt Descriptor Table Register

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 01 /2	LGDT *m16&32*	A	N.E.	Valid	Load *m* into GDTR.
0F 01 /3	LIDT *m16&32*	A	N.E.	Valid	Load *m* into IDTR.
0F 01 /2	LGDT *m16&64*	A	Valid	N.E.	Load *m* into GDTR.
0F 01 /3	LIDT *m16&64*	A	Valid	N.E.	Load *m* into IDTR.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r)	NA	NA	NA

Description

Loads the values in the source operand into the global descriptor table register (GDTR) or the interrupt descriptor table register (IDTR). The source operand specifies a 6-byte memory location that contains the base address (a linear address) and the limit (size of table in bytes) of the global descriptor table (GDT) or the interrupt descriptor table (IDT). If operand-size attribute is 32 bits, a 16-bit limit (lower 2 bytes of the 6-byte data operand) and a 32-bit base address (upper 4 bytes of the data operand) are loaded into the register. If the operand-size attribute is 16 bits, a 16-bit limit (lower 2 bytes) and a 24-bit base address (third, fourth, and fifth byte) are loaded. Here, the high-order byte of the operand is not used and the high-order byte of the base address in the GDTR or IDTR is filled with zeros.

The LGDT and LIDT instructions are used only in operating-system software; they are not used in application programs. They are the only instructions that directly load a linear address (that is, not a segment-relative address) and a limit in protected mode. They are commonly executed in real-address mode to allow processor initialization prior to switching to protected mode.

In 64-bit mode, the instruction's operand size is fixed at 8+2 bytes (an 8-byte base and a 2-byte limit). See the summary chart at the beginning of this section for encoding data and limits.

See "SGDT—Store Global Descriptor Table Register" in Chapter 4, *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 2B*, for information on storing the contents of the GDTR and IDTR.

Operation

```
IF Instruction is LIDT
   THEN
        IF OperandSize = 16
             THEN
```

```
                IDTR(Limit) ← SRC[0:15];
                IDTR(Base) ← SRC[16:47] AND 00FFFFFFH;
            ELSE IF 32-bit Operand Size
                THEN
                    IDTR(Limit) ← SRC[0:15];
                    IDTR(Base) ← SRC[16:47];
                FI;
            ELSE IF 64-bit Operand Size (* In 64-Bit Mode *)
                THEN
                    IDTR(Limit) ← SRC[0:15];
                    IDTR(Base) ← SRC[16:79];
                FI;
        FI;
    ELSE (* Instruction is LGDT *)
        IF OperandSize = 16
            THEN
                GDTR(Limit) ← SRC[0:15];
                GDTR(Base) ← SRC[16:47] AND 00FFFFFFH;
            ELSE IF 32-bit Operand Size
                THEN
                    GDTR(Limit) ← SRC[0:15];
                    GDTR(Base) ← SRC[16:47];
                FI;
            ELSE IF 64-bit Operand Size (* In 64-Bit Mode *)
                THEN
                    GDTR(Limit) ← SRC[0:15];
                    GDTR(Base) ← SRC[16:79];
                FI;
        FI;
FI;
```

Flags Affected

None.

Protected Mode Exceptions

#UD	If source operand is not a memory location.
	If the LOCK prefix is used.
#GP(0)	If the current privilege level is not 0.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.

#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.

Real-Address Mode Exceptions

#UD	If source operand is not a memory location.
	If the LOCK prefix is used.
#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.

Virtual-8086 Mode Exceptions

#UD	If source operand is not a memory location.
	If the LOCK prefix is used.
#GP(0)	The LGDT and LIDT instructions are not recognized in virtual-8086 mode.
#GP	If the current privilege level is not 0.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the current privilege level is not 0.
	If the memory address is in a non-canonical form.
#UD	If source operand is not a memory location.
	If the LOCK prefix is used.
#PF(fault-code)	If a page fault occurs.

LLDT—Load Local Descriptor Table Register

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 00 /2	LLDT *r/m*16	A	Valid	Valid	Load segment selector *r/m*16 into LDTR.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r)	NA	NA	NA

Description

Loads the source operand into the segment selector field of the local descriptor table register (LDTR). The source operand (a general-purpose register or a memory location) contains a segment selector that points to a local descriptor table (LDT). After the segment selector is loaded in the LDTR, the processor uses the segment selector to locate the segment descriptor for the LDT in the global descriptor table (GDT). It then loads the segment limit and base address for the LDT from the segment descriptor into the LDTR. The segment registers DS, ES, SS, FS, GS, and CS are not affected by this instruction, nor is the LDTR field in the task state segment (TSS) for the current task.

If bits 2-15 of the source operand are 0, LDTR is marked invalid and the LLDT instruction completes silently. However, all subsequent references to descriptors in the LDT (except by the LAR, VERR, VERW or LSL instructions) cause a general protection exception (#GP).

The operand-size attribute has no effect on this instruction.

The LLDT instruction is provided for use in operating-system software; it should not be used in application programs. This instruction can only be executed in protected mode or 64-bit mode.

In 64-bit mode, the operand size is fixed at 16 bits.

Operation

```
IF SRC(Offset) > descriptor table limit
    THEN #GP(segment selector); FI;

IF segment selector is valid

    Read segment descriptor;

    IF SegmentDescriptor(Type) ≠ LDT
        THEN #GP(segment selector); FI;
```

```
    IF segment descriptor is not present
        THEN #NP(segment selector); FI;

    LDTR(SegmentSelector) ← SRC;
    LDTR(SegmentDescriptor) ← GDTSegmentDescriptor;
ELSE LDTR ← INVALID
FI;
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If the current privilege level is not 0.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#GP(selector)	If the selector operand does not point into the Global Descriptor Table or if the entry in the GDT is not a Local Descriptor Table.
	Segment selector is beyond GDT limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#NP(selector)	If the LDT descriptor is not present.
#PF(fault-code)	If a page fault occurs.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#UD	The LLDT instruction is not recognized in real-address mode.

Virtual-8086 Mode Exceptions

#UD	The LLDT instruction is not recognized in virtual-8086 mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the current privilege level is not 0.
	If the memory address is in a non-canonical form.

#GP(selector)	If the selector operand does not point into the Global Descriptor Table or if the entry in the GDT is not a Local Descriptor Table.
	Segment selector is beyond GDT limit.
#NP(selector)	If the LDT descriptor is not present.
#PF(fault-code)	If a page fault occurs.
#UD	If the LOCK prefix is used.

LMSW—Load Machine Status Word

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 01 /6	LMSW *r/m*16	A	Valid	Valid	Loads *r/m*16 in machine status word of CR0.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r)	NA	NA	NA

Description

Loads the source operand into the machine status word, bits 0 through 15 of register CR0. The source operand can be a 16-bit general-purpose register or a memory location. Only the low-order 4 bits of the source operand (which contains the PE, MP, EM, and TS flags) are loaded into CR0. The PG, CD, NW, AM, WP, NE, and ET flags of CR0 are not affected. The operand-size attribute has no effect on this instruction.

If the PE flag of the source operand (bit 0) is set to 1, the instruction causes the processor to switch to protected mode. While in protected mode, the LMSW instruction cannot be used to clear the PE flag and force a switch back to real-address mode.

The LMSW instruction is provided for use in operating-system software; it should not be used in application programs. In protected or virtual-8086 mode, it can only be executed at CPL 0.

This instruction is provided for compatibility with the Intel 286 processor; programs and procedures intended to run on the Pentium 4, Intel Xeon, P6 family, Pentium, Intel486, and Intel386 processors should use the MOV (control registers) instruction to load the whole CR0 register. The MOV CR0 instruction can be used to set and clear the PE flag in CR0, allowing a procedure or program to switch between protected and real-address modes.

This instruction is a serializing instruction.

This instruction’s operation is the same in non-64-bit modes and 64-bit mode. Note that the operand size is fixed at 16 bits.

See “Changes to Instruction Behavior in VMX Non-Root Operation” in Chapter 22 of the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 3B*, for more information about the behavior of this instruction in VMX non-root operation.

Operation

```
CR0[0:3] ← SRC[0:3];
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If the current privilege level is not 0.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the current privilege level is not 0.
	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#UD	If the LOCK prefix is used.

LOCK—Assert LOCK# Signal Prefix

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
F0	LOCK	A	Valid	Valid	Asserts LOCK# signal for duration of the accompanying instruction.

NOTES:

* See IA-32 Architecture Compatibility section below.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Causes the processor's LOCK# signal to be asserted during execution of the accompanying instruction (turns the instruction into an atomic instruction). In a multiprocessor environment, the LOCK# signal ensures that the processor has exclusive use of any shared memory while the signal is asserted.

Note that, in later Intel 64 and IA-32 processors (including the Pentium 4, Intel Xeon, and P6 family processors), locking may occur without the LOCK# signal being asserted. See the "IA-32 Architecture Compatibility" section below.

The LOCK prefix can be prepended only to the following instructions and only to those forms of the instructions where the destination operand is a memory operand: ADD, ADC, AND, BTC, BTR, BTS, CMPXCHG, CMPXCH8B, DEC, INC, NEG, NOT, OR, SBB, SUB, XOR, XADD, and XCHG. If the LOCK prefix is used with one of these instructions and the source operand is a memory operand, an undefined opcode exception (#UD) may be generated. An undefined opcode exception will also be generated if the LOCK prefix is used with any instruction not in the above list. The XCHG instruction always asserts the LOCK# signal regardless of the presence or absence of the LOCK prefix.

The LOCK prefix is typically used with the BTS instruction to perform a read-modify-write operation on a memory location in shared memory environment.

The integrity of the LOCK prefix is not affected by the alignment of the memory field. Memory locking is observed for arbitrarily misaligned fields.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

IA-32 Architecture Compatibility

Beginning with the P6 family processors, when the LOCK prefix is prefixed to an instruction and the memory area being accessed is cached internally in the processor, the LOCK# signal is generally not asserted. Instead, only the processor's cache is locked. Here, the processor's cache coherency mechanism ensures that the operation is carried out atomically with regards to memory. See “Effects of a Locked Operation on Internal Processor Caches” in Chapter 8 of *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 3A*, the for more information on locking of caches.

Operation

```
AssertLOCK#(DurationOfAccompaningInstruction);
```

Flags Affected

None.

Protected Mode Exceptions

#UD	If the LOCK prefix is used with an instruction not listed: ADD, ADC, AND, BTC, BTR, BTS, CMPXCHG, CMPXCH8B, DEC, INC, NEG, NOT, OR, SBB, SUB, XOR, XADD, XCHG.
	Other exceptions can be generated by the instruction when the LOCK prefix is applied.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

LODS/LODSB/LODSW/LODSD/LODSQ—Load String

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
AC	LODS *m8*	A	Valid	Valid	For legacy mode, Load byte at address DS:(E)SI into AL. For 64-bit mode load byte at address (R)SI into AL.
AD	LODS *m16*	A	Valid	Valid	For legacy mode, Load word at address DS:(E)SI into AX. For 64-bit mode load word at address (R)SI into AX.
AD	LODS m32	A	Valid	Valid	For legacy mode, Load dword at address DS:(E)SI into EAX. For 64-bit mode load dword at address (R)SI into EAX.
REX.W + AD	LODS *m64*	A	Valid	N.E.	Load qword at address (R)SI into RAX.
AC	LODSB	A	Valid	Valid	For legacy mode, Load byte at address DS:(E)SI into AL. For 64-bit mode load byte at address (R)SI into AL.
AD	LODSW	A	Valid	Valid	For legacy mode, Load word at address DS:(E)SI into AX. For 64-bit mode load word at address (R)SI into AX.
AD	LODSD	A	Valid	Valid	For legacy mode, Load dword at address DS:(E)SI into EAX. For 64-bit mode load dword at address (R)SI into EAX.
REX.W + AD	LODSQ	A	Valid	N.E.	Load qword at address (R)SI into RAX.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Loads a byte, word, or doubleword from the source operand into the AL, AX, or EAX register, respectively. The source operand is a memory location, the address of which is read from the DS:ESI or the DS:SI registers (depending on the address-size attribute of the instruction, 32 or 16, respectively). The DS segment may be overridden with a segment override prefix.

At the assembly-code level, two forms of this instruction are allowed: the "explicit-operands" form and the "no-operands" form. The explicit-operands form (specified with the LODS mnemonic) allows the source operand to be specified explicitly. Here, the source operand should be a symbol that indicates the size and location of the source value. The destination operand is then automatically selected to match the size of the source operand (the AL register for byte operands, AX for word operands, and EAX for doubleword operands). This explicit-operands form is provided to allow documentation; however, note that the documentation provided by this form can be misleading. That is, the source operand symbol must specify the correct **type** (size) of the operand (byte, word, or doubleword), but it does not have to specify the correct **location**. The location is always specified by the DS:(E)SI registers, which must be loaded correctly before the load string instruction is executed.

The no-operands form provides "short forms" of the byte, word, and doubleword versions of the LODS instructions. Here also DS:(E)SI is assumed to be the source operand and the AL, AX, or EAX register is assumed to be the destination operand. The size of the source and destination operands is selected with the mnemonic: LODSB (byte loaded into register AL), LODSW (word loaded into AX), or LODSD (doubleword loaded into EAX).

After the byte, word, or doubleword is transferred from the memory location into the AL, AX, or EAX register, the (E)SI register is incremented or decremented automatically according to the setting of the DF flag in the EFLAGS register. (If the DF flag is 0, the (E)SI register is incremented; if the DF flag is 1, the ESI register is decremented.) The (E)SI register is incremented or decremented by 1 for byte operations, by 2 for word operations, or by 4 for doubleword operations.

In 64-bit mode, use of the REX.W prefix promotes operation to 64 bits. LODS/LODSQ load the quadword at address (R)SI into RAX. The (R)SI register is then incremented or decremented automatically according to the setting of the DF flag in the EFLAGS register.

The LODS, LODSB, LODSW, and LODSD instructions can be preceded by the REP prefix for block loads of ECX bytes, words, or doublewords. More often, however, these instructions are used within a LOOP construct because further processing of the data moved into the register is usually necessary before the next transfer can be made. See "REP/REPE/REPZ /REPNE/REPNZ—Repeat String Operation Prefix" in Chapter 4 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 2B*, for a description of the REP prefix.

Operation

```
IF AL ← SRC; (* Byte load *)
    THEN AL ← SRC; (* Byte load *)
        IF DF = 0
            THEN (E)SI ← (E)SI + 1;
```

```
                ELSE (E)SI ← (E)SI – 1;
          FI;
ELSE IF AX ← SRC; (* Word load *)
    THEN IF DF = 0
                THEN (E)SI ← (E)SI + 2;
                ELSE (E)SI ← (E)SI – 2;
          IF;
    FI;
ELSE IF EAX ← SRC; (* Doubleword load *)
    THEN IF DF = 0
                THEN (E)SI ← (E)SI + 4;
                ELSE (E)SI ← (E)SI – 4;
          FI;
    FI;
ELSE IF RAX ← SRC; (* Quadword load *)
    THEN IF DF = 0
                THEN (R)SI ← (R)SI + 8;
                ELSE (R)SI ← (R)SI – 8;
          FI;
    FI;
FI;
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.

#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

LOOP/LOOP*cc*—Loop According to ECX Counter

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
E2 *cb*	LOOP *rel8*	A	Valid	Valid	Decrement count; jump short if count ≠ 0.
E1 *cb*	LOOPE *rel8*	A	Valid	Valid	Decrement count; jump short if count ≠ 0 and ZF = 1.
E0 *cb*	LOOPNE *rel8*	A	Valid	Valid	Decrement count; jump short if count ≠ 0 and ZF = 0.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	Offset	NA	NA	NA

Description

Performs a loop operation using the RCX, ECX or CX register as a counter (depending on whether address size is 64 bits, 32 bits, or 16 bits). Note that the LOOP instruction ignores REX.W; but 64-bit address size can be over-ridden using a 67H prefix.

Each time the LOOP instruction is executed, the count register is decremented, then checked for 0. If the count is 0, the loop is terminated and program execution continues with the instruction following the LOOP instruction. If the count is not zero, a near jump is performed to the destination (target) operand, which is presumably the instruction at the beginning of the loop.

The target instruction is specified with a relative offset (a signed offset relative to the current value of the instruction pointer in the IP/EIP/RIP register). This offset is generally specified as a label in assembly code, but at the machine code level, it is encoded as a signed, 8-bit immediate value, which is added to the instruction pointer. Offsets of −128 to +127 are allowed with this instruction.

Some forms of the loop instruction (LOOP*cc*) also accept the ZF flag as a condition for terminating the loop before the count reaches zero. With these forms of the instruction, a condition code (*cc*) is associated with each instruction to indicate the condition being tested for. Here, the LOOP*cc* instruction itself does not affect the state of the ZF flag; the ZF flag is changed by other instructions in the loop.

Operation

```
IF (AddressSize = 32)
    THEN Count is ECX;
ELSE IF (AddressSize = 64)
    Count is RCX;
ELSE Count is CX;
FI;

Count ← Count – 1;

IF Instruction is not LOOP
    THEN
        IF (Instruction ← LOOPE) or (Instruction ← LOOPZ)
            THEN IF (ZF = 1) and (Count ≠ 0)
                    THEN BranchCond ← 1;
                    ELSE BranchCond ← 0;
                FI;
            ELSE (Instruction = LOOPNE) or (Instruction = LOOPNZ)
                IF (ZF = 0 ) and (Count ≠ 0)
                    THEN BranchCond ← 1;
                    ELSE BranchCond ← 0;
                FI;
        FI;
    ELSE (* Instruction = LOOP *)
        IF (Count ≠ 0)
            THEN BranchCond ← 1;
            ELSE BranchCond ← 0;
        FI;
FI;

IF BranchCond = 1
    THEN
        IF OperandSize = 32
            THEN EIP ← EIP + SignExtend(DEST);
            ELSE IF OperandSize = 64
                THEN RIP ← RIP + SignExtend(DEST);
                FI;
            ELSE IF OperandSize = 16
                THEN EIP ← EIP AND 0000FFFFH;
                FI;
            ELSE IF OperandSize = (32 or 64)
                THEN IF (R/E)IP < CS.Base or (R/E)IP > CS.Limit
                    #GP; FI;
                FI;
```

```
        FI;
    ELSE
        Terminate loop and continue program execution at (R/E)IP;
FI;
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If the offset being jumped to is beyond the limits of the CS segment.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If the offset being jumped to is beyond the limits of the CS segment or is outside of the effective address space from 0 to FFFFH. This condition can occur if a 32-bit address size override prefix is used.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(0)	If the offset being jumped to is in a non-canonical form.
#UD	If the LOCK prefix is used.

LSL—Load Segment Limit

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 03 /r	LSL *r16, r16/m16*	A	Valid	Valid	Load: *r16* ← segment limit, selector *r16/m16.*
0F 03 /r	LSL *r32, r32/m16**	A	Valid	Valid	Load: *r32* ← segment limit, selector *r32/m16.*
REX.W + 0F 03 /r	LSL *r64, r32/m16**	A	Valid	Valid	Load: *r64* ← segment limit, selector *r32/m16*

NOTES:

* For all loads (regardless of destination sizing), only bits 16-0 are used. Other bits are ignored.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Loads the unscrambled segment limit from the segment descriptor specified with the second operand (source operand) into the first operand (destination operand) and sets the ZF flag in the EFLAGS register. The source operand (which can be a register or a memory location) contains the segment selector for the segment descriptor being accessed. The destination operand is a general-purpose register.

The processor performs access checks as part of the loading process. Once loaded in the destination register, software can compare the segment limit with the offset of a pointer.

The segment limit is a 20-bit value contained in bytes 0 and 1 and in the first 4 bits of byte 6 of the segment descriptor. If the descriptor has a byte granular segment limit (the granularity flag is set to 0), the destination operand is loaded with a byte granular value (byte limit). If the descriptor has a page granular segment limit (the granularity flag is set to 1), the LSL instruction will translate the page granular limit (page limit) into a byte limit before loading it into the destination operand. The translation is performed by shifting the 20-bit "raw" limit left 12 bits and filling the low-order 12 bits with 1s.

When the operand size is 32 bits, the 32-bit byte limit is stored in the destination operand. When the operand size is 16 bits, a valid 32-bit limit is computed; however, the upper 16 bits are truncated and only the low-order 16 bits are loaded into the destination operand.

This instruction performs the following checks before it loads the segment limit into the destination register:

- Checks that the segment selector is not NULL.

- Checks that the segment selector points to a descriptor that is within the limits of the GDT or LDT being accessed
- Checks that the descriptor type is valid for this instruction. All code and data segment descriptors are valid for (can be accessed with) the LSL instruction. The valid special segment and gate descriptor types are given in the following table.
- If the segment is not a conforming code segment, the instruction checks that the specified segment descriptor is visible at the CPL (that is, if the CPL and the RPL of the segment selector are less than or equal to the DPL of the segment selector).

If the segment descriptor cannot be accessed or is an invalid type for the instruction, the ZF flag is cleared and no value is loaded in the destination operand.

Table 3-65. Segment and Gate Descriptor Types

Type	Protected Mode		IA-32e Mode	
	Name	Valid	Name	Valid
0	Reserved	No	Upper 8 byte of a 16-Byte descriptor	Yes
1	Available 16-bit TSS	Yes	Reserved	No
2	LDT	Yes	LDT	Yes
3	Busy 16-bit TSS	Yes	Reserved	No
4	16-bit call gate	No	Reserved	No
5	16-bit/32-bit task gate	No	Reserved	No
6	16-bit interrupt gate	No	Reserved	No
7	16-bit trap gate	No	Reserved	No
8	Reserved	No	Reserved	No
9	Available 32-bit TSS	Yes	64-bit TSS	Yes
A	Reserved	No	Reserved	No
B	Busy 32-bit TSS	Yes	Busy 64-bit TSS	Yes
C	32-bit call gate	No	64-bit call gate	No
D	Reserved	No	Reserved	No
E	32-bit interrupt gate	No	64-bit interrupt gate	No
F	32-bit trap gate	No	64-bit trap gate	No

Operation

```
IF SRC(Offset) > descriptor table limit
    THEN ZF ← 0; FI;

Read segment descriptor;

IF SegmentDescriptor(Type) ≠ conforming code segment
and (CPL > DPL) OR (RPL > DPL)
or Segment type is not valid for instruction
        THEN
            ZF ← 0;
        ELSE
            temp ← SegmentLimit([SRC]);
            IF (G ← 1)
                THEN temp ← ShiftLeft(12, temp) OR 00000FFFH;
            ELSE IF OperandSize = 32
                THEN DEST ← temp; FI;
            ELSE IF OperandSize = 64 (* REX.W used *)
                THEN DEST (* Zero-extended *) ← temp; FI;
            ELSE (* OperandSize = 16 *)
                DEST ← temp AND FFFFH;
            FI;
FI;
```

Flags Affected

The ZF flag is set to 1 if the segment limit is loaded successfully; otherwise, it is set to 0.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and the memory operand effective address is unaligned while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#UD	The LSL instruction cannot be executed in real-address mode.

Virtual-8086 Mode Exceptions

#UD	The LSL instruction cannot be executed in virtual-8086 mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If the memory operand effective address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory operand effective address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and the memory operand effective address is unaligned while the current privilege level is 3.
#UD	If the LOCK prefix is used.

LTR—Load Task Register

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 00 /3	LTR *r/m*16	A	Valid	Valid	Load *r/m*16 into task register.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r)	NA	NA	NA

Description

Loads the source operand into the segment selector field of the task register. The source operand (a general-purpose register or a memory location) contains a segment selector that points to a task state segment (TSS). After the segment selector is loaded in the task register, the processor uses the segment selector to locate the segment descriptor for the TSS in the global descriptor table (GDT). It then loads the segment limit and base address for the TSS from the segment descriptor into the task register. The task pointed to by the task register is marked busy, but a switch to the task does not occur.

The LTR instruction is provided for use in operating-system software; it should not be used in application programs. It can only be executed in protected mode when the CPL is 0. It is commonly used in initialization code to establish the first task to be executed.

The operand-size attribute has no effect on this instruction.

In 64-bit mode, the operand size is still fixed at 16 bits. The instruction references a 16-byte descriptor to load the 64-bit base.

Operation

```
IF SRC is a NULL selector
    THEN #GP(0);

IF SRC(Offset) > descriptor table limit OR IF SRC(type) ≠ global
    THEN #GP(segment selector); FI;

Read segment descriptor;

IF segment descriptor is not for an available TSS
    THEN #GP(segment selector); FI;
IF segment descriptor is not present
    THEN #NP(segment selector); FI;
```

```
TSSsegmentDescriptor(busy) ← 1;
(* Locked read-modify-write operation on the entire descriptor when setting busy flag *)

TaskRegister(SegmentSelector) ← SRC;
TaskRegister(SegmentDescriptor) ← TSSSegmentDescriptor;
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If the current privilege level is not 0.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the source operand contains a NULL segment selector.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
#GP(selector)	If the source selector points to a segment that is not a TSS or to one for a task that is already busy.
	If the selector points to LDT or is beyond the GDT limit.
#NP(selector)	If the TSS is marked not present.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#UD	The LTR instruction is not recognized in real-address mode.

Virtual-8086 Mode Exceptions

#UD	The LTR instruction is not recognized in virtual-8086 mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the current privilege level is not 0.
	If the memory address is in a non-canonical form.
	If the source operand contains a NULL segment selector.

#GP(selector)	If the source selector points to a segment that is not a TSS or to one for a task that is already busy.
	If the selector points to LDT or is beyond the GDT limit.
	If the descriptor type of the upper 8-byte of the 16-byte descriptor is non-zero.
#NP(selector)	If the TSS is marked not present.
#PF(fault-code)	If a page fault occurs.
#UD	If the LOCK prefix is used.

MASKMOVDQU—Store Selected Bytes of Double Quadword

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F F7 /*r* MASKMOVDQU *xmm1*, *xmm2*	A	V/V	SSE2	Selectively write bytes from *xmm1* to memory location using the byte mask in *xmm2*. The default memory location is specified by DS:EDI/RDI.
VEX.128.66.0F.WIG F7 /r VMASKMOVDQU xmm1, xmm2	A	V/V	AVX	Selectively write bytes from xmm1 to memory location using the byte mask in xmm2. The default memory location is specified by DS:DI/EDI/RDI.

Instruction Operand Encoding[1]

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r)	ModRM:r/m (r)	NA	NA

Description

Stores selected bytes from the source operand (first operand) into an 128-bit memory location. The mask operand (second operand) selects which bytes from the source operand are written to memory. The source and mask operands are XMM registers. The location of the first byte of the memory location is specified by DI/EDI and DS registers. The memory location does not need to be aligned on a natural boundary. (The size of the store address depends on the address-size attribute.)

The most significant bit in each byte of the mask operand determines whether the corresponding byte in the source operand is written to the corresponding byte location in memory: 0 indicates no write and 1 indicates write.

The MASKMOVDQU instruction generates a non-temporal hint to the processor to minimize cache pollution. The non-temporal hint is implemented by using a write combining (WC) memory type protocol (see “Caching of Temporal vs. Non-Temporal Data” in Chapter 10, of the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*). Because the WC protocol uses a weakly-ordered memory consistency model, a fencing operation implemented with the SFENCE or MFENCE instruction should be used in conjunction with MASKMOVDQU instructions if multiple processors might use different memory types to read/write the destination memory locations.

1.ModRM.MOD = 011B required

Behavior with a mask of all 0s is as follows:

- No data will be written to memory.
- Signaling of breakpoints (code or data) is not guaranteed; different processor implementations may signal or not signal these breakpoints.
- Exceptions associated with addressing memory and page faults may still be signaled (implementation dependent).
- If the destination memory region is mapped as UC or WP, enforcement of associated semantics for these memory types is not guaranteed (that is, is reserved) and is implementation-specific.

The MASKMOVDQU instruction can be used to improve performance of algorithms that need to merge data on a byte-by-byte basis. MASKMOVDQU should not cause a read for ownership; doing so generates unnecessary bandwidth since data is to be written directly using the byte-mask without allocating old data prior to the store.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

If VMASKMOVDQU is encoded with VEX.L= 1, an attempt to execute the instruction encoded with VEX.L= 1 will cause an #UD exception.

Operation

```
IF (MASK[7] = 1)
    THEN DEST[DI/EDI] ← SRC[7:0] ELSE (* Memory location unchanged *); FI;
IF (MASK[15] = 1)
    THEN DEST[DI/EDI +1] ← SRC[15:8] ELSE (* Memory location unchanged *); FI;
    (* Repeat operation for 3rd through 14th bytes in source operand *)
IF (MASK[127] = 1)
    THEN DEST[DI/EDI +15] ← SRC[127:120] ELSE (* Memory location unchanged *); FI;
```

Intel C/C++ Compiler Intrinsic Equivalent

void _mm_maskmoveu_si128(__m128i d, __m128i n, char * p)

Other Exceptions

See Exceptions Type 4; additionally

#UD	If VEX.L= 1
	If VEX.vvvv != 1111B.

VMASKMOV—Conditional SIMD Packed Loads and Stores

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
VEX.NDS.128.66.0F38.W0 2C /r VMASKMOVPS xmm1, xmm2, m128	A	V/V	AVX	Conditionally load packed single-precision values from m128 using mask in xmm2 and store in xmm1.
VEX.NDS.256.66.0F38.W0 2C /r VMASKMOVPS ymm1, ymm2, m256	A	V/V	AVX	Conditionally load packed single-precision values from m256 using mask in ymm2 and store in ymm1.
VEX.NDS.128.66.0F38.W0 2D /r VMASKMOVPD xmm1, xmm2, m128	A	V/V	AVX	Conditionally load packed double-precision values from m128 using mask in xmm2 and store in xmm1.
VEX.NDS.256.66.0F38.W0 2D /r VMASKMOVPD ymm1, ymm2, m256	A	V/V	AVX	Conditionally load packed double-precision values from m256 using mask in ymm2 and store in ymm1.
VEX.NDS.128.66.0F38.W0 2E /r VMASKMOVPS m128, xmm1, xmm2	B	V/V	AVX	Conditionally store packed single-precision values from xmm2 using mask in xmm1.
VEX.NDS.256.66.0F38.W0 2E /r VMASKMOVPS m256, ymm1, ymm2	B	V/V	AVX	Conditionally store packed single-precision values from ymm2 using mask in ymm1.
VEX.NDS.128.66.0F38.W0 2F /r VMASKMOVPD m128, xmm1, xmm2	B	V/V	AVX	Conditionally store packed double-precision values from xmm2 using mask in xmm1.
VEX.NDS.256.66.0F38.W0 2F /r VMASKMOVPD m256, ymm1, ymm2	B	V/V	AVX	Conditionally store packed double-precision values from ymm2 using mask in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA
B	ModRM:r/m (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Conditionally moves packed data elements from the second source operand into the corresponding data element of the destination operand, depending on the mask bits associated with each data element. The mask bits are specified in the first source operand.

The mask bit for each data element is the most significant bit of that element in the first source operand. If a mask is 1, the corresponding data element is copied from the second source operand to the destination operand. If the mask is 0, the corresponding data element is set to zero in the load form of these instructions, and unmodified in the store form.

The second source operand is a memory address for the load form of these instruction. The destination operand is a memory address for the store form of these instructions. The other operands are both XMM registers (for VEX.128 version) or YMM registers (for VEX.256 version).

Faults occur only due to mask-bit required memory accesses that caused the faults. Faults will not occur due to referencing any memory location if the corresponding mask bit for that memory location is 0. For example, no faults will be detected if the mask bits are all zero.

Unlike previous MASKMOV instructions (MASKMOVQ and MASKMOVDQU), a nontemporal hint is not applied to these instructions

Instruction behavior on alignment check reporting with mask bits of less than all 1s are the same as with mask bits of all 1s.

VMASKMOV should not be used to access memory mapped I/O and un-cached memory as the access and the ordering of the individual loads or stores it does is implementation specific.

In cases where mask bits indicate data should not be loaded or stored paging A and D bits will be set in an implementation dependent way. However, A and D bits are always set for pages where data is actually loaded/stored.

Note: for load forms, the first source (the mask) is encoded in VEX.vvvv; the second source is encoded in rm_field, and the destination register is encoded in reg_field.

Note: for store forms, the first source (the mask) is encoded in VEX.vvvv; the second source register is encoded in reg_field, and the destination memory location is encoded in rm_field.

Operation

VMASKMOVPS -128-bit load

```
DEST[31:0] ← IF (SRC1[31]) Load_32(mem) ELSE 0
DEST[63:32] ← IF (SRC1[63]) Load_32(mem + 4) ELSE 0
DEST[95:64] ← IF (SRC1[95]) Load_32(mem + 8) ELSE 0
DEST[127:97] ← IF (SRC1[127]) Load_32(mem + 12) ELSE 0
DEST[VLMAX-1:128] ← 0
```

VMASKMOVPS - 256-bit load

DEST[31:0] ← IF (SRC1[31]) Load_32(mem) ELSE 0
DEST[63:32] ← IF (SRC1[63]) Load_32(mem + 4) ELSE 0
DEST[95:64] ← IF (SRC1[95]) Load_32(mem + 8) ELSE 0
DEST[127:96] ← IF (SRC1[127]) Load_32(mem + 12) ELSE 0
DEST[159:128] ← IF (SRC1[159]) Load_32(mem + 16) ELSE 0
DEST[191:160] ← IF (SRC1[191]) Load_32(mem + 20) ELSE 0
DEST[223:192] ← IF (SRC1[223]) Load_32(mem + 24) ELSE 0
DEST[255:224] ← IF (SRC1[255]) Load_32(mem + 28) ELSE 0

VMASKMOVPD - 128-bit load
DEST[63:0] ← IF (SRC1[63]) Load_64(mem) ELSE 0
DEST[127:64] ← IF (SRC1[127]) Load_64(mem + 16) ELSE 0
DEST[VLMAX-1:128] ← 0

VMASKMOVPD - 256-bit load
DEST[63:0] ← IF (SRC1[63]) Load_64(mem) ELSE 0
DEST[127:64] ← IF (SRC1[127]) Load_64(mem + 8) ELSE 0
DEST[195:128] ← IF (SRC1[191]) Load_64(mem + 16) ELSE 0
DEST[255:196] ← IF (SRC1[255]) Load_64(mem + 24) ELSE 0

VMASKMOVPS - 128-bit store
IF (SRC1[31]) DEST[31:0] ← SRC2[31:0]
IF (SRC1[63]) DEST[63:32] ← SRC2[63:32]
IF (SRC1[95]) DEST[95:64] ← SRC2[95:64]
IF (SRC1[127]) DEST[127:96] ← SRC2[127:96]

VMASKMOVPS - 256-bit store
IF (SRC1[31]) DEST[31:0] ← SRC2[31:0]
IF (SRC1[63]) DEST[63:32] ← SRC2[63:32]
IF (SRC1[95]) DEST[95:64] ← SRC2[95:64]
IF (SRC1[127]) DEST[127:96] ← SRC2[127:96]
IF (SRC1[159]) DEST[159:128] ←SRC2[159:128]
IF (SRC1[191]) DEST[191:160] ← SRC2[191:160]
IF (SRC1[223]) DEST[223:192] ← SRC2[223:192]
IF (SRC1[255]) DEST[255:224] ← SRC2[255:224]

VMASKMOVPD - 128-bit store
IF (SRC1[63]) DEST[63:0] ← SRC2[63:0]
IF (SRC1[127]) DEST[127:64] ←SRC2[127:64]

VMASKMOVPD - 256-bit store
IF (SRC1[63]) DEST[63:0] ← SRC2[63:0]
IF (SRC1[127]) DEST[127:64] ←SRC2[127:64]

```
IF (SRC1[191]) DEST[191:128] ← SRC2[191:128]
IF (SRC1[255]) DEST[255:192] ← SRC2[255:192]
```

Intel C/C++ Compiler Intrinsic Equivalent

```
__m256  _mm256_maskload_ps(float const *a, __m256i mask)

void    _mm256_maskstore_ps(float *a, __m256i mask, __m256 b)

__m256d _mm256_maskload_pd(double *a, __m256i mask);

void    _mm256_maskstore_pd(double *a, __m256i mask, __m256d b);

__m128 _mm256_maskload_ps(float const *a, __m128i mask)

void    _mm256_maskstore_ps(float *a, __m128i mask, __m128 b)

__m128d _mm256_maskload_pd(double *a, __m128i mask);

void    _mm256_maskstore_pd(double *a, __m128i mask, __m128d b);
```

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 6 (No AC# reported for any mask bit combinations); additionally

#UD	If VEX.W = 1.

MASKMOVQ—Store Selected Bytes of Quadword

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F F7 */r*	MASKMOVQ *mm1, mm2*	A	Valid	Valid	Selectively write bytes from *mm1* to memory location using the byte mask in *mm2.* The default memory location is specified by DS:EDI/RDI.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r)	ModRM:r/m (r)	NA	NA

Description

Stores selected bytes from the source operand (first operand) into a 64-bit memory location. The mask operand (second operand) selects which bytes from the source operand are written to memory. The source and mask operands are MMX technology registers. The location of the first byte of the memory location is specified by DI/EDI and DS registers. (The size of the store address depends on the address-size attribute.)

The most significant bit in each byte of the mask operand determines whether the corresponding byte in the source operand is written to the corresponding byte location in memory: 0 indicates no write and 1 indicates write.

The MASKMOVQ instruction generates a non-temporal hint to the processor to minimize cache pollution. The non-temporal hint is implemented by using a write combining (WC) memory type protocol (see “Caching of Temporal vs. Non-Temporal Data” in Chapter 10, of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*). Because the WC protocol uses a weakly-ordered memory consistency model, a fencing operation implemented with the SFENCE or MFENCE instruction should be used in conjunction with MASKMOVQ instructions if multiple processors might use different memory types to read/write the destination memory locations.

This instruction causes a transition from x87 FPU to MMX technology state (that is, the x87 FPU top-of-stack pointer is set to 0 and the x87 FPU tag word is set to all 0s [valid]).

The behavior of the MASKMOVQ instruction with a mask of all 0s is as follows:

- No data will be written to memory.
- Transition from x87 FPU to MMX technology state will occur.
- Exceptions associated with addressing memory and page faults may still be signaled (implementation dependent).

- Signaling of breakpoints (code or data) is not guaranteed (implementation dependent).
- If the destination memory region is mapped as UC or WP, enforcement of associated semantics for these memory types is not guaranteed (that is, is reserved) and is implementation-specific.

The MASKMOVQ instruction can be used to improve performance for algorithms that need to merge data on a byte-by-byte basis. It should not cause a read for ownership; doing so generates unnecessary bandwidth since data is to be written directly using the byte-mask without allocating old data prior to the store.

In 64-bit mode, the memory address is specified by DS:RDI.

Operation

```
IF (MASK[7] = 1)
    THEN DEST[DI/EDI] ← SRC[7:0] ELSE (* Memory location unchanged *); FI;
IF (MASK[15] = 1)
    THEN DEST[DI/EDI +1] ← SRC[15:8] ELSE (* Memory location unchanged *); FI;
    (* Repeat operation for 3rd through 6th bytes in source operand *)
IF (MASK[63] = 1)
    THEN DEST[DI/EDI +15] ← SRC[63:56] ELSE (* Memory location unchanged *); FI;
```

Intel C/C++ Compiler Intrinsic Equivalent

void _mm_maskmove_si64(__m64d, __m64n, char * p)

Protected Mode Exceptions

#GP(0)	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments (even if mask is all 0s).
	If the destination operand is in a nonwritable segment.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	For an illegal address in the SS segment (even if mask is all 0s).
#PF(fault-code)	For a page fault (implementation specific).
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending FPU exception.
#UD	If CR0.EM[bit 2] = 1.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If Mod field of the ModR/M byte not 11B.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

Real-Address Mode Exceptions

GP	If any part of the operand lies outside the effective address space from 0 to FFFFH. (even if mask is all 0s).
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending FPU exception.
#UD	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	For a page fault (implementation specific).
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(0)	If the memory address is in a non-canonical form.
#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#PF(fault-code)	For a page fault (implementation specific).
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending FPU exception.
#UD	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If Mod field of the ModR/M byte not 11B.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

MAXPD—Return Maximum Packed Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 5F /*r* MAXPD *xmm1, xmm2/m128*	A	V/V	SSE2	Return the maximum double-precision floating-point values between *xmm2/m128* and *xmm1*.
VEX.NDS.128.66.0F.WIG 5F /r VMAXPD xmm1,xmm2, xmm3/m128	B	V/V	AVX	Return the maximum double-precision floating-point values between xmm2 and xmm3/mem.
VEX.NDS.256.66.0F.WIG 5F /r VMAXPD ymm1, ymm2, ymm3/m256	B	V/V	AVX	Return the maximum packed double-precision floating-point values between ymm2 and ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs an SIMD compare of the packed double-precision floating-point values in the first source operand and the second source operand and returns the maximum value for each pair of values to the destination operand.

If the values being compared are both 0.0s (of either sign), the value in the second operand (source operand) is returned. If a value in the second operand is an SNaN, that SNaN is forwarded unchanged to the destination (that is, a QNaN version of the SNaN is not returned).

If only one value is a NaN (SNaN or QNaN) for this instruction, the second operand (source operand), either a NaN or a valid floating-point value, is written to the result. If instead of this behavior, it is required that the NaN source operand (from either the first or second operand) be returned, the action of MAXPD can be emulated using a sequence of instructions, such as, a comparison followed by AND, ANDN and OR.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register

and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

```
MAX(SRC1, SRC2)
{
    IF ((SRC1 = 0.0) and (SRC2 = 0.0)) THEN DEST ← SRC2;
        ELSE IF (SRC1 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF (SRC2 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF (SRC1 > SRC2) THEN DEST ← SRC1;
        ELSE DEST ← SRC2;
    FI;
}

MAXPD (128-bit Legacy SSE version)
DEST[63:0] ← MAX(DEST[63:0], SRC[63:0])
DEST[127:64] ← MAX(DEST[127:64], SRC[127:64])
DEST[VLMAX-1:128] (Unmodified)

VMAXPD (VEX.128 encoded version)
DEST[63:0] ← MAX(SRC1[63:0], SRC2[63:0])
DEST[127:64] ← MAX(SRC1[127:64], SRC2[127:64])
DEST[VLMAX-1:128] ← 0

VMAXPD (VEX.256 encoded version)
DEST[63:0] ← MAX(SRC1[63:0], SRC2[63:0])
DEST[127:64] ← MAX(SRC1[127:64], SRC2[127:64])
DEST[191:128] ← MAX(SRC1[191:128], SRC2[191:128])
DEST[255:192] ← MAX(SRC1[255:192], SRC2[255:192])
```

Intel C/C++ Compiler Intrinsic Equivalent

MAXPD __m128d _mm_max_pd(__m128d a, __m128d b);

VMAXPD __m256d _mm256_max_pd (__m256d a, __m256d b);

SIMD Floating-Point Exceptions

Invalid (including QNaN source operand), Denormal.

Other Exceptions

See Exceptions Type 2.

MAXPS—Return Maximum Packed Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 5F /*r* MAXPS *xmm1*, *xmm2/m128*	A	V/V	SSE	Return the maximum single-precision floating-point values between *xmm2/m128* and *xmm1*.
VEX.NDS.128.0F.WIG 5F /r VMAXPS xmm1,xmm2, xmm3/m128	B	V/V	AVX	Return the maximum single-precision floating-point values between xmm2 and xmm3/mem.
VEX.NDS.256.0F.WIG 5F /r VMAXPS ymm1, ymm2, ymm3/m256	B	V/V	AVX	Return the maximum single double-precision floating-point values between ymm2 and ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs an SIMD compare of the packed single-precision floating-point values in the first source operand and the second source operand and returns the maximum value for each pair of values to the destination operand.

If the values being compared are both 0.0s (of either sign), the value in the second operand (source operand) is returned. If a value in the second operand is an SNaN, that SNaN is forwarded unchanged to the destination (that is, a QNaN version of the SNaN is not returned).

If only one value is a NaN (SNaN or QNaN) for this instruction, the second operand (source operand), either a NaN or a valid floating-point value, is written to the result. If instead of this behavior, it is required that the NaN source operand (from either the first or second operand) be returned, the action of MAXPS can be emulated using a sequence of instructions, such as, a comparison followed by AND, ANDN and OR.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register

and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

MAX(SRC1, SRC2)

```
{
    IF ((SRC1 = 0.0) and (SRC2 = 0.0)) THEN DEST ← SRC2;
        ELSE IF (SRC1 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF SRC2 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF (SRC1 > SRC2) THEN DEST ← SRC1;
        ELSE DEST ← SRC2;
    FI;
}
```

MAXPS (128-bit Legacy SSE version)

```
DEST[31:0] ← MAX(DEST[31:0], SRC[31:0])
DEST[63:32] ← MAX(DEST[63:32], SRC[63:32])
DEST[95:64] ← MAX(DEST[95:64], SRC[95:64])
DEST[127:96] ← MAX(DEST[127:96], SRC[127:96])
DEST[VLMAX-1:128] (Unmodified)
```

VMAXPS (VEX.128 encoded version)

```
DEST[31:0] ← MAX(SRC1[31:0], SRC2[31:0])
DEST[63:32] ← MAX(SRC1[63:32], SRC2[63:32])
DEST[95:64] ← MAX(SRC1[95:64], SRC2[95:64])
DEST[127:96] ← MAX(SRC1[127:96], SRC2[127:96])
DEST[VLMAX-1:128] ← 0
```

VMAXPS (VEX.256 encoded version)

```
DEST[31:0] ← MAX(SRC1[31:0], SRC2[31:0])
DEST[63:32] ← MAX(SRC1[63:32], SRC2[63:32])
DEST[95:64] ← MAX(SRC1[95:64], SRC2[95:64])
DEST[127:96] ← MAX(SRC1[127:96], SRC2[127:96])
DEST[159:128] ← MAX(SRC1[159:128], SRC2[159:128])
DEST[191:160] ← MAX(SRC1[191:160], SRC2[191:160])
DEST[223:192] ← MAX(SRC1[223:192], SRC2[223:192])
```

DEST[255:224] ← MAX(SRC1[255:224], SRC2[255:224])

Intel C/C++ Compiler Intrinsic Equivalent

MAXPS __m128 _mm_max_ps (__m128 a, __m128 b);

VMAXPS __m256 _mm256_max_ps (__m256 a, __m256 b);

SIMD Floating-Point Exceptions

Invalid (including QNaN source operand), Denormal.

Other Exceptions

See Exceptions Type 2.

MAXSD—Return Maximum Scalar Double-Precision Floating-Point Value

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 5F /*r* MAXSD *xmm1, xmm2/m64*	A	V/V	SSE2	Return the maximum scalar double-precision floating-point value between *xmm2/mem64* and *xmm1*.
VEX.NDS.LIG.F2.0F.WIG 5F /r VMAXSD xmm1, xmm2, xmm3/m64	B	V/V	AVX	Return the maximum scalar double-precision floating-point value between xmm3/mem64 and xmm2.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Compares the low double-precision floating-point values in the first source operand and second the source operand, and returns the maximum value to the low quadword of the destination operand. The second source operand can be an XMM register or a 64-bit memory location. The first source and destination operands are XMM registers. When the second source operand is a memory operand, only 64 bits are accessed. The high quadword of the destination operand is copied from the same bits of first source operand.

If the values being compared are both 0.0s (of either sign), the value in the second source operand is returned. If a value in the second source operand is an SNaN, that SNaN is returned unchanged to the destination (that is, a QNaN version of the SNaN is not returned).

If only one value is a NaN (SNaN or QNaN) for this instruction, the second source operand, either a NaN or a valid floating-point value, is written to the result. If instead of this behavior, it is required that the NaN of either source operand be returned, the action of MAXSD can be emulated using a sequence of instructions, such as, a comparison followed by AND, ANDN and OR.

The second source operand can be an XMM register or a 64-bit memory location. The first source and destination operands are XMM registers.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The destination and first source operand are the same. Bits (VLMAX-1:64) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (127:64) of the XMM register destination are copied from corresponding bits in the first source operand. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

```
MAX(SRC1, SRC2)
{
    IF ((SRC1 = 0.0) and (SRC2 = 0.0)) THEN DEST ← SRC2;
        ELSE IF (SRC1 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF SRC2 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF (SRC1 > SRC2) THEN DEST ← SRC1;
        ELSE DEST ← SRC2;
    FI;
}
```

MAXSD (128-bit Legacy SSE version)

```
DEST[63:0] ←MAX(DEST[63:0], SRC[63:0])
DEST[VLMAX-1:64] (Unmodified)
```

VMAXSD (VEX.128 encoded version)

```
DEST[63:0] ←MAX(SRC1[63:0], SRC2[63:0])
DEST[127:64] ←SRC1[127:64]
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

MAXSD __m128d _mm_max_sd(__m128d a, __m128d b)

SIMD Floating-Point Exceptions

Invalid (including QNaN source operand), Denormal.

Other Exceptions

See Exceptions Type 3.

MAXSS—Return Maximum Scalar Single-Precision Floating-Point Value

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 5F /*r* MAXSS *xmm1*, *xmm2/m32*	A	V/V	SSE	Return the maximum scalar single-precision floating-point value between *xmm2/mem32* and *xmm1*.
VEX.NDS.LIG.F3.0F.WIG 5F /r VMAXSS xmm1, xmm2, xmm3/m32	B	V/V	AVX	Return the maximum scalar single-precision floating-point value between xmm3/mem32 and xmm2.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Compares the low single-precision floating-point values in the first source operand and the second source operand, and returns the maximum value to the low doubleword of the destination operand.

If the values being compared are both 0.0s (of either sign), the value in the second source operand is returned. If a value in the second source operand is an SNaN, that SNaN is returned unchanged to the destination (that is, a QNaN version of the SNaN is not returned).

If only one value is a NaN (SNaN or QNaN) for this instruction, the second source operand, either a NaN or a valid floating-point value, is written to the result. If instead of this behavior, it is required that the NaN from either source operand be returned, the action of MAXSS can be emulated using a sequence of instructions, such as, a comparison followed by AND, ANDN and OR.

The second source operand can be an XMM register or a 32-bit memory location. The first source and destination operands are XMM registers.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The destination and first source operand are the same. Bits (VLMAX-1:32) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (127:32) of the XMM register destination are copied from corresponding bits in the first source operand. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

```
MAX(SRC1, SRC2)
{
    IF ((SRC1 = 0.0) and (SRC2 = 0.0)) THEN DEST ← SRC2;
        ELSE IF (SRC1 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF SRC2 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF (SRC1 > SRC2) THEN DEST ← SRC1;
        ELSE DEST ← SRC2;
    FI;
}
```

MAXSS (128-bit Legacy SSE version)

```
DEST[31:0] ←MAX(DEST[31:0], SRC[31:0])
DEST[VLMAX-1:32] (Unmodified)
```

VMAXSS (VEX.128 encoded version)

```
DEST[31:0] ←MAX(SRC1[31:0], SRC2[31:0])
DEST[127:32] ←SRC1[127:32]
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

__m128d _mm_max_ss(__m128d a, __m128d b)

SIMD Floating-Point Exceptions

Invalid (including QNaN source operand), Denormal.

Other Exceptions

See Exceptions Type 3.

MFENCE—Memory Fence

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F AE /6	MFENCE	A	Valid	Valid	Serializes load and store operations.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Performs a serializing operation on all load-from-memory and store-to-memory instructions that were issued prior the MFENCE instruction. This serializing operation guarantees that every load and store instruction that precedes the MFENCE instruction in program order becomes globally visible before any load or store instruction that follows the MFENCE instruction.[1] The MFENCE instruction is ordered with respect to all load and store instructions, other MFENCE instructions, any LFENCE and SFENCE instructions, and any serializing instructions (such as the CPUID instruction). MFENCE does not serialize the instruction stream.

Weakly ordered memory types can be used to achieve higher processor performance through such techniques as out-of-order issue, speculative reads, write-combining, and write-collapsing. The degree to which a consumer of data recognizes or knows that the data is weakly ordered varies among applications and may be unknown to the producer of this data. The MFENCE instruction provides a performance-efficient way of ensuring load and store ordering between routines that produce weakly-ordered results and routines that consume that data.

Processors are free to fetch and cache data speculatively from regions of system memory that use the WB, WC, and WT memory types. This speculative fetching can occur at any time and is not tied to instruction execution. Thus, it is not ordered with respect to executions of the MFENCE instruction; data can be brought into the caches speculatively just before, during, or after the execution of an MFENCE instruction.Processors are free to fetch and cache data speculatively from regions of system memory that use the WB, WC, and WT memory types. This speculative fetching can occur at any time and is not tied to instruction execution. Thus, it is not ordered with respect to executions of the MFENCE instruction; data can be brought into the caches speculatively just before, during, or after the execution of an MFENCE instruction.

This instruction’s operation is the same in non-64-bit modes and 64-bit mode.

1. A load instruction is considered to become globally visible when the value to be loaded into its destination register is determined.

Operation

```
Wait_On_Following_Loads_And_Stores_Until(preceding_loads_and_stores_globally_visible);
```

Intel C/C++ Compiler Intrinsic Equivalent

```
void _mm_mfence(void)
```

Exceptions (All Modes of Operation)

#UD	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.

MINPD—Return Minimum Packed Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 5D /*r* MINPD *xmm1*, *xmm2/m128*	A	V/V	SSE2	Return the minimum double-precision floating-point values between *xmm2/m128* and *xmm1*.
VEX.NDS.128.66.0F.WIG 5D /r VMINPD xmm1,xmm2, xmm3/m128	B	V/V	AVX	Return the minimum double-precision floating-point values between xmm2 and xmm3/mem.
VEX.NDS.256.66.0F.WIG 5D /r VMINPD ymm1, ymm2, ymm3/m256	B	V/V	AVX	Return the minimum packed double-precision floating-point values between ymm2 and ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs an SIMD compare of the packed double-precision floating-point values in the first source operand and the second source operand and returns the minimum value for each pair of values to the destination operand.

If the values being compared are both 0.0s (of either sign), the value in the second operand (source operand) is returned. If a value in the second operand is an SNaN, that SNaN is forwarded unchanged to the destination (that is, a QNaN version of the SNaN is not returned).

If only one value is a NaN (SNaN or QNaN) for this instruction, the second operand (source operand), either a NaN or a valid floating-point value, is written to the result. If instead of this behavior, it is required that the NaN source operand (from either the first or second operand) be returned, the action of MINPD can be emulated using a sequence of instructions, such as, a comparison followed by AND, ANDN and OR.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register

and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

Operation

```
MIN(SRC1, SRC2)
{
    IF ((SRC1 = 0.0) and (SRC2 = 0.0)) THEN DEST ← SRC2;
        ELSE IF (SRC1 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF (SRC2 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF (SRC1 < SRC2) THEN DEST ← SRC1;
        ELSE DEST ← SRC2;
    FI;
}

MINPD (128-bit Legacy SSE version)
DEST[63:0] ← MIN(SRC1[63:0], SRC2[63:0])
DEST[127:64] ← MIN(SRC1[127:64], SRC2[127:64])
DEST[VLMAX-1:128] (Unmodified)

VMINPD (VEX.128 encoded version)
DEST[63:0] ← MIN(SRC1[63:0], SRC2[63:0])
DEST[127:64] ← MIN(SRC1[127:64], SRC2[127:64])
DEST[VLMAX-1:128] ← 0

VMINPD (VEX.256 encoded version)
DEST[63:0] ← MIN(SRC1[63:0], SRC2[63:0])
DEST[127:64] ← MIN(SRC1[127:64], SRC2[127:64])
DEST[191:128] ← MIN(SRC1[191:128], SRC2[191:128])
DEST[255:192] ← MIN(SRC1[255:192], SRC2[255:192])
```

Intel C/C++ Compiler Intrinsic Equivalent

MINPD __m128d _mm_min_pd(__m128d a, __m128d b);

VMINPD __m256d _mm256_min_pd (__m256d a, __m256d b);

SIMD Floating-Point Exceptions

Invalid (including QNaN source operand), Denormal.

Other Exceptions

See Exceptions Type 2.

MINPS—Return Minimum Packed Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 5D /*r* MINPS *xmm1*, *xmm2/m128*	A	V/V	SSE	Return the minimum single-precision floating-point values between *xmm2/m128* and *xmm1*.
VEX.NDS.128.0F.WIG 5D /r VMINPS xmm1,xmm2, xmm3/m128	B	V/V	AVX	Return the minimum single-precision floating-point values between xmm2 and xmm3/mem.
VEX.NDS.256.0F.WIG 5D /r VMINPS ymm1, ymm2, ymm3/m256	B	V/V	AVX	Return the minimum single double-precision floating-point values between ymm2 and ymm3/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs an SIMD compare of the packed single-precision floating-point values in the first source operand and the second source operand and returns the minimum value for each pair of values to the destination operand.

If the values being compared are both 0.0s (of either sign), the value in the second operand (source operand) is returned. If a value in the second operand is an SNaN, that SNaN is forwarded unchanged to the destination (that is, a QNaN version of the SNaN is not returned).

If only one value is a NaN (SNaN or QNaN) for this instruction, the second operand (source operand), either a NaN or a valid floating-point value, is written to the result. If instead of this behavior, it is required that the NaN source operand (from either the first or second operand) be returned, the action of MINPS can be emulated using a sequence of instructions, such as, a comparison followed by AND, ANDN and OR.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register

and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the corresponding YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

```
MIN(SRC1, SRC2)
{
    IF ((SRC1 = 0.0) and (SRC2 = 0.0)) THEN DEST ← SRC2;
        ELSE IF (SRC1 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF (SRC2 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF (SRC1 < SRC2) THEN DEST ← SRC1;
        ELSE DEST ← SRC2;
    FI;
}

MINPS (128-bit Legacy SSE version)
DEST[31:0] ← MIN(SRC1[31:0], SRC2[31:0])
DEST[63:32] ← MIN(SRC1[63:32], SRC2[63:32])
DEST[95:64] ← MIN(SRC1[95:64], SRC2[95:64])
DEST[127:96] ← MIN(SRC1[127:96], SRC2[127:96])
DEST[VLMAX-1:128] (Unmodified)

VMINPS (VEX.128 encoded version)
DEST[31:0] ← MIN(SRC1[31:0], SRC2[31:0])
DEST[63:32] ← MIN(SRC1[63:32], SRC2[63:32])
DEST[95:64] ← MIN(SRC1[95:64], SRC2[95:64])
DEST[127:96] ← MIN(SRC1[127:96], SRC2[127:96])
DEST[VLMAX-1:128] ← 0

VMINPS (VEX.256 encoded version)
DEST[31:0] ← MIN(SRC1[31:0], SRC2[31:0])
DEST[63:32] ← MIN(SRC1[63:32], SRC2[63:32])
DEST[95:64] ← MIN(SRC1[95:64], SRC2[95:64])
DEST[127:96] ← MIN(SRC1[127:96], SRC2[127:96])
DEST[159:128] ← MIN(SRC1[159:128], SRC2[159:128])
DEST[191:160] ← MIN(SRC1[191:160], SRC2[191:160])
DEST[223:192] ← MIN(SRC1[223:192], SRC2[223:192])
```

DEST[255:224] ← MIN(SRC1[255:224], SRC2[255:224])

Intel C/C++ Compiler Intrinsic Equivalent

MINPS __m128d _mm_min_ps(__m128d a, __m128d b);

VMINPS __m256 _mm256_min_ps (__m256 a, __m256 b);

SIMD Floating-Point Exceptions

Invalid (including QNaN source operand), Denormal.

Other Exceptions

See Exceptions Type 2.

MINSD—Return Minimum Scalar Double-Precision Floating-Point Value

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 5D /*r* MINSD *xmm1*, *xmm2/m64*	A	V/V	SSE2	Return the minimum scalar double-precision floating-point value between *xmm2/mem64* and *xmm1*.
VEX.NDS.LIG.F2.0F.WIG 5D /r VMINSD xmm1, xmm2, xmm3/m64	B	V/V	AVX	Return the minimum scalar double precision floating-point value between xmm3/mem64 and xmm2.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Compares the low double-precision floating-point values in the first source operand and the second source operand, and returns the minimum value to the low quadword of the destination operand. When the source operand is a memory operand, only the 64 bits are accessed. The high quadword of the destination operand is copied from the same bits in the first source operand.

If the values being compared are both 0.0s (of either sign), the value in the second source operand is returned. If a value in the second source operand is an SNaN, that SNaN is returned unchanged to the destination (that is, a QNaN version of the SNaN is not returned).

If only one value is a NaN (SNaN or QNaN) for this instruction, the second source operand, either a NaN or a valid floating-point value, is written to the result. If instead of this behavior, it is required that the NaN source operand (from either the first or second source) be returned, the action of MINSD can be emulated using a sequence of instructions, such as, a comparison followed by AND, ANDN and OR.

The second source operand can be an XMM register or a 64-bit memory location. The first source and destination operands are XMM registers.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The destination and first source operand are the same. Bits (VLMAX-1:64) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (127:64) of the XMM register destination are copied from corresponding bits in the first source operand. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

MIN(SRC1, SRC2)

```
{
    IF ((SRC1 = 0.0) and (SRC2 = 0.0)) THEN DEST ← SRC2;
        ELSE IF (SRC1 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF SRC2 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF (SRC1 < SRC2) THEN DEST ← SRC1;
        ELSE DEST ← SRC2;
    FI;
}
```

MINSD (128-bit Legacy SSE version)

```
DEST[63:0] ← MIN(SRC1[63:0], SRC2[63:0])
DEST[VLMAX-1:64] (Unmodified)
```

MINSD (VEX.128 encoded version)

```
DEST[63:0] ← MIN(SRC1[63:0], SRC2[63:0])
DEST[127:64] ← SRC1[127:64]
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

MINSD __m128d _mm_min_sd(__m128d a, __m128d b)

SIMD Floating-Point Exceptions

Invalid (including QNaN source operand), Denormal.

Other Exceptions

See Exceptions Type 3.

MINSS—Return Minimum Scalar Single-Precision Floating-Point Value

Opcode/ Instruction		Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 5D /*r* MINSS *xmm1*, *xmm2/m32*		A	V/V	SSE	Return the minimum scalar single-precision floating-point value between *xmm2/mem32* and *xmm1*.
VEX.NDS.LIG.F3.0F.WIG 5D /r	VMINSS xmm1,xmm2, xmm3/m32	B	V/V	AVX	Return the minimum scalar single precision floating-point value between xmm3/mem32 and xmm2.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Compares the low single-precision floating-point values in the first source operand and the second source operand and returns the minimum value to the low doubleword of the destination operand.

If the values being compared are both 0.0s (of either sign), the value in the second source operand is returned. If a value in the second operand is an SNaN, that SNaN is returned unchanged to the destination (that is, a QNaN version of the SNaN is not returned).

If only one value is a NaN (SNaN or QNaN) for this instruction, the second source operand, either a NaN or a valid floating-point value, is written to the result. If instead of this behavior, it is required that the NaN in either source operand be returned, the action of MINSD can be emulated using a sequence of instructions, such as, a comparison followed by AND, ANDN and OR.

The second source operand can be an XMM register or a 32-bit memory location. The first source and destination operands are XMM registers.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The destination and first source operand are the same. Bits (VLMAX-1:32) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (127:32) of the XMM register destination are copied from corresponding bits in the first source operand. Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

MIN(SRC1, SRC2)

```
{
    IF ((SRC1 = 0.0) and (SRC2 = 0.0)) THEN DEST ← SRC2;
        ELSE IF (SRC1 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF SRC2 = SNaN) THEN DEST ← SRC2; FI;
        ELSE IF (SRC1 < SRC2) THEN DEST ← SRC1;
        ELSE DEST ← SRC2;
    FI;
}
```

MINSS (128-bit Legacy SSE version)

```
DEST[31:0] ← MIN(SRC1[31:0], SRC2[31:0])
DEST[VLMAX-1:32] (Unmodified)
```

VMINSS (VEX.128 encoded version)

```
DEST[31:0] ← MIN(SRC1[31:0], SRC2[31:0])
DEST[127:32] ← SRC1[127:32]
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

MINSS __m128d _mm_min_ss(__m128d a, __m128d b)

SIMD Floating-Point Exceptions

Invalid (including QNaN source operand), Denormal.

Other Exceptions

See Exceptions Type 3.

MONITOR—Set Up Monitor Address

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 01 *C8*	MONITOR	A	Valid	Valid	Sets up a linear address range to be monitored by hardware and activates the monitor. The address range should be a write-back memory caching type. The address is DS:EAX (DS:RAX in 64-bit mode).

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

The MONITOR instruction arms address monitoring hardware using an address specified in EAX (the address range that the monitoring hardware checks for store operations can be determined by using CPUID). A store to an address within the specified address range triggers the monitoring hardware. The state of monitor hardware is used by MWAIT.

The content of EAX is an effective address (in 64-bit mode, RAX is used). By default, the DS segment is used to create a linear address that is monitored. Segment overrides can be used.

ECX and EDX are also used. They communicate other information to MONITOR. ECX specifies optional extensions. EDX specifies optional hints; it does not change the architectural behavior of the instruction. For the Pentium 4 processor (family 15, model 3), no extensions or hints are defined. Undefined hints in EDX are ignored by the processor; undefined extensions in ECX raises a general protection fault.

The address range must use memory of the write-back type. Only write-back memory will correctly trigger the monitoring hardware. Additional information on determining what address range to use in order to prevent false wake-ups is described in Chapter 8, "Multiple-Processor Management" of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*.

The MONITOR instruction is ordered as a load operation with respect to other memory transactions. The instruction is subject to the permission checking and faults associated with a byte load. Like a load, MONITOR sets the A-bit but not the D-bit in page tables.

The MONITOR CPUID feature flag (ECX bit 3; CPUID executed EAX = 1) indicates the availability of MONITOR and MWAIT in the processor. When set, MONITOR may be executed only at privilege level 0 (use at any other privilege level results in an invalid-opcode exception). The operating system or system BIOS may disable this instruction by using the IA32_MISC_ENABLE MSR; disabling MONITOR clears the CPUID feature flag and causes execution to generate an illegal opcode exception.

The instruction's operation is the same in non-64-bit modes and 64-bit mode.

Operation

MONITOR sets up an address range for the monitor hardware using the content of EAX (RAX in 64-bit mode) as an effective address and puts the monitor hardware in armed state. Always use memory of the write-back caching type. A store to the specified address range will trigger the monitor hardware. The content of ECX and EDX are used to communicate other information to the monitor hardware.

Intel C/C++ Compiler Intrinsic Equivalent

MONITOR void _mm_monitor(void const *p, unsigned extensions,unsigned hints)

Numeric Exceptions

None

Protected Mode Exceptions

#GP(0)	If the value in EAX is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register is used to access memory and it contains a NULL segment selector.
	If ECX $\neq$ 0.
#SS(0)	If the value in EAX is outside the SS segment limit.
#PF(fault-code)	For a page fault.
#UD	If CPUID.01H:ECX.MONITOR[bit 3] = 0.
	If current privilege level is not 0.

Real Address Mode Exceptions

#GP	If the CS, DS, ES, FS, or GS register is used to access memory and the value in EAX is outside of the effective address space from 0 to FFFFH.
	If ECX $\neq$ 0.
#SS	If the SS register is used to access memory and the value in EAX is outside of the effective address space from 0 to FFFFH.
#UD	If CPUID.01H:ECX.MONITOR[bit 3] = 0.

Virtual 8086 Mode Exceptions

#UD	The MONITOR instruction is not recognized in virtual-8086 mode (even if CPUID.01H:ECX.MONITOR[bit 3] = 1).

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(0)	If the linear address of the operand in the CS, DS, ES, FS, or GS segment is in a non-canonical form.
	If RCX ≠ 0.
#SS(0)	If the SS register is used to access memory and the value in EAX is in a non-canonical form.
#PF(fault-code)	For a page fault.
#UD	If the current privilege level is not 0.
	If CPUID.01H:ECX.MONITOR[bit 3] = 0.

MOV—Move

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
88 */r*	MOV *r/m8,r8*	A	Valid	Valid	Move *r8* to *r/m8.*
REX + 88 */r*	MOV *r/m8****,*r8****	A	Valid	N.E.	Move *r8* to *r/m8.*
89 */r*	MOV *r/m16,r16*	A	Valid	Valid	Move *r16* to *r/m16.*
89 */r*	MOV *r/m32,r32*	A	Valid	Valid	Move *r32* to *r/m32.*
REX.W + 89 */r*	MOV *r/m64,r64*	A	Valid	N.E.	Move *r64* to *r/m64.*
8A */r*	MOV *r8,r/m8*	B	Valid	Valid	Move *r/m8* to *r8.*
REX + 8A */r*	MOV *r8***,r/m8****	B	Valid	N.E.	Move *r/m8* to *r8.*
8B */r*	MOV *r16,r/m16*	B	Valid	Valid	Move *r/m16* to *r16.*
8B */r*	MOV *r32,r/m32*	B	Valid	Valid	Move *r/m32* to *r32.*
REX.W + 8B */r*	MOV *r64,r/m64*	B	Valid	N.E.	Move *r/m64* to *r64.*
8C */r*	MOV *r/m16,Sreg***	A	Valid	Valid	Move segment register to *r/m16.*
REX.W + 8C */r*	MOV *r/m64,Sreg***	A	Valid	Valid	Move zero extended 16-bit segment register to *r/m64.*
8E */r*	MOV *Sreg,r/m16***	B	Valid	Valid	Move *r/m16* to segment register.
REX.W + 8E */r*	MOV *Sreg,r/m64***	B	Valid	Valid	Move *lower 16 bits of r/m64* to segment register.
A0	MOV AL,*moffs8**	C	Valid	Valid	Move byte at (*seg:offset*) to AL.
REX.W + A0	MOV AL,*moffs8**	C	Valid	N.E.	Move byte at (*offset*) to AL.
A1	MOV AX,*moffs16**	C	Valid	Valid	Move word at (*seg:offset*) to AX.
A1	MOV EAX,*moffs32**	C	Valid	Valid	Move doubleword at (*seg:offset*) to EAX.
REX.W + A1	MOV RAX,*moffs64**	C	Valid	N.E.	Move quadword at (*offset*) to RAX.
A2	MOV *moffs8*,AL	D	Valid	Valid	Move AL to (*seg:offset*).
REX.W + A2	MOV *moffs8****,AL	D	Valid	N.E.	Move AL to (*offset*).
A3	MOV *moffs16**,AX	D	Valid	Valid	Move AX to (*seg:offset*).
A3	MOV *moffs32**,EAX	D	Valid	Valid	Move EAX to (*seg:offset*).

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
REX.W + A3	MOV *moffs64**,RAX	D	Valid	N.E.	Move RAX to (*offset*).
B0+ *rb*	MOV *r8, imm8*	E	Valid	Valid	Move *imm8* to *r8.*
REX + B0+ *rb*	MOV *r8****, *imm8*	E	Valid	N.E.	Move *imm8* to *r8.*
B8+ *rw*	MOV *r16, imm16*	E	Valid	Valid	Move *imm16* to *r16.*
B8+ *rd*	MOV *r32, imm32*	E	Valid	Valid	Move *imm32* to *r32.*
REX.W + B8+ *rd*	MOV *r64, imm64*	E	Valid	N.E.	Move *imm64* to *r64.*
C6 */0*	MOV *r/m8, imm8*	F	Valid	Valid	Move *imm8* to *r/m8.*
REX + C6 */0*	MOV *r/m8****, imm8*	F	Valid	N.E.	Move *imm8* to *r/m8.*
C7 */0*	MOV *r/m16, imm16*	F	Valid	Valid	Move *imm16* to *r/m16.*
C7 */0*	MOV *r/m32, imm32*	F	Valid	Valid	Move *imm32* to *r/m32.*
REX.W + C7 */0*	MOV *r/m64, imm32*	F	Valid	N.E.	Move *imm32 sign extended to 64-bits* to *r/m64.*

NOTES:

* The *moffs8*, *moffs16*, *moffs32* and *moffs64* operands specify a simple offset relative to the segment base, where 8, 16, 32 and 64 refer to the size of the data. The address-size attribute of the instruction determines the size of the offset, either 16, 32 or 64 bits.

** In 32-bit mode, the assembler may insert the 16-bit operand-size prefix with this instruction (see the following "Description" section for further information).

***In 64-bit mode, *r/m8* can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (w)	ModRM:reg (r)	NA	NA
B	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
C	AL/AX/EAX/RAX	Displacement	NA	NA
D	Displacement	AL/AX/EAX/RAX	NA	NA
E	reg (w)	imm8/16/32/64	NA	NA
F	ModRM:r/m (w)	imm8/16/32/64	NA	NA

Description

Copies the second operand (source operand) to the first operand (destination operand). The source operand can be an immediate value, general-purpose register, segment register, or memory location; the destination register can be a general-purpose register, segment register, or memory location. Both operands must be the same size, which can be a byte, a word, a doubleword, or a quadword.

The MOV instruction cannot be used to load the CS register. Attempting to do so results in an invalid opcode exception (#UD). To load the CS register, use the far JMP, CALL, or RET instruction.

If the destination operand is a segment register (DS, ES, FS, GS, or SS), the source operand must be a valid segment selector. In protected mode, moving a segment selector into a segment register automatically causes the segment descriptor information associated with that segment selector to be loaded into the hidden (shadow) part of the segment register. While loading this information, the segment selector and segment descriptor information is validated (see the “Operation” algorithm below). The segment descriptor data is obtained from the GDT or LDT entry for the specified segment selector.

A NULL segment selector (values 0000-0003) can be loaded into the DS, ES, FS, and GS registers without causing a protection exception. However, any subsequent attempt to reference a segment whose corresponding segment register is loaded with a NULL value causes a general protection exception (#GP) and no memory reference occurs.

Loading the SS register with a MOV instruction inhibits all interrupts until after the execution of the next instruction. This operation allows a stack pointer to be loaded into the ESP register with the next instruction (MOV ESP, **stack-pointer value**) before an interrupt occurs[1]. Be aware that the LSS instruction offers a more efficient method of loading the SS and ESP registers.

When operating in 32-bit mode and moving data between a segment register and a general-purpose register, the 32-bit IA-32 processors do not require the use of the 16-bit operand-size prefix (a byte with the value 66H) with this instruction, but most assemblers will insert it if the standard form of the instruction is used (for example, MOV DS, AX). The processor will execute this instruction correctly, but it will usually require an extra clock. With most assemblers, using the instruction form MOV DS, EAX will avoid this unneeded 66H prefix. When the processor executes the instruction with a 32-bit general-purpose register, it assumes that the 16 least-significant bits of the general-purpose register are the destination or source operand. If the register is a destination operand, the resulting value in the two high-order bytes of the register is implementation dependent. For the Pentium 4, Intel Xeon, and P6 family processors, the two high-order bytes are filled with zeros; for earlier 32-bit IA-32 processors, the two high order bytes are undefined.

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
DEST ← SRC;
```

Loading a segment register while in protected mode results in special checks and actions, as described in the following listing. These checks are performed on the segment selector and the segment descriptor to which it points.

```
IF SS is loaded
   THEN
      IF segment selector is NULL
         THEN #GP(0); FI;
      IF segment selector index is outside descriptor table limits
      or segment selector's RPL ≠ CPL
      or segment is not a writable data segment
      or DPL ≠ CPL
         THEN #GP(selector); FI;
      IF segment not marked present
         THEN #SS(selector);
         ELSE
            SS ← segment selector;
            SS ← segment descriptor; FI;
```

1. If a code instruction breakpoint (for debug) is placed on an instruction located immediately after a MOV SS instruction, the breakpoint may not be triggered. However, in a sequence of instructions that load the SS register, only the first instruction in the sequence is guaranteed to delay an interrupt.

   In the following sequence, interrupts may be recognized before MOV ESP, EBP executes:

   MOV SS, EDX
   MOV SS, EAX
   MOV ESP, EBP

```
FI;

IF DS, ES, FS, or GS is loaded with non-NULL selector
THEN
    IF segment selector index is outside descriptor table limits
    or segment is not a data or readable code segment
    or ((segment is a data or nonconforming code segment)
    and (both RPL and CPL > DPL))
        THEN #GP(selector); FI;
    IF segment not marked present
        THEN #NP(selector);
        ELSE
            SegmentRegister ← segment selector;
            SegmentRegister ← segment descriptor; FI;
FI;

IF DS, ES, FS, or GS is loaded with NULL selector
    THEN
        SegmentRegister ← segment selector;
        SegmentRegister ← segment descriptor;
FI;
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If attempt is made to load SS register with NULL segment selector.
	If the destination operand is in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#GP(selector)	If segment selector index is outside descriptor table limits.
	If the SS register is being loaded and the segment selector's RPL and the segment descriptor’s DPL are not equal to the CPL.
	If the SS register is being loaded and the segment pointed to is a non-writable data segment.
	If the DS, ES, FS, or GS register is being loaded and the segment pointed to is not a data or readable code segment.
	If the DS, ES, FS, or GS register is being loaded and the segment pointed to is a data or nonconforming code segment, but both the RPL and the CPL are greater than the DPL.

#SS(0)	If a memory operand effective address is outside the SS segment limit.
#SS(selector)	If the SS register is being loaded and the segment pointed to is marked not present.
#NP	If the DS, ES, FS, or GS register is being loaded and the segment pointed to is marked not present.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If attempt is made to load the CS register.
	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If attempt is made to load the CS register.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If attempt is made to load the CS register.
	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(0)	If the memory address is in a non-canonical form.
	If an attempt is made to load SS register with NULL segment selector when CPL = 3.
	If an attempt is made to load SS register with NULL segment selector when CPL < 3 and CPL ≠ RPL.

#GP(selector)	If segment selector index is outside descriptor table limits.
	If the memory access to the descriptor table is non-canonical.
	If the SS register is being loaded and the segment selector's RPL and the segment descriptor’s DPL are not equal to the CPL.
	If the SS register is being loaded and the segment pointed to is a nonwritable data segment.
	If the DS, ES, FS, or GS register is being loaded and the segment pointed to is not a data or readable code segment.
	If the DS, ES, FS, or GS register is being loaded and the segment pointed to is a data or nonconforming code segment, but both the RPL and the CPL are greater than the DPL.
#SS(0)	If the stack address is in a non-canonical form.
#SS(selector)	If the SS register is being loaded and the segment pointed to is marked not present.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If attempt is made to load the CS register.
	If the LOCK prefix is used.

MOV—Move to/from Control Registers

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 20/*r*	MOV *r32*, CR0–CR7	A	N.E.	Valid	Move control register to *r32*
0F 20/*r*	MOV *r64*, CR0–CR7	A	Valid	N.E.	Move extended control register to *r64*.
REX.R + 0F 20 /0	MOV *r64*, CR8	A	Valid	N.E.	Move extended CR8 to *r64*.[1]
0F 22 /*r*	MOV CR0–CR7, *r32*	A	N.E.	Valid	Move *r32* to control register
0F 22 /*r*	MOV CR0–CR7, *r64*	A	Valid	N.E.	Move *r64* to extended control register.
REX.R + 0F 22 /0	MOV CR8, *r64*	A	Valid	N.E.	Move *r64* to extended CR8.[1]

NOTE:

1. MOV CR* instructions, except for MOV CR8, are serializing instructions. MOV CR8 is not architecturally defined as a serializing instruction. For more information, see Chapter 8 in *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Moves the contents of a control register (CR0, CR2, CR3, CR4, or CR8) to a general-purpose register or the contents of a general purpose register to a control register. The operand size for these instructions is always 32 bits in non-64-bit modes, regardless of the operand-size attribute. (See "Control Registers" in Chapter 2 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*, for a detailed description of the flags and fields in the control registers.) This instruction can be executed only when the current privilege level is 0.

At the opcode level, the *reg* field within the ModR/M byte specifies which of the control registers is loaded or read. The 2 bits in the *mod* field are ignored. The *r/m* field specifies the general-purpose register loaded or read. Attempts to reference CR1, CR5, CR6, CR7, and CR9–CR15 result in undefined opcode (#UD) exceptions.

When loading control registers, programs should not attempt to change the reserved bits; that is, always set reserved bits to the value previously read. An attempt to change CR4's reserved bits will cause a general protection fault. Reserved bits in CR0 and CR3 remain clear after any load of those registers; attempts to set them have no impact. On Pentium 4, Intel Xeon and P6 family processors, CR0.ET remains set after any load of CR0; attempts to clear this bit have no impact.

In certain cases, these instructions have the side effect of invalidating entries in the TLBs and the paging-structure caches. See Section 4.10.4.1, "Operations that Invalidate TLBs and Paging-Structure Caches," in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A* for details.

The following side effects are implementation-specific for the Pentium 4, Intel Xeon, and P6 processor family: when modifying PE or PG in register CR0, or PSE or PAE in register CR4, all TLB entries are flushed, including global entries. Software should not depend on this functionality in all Intel 64 or IA-32 processors.

In 64-bit mode, the instruction's default operation size is 64 bits. The REX.R prefix must be used to access CR8. Use of REX.B permits access to additional registers (R8-R15). Use of the REX.W prefix or 66H prefix is ignored. Use of the REX.R prefix to specify a register other than CR8 causes an invalid-opcode exception. See the summary chart at the beginning of this section for encoding data and limits.

If CR4.PCIDE = 1, bit 63 of the source operand to MOV to CR3 determines whether the instruction invalidates entries in the TLBs and the paging-structure caches (see Section 4.10.4.1, "Operations that Invalidate TLBs and Paging-Structure Caches," in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*). The instruction does not modify bit 63 of CR3, which is reserved and always 0.

See "Changes to Instruction Behavior in VMX Non-Root Operation" in Chapter 22 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3B*, for more information about the behavior of this instruction in VMX non-root operation.

Operation

```
DEST ← SRC;
```

Flags Affected

The OF, SF, ZF, AF, PF, and CF flags are undefined.

Protected Mode Exceptions

#GP(0)	If the current privilege level is not 0.
	If an attempt is made to write invalid bit combinations in CR0 (such as setting the PG flag to 1 when the PE flag is set to 0, or setting the CD flag to 0 when the NW flag is set to 1).
	If an attempt is made to write a 1 to any reserved bit in CR4.
	If an attempt is made to write 1 to CR4.PCIDE.
	If any of the reserved bits are set in the page-directory pointers table (PDPT) and the loading of a control register causes the PDPT to be loaded into the processor.
#UD	If the LOCK prefix is used.

If an attempt is made to access CR1, CR5, CR6, or CR7.

Real-Address Mode Exceptions

#GP	If an attempt is made to write a 1 to any reserved bit in CR4.
	If an attempt is made to write 1 to CR4.PCIDE.
	If an attempt is made to write invalid bit combinations in CR0 (such as setting the PG flag to 1 when the PE flag is set to 0).
#UD	If the LOCK prefix is used.
	If an attempt is made to access CR1, CR5, CR6, or CR7.

Virtual-8086 Mode Exceptions

#GP(0)	These instructions cannot be executed in virtual-8086 mode.

Compatibility Mode Exceptions

#GP(0)	If the current privilege level is not 0.
	If an attempt is made to write invalid bit combinations in CR0 (such as setting the PG flag to 1 when the PE flag is set to 0, or setting the CD flag to 0 when the NW flag is set to 1).
	If an attempt is made to change CR4.PCIDE from 0 to 1 while CR3[11:0] ≠ 000H.
	If an attempt is made to clear CR0.PG[bit 31] while CR4.PCIDE = 1.
	If an attempt is made to write a 1 to any reserved bit in CR3.
	If an attempt is made to leave IA-32e mode by clearing CR4.PAE[bit 5].
#UD	If the LOCK prefix is used.
	If an attempt is made to access CR1, CR5, CR6, or CR7.

64-Bit Mode Exceptions

#GP(0)	If the current privilege level is not 0.
	If an attempt is made to write invalid bit combinations in CR0 (such as setting the PG flag to 1 when the PE flag is set to 0, or setting the CD flag to 0 when the NW flag is set to 1).
	If an attempt is made to change CR4.PCIDE from 0 to 1 while CR3[11:0] ≠ 000H.
	If an attempt is made to clear CR0.PG[bit 31].
	If an attempt is made to write a 1 to any reserved bit in CR4.
	If an attempt is made to write a 1 to any reserved bit in CR8.
	If an attempt is made to write a 1 to any reserved bit in CR3.

	If an attempt is made to leave IA-32e mode by clearing CR4.PAE[bit 5].
#UD	If the LOCK prefix is used.
	If an attempt is made to access CR1, CR5, CR6, or CR7.
	If the REX.R prefix is used to specify a register other than CR8.

MOV—Move to/from Debug Registers

Opcode	Instruction	Op/En	64-Bit Mode	Compat/Leg Mode	Description
0F 21/*r*	MOV *r32,* DR0–DR7	A	N.E.	Valid	Move debug register to *r32*
0F 21/*r*	MOV *r64,* DR0–DR7	A	Valid	N.E.	Move extended debug register to *r64*.
0F 23 /*r*	MOV DR0–DR7, *r32*	A	N.E.	Valid	Move *r32* to debug register
0F 23 /*r*	MOV DR0–DR7, *r64*	A	Valid	N.E.	Move *r64* to extended debug register.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Moves the contents of a debug register (DR0, DR1, DR2, DR3, DR4, DR5, DR6, or DR7) to a general-purpose register or vice versa. The operand size for these instructions is always 32 bits in non-64-bit modes, regardless of the operand-size attribute. (See Chapter 20, "Introduction to Virtual-Machine Extensions", of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*, for a detailed description of the flags and fields in the debug registers.)

The instructions must be executed at privilege level 0 or in real-address mode.

When the debug extension (DE) flag in register CR4 is clear, these instructions operate on debug registers in a manner that is compatible with Intel386 and Intel486 processors. In this mode, references to DR4 and DR5 refer to DR6 and DR7, respectively. When the DE flag in CR4 is set, attempts to reference DR4 and DR5 result in an undefined opcode (#UD) exception. (The CR4 register was added to the IA-32 Architecture beginning with the Pentium processor.)

At the opcode level, the *reg* field within the ModR/M byte specifies which of the debug registers is loaded or read. The two bits in the *mod* field are ignored. The *r/m* field specifies the general-purpose register loaded or read.

In 64-bit mode, the instruction's default operation size is 64 bits. Use of the REX.B prefix permits access to additional registers (R8–R15). Use of the REX.W or 66H prefix is ignored. Use of the REX.R prefix causes an invalid-opcode exception. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
IF ((DE = 1) and (SRC or DEST = DR4 or DR5))
    THEN
        #UD;
    ELSE
        DEST ← SRC;

FI;
```

Flags Affected

The OF, SF, ZF, AF, PF, and CF flags are undefined.

Protected Mode Exceptions

#GP(0)	If the current privilege level is not 0.
#UD	If CR4.DE[bit 3] = 1 (debug extensions) and a MOV instruction is executed involving DR4 or DR5.
	If the LOCK prefix is used.
#DB	If any debug register is accessed while the DR7.GD[bit 13] = 1.

Real-Address Mode Exceptions

#UD	If CR4.DE[bit 3] = 1 (debug extensions) and a MOV instruction is executed involving DR4 or DR5.
	If the LOCK prefix is used.
#DB	If any debug register is accessed while the DR7.GD[bit 13] = 1.

Virtual-8086 Mode Exceptions

#GP(0)	The debug registers cannot be loaded or read when in virtual-8086 mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(0)	If the current privilege level is not 0.
#UD	If CR4.DE[bit 3] = 1 (debug extensions) and a MOV instruction is executed involving DR4 or DR5.
	If the LOCK prefix is used.
	If the REX.R prefix is used.
#DB	If any debug register is accessed while the DR7.GD[bit 13] = 1.

MOVAPD—Move Aligned Packed Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 28 /*r* MOVAPD *xmm1*, *xmm2/m128*	A	V/V	SSE2	Move packed double-precision floating-point values from *xmm2/m128* to *xmm1*.
66 0F 29 /*r* MOVAPD *xmm2/m128*, *xmm1*	B	V/V	SSE2	Move packed double-precision floating-point values from *xmm1* to *xmm2/m128*.
VEX.128.66.0F.WIG 28 /r VMOVAPD xmm1, xmm2/m128	A	V/V	AVX	Move aligned packed double-precision floating-point values from xmm2/mem to xmm1.
VEX.128.66.0F.WIG 29 /r VMOVAPD xmm2/m128, xmm1	B	V/V	AVX	Move aligned packed double-precision floating-point values from xmm1 to xmm2/mem.
VEX.256.66.0F.WIG 28 /r VMOVAPD ymm1, ymm2/m256	A	V/V	AVX	Move aligned packed double-precision floating-point values from ymm2/mem to ymm1.
VEX.256.66.0F.WIG 29 /r VMOVAPD ymm2/m256, ymm1	B	V/V	AVX	Move aligned packed double-precision floating-point values from ymm1 to ymm2/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

Moves 2 or 4 double-precision floating-point values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load an XMM or YMM register from an 128-bit or 256-bit memory location, to store the contents of an XMM or YMM register into a 128-bit or 256-bit memory location, or to move data between two XMM or two YMM registers. When the source or destination operand is a memory operand, the operand must be aligned on a 16-byte (128-bit version) or 32-byte (VEX.256 encoded version) boundary or a general-protection exception (#GP) will be generated.

To move double-precision floating-point values to and from unaligned memory locations, use the (V)MOVUPD instruction.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

128-bit versions:

Moves 128 bits of packed double-precision floating-point values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load an XMM register from a 128-bit memory location, to store the contents of an XMM register into a 128-bit memory location, or to move data between two XMM registers. When the source or destination operand is a memory operand, the operand must be aligned on a 16-byte boundary or a general-protection exception (#GP) will be generated. To move single-precision floating-point values to and from unaligned memory locations, use the VMOVUPD instruction.

128-bit Legacy SSE version: Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register destination are zeroed.

VEX.256 encoded version:

Moves 256 bits of packed double-precision floating-point values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load a YMM register from a 256-bit memory location, to store the contents of a YMM register into a 256-bit memory location, or to move data between two YMM registers. When the source or destination operand is a memory operand, the operand must be aligned on a 32-byte boundary or a general-protection exception (#GP) will be generated. To move single-precision floating-point values to and from unaligned memory locations, use the VMOVUPD instruction.

Operation

MOVAPD (128-bit load- and register-copy- form Legacy SSE version)

DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] (Unmodified)

(V)MOVAPD (128-bit store-form version)

DEST[127:0] ← SRC[127:0]

VMOVAPD (VEX.128 encoded version)

DEST[127:0] ← SRC[127:0]

DEST[VLMAX-1:128] ← 0

VMOVAPD (VEX.256 encoded version)
DEST[255:0] ← SRC[255:0]

Intel C/C++ Compiler Intrinsic Equivalent

MOVAPD __m128d _mm_load_pd (double const * p);

MOVAPD _mm_store_pd(double * p, __m128d a);

VMOVAPD __m256d _mm256_load_pd (double const * p);

VMOVAPD _mm256_store_pd(double * p, __m256d a);

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 1.SSE2; additionally

#UD	If VEX.vvvv != 1111B.

MOVAPS—Move Aligned Packed Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 28 /*r* MOVAPS *xmm1*, *xmm2/m128*	A	V/V	SSE	Move packed single-precision floating-point values from *xmm2/m128* to *xmm1*.
0F 29 /*r* MOVAPS *xmm2/m128*, *xmm1*	B	V/V	SSE	Move packed single-precision floating-point values from *xmm1* to *xmm2/m128*.
VEX.128.0F.WIG 28 /r VMOVAPS xmm1, xmm2/m128	A	V/V	AVX	Move aligned packed single-precision floating-point values from xmm2/mem to xmm1.
VEX.128.0F.WIG 29 /r VMOVAPS xmm2/m128, xmm1	B	V/V	AVX	Move aligned packed single-precision floating-point values from xmm1 to xmm2/mem.
VEX.256.0F.WIG 28 /r VMOVAPS ymm1, ymm2/m256	A	V/V	AVX	Move aligned packed single-precision floating-point values from ymm2/mem to ymm1.
VEX.256.0F.WIG 29 /r VMOVAPS ymm2/m256, ymm1	B	V/V	AVX	Move aligned packed single-precision floating-point values from ymm1 to ymm2/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

Moves 4 or8 single-precision floating-point values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load an XMM or YMM register from an 128-bit or 256-bit memory location, to store the contents of an XMM or YMM register into a 128-bit or 256-bit memory location, or to move data between two XMM or two YMM registers. When the source or destination operand is a memory operand, the operand must be aligned on a 16-byte (128-

bit version) or 32-byte (VEX.256 encoded version) boundary or a general-protection exception (#GP) will be generated.

To move single-precision floating-point values to and from unaligned memory locations, use the (V)MOVUPS instruction.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

128-bit versions:

Moves 128 bits of packed single-precision floating-point values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load an XMM register from a 128-bit memory location, to store the contents of an XMM register into a 128-bit memory location, or to move data between two XMM registers. When the source or destination operand is a memory operand, the operand must be aligned on a 16-byte boundary or a general-protection exception (#GP) will be generated. To move single-precision floating-point values to and from unaligned memory locations, use the VMOVUPS instruction.

128-bit Legacy SSE version: Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

VEX.256 encoded version:

Moves 256 bits of packed single-precision floating-point values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load a YMM register from a 256-bit memory location, to store the contents of a YMM register into a 256-bit memory location, or to move data between two YMM registers.

Operation

MOVAPS (128-bit load- and register-copy- form Legacy SSE version)

DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] (Unmodified)

(V)MOVAPS (128-bit store form)

DEST[127:0] ← SRC[127:0]

VMOVAPS (VEX.128 encoded version)

DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] ← 0

VMOVAPS (VEX.256 encoded version)

DEST[255:0] ← SRC[255:0]

Intel C/C++ Compiler Intrinsic Equivalent

MOVAPS __m128 _mm_load_ps (float const * p);

MOVAPS _mm_store_ps(float * p, __m128 a);

VMOVAPS __m256 _mm256_load_ps (float const * p);

VMOVAPS _mm256_store_ps(float * p, __m256 a);

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 1.SSE; additionally

#UD	If VEX.vvvv != 1111B.

MOVBE—Move Data After Swapping Bytes

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 38 F0 /*r*	MOVBE *r16*, m16	A	Valid	Valid	Reverse byte order in *m16 and move to r16*
0F 38 F0 /*r*	MOVBE *r32*, m32	A	Valid	Valid	Reverse byte order in *m32 and move to r32*
REX.W + 0F 38 F0 /*r*	MOVBE *r64*, m64	A	Valid	N.E.	Reverse byte order in *m64 and move to r64*.
0F 38 F1 /*r*	MOVBE *m16*, r16	B	Valid	Valid	Reverse byte order in *r16 and move to m16*
0F 38 F1 /*r*	MOVBE *m32*, r32	B	Valid	Valid	Reverse byte order in *r32 and move to m32*
REX.W + 0F 38 F1 /*r*	MOVBE *m64*, r64	B	Valid	N.E.	Reverse byte order in *r64 and move to m64*.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

Performs a byte swap operation on the data copied from the second operand (source operand) and store the result in the first operand (destination operand). The source operand can be a general-purpose register, or memory location; the destination register can be a general-purpose register, or a memory location; however, both operands can not be registers, and only one operand can be a memory location. Both operands must be the same size, which can be a word, a doubleword or quadword.

The MOVBE instruction is provided for swapping the bytes on a read from memory or on a write to memory; thus providing support for converting little-endian values to big-endian format and vice versa.

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

TEMP ← SRC

```
IF ( OperandSize = 16)
    THEN
        DEST[7:0] ← TEMP[15:8];
        DEST[15:8] ← TEMP[7:0];
    ELES IF    ( OperandSize = 32)
        DEST[7:0] ← TEMP[31:24];
        DEST[15:8] ← TEMP[23:16];
        DEST[23:16] ← TEMP[15:8];
        DEST[31:23] ← TEMP[7:0];
    ELSE IF ( OperandSize = 64)
        DEST[7:0] ← TEMP[63:56];
        DEST[15:8] ← TEMP[55:48];
        DEST[23:16] ← TEMP[47:40];
        DEST[31:24] ← TEMP[39:32];
        DEST[39:32] ← TEMP[31:24];
        DEST[47:40] ← TEMP[23:16];
        DEST[55:48] ← TEMP[15:8];
        DEST[63:56] ← TEMP[7:0];

FI;
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If the destination operand is in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If CPUID.01H:ECX.MOVBE[bit 22] = 0 .
	If the LOCK prefix is used.
	If REP (F3H) prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.

#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If CPUID.01H:ECX.MOVBE[bit 22] = 0 .
	If the LOCK prefix is used.
	If REP (F3H) prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#UD	If CPUID.01H:ECX.MOVBE[bit 22] = 0 .
	If the LOCK prefix is used.
	If REP (F3H) prefix is used.
	If REPNE (F2H) prefix is used and CPUID.01H:ECX.SSE4_2[bit 20] = 0.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(0)	If the memory address is in a non-canonical form.
#SS(0)	If the stack address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If CPUID.01H:ECX.MOVBE[bit 22] = 0 .
	If the LOCK prefix is used.
	If REP (F3H) prefix is used.

MOVD/MOVQ—Move Doubleword/Move Quadword

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 6E /*r* MOVD *mm, r/m32*	A	V/V	SSE2	Move doubleword from *r/m32* to *mm*.
REX.W + 0F 6E /*r* MOVQ *mm, r/m64*	A	V/N.E.	SSE2	Move quadword from *r/m64* to *mm*.
0F 7E /*r* MOVD *r/m32, mm*	B	V/V	SSE2	Move doubleword from *mm* to *r/m32*.
REX.W + 0F 7E /*r* MOVQ *r/m64, mm*	B	V/N.E.	SSE2	Move quadword from *mm* to *r/m64*.
VEX.128.66.0F.W0 6E / VMOVD xmm1, r32/m32	A	V/V	AVX	Move doubleword from r/m32 to xmm1.
VEX.128.66.0F.W1 6E /r VMOVQ xmm1, r64/m64	A	V/N.E.	AVX	Move quadword from r/m64 to xmm1.
66 0F 6E /*r* MOVD *xmm, r/m32*	A	V/V	SSE2	Move doubleword from *r/m32* to *xmm*.
66 REX.W 0F 6E /*r* MOVQ *xmm, r/m64*	A	V/N.E.	SSE2	Move quadword from *r/m64* to *xmm*.
66 0F 7E /*r* MOVD *r/m32, xmm*	B	V/V	SSE2	Move doubleword from *xmm* register to *r/m32*.
66 REX.W 0F 7E /*r* MOVQ *r/m64, xmm*	B	V/N.E.	SSE2	Move quadword from *xmm* register to *r/m64*.
VEX.128.66.0F.W0 7E /r VMOVD r32/m32, xmm1	B	V/V	AVX	Move doubleword from xmm1 register to r/m32.
VEX.128.66.0F.W1 7E /r VMOVQ r64/m64, xmm1	B	V/N.E.	AVX	Move quadword from xmm1 register to r/m64.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

Copies a doubleword from the source operand (second operand) to the destination operand (first operand). The source and destination operands can be general-purpose registers, MMX technology registers, XMM registers, or 32-bit memory locations. This instruction can be used to move a doubleword to and from the low doubleword of an MMX technology register and a general-purpose register or a 32-bit memory location, or to and from the low doubleword of an XMM register and a general-purpose register or a 32-bit memory location. The instruction cannot be used to transfer data between MMX technology registers, between XMM registers, between general-purpose registers, or between memory locations.

When the destination operand is an MMX technology register, the source operand is written to the low doubleword of the register, and the register is zero-extended to 64 bits. When the destination operand is an XMM register, the source operand is written to the low doubleword of the register, and the register is zero-extended to 128 bits.

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

MOVD (when destination operand is MMX technology register)

```
DEST[31:0] ← SRC;
DEST[63:32] ← 00000000H;
```

MOVD (when destination operand is XMM register)

```
DEST[31:0] ← SRC;
DEST[127:32] ← 000000000000000000000000H;
DEST[VLMAX-1:128] (Unmodified)
```

MOVD (when source operand is MMX technology or XMM register)

```
DEST ← SRC[31:0];
```

VMOVD (VEX-encoded version when destination is an XMM register)

```
DEST[31:0] ← SRC[31:0]
DEST[VLMAX-1:32] ← 0
```

MOVQ (when destination operand is XMM register)

```
DEST[63:0] ← SRC[63:0];
DEST[127:64] ← 0000000000000000H;
DEST[VLMAX-1:128] (Unmodified)
```

MOVQ (when destination operand is r/m64)

```
DEST[63:0] ← SRC[63:0];
```

MOVQ (when source operand is XMM register or r/m64)

```
DEST ← SRC[63:0];
```

VMOVQ (VEX-encoded version when destination is an XMM register)

DEST[63:0] ← SRC[63:0]
DEST[VLMAX-1:64] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

MOVD	__m64 _mm_cvtsi32_si64 (int i)
MOVD	int _mm_cvtsi64_si32 (__m64m)
MOVD	__m128i _mm_cvtsi32_si128 (int a)
MOVD	int _mm_cvtsi128_si32 (__m128i a)

Flags Affected

None.

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 5; additionally

#UD	If VEX.L = 1.
	If VEX.vvvv != 1111B.

MOVDDUP—Move One Double-FP and Duplicate

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 12 /*r* MOVDDUP *xmm1*, *xmm2/m64*	A	V/V	SSE3	Move one double-precision floating-point value from the lower 64-bit operand in *xmm2/m64* to *xmm1* and duplicate.
VEX.128.F2.0F.WIG 12 /r VMOVDDUP xmm1, xmm2/m64	A	V/V	AVX	Move double-precision floating-point values from xmm2/mem and duplicate into xmm1.
VEX.256.F2.0F.WIG 12 /r VMOVDDUP ymm1, ymm2/m256	A	V/V	AVX	Move even index double-precision floating-point values from ymm2/mem and duplicate each element into ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

The linear address corresponds to the address of the least-significant byte of the referenced memory data. When a memory address is indicated, the 8 bytes of data at memory location m64 are loaded. When the register-register form of this operation is used, the lower half of the 128-bit source register is duplicated and copied into the 128-bit destination register. See Figure 3-27.

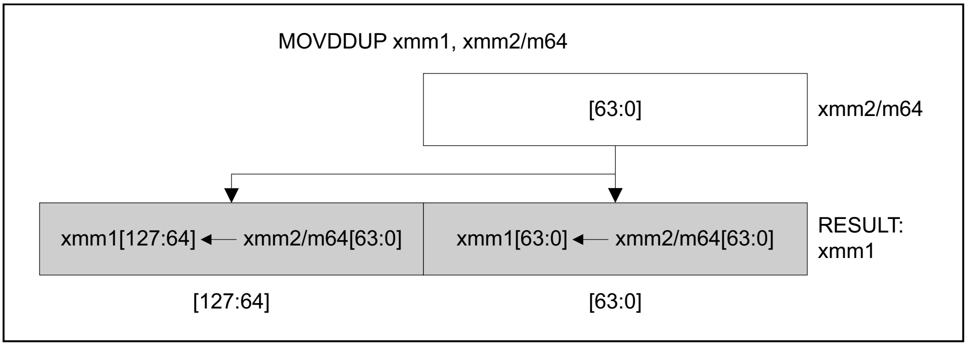


Figure 3-27. MOVDDUP—Move One Double-FP and Duplicate

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Operation

```
IF (Source = m64)
    THEN
        (* Load instruction *)
        xmm1[63:0] = m64;
        xmm1[127:64] = m64;
    ELSE
        (* Move instruction *)
        xmm1[63:0] = xmm2[63:0];
        xmm1[127:64] = xmm2[63:0];
FI;
```

MOVDDUP (128-bit Legacy SSE version)

DEST[63:0] ← SRC[63:0]
DEST[127:64] ← SRC[63:0]
DEST[VLMAX-1:128] (Unmodified)

VMOVDDUP (VEX.128 encoded version)

DEST[63:0] ← SRC[63:0]
DEST[127:64] ← SRC[63:0]

DEST[VLMAX-1:128] ← 0

VMOVDDUP (VEX.256 encoded version)
DEST[63:0] ← SRC[63:0]
DEST[127:64] ← SRC[63:0]
DEST[191:128] ← SRC[191:128]
DEST[255:192] ← SRC[191:128]

Intel C/C++ Compiler Intrinsic Equivalent

MOVDDUP __m128d _mm_movedup_pd(__m128d a)

MOVDDUP __m128d _mm_loaddup_pd(double const * dp)

SIMD Floating-Point Exceptions

None

Other Exceptions

See Exceptions Type 5; additionally

#UD If VEX.vvvv != 1111B.

MOVDQA—Move Aligned Double Quadword

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 6F /*r* MOVDQA *xmm1*, *xmm2/m128*	A	V/V	SSE2	Move aligned double quadword from *xmm2/m128* to *xmm1*.
66 0F 7F /*r* MOVDQA *xmm2/m128*, *xmm1*	B	V/V	SSE2	Move aligned double quadword from *xmm1* to *xmm2/m128*.
VEX.128.66.0F.WIG 6F /r VMOVDQA xmm1, xmm2/m128	A	V/V	AVX	Move aligned packed integer values from xmm2/mem to xmm1.
VEX.128.66.0F.WIG 7F /r VMOVDQA xmm2/m128, xmm1	B	V/V	AVX	Move aligned packed integer values from xmm1 to xmm2/mem.
VEX.256.66.0F.WIG 6F /r VMOVDQA ymm1, ymm2/m256	A	V/V	AVX	Move aligned packed integer values from ymm2/mem to ymm1.
VEX.256.66.0F.WIG 7F /r VMOVDQA ymm2/m256, ymm1	B	V/V	AVX	Move aligned packed integer values from ymm1 to ymm2/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

128-bit versions:

Moves 128 bits of packed integer values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load an XMM register from a 128-bit memory location, to store the contents of an XMM register into a 128-bit memory location, or to move data between two XMM registers.

When the source or destination operand is a memory operand, the operand must be aligned on a 16-byte boundary or a general-protection exception (#GP) will be generated. To move integer data to and from unaligned memory locations, use the VMOVDQU instruction.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

VEX.256 encoded version:

Moves 256 bits of packed integer values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load a YMM register from a 256-bit memory location, to store the contents of a YMM register into a 256-bit memory location, or to move data between two YMM registers.

When the source or destination operand is a memory operand, the operand must be aligned on a 32-byte boundary or a general-protection exception (#GP) will be generated. To move integer data to and from unaligned memory locations, use the VMOVDQU instruction.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

Operation

MOVDQA (128-bit load- and register- form Legacy SSE version)

```
DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] (Unmodified)
(* #GP if SRC or DEST unaligned memory operand *)
```

(V)MOVDQA (128-bit store forms)

```
DEST[127:0] ← SRC[127:0]
```

VMOVDQA (VEX.128 encoded version)

```
DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] ← 0
```

VMOVDQA (VEX.256 encoded version)

```
DEST[255:0] ← SRC[255:0]
```

Intel C/C++ Compiler Intrinsic Equivalent

MOVDQA __m128i _mm_load_si128 (__m128i *p)

MOVDQA void _mm_store_si128 (__m128i *p, __m128i a)

VMOVDQA __m256i _mm256_load_si256 (__m256i * p);

VMOVDQA _mm256_store_si256(_m256i *p, __m256i a);

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 1.SSE2; additionally

#UD	If VEX.vvvv != 1111B.

MOVDQU—Move Unaligned Double Quadword

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 6F /*r* MOVDQU *xmm1*, *xmm2/m128*	A	V/V	SSE2	Move unaligned double quadword from *xmm2/m128* to *xmm1*.
F3 0F 7F /*r* MOVDQU *xmm2/m128*, *xmm1*	B	V/V	SSE2	Move unaligned double quadword from *xmm1* to *xmm2/m128*.
VEX.128.F3.0F.WIG 6F /r VMOVDQU xmm1, xmm2/m128	A	V/V	AVX	Move unaligned packed integer values from xmm2/mem to xmm1.
VEX.128.F3.0F.WIG 7F /r VMOVDQU xmm2/m128, xmm1	B	V/V	AVX	Move unaligned packed integer values from xmm1 to xmm2/mem.
VEX.256.F3.0F.WIG 6F /r VMOVDQU ymm1, ymm2/m256	A	V/V	AVX	Move unaligned packed integer values from ymm2/mem to ymm1.
VEX.256.F3.0F.WIG 7F /r VMOVDQU ymm2/m256, ymm1	B	V/V	AVX	Move unaligned packed integer values from ymm1 to ymm2/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

128-bit versions:

Moves 128 bits of packed integer values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load an XMM register from a 128-bit memory location, to store the contents of an XMM register into a 128-bit memory location, or to move data between two XMM registers. When the source or destination operand is a memory operand, the operand may be unaligned on a 16-byte boundary without causing a general-protection exception (#GP) to be generated.[1]

To move a double quadword to or from memory locations that are known to be aligned on 16-byte boundaries, use the MOVDQA instruction.

While executing in 16-bit addressing mode, a linear address for a 128-bit data access that overlaps the end of a 16-bit segment is not allowed and is defined as reserved behavior. A specific processor implementation may or may not generate a general-protection exception (#GP) in this situation, and the address that spans the end of the segment may or may not wrap around to the beginning of the segment.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

When the source or destination operand is a memory operand, the operand may be unaligned to any alignment without causing a general-protection exception (#GP) to be generated

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

VEX.256 encoded version:

Moves 256 bits of packed integer values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load a YMM register from a 256-bit memory location, to store the contents of a YMM register into a 256-bit memory location, or to move data between two YMM registers.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

Operation

MOVDQU load and register copy (128-bit Legacy SSE version)

DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] (Unmodified)

(V)MOVDQU 128-bit store-form versions

DEST[127:0] ← SRC[127:0]

VMOVDQU (VEX.128 encoded version)

DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] ← 0

VMOVDQU (VEX.256 encoded version)

DEST[255:0] ← SRC[255:0]

---

1. If alignment checking is enabled (CR0.AM = 1, RFLAGS.AC = 1, and CPL = 3), an alignment-check exception (#AC) may or may not be generated (depending on processor implementation) when the operand is not aligned on an 8-byte boundary.

Intel C/C++ Compiler Intrinsic Equivalent

MOVDQU void _mm_storeu_si128 (__m128i *p, __m128i a)

MOVDQU __m128i _mm_loadu_si128 (__m128i *p)

VMOVDQU __m256i _mm256_loadu_si256 (__m256i * p);

VMOVDQU _mm256_storeu_si256(_m256i *p, __m256i a);

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 4; additionally

#UD	If VEX.vvvv != 1111B.

MOVDQ2Q—Move Quadword from XMM to MMX Technology Register

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
F2 0F D6	MOVDQ2Q *mm, xmm*	A	Valid	Valid	Move low quadword from *xmm* to *mmx* register.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:reg (r)	NA	NA

Description

Moves the low quadword from the source operand (second operand) to the destination operand (first operand). The source operand is an XMM register and the destination operand is an MMX technology register.

This instruction causes a transition from x87 FPU to MMX technology operation (that is, the x87 FPU top-of-stack pointer is set to 0 and the x87 FPU tag word is set to all 0s [valid]). If this instruction is executed while an x87 FPU floating-point exception is pending, the exception is handled before the MOVDQ2Q instruction is executed.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Operation

```
DEST ← SRC[63:0];
```

Intel C/C++ Compiler Intrinsic Equivalent

MOVDQ2Q __m64 _mm_movepi64_pi64 (__m128i a)

SIMD Floating-Point Exceptions

None.

Protected Mode Exceptions

#NM	If CR0.TS[bit 3] = 1.
#UD	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.
#MF	If there is a pending x87 FPU exception.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

MOVHLPS— Move Packed Single-Precision Floating-Point Values High to Low

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 12 /*r* MOVHLPS *xmm1, xmm2*	A	V/V	SSE3	Move two packed single-precision floating-point values from high quadword of *xmm2* to low quadword of *xmm1*.
VEX.NDS.128.0F.WIG 12 /r VMOVHLPS xmm1, xmm2, xmm3	B	V/V	AVX	Merge two packed single-precision floating-point values from high quadword of xmm3 and low quadword of xmm2.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:reg (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

This instruction cannot be used for memory to register moves.

128-bit two-argument form:

Moves two packed single-precision floating-point values from the high quadword of the second XMM argument (second operand) to the low quadword of the first XMM register (first argument). The high quadword of the destination operand is left unchanged. Bits (VLMAX-1:64) of the corresponding YMM destination register are unmodified.

128-bit three-argument form

Moves two packed single-precision floating-point values from the high quadword of the third XMM argument (third operand) to the low quadword of the destination (first operand). Copies the high quadword from the second XMM argument (second operand) to the high quadword of the destination (first operand). Bits (VLMAX-1:128) of the destination YMM register are zeroed.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

If VMOVHLPS is encoded with VEX.L= 1, an attempt to execute the instruction encoded with VEX.L= 1 will cause an #UD exception.

Operation

MOVHLPS (128-bit two-argument form)
DEST[63:0] ← SRC[127:64]
DEST[VLMAX-1:64] (Unmodified)

VMOVHLPS (128-bit three-argument form)
DEST[63:0] ← SRC2[127:64]
DEST[127:64] ← SRC1[127:64]
DEST[VLMAX-1:128] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

MOVHLPS __m128 _mm_movehl_ps(__m128 a, __m128 b)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 7; additionally

#UD	If VEX.L= 1.

MOVHPD—Move High Packed Double-Precision Floating-Point Value

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 16 /*r* MOVHPD *xmm, m64*	A	V/V	SSE2	Move double-precision floating-point value from *m64* to high quadword of *xmm*.
66 0F 17 /*r* MOVHPD *m64, xmm*	B	V/V	SSE2	Move double-precision floating-point value from high quadword of *xmm* to *m64*.
VEX.NDS.128.66.0F.WIG 16 /r VMOVHPD xmm2, xmm1, m64	C	V/V	AVX	Merge double-precision floating-point value from m64 and the low quadword of xmm1.
VEX128.66.0F.WIG 17/r VMOVHPD m64, xmm1	B	V/V	AVX	Move double-precision floating-point values from high quadword of xmm1 to m64.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA
C	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

This instruction cannot be used for register to register or memory to memory moves.

128-bit Legacy SSE load:

Moves a double-precision floating-point value from the source 64-bit memory operand and stores it in the high 64-bits of the destination XMM register. The lower 64bits of the XMM register are preserved. The upper 128-bits of the corresponding YMM destination register are preserved.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

VEX.128 encoded load:

Loads a double-precision floating-point value from the source 64-bit memory operand (third operand) and stores it in the upper 64-bits of the destination XMM register (first operand). The low 64-bits from second XMM register (second operand) are stored in the lower 64-bits of the destination. The upper 128-bits of the destination YMM register are zeroed.

128-bit store:

Stores a double-precision floating-point value from the high 64-bits of the XMM register source (second operand) to the 64-bit memory location (first operand).

Note: VMOVHPD (store) (VEX.128.66.0F 17 /r) is legal and has the same behavior as the existing 66 0F 17 store. For VMOVHPD (store) (VEX.128.66.0F 17 /r) instruction version, VEX.vvvv is reserved and must be 1111b otherwise instruction will #UD.

If VMOVHPD is encoded with VEX.L= 1, an attempt to execute the instruction encoded with VEX.L= 1 will cause an #UD exception.

Operation

MOVHPD (128-bit Legacy SSE load)

```
DEST[63:0] (Unmodified)
DEST[127:64] ← SRC[63:0]
DEST[VLMAX-1:128] (Unmodified)
```

VMOVHPD (VEX.128 encoded load)

```
DEST[63:0] ← SRC1[63:0]
DEST[127:64] ← SRC2[63:0]
DEST[VLMAX-1:128] ← 0
```

VMOVHPD (store)

```
DEST[63:0] ← SRC[127:64]
```

Intel C/C++ Compiler Intrinsic Equivalent

MOVHPD __m128d _mm_loadh_pd (__m128d a, double *p)

MOVHPD void _mm_storeh_pd (double *p, __m128d a)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 5; additionally

#UD	If VEX.L= 1.

MOVHPS—Move High Packed Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 16 /*r* MOVHPS *xmm, m64*	A	V/V	SSE	Move two packed single-precision floating-point values from *m64* to high quadword of *xmm*.
0F 17 /*r* MOVHPS *m64, xmm*	B	V/V	SSE	Move two packed single-precision floating-point values from high quadword of *xmm* to *m64*.
VEX.NDS.128.0F.WIG 16 /r VMOVHPS xmm2, xmm1, m64	C	V/V	AVX	Merge two packed single-precision floating-point values from m64 and the low quadword of xmm1.
VEX.128.0F.WIG 17/r VMOVHPS m64, xmm1	B	V/V	AVX	Move two packed single-precision floating-point values from high quadword of xmm1to m64.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA
C	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

This instruction cannot be used for register to register or memory to memory moves.

128-bit Legacy SSE load:

Moves two packed single-precision floating-point values from the source 64-bit memory operand and stores them in the high 64-bits of the destination XMM register. The lower 64bits of the XMM register are preserved. The upper 128-bits of the corresponding YMM destination register are preserved.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

VEX.128 encoded load:

Loads two single-precision floating-point values from the source 64-bit memory operand (third operand) and stores it in the upper 64-bits of the destination XMM register (first operand). The low 64-bits from second XMM register (second operand) are stored in the lower 64-bits of the destination. The upper 128-bits of the destination YMM register are zeroed.

128-bit store:

Stores two packed single-precision floating-point values from the high 64-bits of the XMM register source (second operand) to the 64-bit memory location (first operand).

Note: VMOVHPS (store) (VEX.NDS.128.0F 17 /r) is legal and has the same behavior as the existing 0F 17 store. For VMOVHPS (store) (VEX.NDS.128.0F 17 /r) instruction version, VEX.vvvv is reserved and must be 1111b otherwise instruction will #UD.

If VMOVHPS is encoded with VEX.L= 1, an attempt to execute the instruction encoded with VEX.L= 1 will cause an #UD exception.

Operation

MOVHPS (128-bit Legacy SSE load)

```
DEST[63:0] (Unmodified)
DEST[127:64] ← SRC[63:0]
DEST[VLMAX-1:128] (Unmodified)
```

VMOVHPS (VEX.128 encoded load)

```
DEST[63:0] ← SRC1[63:0]
DEST[127:64] ← SRC2[63:0]
DEST[VLMAX-1:128] ← 0
```

VMOVHPS (store)

```
DEST[63:0] ← SRC[127:64]
```

Intel C/C++ Compiler Intrinsic Equivalent

MOVHPS	__m128d _mm_loadh_pi (__m128d a, __m64 *p)
MOVHPS	void _mm_storeh_pi (__m64 *p, __m128d a)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 5; additionally

#UD	If VEX.L= 1.

MOVLHPS—Move Packed Single-Precision Floating-Point Values Low to High

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 16 /r MOVLHPS *xmm1, xmm2*	A	V/V	SSE	Move two packed single-precision floating-point values from low quadword of *xmm2* to high quadword of *xmm1*.
VEX.NDS.128.0F.WIG 16 /r VMOVLHPS xmm1, xmm2, xmm3	B	V/V	AVX	Merge two packed single-precision floating-point values from low quadword of xmm3 and low quadword of xmm2.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:reg (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

This instruction cannot be used for memory to register moves.

128-bit two-argument form:

Moves two packed single-precision floating-point values from the low quadword of the second XMM argument (second operand) to the high quadword of the first XMM register (first argument). The low quadword of the destination operand is left unchanged. The upper 128 bits of the corresponding YMM destination register are unmodified.

128-bit three-argument form

Moves two packed single-precision floating-point values from the low quadword of the third XMM argument (third operand) to the high quadword of the destination (first operand). Copies the low quadword from the second XMM argument (second operand) to the low quadword of the destination (first operand). The upper 128-bits of the destination YMM register are zeroed.

If VMOVLHPS is encoded with VEX.L= 1, an attempt to execute the instruction encoded with VEX.L= 1 will cause an #UD exception.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Operation

MOVLHPS (128-bit two-argument form)
DEST[63:0] (Unmodified)
DEST[127:64] ← SRC[63:0]
DEST[VLMAX-1:128] (Unmodified)

VMOVLHPS (128-bit three-argument form)
DEST[63:0] ← SRC1[63:0]
DEST[127:64] ← SRC2[63:0]
DEST[VLMAX-1:128] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

MOVHLPS __m128 _mm_movelh_ps(__m128 a, __m128 b)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 7; additionally

#UD	If VEX.L= 1.

MOVLPD—Move Low Packed Double-Precision Floating-Point Value

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 12 /*r* MOVLPD *xmm, m64*	A	V/V	SSE2	Move double-precision floating-point value from *m64* to low quadword of *xmm* register.
66 0F 13 /*r* MOVLPD *m64, xmm*	B	V/V	SSE2	Move double-precision floating-point nvalue from low quadword of *xmm* register to *m64*.
VEX.NDS.128.66.0F.WIG 12 /r VMOVLPD xmm2, xmm1, m64	C	V/V	AVX	Merge double-precision floating-point value from m64 and the high quadword of xmm1.
VEX.128.66.0F.WIG 13/r VMOVLPD m64, xmm1	B	V/V	AVX	Move double-precision floating-point values from low quadword of xmm1 to m64.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA
C	ModRM:r/m (r)	ModRM:reg (r, w)	VEX.vvvv (r)	NA

Description

This instruction cannot be used for register to register or memory to memory moves.

128-bit Legacy SSE load:

Moves a double-precision floating-point value from the source 64-bit memory operand and stores it in the low 64-bits of the destination XMM register. The upper 64bits of the XMM register are preserved. The upper 128-bits of the corresponding YMM destination register are preserved.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

VEX.128 encoded load:

Loads a double-precision floating-point value from the source 64-bit memory operand (third operand), merges it with the upper 64-bits of the first source XMM register (second operand), and stores it in the low 128-bits of the destination XMM register (first operand). The upper 128-bits of the destination YMM register are zeroed.

128-bit store:

Stores a double-precision floating-point value from the low 64-bits of the XMM register source (second operand) to the 64-bit memory location (first operand).

Note: VMOVLPD (store) (VEX.128.66.0F 13 /r) is legal and has the same behavior as the existing 66 0F 13 store. For VMOVLPD (store) (VEX.128.66.0F 13 /r) instruction version, VEX.vvvv is reserved and must be 1111b otherwise instruction will #UD.

If VMOVLPD is encoded with VEX.L= 1, an attempt to execute the instruction encoded with VEX.L= 1 will cause an #UD exception.

Operation

MOVLPD (128-bit Legacy SSE load)

DEST[63:0] ← SRC[63:0]
DEST[VLMAX-1:64] (Unmodified)

VMOVLPD (VEX.128 encoded load)

DEST[63:0] ← SRC2[63:0]
DEST[127:64] ← SRC1[127:64]
DEST[VLMAX-1:128] ← 0

VMOVLPD (store)

DEST[63:0] ← SRC[63:0]

Intel C/C++ Compiler Intrinsic Equivalent

MOVLPD __m128d _mm_loadl_pd (__m128d a, double *p)

MOVLPD void _mm_storel_pd (double *p, __m128d a)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 5; additionally

#UD	If VEX.L= 1.
	If VEX.vvvv != 1111B.

MOVLPS—Move Low Packed Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 12 /*r* MOVLPS *xmm*, *m64*	A	V/V	SSE	Move two packed single-precision floating-point values from *m64* to low quadword of *xmm*.
0F 13 /*r* MOVLPS *m64*, *xmm*	B	V/V	SSE	Move two packed single-precision floating-point values from low quadword of *xmm* to *m64*.
VEX.NDS.128.0F.WIG 12 /r VMOVLPS xmm2, xmm1, m64	C	V/V	AVX	Merge two packed single-precision floating-point values from m64 and the high quadword of xmm1.
VEX.128.0F.WIG 13/r VMOVLPS m64, xmm1	B	V/V	AVX	Move two packed single-precision floating-point values from low quadword of xmm1 to m64.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA
C	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

This instruction cannot be used for register to register or memory to memory moves.

128-bit Legacy SSE load:

Moves two packed single-precision floating-point values from the source 64-bit memory operand and stores them in the low 64-bits of the destination XMM register. The upper 64bits of the XMM register are preserved. The upper 128-bits of the corresponding YMM destination register are preserved.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

VEX.128 encoded load:

Loads two packed single-precision floating-point values from the source 64-bit memory operand (third operand), merges them with the upper 64-bits of the first source XMM register (second operand), and stores them in the low 128-bits of the destination XMM register (first operand). The upper 128-bits of the destination YMM register are zeroed.

128-bit store:

Loads two packed single-precision floating-point values from the low 64-bits of the XMM register source (second operand) to the 64-bit memory location (first operand).

Note: VMOVLPS (store) (VEX.128.0F 13 /r) is legal and has the same behavior as the existing 0F 13 store. For VMOVLPS (store) (VEX.128.0F 13 /r) instruction version, VEX.vvvv is reserved and must be 1111b otherwise instruction will #UD.

If VMOVLPS is encoded with VEX.L= 1, an attempt to execute the instruction encoded with VEX.L= 1 will cause an #UD exception.

Operation

MOVLPS (128-bit Legacy SSE load)

DEST[63:0] ← SRC[63:0]
DEST[VLMAX-1:64] (Unmodified)

VMOVLPS (VEX.128 encoded load)

DEST[63:0] ← SRC2[63:0]
DEST[127:64] ← SRC1[127:64]
DEST[VLMAX-1:128] ← 0

VMOVLPS (store)

DEST[63:0] ← SRC[63:0]

Intel C/C++ Compiler Intrinsic Equivalent

MOVLPS	__m128 _mm_loadl_pi (__m128 a, __m64 *p)
MOVLPS	void _mm_storel_pi (__m64 *p, __m128 a)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 5; additionally

#UD	If VEX.L= 1.
	If VEX.vvvv != 1111B.

MOVMSKPD—Extract Packed Double-Precision Floating-Point Sign Mask

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 50 /*r* MOVMSKPD *reg, xmm*	A	V/V	SSE2	Extract 2-bit sign mask from *xmm* and store in *reg*. The upper bits of r32 or r64 are filled with zeros.
VEX.128.66.0F.WIG 50 /r VMOVMSKPD reg, xmm2	A	V/V	AVX	Extract 2-bit sign mask from xmm2 and store in reg. The upper bits of r32 or r64 are zeroed.
VEX.256.66.0F.WIG 50 /r VMOVMSKPD reg, ymm2	A	V/V	AVX	Extract 4-bit sign mask from ymm2 and store in reg. The upper bits of r32 or r64 are zeroed.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:reg (r)	NA	NA

Description

Extracts the sign bits from the packed double-precision floating-point values in the source operand (second operand), formats them into a 2-bit mask, and stores the mask in the destination operand (first operand). The source operand is an XMM register, and the destination operand is a general-purpose register. The mask is stored in the 2 low-order bits of the destination operand. Zero-extend the upper bits of the destination.

In 64-bit mode, the instruction can access additional registers (XMM8-XMM15, R8-R15) when used with a REX.R prefix. The default operand size is 64-bit in 64-bit mode.

128-bit versions: The source operand is a YMM register. The destination operand is a general purpose register.

VEX.256 encoded version: The source operand is a YMM register. The destination operand is a general purpose register.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b, otherwise instructions will #UD.

Operation

(V)MOVMSKPD (128-bit versions)

```
DEST[0] ← SRC[63]
DEST[1] ← SRC[127]
IF DEST = r32
    THEN DEST[31:2] ← 0;
    ELSE DEST[63:2] ← 0;
FI
```

VMOVMSKPD (VEX.256 encoded version)

```
DEST[0] ← SRC[63]
DEST[1] ← SRC[127]
DEST[2] ← SRC[191]
DEST[3] ← SRC[255]
IF DEST = r32
    THEN DEST[31:4] ← 0;
    ELSE DEST[63:4] ← 0;
FI
```

Intel C/C++ Compiler Intrinsic Equivalent

MOVMSKPD	int _mm_movemask_pd (__m128d a)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 7; additionally

#UD	If VEX.vvvv != 1111B.

MOVMSKPS—Extract Packed Single-Precision Floating-Point Sign Mask

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 50 /r MOVMSKPS *reg, xmm*	A	V/V	SSE	Extract 4-bit sign mask from *xmm* and store in *reg*. The upper bits of r32 or r64 are filled with zeros.
VEX.128.0F.WIG 50 /r VMOVMSKPS reg, xmm2	A	V/V	AVX	Extract 4-bit sign mask from xmm2 and store in reg. The upper bits of r32 or r64 are zeroed.
VEX.256.0F.WIG 50 /r VMOVMSKPS reg, ymm2	A	V/V	AVX	Extract 8-bit sign mask from ymm2 and store in reg. The upper bits of r32 or r64 are zeroed.

Instruction Operand Encoding[1]

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Extracts the sign bits from the packed single-precision floating-point values in the source operand (second operand), formats them into a 4- or 8-bit mask, and stores the mask in the destination operand (first operand). The source operand is an XMM or YMM register, and the destination operand is a general-purpose register. The mask is stored in the 4 or 8 low-order bits of the destination operand. The upper bits of the destination operand beyond the mask are filled with zeros.

In 64-bit mode, the instruction can access additional registers (XMM8-XMM15, R8-R15) when used with a REX.R prefix. The default operand size is 64-bit in 64-bit mode.

128-bit versions: The source operand is a YMM register. The destination operand is a general purpose register.

VEX.256 encoded version: The source operand is a YMM register. The destination operand is a general purpose register.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b, otherwise instructions will #UD.

1. ModRM.MOD = 011B required

Operation

```
DEST[0] ← SRC[31];
DEST[1] ← SRC[63];
DEST[2] ← SRC[95];
DEST[3] ← SRC[127];

IF DEST = r32
    THEN DEST[31:4] ← ZeroExtend;
    ELSE DEST[63:4] ← ZeroExtend;
FI;
```

(V)MOVMSKPS (128-bit version)

```
DEST[0] ← SRC[31]
DEST[1] ← SRC[63]
DEST[2] ← SRC[95]
DEST[3] ← SRC[127]
IF DEST = r32
    THEN DEST[31:4] ← 0;
    ELSE DEST[63:4] ← 0;
FI
```

VMOVMSKPS (VEX.256 encoded version)

```
DEST[0] ← SRC[31]
DEST[1] ← SRC[63]
DEST[2] ← SRC[95]
DEST[3] ← SRC[127]
DEST[4] ← SRC[159]
DEST[5] ← SRC[191]
DEST[6] ← SRC[223]
DEST[7] ← SRC[255]
IF DEST = r32
    THEN DEST[31:8] ← 0;
    ELSE DEST[63:8] ← 0;
FI
```

Intel C/C++ Compiler Intrinsic Equivalent

int _mm_movemask_ps(__m128 a)

int _mm256_movemask_ps(__m256 a)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 7; additionally

#UD	If VEX.vvvv != 1111B.

MOVNTDQA — Load Double Quadword Non-Temporal Aligned Hint

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 38 2A /r MOVNTDQA *xmm1, m128*	A	V/V	SSE4_1	Move double quadword from *m128* to *xmm* using non-temporal hint if WC memory type.
VEX.128.66.0F38.WIG 2A /r VMOVNTDQA xmm1, m128	A	V/V	AVX	Move double quadword from m128 to xmm using non-temporal hint if WC memory type.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

MOVNTDQA loads a double quadword from the source operand (second operand) to the destination operand (first operand) using a non-temporal hint. A processor implementation may make use of the non-temporal hint associated with this instruction if the memory source is WC (write combining) memory type. An implementation may also make use of the non-temporal hint associated with this instruction if the memory source is WB (write back) memory type.

A processor's implementation of the non-temporal hint does not override the effective memory type semantics, but the implementation of the hint is processor dependent. For example, a processor implementation may choose to ignore the hint and process the instruction as a normal MOVDQA for any memory type. Another implementation of the hint for WC memory type may optimize data transfer throughput of WC reads. A third implementation may optimize cache reads generated by MOVNTDQA on WB memory type to reduce cache evictions.

WC Streaming Load Hint

For WC memory type in particular, the processor never appears to read the data into the cache hierarchy. Instead, the non-temporal hint may be implemented by loading a temporary internal buffer with the equivalent of an aligned cache line without filling this data to the cache. Any memory-type aliased lines in the cache will be snooped and flushed. Subsequent MOVNTDQA reads to unread portions of the WC cache line will receive data from the temporary internal buffer if data is available. The temporary internal buffer may be flushed by the processor at any time for any reason, for example:

- A load operation other than a MOVNTDQA which references memory already resident in a temporary internal buffer.

- A non-WC reference to memory already resident in a temporary internal buffer.
- Interleaving of reads and writes to memory currently residing in a single temporary internal buffer.
- Repeated (V)MOVNTDQA loads of a particular 16-byte item in a streaming line.
- Certain micro-architectural conditions including resource shortages, detection of a mis-speculation condition, and various fault conditions

The memory type of the region being read can override the non-temporal hint, if the memory address specified for the non-temporal read is not a WC memory region. Information on non-temporal reads and writes can be found in Chapter 11, "Memory Cache Control" of *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 3A*.

Because the WC protocol uses a weakly-ordered memory consistency model, an MFENCE or locked instruction should be used in conjunction with MOVNTDQA instructions if multiple processors might reference the same WC memory locations or in order to synchronize reads of a processor with writes by other agents in the system. Because of the speculative nature of fetching due to MOVNTDQA, Streaming loads must not be used to reference memory addresses that are mapped to I/O devices having side effects or when reads to these devices are destructive. For additional information on MOVNTDQA usages, see Section 12.10.3 in Chapter 12, "Programming with SSE3, SSSE3 and SSE4" of *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*.

The 128-bit (V)MOVNTDQA addresses must be 16-byte aligned or the instruction will cause a #GP.

Note: In VEX-128 encoded versions, VEX.vvvv is reserved and must be 1111b, VEX.L must be 0; otherwise instructions will #UD.

Operation

MOVNTDQA (128bit- Legacy SSE form)

DEST ← SRC
DEST[VLMAX-1:128] (Unmodified)

VMOVNTDQA (VEX.128 encoded form)

DEST ← SRC
DEST[VLMAX-1:128] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

MOVNTDQA __m128i _mm_stream_load_si128 (__m128i *p);

Flags Affected

None

Other Exceptions

See Exceptions Type 1.SSE4.1; additionally

#UD	If VEX.L= 1.
	If VEX.vvvv != 1111B.

MOVNTDQ—Store Double Quadword Using Non-Temporal Hint

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F E7 /r MOVNTDQ *m128*, xmm	A	V/V	SSE2	Move double quadword from *xmm* to *m128* using non-temporal hint.
VEX.128.66.0F.WIG E7 /r VMOVNTDQ m128, xmm1	A	V/V	AVX	Move packed integer values in xmm1 to m128 using non-temporal hint.
VEX.256.66.0F.WIG E7 /r VMOVNTDQ m256, ymm1	A	V/V	AVX	Move packed integer values in ymm1 to m256 using non-temporal hint.

Instruction Operand Encoding[1]

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

Moves the packed integers in the source operand (second operand) to the destination operand (first operand) using a non-temporal hint to prevent caching of the data during the write to memory. The source operand is an XMM register or YMM register, which is assumed to contain integer data (packed bytes, words, doublewords, or quadwords). The destination operand is a 128-bit or 256-bit memory location. The memory operand must be aligned on a 16-byte (128-bit version) or 32-byte (VEX.256 encoded version) boundary otherwise a general-protection exception (#GP) will be generated.

The non-temporal hint is implemented by using a write combining (WC) memory type protocol when writing the data to memory. Using this protocol, the processor does not write the data into the cache hierarchy, nor does it fetch the corresponding cache line from memory into the cache hierarchy. The memory type of the region being written to can override the non-temporal hint, if the memory address specified for the non-temporal store is in an uncacheable (UC) or write protected (WP) memory region. For more information on non-temporal stores, see "Caching of Temporal vs. Non-Temporal Data" in Chapter 10 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*.

Because the WC protocol uses a weakly-ordered memory consistency model, a fencing operation implemented with the SFENCE or MFENCE instruction should be used in conjunction with MOVNTDQ instructions if multiple processors might use different memory types to read/write the destination memory locations.

1. ModRM.MOD = 011B required

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Note: In VEX-128 encoded versions, VEX.vvvv is reserved and must be 1111b, VEX.L must be 0; otherwise instructions will #UD.

Operation

DEST ← SRC;

Intel C/C++ Compiler Intrinsic Equivalent

MOVNTDQ void _mm_stream_pd(double* p, __m128d a)

VMOVNTDQ void _mm256_stream_si256 (__m256i * p, __m256i a);

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 1.SSE2; additionally

#UD	If VEX.vvvv != 1111B.

MOVNTI—Store Doubleword Using Non-Temporal Hint

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F C3 /*r*	MOVNTI *m32, r32*	A	Valid	Valid	Move doubleword from *r32* to *m32* using non-temporal hint.
REX.W + 0F C3 /*r*	MOVNTI *m64, r64*	A	Valid	N.E.	Move quadword from *r64* to *m64* using non-temporal hint.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

Moves the doubleword integer in the source operand (second operand) to the destination operand (first operand) using a non-temporal hint to minimize cache pollution during the write to memory. The source operand is a general-purpose register. The destination operand is a 32-bit memory location.

The non-temporal hint is implemented by using a write combining (WC) memory type protocol when writing the data to memory. Using this protocol, the processor does not write the data into the cache hierarchy, nor does it fetch the corresponding cache line from memory into the cache hierarchy. The memory type of the region being written to can override the non-temporal hint, if the memory address specified for the non-temporal store is in an uncacheable (UC) or write protected (WP) memory region. For more information on non-temporal stores, see “Caching of Temporal vs. Non-Temporal Data” in Chapter 10 in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*.

Because the WC protocol uses a weakly-ordered memory consistency model, a fencing operation implemented with the SFENCE or MFENCE instruction should be used in conjunction with MOVNTI instructions if multiple processors might use different memory types to read/write the destination memory locations.

In 64-bit mode, the instruction’s default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

DEST ← SRC;

Intel C/C++ Compiler Intrinsic Equivalent

MOVNTI void _mm_stream_si32 (int *p, int a)

SIMD Floating-Point Exceptions

None.

Protected Mode Exceptions

#GP(0)	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
#SS(0)	For an illegal address in the SS segment.
#PF(fault-code)	For a page fault.
#UD	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand is not aligned on a 16-byte boundary, regardless of segment.
	If any part of the operand lies outside the effective address space from 0 to FFFFH.
#UD	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	For a page fault.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	For a page fault.
#UD	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

MOVNTPD—Store Packed Double-Precision Floating-Point Values Using Non-Temporal Hint

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 2B /*r* MOVNTPD *m128, xmm*	A	V/V	SSE2	Move packed double-precision floating-point values from *xmm* to *m128* using non-temporal hint.
VEX.128.66.0F.WIG 2B /r VMOVNTPD m128, xmm1	A	V/V	AVX	Move packed double-precision values in xmm1 to m128 using non-temporal hint.
VEX.256.66.0F.WIG 2B /r VMOVNTPD m256, ymm1	A	V/V	AVX	Move packed double-precision values in ymm1 to m256 using non-temporal hint.

Instruction Operand Encoding[1]

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

Moves the packed double-precision floating-point values in the source operand (second operand) to the destination operand (first operand) using a non-temporal hint to prevent caching of the data during the write to memory. The source operand is an XMM register or YMM register, which is assumed to contain packed double-precision, floating-pointing data. The destination operand is a 128-bit or 256-bit memory location. The memory operand must be aligned on a 16-byte (128-bit version) or 32-byte (VEX.256 encoded version) boundary otherwise a general-protection exception (#GP) will be generated.

The non-temporal hint is implemented by using a write combining (WC) memory type protocol when writing the data to memory. Using this protocol, the processor does not write the data into the cache hierarchy, nor does it fetch the corresponding cache line from memory into the cache hierarchy. The memory type of the region being written to can override the non-temporal hint, if the memory address specified for the non-temporal store is in an uncacheable (UC) or write protected (WP) memory region. For more information on non-temporal stores, see “Caching of Temporal vs. Non-Temporal Data” in Chapter 10 in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*.

1. ModRM.MOD = 011B required

Because the WC protocol uses a weakly-ordered memory consistency model, a fencing operation implemented with the SFENCE or MFENCE instruction should be used in conjunction with MOVNTPD instructions if multiple processors might use different memory types to read/write the destination memory locations.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Note: In VEX-128 encoded versions, VEX.vvvv is reserved and must be 1111b, VEX.L must be 0; otherwise instructions will #UD.

Operation

```
DEST ← SRC;
```

Intel C/C++ Compiler Intrinsic Equivalent

MOVNTPD void _mm_stream_pd(double *p, __m128d a)

VMOVNTPD void _mm256_stream_pd (double * p, __m256d a);

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 1.SSE2; additionally

#UD	If VEX.vvvv != 1111B.

MOVNTPS—Store Packed Single-Precision Floating-Point Values Using Non-Temporal Hint

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 2B /*r* MOVNTPS *m128, xmm*	A	V/V	SSE	Move packed single-precision floating-point values from *xmm* to *m128* using non-temporal hint.
VEX.128.0F.WIG 2B /r VMOVNTPS m128, xmm1	A	V/V	AVX	Move packed single-precision values xmm1 to mem using non-temporal hint.
VEX.256.0F.WIG 2B /r VMOVNTPS m256, ymm1	A	V/V	AVX	Move packed single-precision values ymm1 to mem using non-temporal hint.

Instruction Operand Encoding[1]

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

Moves the packed single-precision floating-point values in the source operand (second operand) to the destination operand (first operand) using a non-temporal hint to prevent caching of the data during the write to memory. The source operand is an XMM register or YMM register, which is assumed to contain packed single-precision, floating-pointing. The destination operand is a 128-bit or 256-bitmemory location. The memory operand must be aligned on a 16-byte (128-bit version) or 32-byte (VEX.256 encoded version) boundary otherwise a general-protection exception (#GP) will be generated.

The non-temporal hint is implemented by using a write combining (WC) memory type protocol when writing the data to memory. Using this protocol, the processor does not write the data into the cache hierarchy, nor does it fetch the corresponding cache line from memory into the cache hierarchy. The memory type of the region being written to can override the non-temporal hint, if the memory address specified for the non-temporal store is in an uncacheable (UC) or write protected (WP) memory region. For more information on non-temporal stores, see “Caching of Temporal vs. Non-Temporal Data” in Chapter 10 in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*.

1. ModRM.MOD = 011B required

Because the WC protocol uses a weakly-ordered memory consistency model, a fencing operation implemented with the SFENCE or MFENCE instruction should be used in conjunction with MOVNTPS instructions if multiple processors might use different memory types to read/write the destination memory locations.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

Operation

```
DEST ← SRC;
```

Intel C/C++ Compiler Intrinsic Equivalent

MOVNTDQ void _mm_stream_ps(float * p, __m128 a)

VMOVNTPS void _mm256_stream_ps (float * p, __m256 a);

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 1.SSE; additionally

#UD	If VEX.vvvv != 1111B.

MOVNTQ—Store of Quadword Using Non-Temporal Hint

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F E7 /*r*	MOVNTQ *m64, mm*	A	Valid	Valid	Move quadword from *mm* to *m64* using non-temporal hint.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

Moves the quadword in the source operand (second operand) to the destination operand (first operand) using a non-temporal hint to minimize cache pollution during the write to memory. The source operand is an MMX technology register, which is assumed to contain packed integer data (packed bytes, words, or doublewords). The destination operand is a 64-bit memory location.

The non-temporal hint is implemented by using a write combining (WC) memory type protocol when writing the data to memory. Using this protocol, the processor does not write the data into the cache hierarchy, nor does it fetch the corresponding cache line from memory into the cache hierarchy. The memory type of the region being written to can override the non-temporal hint, if the memory address specified for the non-temporal store is in an uncacheable (UC) or write protected (WP) memory region. For more information on non-temporal stores, see “Caching of Temporal vs. Non-Temporal Data” in Chapter 10 in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*.

Because the WC protocol uses a weakly-ordered memory consistency model, a fencing operation implemented with the SFENCE or MFENCE instruction should be used in conjunction with MOVNTQ instructions if multiple processors might use different memory types to read/write the destination memory locations.

This instruction’s operation is the same in non-64-bit modes and 64-bit mode.

Operation

```
DEST ← SRC;
```

Intel C/C++ Compiler Intrinsic Equivalent

MOVNTQ void _mm_stream_pi(__m64 * p, __m64 a)

SIMD Floating-Point Exceptions

None.

Protected Mode Exceptions

#GP(0)	For an illegal memory operand effective address in the CS, DS, ES, FS or GS segments.
#SS(0)	For an illegal address in the SS segment.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If CR0.EM[bit 2] = 1.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

Real-Address Mode Exceptions

#GP	If a memory operand is not aligned on a 16-byte boundary, regardless of segment.
	If any part of the operand lies outside the effective address space from 0 to FFFFH.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If CR0.EM[bit 2] = 1.
	If CPUID.01H:EDX.SSE2[bit 26] = 0.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	For a page fault.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.

#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	For a page fault.
#NM	If CR0.TS[bit 3] = 1.
#MF	If there is a pending x87 FPU exception.
#UD	If CR0.EM[bit 2] = 1.
	If CPUID.01H:EDX.SSE[bit 25] = 0.
	If the LOCK prefix is used.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

MOVQ—Move Quadword

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 6F /*r*	MOVQ *mm, mm/m64*	A	Valid	Valid	Move quadword from *mm/m64* to *mm*.
0F 7F /*r*	MOVQ *mm/m64, mm*	B	Valid	Valid	Move quadword from *mm* to *mm/m64*.
F3 0F 7E	MOVQ *xmm1, xmm2/m64*	A	Valid	Valid	Move quadword from *xmm2/mem64* to *xmm1*.
66 0F D6	MOVQ *xmm2/m64, xmm1*	B	Valid	Valid	Move quadword from *xmm1* to *xmm2/mem64*.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

Copies a quadword from the source operand (second operand) to the destination operand (first operand). The source and destination operands can be MMX technology registers, XMM registers, or 64-bit memory locations. This instruction can be used to move a quadword between two MMX technology registers or between an MMX technology register and a 64-bit memory location, or to move data between two XMM registers or between an XMM register and a 64-bit memory location. The instruction cannot be used to transfer data between memory locations.

When the source operand is an XMM register, the low quadword is moved; when the destination operand is an XMM register, the quadword is stored to the low quadword of the register, and the high quadword is cleared to all 0s.

In 64-bit mode, use of the REX prefix in the form of REX.R permits this instruction to access additional registers (XMM8-XMM15).

Operation

```
MOVQ instruction when operating on MMX technology registers and memory locations:
    DEST ← SRC;

MOVQ instruction when source and destination operands are XMM registers:
    DEST[63:0] ← SRC[63:0];
    DEST[127:64] ← 0000000000000000H;
```

```
MOVQ instruction when source operand is XMM register and destination
operand is memory location:
    DEST ← SRC[63:0];

MOVQ instruction when source operand is memory location and destination
operand is XMM register:
    DEST[63:0] ← SRC;
    DEST[127:64] ← 0000000000000000H;
```

Flags Affected

None.

Intel C/C++ Compiler Intrinsic Equivalent

MOVQ m128i _mm_mov_epi64(__m128i a)

SIMD Floating-Point Exceptions

None.

Protected Mode Exceptions

#GP(0)	If the destination operand is in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#UD	If CR0.EM[bit 2] = 1.
	128-bit operations will generate #UD only if CR4.OSFXSR[bit 9] = 0. Execution of 128-bit instructions on a non-SSE2 capable processor (one that is MMX technology capable) will result in the instruction operating on the mm registers, not #UD.
	If the LOCK prefix is used.
#NM	If CR0.TS[bit 3] = 1.
#MF	(MMX register operations only) If there is a pending FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

Real-Address Mode Exceptions

#GP	If any part of the operand lies outside of the effective address space from 0 to FFFFH.
#UD	If CR0.EM[bit 2] = 1.
	128-bit operations will generate #UD only if CR4.OSFXSR[bit 9] = 0. Execution of 128-bit instructions on a non-SSE2 capable processor (one that is MMX technology capable) will result in the instruction operating on the mm registers, not #UD.
	If the LOCK prefix is used.
#NM	If CR0.TS[bit 3] = 1.
#MF	(MMX register operations only) If there is a pending FPU exception.

Virtual-8086 Mode Exceptions

Same exceptions as in real address mode.

#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#UD	If CR0.EM[bit 2] = 1.
	(XMM register operations only) If CR4.OSFXSR[bit 9] = 0.
	(XMM register operations only) If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.
#NM	If CR0.TS[bit 3] = 1.
#MF	(MMX register operations only) If there is a pending FPU exception.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

MOVQ2DQ—Move Quadword from MMX Technology to XMM Register

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
F3 0F D6	MOVQ2DQ *xmm, mm*	A	Valid	Valid	Move quadword from *mmx* to low quadword of *xmm*.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:reg (r)	NA	NA

Description

Moves the quadword from the source operand (second operand) to the low quadword of the destination operand (first operand). The source operand is an MMX technology register and the destination operand is an XMM register.

This instruction causes a transition from x87 FPU to MMX technology operation (that is, the x87 FPU top-of-stack pointer is set to 0 and the x87 FPU tag word is set to all 0s [valid]). If this instruction is executed while an x87 FPU floating-point exception is pending, the exception is handled before the MOVQ2DQ instruction is executed.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Operation

```
DEST[63:0] ← SRC[63:0];
DEST[127:64] ← 00000000000000000H;
```

Intel C/C++ Compiler Intrinsic Equivalent

MOVQ2DQ __128i _mm_movpi64_pi64 (__m64 a)

SIMD Floating-Point Exceptions

None.

Protected Mode Exceptions

#NM	If CR0.TS[bit 3] = 1.
#UD	If CR0.EM[bit 2] = 1.
	If CR4.OSFXSR[bit 9] = 0.
	If CPUID.01H:EDX.SSE2[bit 26] = 0.
	If the LOCK prefix is used.
#MF	If there is a pending x87 FPU exception.

Real-Address Mode Exceptions

Same exceptions as in protected mode.

Virtual-8086 Mode Exceptions

Same exceptions as in protected mode.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

Same exceptions as in protected mode.

MOVS/MOVSB/MOVSW/MOVSD/MOVSQ—Move Data from String to String

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
A4	MOVS *m8, m8*	A	Valid	Valid	For legacy mode, Move byte from address DS:(E)SI to ES:(E)DI. For 64-bit mode move byte from address (R|E)SI to (R|E)DI.
A5	MOVS *m16, m16*	A	Valid	Valid	For legacy mode, move word from address DS:(E)SI to ES:(E)DI. For 64-bit mode move word at address (R|E)SI to (R|E)DI.
A5	MOVS *m32, m32*	A	Valid	Valid	For legacy mode, move dword from address DS:(E)SI to ES:(E)DI. For 64-bit mode move dword from address (R|E)SI to (R|E)DI.
REX.W + A5	MOVS *m64, m64*	A	Valid	N.E.	Move qword from address (R|E)SI to (R|E)DI.
A4	MOVSB	A	Valid	Valid	For legacy mode, Move byte from address DS:(E)SI to ES:(E)DI. For 64-bit mode move byte from address (R|E)SI to (R|E)DI.
A5	MOVSW	A	Valid	Valid	For legacy mode, move word from address DS:(E)SI to ES:(E)DI. For 64-bit mode move word at address (R|E)SI to (R|E)DI.
A5	MOVSD	A	Valid	Valid	For legacy mode, move dword from address DS:(E)SI to ES:(E)DI. For 64-bit mode move dword from address (R|E)SI to (R|E)DI.
REX.W + A5	MOVSQ	A	Valid	N.E.	Move qword from address (R|E)SI to (R|E)DI.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

Moves the byte, word, or doubleword specified with the second operand (source operand) to the location specified with the first operand (destination operand). Both the source and destination operands are located in memory. The address of the source operand is read from the DS:ESI or the DS:SI registers (depending on the address-size attribute of the instruction, 32 or 16, respectively). The address of the destination operand is read from the ES:EDI or the ES:DI registers (again depending on the address-size attribute of the instruction). The DS segment may be overridden with a segment override prefix, but the ES segment cannot be overridden.

At the assembly-code level, two forms of this instruction are allowed: the "explicit-operands" form and the "no-operands" form. The explicit-operands form (specified with the MOVS mnemonic) allows the source and destination operands to be specified explicitly. Here, the source and destination operands should be symbols that indicate the size and location of the source value and the destination, respectively. This explicit-operands form is provided to allow documentation; however, note that the documentation provided by this form can be misleading. That is, the source and destination operand symbols must specify the correct **type** (size) of the operands (bytes, words, or doublewords), but they do not have to specify the correct **location**. The locations of the source and destination operands are always specified by the DS:(E)SI and ES:(E)DI registers, which must be loaded correctly before the move string instruction is executed.

The no-operands form provides "short forms" of the byte, word, and doubleword versions of the MOVS instructions. Here also DS:(E)SI and ES:(E)DI are assumed to be the source and destination operands, respectively. The size of the source and destination operands is selected with the mnemonic: MOVSB (byte move), MOVSW (word move), or MOVSD (doubleword move).

After the move operation, the (E)SI and (E)DI registers are incremented or decremented automatically according to the setting of the DF flag in the EFLAGS register. (If the DF flag is 0, the (E)SI and (E)DI register are incremented; if the DF flag is 1, the (E)SI and (E)DI registers are decremented.) The registers are incremented or decremented by 1 for byte operations, by 2 for word operations, or by 4 for doubleword operations.

The MOVS, MOVSB, MOVSW, and MOVSD instructions can be preceded by the REP prefix (see "REP/REPE/REPZ /REPNE/REPNZ—Repeat String Operation Prefix" in Chapter 4 of the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 2B*, for a description of the REP prefix) for block moves of ECX bytes, words, or doublewords.

In 64-bit mode, the instruction's default address size is 64 bits, 32-bit address size is supported using the prefix 67H. The 64-bit addresses are specified by RSI and RDI;

32-bit address are specified by ESI and EDI. Use of the REX.W prefix promotes doubleword operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

```
DEST ← SRC;

Non-64-bit Mode:

IF (Byte move)
   THEN IF DF = 0
        THEN
             (E)SI ← (E)SI + 1;
             (E)DI ← (E)DI + 1;
        ELSE
             (E)SI ← (E)SI – 1;
             (E)DI ← (E)DI – 1;
        FI;
   ELSE IF (Word move)
        THEN IF DF = 0
             (E)SI ← (E)SI + 2;
             (E)DI ← (E)DI + 2;
             FI;
        ELSE
             (E)SI ← (E)SI – 2;
             (E)DI ← (E)DI – 2;
        FI;
   ELSE IF (Doubleword move)
        THEN IF DF = 0
             (E)SI ← (E)SI + 4;
             (E)DI ← (E)DI + 4;
             FI;
        ELSE
             (E)SI ← (E)SI – 4;
             (E)DI ← (E)DI – 4;
        FI;
FI;
64-bit Mode:

IF (Byte move)
   THEN IF DF = 0
        THEN
             (R|E)SI ← (R|E)SI + 1;
             (R|E)DI ← (R|E)DI + 1;
```

```
        ELSE
            (R|E)SI ← (R|E)SI – 1;
            (R|E)DI ← (R|E)DI – 1;
        FI;
    ELSE IF (Word move)
        THEN IF DF = 0
            (R|E)SI ← (R|E)SI + 2;
            (R|E)DI ← (R|E)DI + 2;
            FI;
        ELSE
            (R|E)SI ← (R|E)SI – 2;
            (R|E)DI ← (R|E)DI – 2;
        FI;
    ELSE IF (Doubleword move)
        THEN IF DF = 0
            (R|E)SI ← (R|E)SI + 4;
            (R|E)DI ← (R|E)DI + 4;
            FI;
        ELSE
            (R|E)SI ← (R|E)SI – 4;
            (R|E)DI ← (R|E)DI – 4;
        FI;
    ELSE IF (Quadword move)
        THEN IF DF = 0
            (R|E)SI ← (R|E)SI + 8;
            (R|E)DI ← (R|E)DI + 8;
            FI;
        ELSE
            (R|E)SI ← (R|E)SI – 8;
            (R|E)DI ← (R|E)DI – 8;
        FI;
FI;
```

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If the destination is located in a non-writable segment.
	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.

#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

MOVSD—Move Scalar Double-Precision Floating-Point Value

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 10 /*r* MOVSD *xmm1*, *xmm2/m64*	A	V/V	SSE2	Move scalar double-precision floating-point value from *xmm2/m64* to *xmm1* register.
VEX.NDS.LIG.F2.0F.WIG 10 /r VMOVSD xmm1, xmm2, xmm3	B	V/V	AVX	Merge scalar double-precision floating-point value from xmm2 and xmm3 to xmm1 register.
VEX.LIG.F2.0F.WIG 10 /r VMOVSD xmm1, m64	D	V/V	AVX	Load scalar double-precision floating-point value from m64 to xmm1 register.
F2 0F 11 /*r* MOVSD *xmm2/m64*, *xmm1*	C	V/V	SSE2	Move scalar double-precision floating-point value from *xmm1* register to *xmm2/m64*.
VEX.NDS.LIG.F2.0F.WIG 11 /r VMOVSD xmm1, xmm2, xmm3	E	V/V	AVX	Merge scalar double-precision floating-point value from xmm2 and xmm3 registers to xmm1.
VEX.LIG.F2.0F.WIG 11 /r VMOVSD m64, xmm1	C	V/V	AVX	Move scalar double-precision floating-point value from xmm1 register to m64.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA
C	ModRM:r/m (w)	ModRM:reg (r)	NA	NA
D	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
E	ModRM:r/m (w)	VEX.vvvv (r)	ModRM:reg (r)	NA

Description

MOVSD moves a scalar double-precision floating-point value from the source operand (second operand) to the destination operand (first operand). The source and destination operands can be XMM registers or 64-bit memory locations. This instruction can be used to move a double-precision floating-point value to and from the low quadword of an XMM register and a 64-bit memory location, or to move a double-precision floating-point value between the low quadwords of two XMM registers. The instruction cannot be used to transfer data between memory locations.

For non-VEX encoded instruction syntax and when the source and destination operands are XMM registers, the high quadword of the destination operand remains unchanged. When the source operand is a memory location and destination operand is an XMM registers, the high quadword of the destination operand is cleared to all 0s.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

Note: For the "VMOVSD m64, xmm1" (memory store form) instruction version, VEX.vvvv is reserved and must be 1111b, otherwise instruction will #UD.

Note: For the "VMOVSD xmm1, m64" (memory load form) instruction version, VEX.vvvv is reserved and must be 1111b otherwise instruction will #UD.

VEX encoded instruction syntax supports two source operands and a destination operand if ModR/M.mod field is 11B. VEX.vvvv is used to encode the first source operand (the second operand). The low 128 bits of the destination operand stores the result of merging the low quadword of the second source operand with the quad word in bits 127:64 of the first source operand. The upper bits of the destination operand are cleared.

Operation

MOVSD (128-bit Legacy SSE version: MOVSD XMM1, XMM2)

DEST[63:0] ← SRC[63:0]
DEST[VLMAX-1:64] (Unmodified)

MOVSD/VMOVSD (128-bit versions: MOVSD m64, xmm1 or VMOVSD m64, xmm1)

DEST[63:0] ← SRC[63:0]

MOVSD (128-bit Legacy SSE version: MOVSD XMM1, m64)

DEST[63:0] ← SRC[63:0]
DEST[127:64] ← 0
DEST[VLMAX-1:128] (Unmodified)

VMOVSD (VEX.NDS.128.F2.0F 11 /r: VMOVSD xmm1, xmm2, xmm3)

DEST[63:0] ← SRC2[63:0]
DEST[127:64] ← SRC1[127:64]
DEST[VLMAX-1:128] ← 0

VMOVSD (VEX.NDS.128.F2.0F 10 /r: VMOVSD xmm1, xmm2, xmm3)

DEST[63:0] ← SRC2[63:0]
DEST[127:64] ← SRC1[127:64]

DEST[VLMAX-1:128] ← 0

VMOVSD (VEX.NDS.128.F2.0F 10 /r: VMOVSD xmm1, m64)
DEST[63:0] ← SRC[63:0]
DEST[VLMAX-1:64] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

MOVSD	__m128d _mm_load_sd (double *p)
MOVSD	void _mm_store_sd (double *p, __m128d a)
MOVSD	__m128d _mm_store_sd (__m128d a, __m128d b)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 5; additionally

#UD	If VEX.vvvv != 1111B.

MOVSHDUP—Move Packed Single-FP High and Duplicate

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 16 /*r* MOVSHDUP *xmm1*, *xmm2/m128*	A	V/V	SSE3	Move two single-precision floating-point values from the higher 32-bit operand of each qword in *xmm2/m128* to *xmm1* and duplicate each 32-bit operand to the lower 32-bits of each qword.
VEX.128.F3.0F.WIG 16 /r VMOVSHDUP xmm1, xmm2/m128	A	V/V	AVX	Move odd index single-precision floating-point values from xmm2/mem and duplicate each element into xmm1.
VEX.256.F3.0F.WIG 16 /r VMOVSHDUP ymm1, ymm2/m256	A	V/V	AVX	Move odd index single-precision floating-point values from ymm2/mem and duplicate each element into ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

The linear address corresponds to the address of the least-significant byte of the referenced memory data. When a memory address is indicated, the 16 bytes of data at memory location m128 are loaded and the single-precision elements in positions 1 and 3 are duplicated. When the register-register form of this operation is used, the same operation is performed but with data coming from the 128-bit source register. See Figure 3-28.

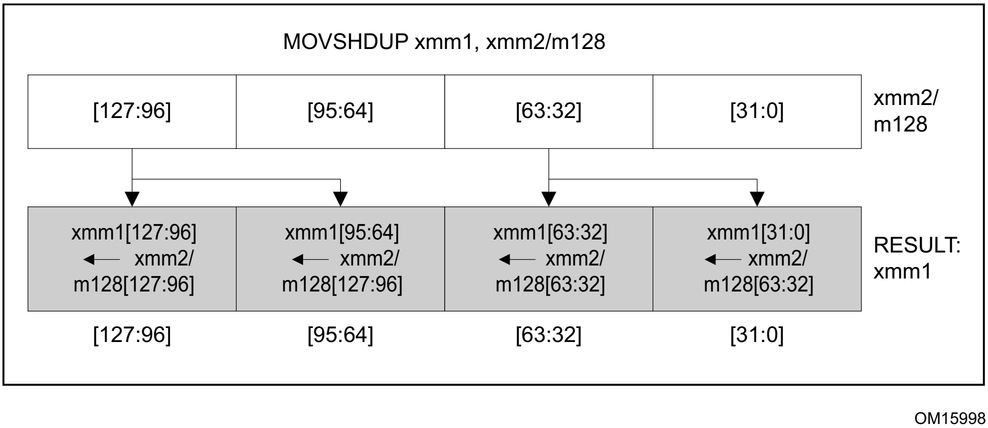


Figure 3-28. MOVSHDUP—Move Packed Single-FP High and Duplicate

In 64-bit mode, use of the REX prefix in the form of REX.R permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

Operation

MOVSHDUP (128-bit Legacy SSE version)
DEST[31:0] ← SRC[63:32]
DEST[63:32] ← SRC[63:32]
DEST[95:64] ← SRC[127:96]
DEST[127:96] ← SRC[127:96]
DEST[VLMAX-1:128] (Unmodified)

VMOVSHDUP (VEX.128 encoded version)
DEST[31:0] ← SRC[63:32]
DEST[63:32] ← SRC[63:32]
DEST[95:64] ← SRC[127:96]
DEST[127:96] ← SRC[127:96]

DEST[VLMAX-1:128] ← 0

VMOVSHDUP (VEX.256 encoded version)
DEST[31:0] ← SRC[63:32]
DEST[63:32] ← SRC[63:32]
DEST[95:64] ← SRC[127:96]
DEST[127:96] ← SRC[127:96]
DEST[159:128] ← SRC[191:160]
DEST[191:160] ← SRC[191:160]
DEST[223:192] ← SRC[255:224]
DEST[255:224] ← SRC[255:224]

Intel C/C++ Compiler Intrinsic Equivalent

(V)MOVSHDUP __m128 _mm_movehdup_ps(__m128 a)

VMOVSHDUP __m256 _mm256_movehdup_ps (__m256 a);

Exceptions

General protection exception if not aligned on 16-byte boundary, regardless of segment.

Numeric Exceptions

None

Other Exceptions

See Exceptions Type 2.

MOVSLDUP—Move Packed Single-FP Low and Duplicate

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 12 /*r* MOVSLDUP *xmm1*, *xmm2/m128*	A	V/V	SSE3	Move two single-precision floating-point values from the lower 32-bit operand of each qword in *xmm2/m128* to *xmm1* and duplicate each 32-bit operand to the higher 32-bits of each qword.
VEX.128.F3.0F.WIG 12 /r VMOVSLDUP xmm1, xmm2/m128	A	V/V	AVX	Move even index single-precision floating-point values from xmm2/mem and duplicate each element into xmm1.
VEX.256.F3.0F.WIG 12 /r VMOVSLDUP ymm1, ymm2/m256	A	V/V	AVX	Move even index single-precision floating-point values from ymm2/mem and duplicate each element into ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

The linear address corresponds to the address of the least-significant byte of the referenced memory data. When a memory address is indicated, the 16 bytes of data at memory location m128 are loaded and the single-precision elements in positions 0 and 2 are duplicated. When the register-register form of this operation is used, the same operation is performed but with data coming from the 128-bit source register.

See Figure 3-29.

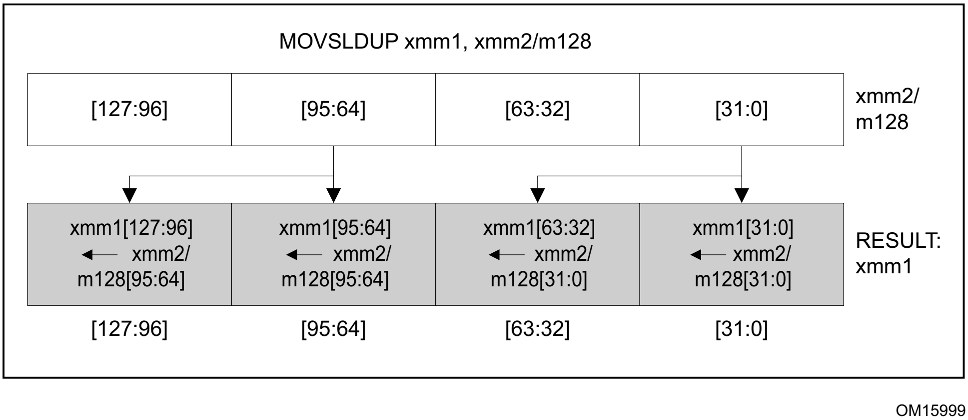


Figure 3-29. MOVSLDUP—Move Packed Single-FP Low and Duplicate

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

Operation

MOVSLDUP (128-bit Legacy SSE version)

DEST[31:0] ← SRC[31:0]
DEST[63:32] ← SRC[31:0]
DEST[95:64] ← SRC[95:64]
DEST[127:96] ← SRC[95:64]
DEST[VLMAX-1:128] (Unmodified)

VMOVSLDUP (VEX.128 encoded version)

DEST[31:0] ← SRC[31:0]
DEST[63:32] ← SRC[31:0]
DEST[95:64] ← SRC[95:64]
DEST[127:96] ← SRC[95:64]

```
DEST[VLMAX-1:128] ← 0
```

VMOVSLDUP (VEX.256 encoded version)

```
DEST[31:0] ← SRC[31:0]
DEST[63:32] ← SRC[31:0]
DEST[95:64] ← SRC[95:64]
DEST[127:96] ← SRC[95:64]
DEST[159:128] ← SRC[159:128]
DEST[191:160] ← SRC[159:128]
DEST[223:192] ← SRC[223:192]
DEST[255:224] ← SRC[223:192]
```

Intel C/C++ Compiler Intrinsic Equivalent

```
(V)MOVSLDUP     __m128 _mm_moveldup_ps(__m128 a)
VMOVSLDUP __m256 _mm256_moveldup_ps (__m256 a);
```

Exceptions

General protection exception if not aligned on 16-byte boundary, regardless of segment.

Numeric Exceptions

None.

Other Exceptions

See Exceptions Type 4; additionally

#UD	If VEX.vvvv != 1111B.

MOVSS—Move Scalar Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 10 /*r* MOVSS *xmm1*, *xmm2/m32*	A	V/V	SSE	Move scalar single-precision floating-point value from *xmm2/m32* to *xmm1* register.
VEX.NDS.LIG.F3.0F.WIG 10 /r VMOVSS xmm1, xmm2, xmm3	B	V/V	AVX	Merge scalar single-precision floating-point value from xmm2 and xmm3 to xmm1 register.
VEX.LIG.F3.0F.WIG 10 /r VMOVSS xmm1, m32	D	V/V	AVX	Load scalar single-precision floating-point value from m32 to xmm1 register.
F3 0F 11 /*r* MOVSS *xmm2/m32*, *xmm*	C	V/V	SSE	Move scalar single-precision floating-point value from *xmm1* register to *xmm2/m32*.
VEX.NDS.LIG.F3.0F.WIG 11 /r VMOVSS xmm1, xmm2, xmm3	E	V/V	AVX	Move scalar single-precision floating-point value from xmm2 and xmm3 to xmm1 register.
VEX.LIG.F3.0F.WIG 11 /r VMOVSS m32, xmm1	C	V/V	AVX	Move scalar single-precision floating-point value from xmm1 register to m32.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA
C	ModRM:r/m (w)	ModRM:reg (r)	NA	NA
D	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
E	ModRM:r/m (w)	VEX.vvvv (r)	ModRM:reg (r)	NA

Description

Moves a scalar single-precision floating-point value from the source operand (second operand) to the destination operand (first operand). The source and destination operands can be XMM registers or 32-bit memory locations. This instruction can be used to move a single-precision floating-point value to and from the low doubleword of an XMM register and a 32-bit memory location, or to move a single-precision floating-point value between the low doublewords of two XMM registers. The instruction cannot be used to transfer data between memory locations.

For non-VEX encoded syntax and when the source and destination operands are XMM registers, the high doublewords of the destination operand remains unchanged. When the source operand is a memory location and destination operand is an XMM registers, the high doublewords of the destination operand is cleared to all 0s.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

VEX encoded instruction syntax supports two source operands and a destination operand if ModR/M.mod field is 11B. VEX.vvvv is used to encode the first source operand (the second operand). The low 128 bits of the destination operand stores the result of merging the low dword of the second source operand with three dwords in bits 127:32 of the first source operand. The upper bits of the destination operand are cleared.

Note: For the "VMOVSS m32, xmm1" (memory store form) instruction version, VEX.vvvv is reserved and must be 1111b otherwise instruction will #UD.

Note: For the "VMOVSS xmm1, m32" (memory load form) instruction version, VEX.vvvv is reserved and must be 1111b otherwise instruction will #UD.

Operation

MOVSS (Legacy SSE version when the source and destination operands are both XMM registers)

```
DEST[31:0] ← SRC[31:0]
DEST[VLMAX-1:32] (Unmodified)
```

MOVSS/VMOVSS (when the source operand is an XMM register and the destination is memory)

```
DEST[31:0] ← SRC[31:0]
```

MOVSS (Legacy SSE version when the source operand is memory and the destination is an XMM register)

```
DEST[31:0] ← SRC[31:0]
DEST[127:32] ← 0
DEST[VLMAX-1:128] (Unmodified)
```

VMOVSS (VEX.NDS.128.F3.0F 11 /r where the destination is an XMM register)

```
DEST[31:0] ← SRC2[31:0]
DEST[127:32] ← SRC1[127:32]
DEST[VLMAX-1:128] ← 0
```

VMOVSS (VEX.NDS.128.F3.0F 10 /r where the source and destination are XMM registers)

DEST[31:0] ← SRC2[31:0]
DEST[127:32] ← SRC1[127:32]
DEST[VLMAX-1:128] ← 0

VMOVSS (VEX.NDS.128.F3.0F 10 /r when the source operand is memory and the destination is an XMM register)
DEST[31:0] ← SRC[31:0]
DEST[VLMAX-1:32] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

MOVSS	__m128 _mm_load_ss(float * p)
MOVSS	void _mm_store_ss(float * p, __m128 a)
MOVSS	__m128 _mm_move_ss(__m128 a, __m128 b)

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 5; additionally

#UD	If VEX.vvvv != 1111B.

MOVSX/MOVSXD—Move with Sign-Extension

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F BE /r	MOVSX *r16, r/m8*	A	Valid	Valid	Move byte to word with sign-extension.
0F BE /r	MOVSX *r32, r/m8*	A	Valid	Valid	Move byte to doubleword with sign-extension.
REX + 0F BE /r	MOVSX *r64, r/m8**	A	Valid	N.E.	Move byte to quadword with sign-extension.
0F BF /r	MOVSX *r32, r/m16*	A	Valid	Valid	Move word to doubleword, with sign-extension.
REX.W + 0F BF /r	MOVSX *r64, r/m16*	A	Valid	N.E.	Move word to quadword with sign-extension.
REX.W** + 63 /r	MOVSXD *r64, r/m32*	A	Valid	N.E.	Move doubleword to quadword with sign-extension.

NOTES:

* In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

** The use of MOVSXD without REX.W in 64-bit mode is discouraged, Regular MOV should be used instead of using MOVSXD without REX.W.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Copies the contents of the source operand (register or memory location) to the destination operand (register) and sign extends the value to 16 or 32 bits (see Figure 7-6 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*). The size of the converted value depends on the operand-size attribute.

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits. See the summary chart at the beginning of this section for encoding data and limits.

Operation

DEST ← SignExtend(SRC);

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

#UD	If the LOCK prefix is used.

MOVUPD—Move Unaligned Packed Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 10 /*r* MOVUPD *xmm1*, *xmm2/m128*	A	V/V	SSE2	Move packed double-precision floating-point values from *xmm2/m128* to xmm1.
VEX.128.66.0F.WIG 10 /r VMOVUPD xmm1, xmm2/m128	A	V/V	AVX	Move unaligned packed double-precision floating-point from xmm2/mem to xmm1.
VEX.256.66.0F.WIG 10 /r VMOVUPD ymm1, ymm2/m256	A	V/V	AVX	Move unaligned packed double-precision floating-point from ymm2/mem to ymm1.
66 0F 11 /*r* MOVUPD *xmm2/m128, xmm*	B	V/V	SSE2	Move packed double-precision floating-point values from xmm1 to *xmm2/m128*.
VEX.128.66.0F.WIG 11 /r VMOVUPD xmm2/m128, xmm1	B	V/V	AVX	Move unaligned packed double-precision floating-point from xmm1 to xmm2/mem.
VEX.256.66.0F.WIG 11 /r VMOVUPD ymm2/m256, ymm1	B	V/V	AVX	Move unaligned packed double-precision floating-point from ymm1 to ymm2/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

128-bit versions:

Moves a double quadword containing two packed double-precision floating-point values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load an XMM register from a 128-bit memory location, store the contents of an XMM register into a 128-bit memory location, or move data between two XMM registers.

128-bit Legacy SSE version: Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

When the source or destination operand is a memory operand, the operand may be unaligned on a 16-byte boundary without causing a general-protection exception (#GP) to be generated.[1]

To move double-precision floating-point values to and from memory locations that are known to be aligned on 16-byte boundaries, use the MOVAPD instruction.

While executing in 16-bit addressing mode, a linear address for a 128-bit data access that overlaps the end of a 16-bit segment is not allowed and is defined as reserved behavior. A specific processor implementation may or may not generate a general-protection exception (#GP) in this situation, and the address that spans the end of the segment may or may not wrap around to the beginning of the segment.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

VEX.256 encoded version:

Moves 256 bits of packed double-precision floating-point values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load a YMM register from a 256-bit memory location, to store the contents of a YMM register into a 256-bit memory location, or to move data between two YMM registers.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

Operation

MOVUPD (128-bit load and register-copy form Legacy SSE version)

DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] (Unmodified)

(V)MOVUPD (128-bit store form)

DEST[127:0] ← SRC[127:0]

1. If alignment checking is enabled (CR0.AM = 1, RFLAGS.AC = 1, and CPL = 3), an alignment-check exception (#AC) may or may not be generated (depending on processor implementation) when the operand is not aligned on an 8-byte boundary.

VMOVUPD (VEX.128 encoded version)
DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] ← 0

VMOVUPD (VEX.256 encoded version)
DEST[255:0] ← SRC[255:0]

Intel C/C++ Compiler Intrinsic Equivalent

MOVUPD __m128 _mm_loadu_pd(double * p)

MOVUPD void _mm_storeu_pd(double *p, __m128 a)

VMOVUPD __m256d _mm256_loadu_pd (__m256d * p);

VMOVUPD _mm256_storeu_pd(_m256d *p, __m256d a);

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 4

Note treatment of #AC varies; additionally

#UD	If VEX.vvvv != 1111B.

MOVUPS—Move Unaligned Packed Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 10 /*r* MOVUPS *xmm1, xmm2/m128*	A	V/V	SSE	Move packed single-precision floating-point values from *xmm2/m128* to xmm1.
VEX.128.0F.WIG 10 /r VMOVUPS xmm1, xmm2/m128	A	V/V	AVX	Move unaligned packed single-precision floating-point from xmm2/mem to xmm1.
VEX.256.0F.WIG 10 /r VMOVUPS ymm1, ymm2/m256	A	V/V	AVX	Move unaligned packed single-precision floating-point from ymm2/mem to ymm1.
0F 11 /*r* MOVUPS *xmm2/m128, xmm1*	B	V/V	SSE	Move packed single-precision floating-point values from xmm1 to *xmm2/m128*.
VEX.128.0F.WIG 11 /r VMOVUPS xmm2/m128, xmm1	B	V/V	AVX	Move unaligned packed single-precision floating-point from xmm1 to xmm2/mem.
VEX.256.0F.WIG 11 /r VMOVUPS ymm2/m256, ymm1	B	V/V	AVX	Move unaligned packed single-precision floating-point from ymm1 to ymm2/mem.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA
B	ModRM:r/m (w)	ModRM:reg (r)	NA	NA

Description

128-bit versions: Moves a double quadword containing four packed single-precision floating-point values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load an XMM register from a 128-bit memory location, store the contents of an XMM register into a 128-bit memory location, or move data between two XMM registers.

128-bit Legacy SSE version: Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

When the source or destination operand is a memory operand, the operand may be unaligned on a 16-byte boundary without causing a general-protection exception (#GP) to be generated.[1]

To move packed single-precision floating-point values to and from memory locations that are known to be aligned on 16-byte boundaries, use the MOVAPS instruction.

While executing in 16-bit addressing mode, a linear address for a 128-bit data access that overlaps the end of a 16-bit segment is not allowed and is defined as reserved behavior. A specific processor implementation may or may not generate a general-protection exception (#GP) in this situation, and the address that spans the end of the segment may or may not wrap around to the beginning of the segment.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

VEX.256 encoded version: Moves 256 bits of packed single-precision floating-point values from the source operand (second operand) to the destination operand (first operand). This instruction can be used to load a YMM register from a 256-bit memory location, to store the contents of a YMM register into a 256-bit memory location, or to move data between two YMM registers.

Note: In VEX-encoded versions, VEX.vvvv is reserved and must be 1111b otherwise instructions will #UD.

Operation

MOVUPS (128-bit load and register-copy form Legacy SSE version)

DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] (Unmodified)

(V)MOVUPS (128-bit store form)

DEST[127:0] ← SRC[127:0]

VMOVUPS (VEX.128 encoded load-form)

DEST[127:0] ← SRC[127:0]
DEST[VLMAX-1:128] ← 0

VMOVUPS (VEX.256 encoded version)

1. If alignment checking is enabled (CR0.AM = 1, RFLAGS.AC = 1, and CPL = 3), an alignment-check exception (#AC) may or may not be generated (depending on processor implementation) when the operand is not aligned on an 8-byte boundary.

DEST[255:0] ← SRC[255:0]

Intel C/C++ Compiler Intrinsic Equivalent

MOVUPS __m128 _mm_loadu_ps(double * p)

MOVUPS void _mm_storeu_ps(double *p, __m128 a)

VMOVUPS __m256 _mm256_loadu_ps (__m256 * p);

VMOVUPS _mm256_storeu_ps(_m256 *p, __m256 a);

SIMD Floating-Point Exceptions

None.

Other Exceptions

See Exceptions Type 4

Note treatment of #AC varies; additionally

#UD	If VEX.vvvv != 1111B.

MOVZX—Move with Zero-Extend

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F B6 /*r*	MOVZX *r16, r/m8*	A	Valid	Valid	Move byte to word with zero-extension.
0F B6 /*r*	MOVZX *r32, r/m8*	A	Valid	Valid	Move byte to doubleword, zero-extension.
REX.W + 0F B6 /*r*	MOVZX *r64, r/m8**	A	Valid	N.E.	Move byte to quadword, zero-extension.
0F B7 /*r*	MOVZX *r32, r/m16*	A	Valid	Valid	Move word to doubleword, zero-extension.
REX.W + 0F B7 /*r*	MOVZX *r64, r/m16*	A	Valid	N.E.	Move word to quadword, zero-extension.

NOTES:

* In 64-bit mode, r/m8 can not be encoded to access the following byte registers if the REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (w)	ModRM:r/m (r)	NA	NA

Description

Copies the contents of the source operand (register or memory location) to the destination operand (register) and zero extends the value. The size of the converted value depends on the operand-size attribute.

In 64-bit mode, the instruction's default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bit operands. See the summary chart at the beginning of this section for encoding data and limits.

Operation

DEST ← ZeroExtend(SRC);

Flags Affected

None.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

MPSADBW — Compute Multiple Packed Sums of Absolute Difference

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 3A 42 /r ib MPSADBW *xmm1, xmm2/m128, imm8*	A	V/V	SSE4_1	Sums absolute 8-bit integer difference of adjacent groups of 4 byte integers in *xmm1* and *xmm2/m128* and writes the results in *xmm1*. Starting offsets within *xmm1* and *xmm2/m128* are determined by *imm8*.
VEX.NDS.128.66.0F3A.WIG 42 /r ib VMPSADBW xmm1, xmm2, xmm3/m128, imm8	B	V/V	AVX	Sums absolute 8-bit integer difference of adjacent groups of 4 byte integers in xmm2 and xmm3/m128 and writes the results in xmm1. Starting offsets within xmm2 and xmm3/m128 are determined by imm8.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	imm8	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

MPSADBW sums the absolute difference (SAD) of a pair of unsigned bytes for a group of 4 byte pairs, and produces 8 SAD results (one for each 4 byte-pairs) stored as 8 word integers in the destination operand (first operand). Each 4 byte pairs are selected from the source operand (first opeand) and the destination according to the bit fields specified in the immediate byte (third operand).

The immediate byte provides two bit fields:

SRC_OFFSET: the value of Imm8[1:0]*32 specifies the offset of the 4 sequential source bytes in the source operand.

DEST_OFFSET: the value of Imm8[2]*32 specifies the offset of the first of 8 groups of 4 sequential destination bytes in the destination operand. The next four destination bytes starts at DEST_OFFSET + 8, etc.

The SAD operation is repeated 8 times, each time using the same 4 source bytes but selecting the next group of 4 destination bytes starting at the next higher byte in the destination. Each 16-bit sum is written to destination.

128-bit Legacy SSE version: The first source and destination are the same. Bits (VLMAX-1:128) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

If VMPSADBW is encoded with VEX.L= 1, an attempt to execute the instruction encoded with VEX.L= 1 will cause an #UD exception.

Operation

MPSADBW (128-bit Legacy SSE version)

```
SRC_OFFSET ← imm8[1:0]*32
DEST_OFFSET ← imm8[2]*32
DEST_BYTE0 ← DEST[DEST_OFFSET+7:DEST_OFFSET]
DEST_BYTE1 ← DEST[DEST_OFFSET+15:DEST_OFFSET+8]
DEST_BYTE2 ← DEST[DEST_OFFSET+23:DEST_OFFSET+16]
DEST_BYTE3 ← DEST[DEST_OFFSET+31:DEST_OFFSET+24]
DEST_BYTE4 ← DEST[DEST_OFFSET+39:DEST_OFFSET+32]
DEST_BYTE5 ← DEST[DEST_OFFSET+47:DEST_OFFSET+40]
DEST_BYTE6 ← DEST[DEST_OFFSET+55:DEST_OFFSET+48]
DEST_BYTE7 ← DEST[DEST_OFFSET+63:DEST_OFFSET+56]
DEST_BYTE8 ← DEST[DEST_OFFSET+71:DEST_OFFSET+64]
DEST_BYTE9 ← DEST[DEST_OFFSET+79:DEST_OFFSET+72]
DEST_BYTE10 ← DEST[DEST_OFFSET+87:DEST_OFFSET+80]

SRC_BYTE0 ← SRC[SRC_OFFSET+7:SRC_OFFSET]
SRC_BYTE1 ← SRC[SRC_OFFSET+15:SRC_OFFSET+8]
SRC_BYTE2 ← SRC[SRC_OFFSET+23:SRC_OFFSET+16]
SRC_BYTE3 ← SRC[SRC_OFFSET+31:SRC_OFFSET+24]

TEMP0 ← ABS( DEST_BYTE0 - SRC_BYTE0)
TEMP1 ← ABS( DEST_BYTE1 - SRC_BYTE1)
TEMP2 ← ABS( DEST_BYTE2 - SRC_BYTE2)
TEMP3 ← ABS( DEST_BYTE3 - SRC_BYTE3)
DEST[15:0] ← TEMP0 + TEMP1 + TEMP2 + TEMP3

TEMP0 ← ABS( DEST_BYTE1 - SRC_BYTE0)
TEMP1 ← ABS( DEST_BYTE2 - SRC_BYTE1)
TEMP2 ← ABS( DEST_BYTE3 - SRC_BYTE2)
TEMP3 ← ABS( DEST_BYTE4 - SRC_BYTE3)
DEST[31:16] ← TEMP0 + TEMP1 + TEMP2 + TEMP3
```

```
TEMP0 ← ABS( DEST_BYTE2 - SRC_BYTE0)
TEMP1 ← ABS( DEST_BYTE3 - SRC_BYTE1)
TEMP2 ← ABS( DEST_BYTE4 - SRC_BYTE2)
TEMP3 ← ABS( DEST_BYTE5 - SRC_BYTE3)
DEST[47:32] ← TEMP0 + TEMP1 + TEMP2 + TEMP3

TEMP0 ← ABS( DEST_BYTE3 - SRC_BYTE0)
TEMP1 ← ABS( DEST_BYTE4 - SRC_BYTE1)
TEMP2 ← ABS( DEST_BYTE5 - SRC_BYTE2)
TEMP3 ← ABS( DEST_BYTE6 - SRC_BYTE3)
DEST[63:48] ← TEMP0 + TEMP1 + TEMP2 + TEMP3

TEMP0 ← ABS( DEST_BYTE4 - SRC_BYTE0)
TEMP1 ← ABS( DEST_BYTE5 - SRC_BYTE1)
TEMP2 ← ABS( DEST_BYTE6 - SRC_BYTE2)
TEMP3 ← ABS( DEST_BYTE7 - SRC_BYTE3)
DEST[79:64] ← TEMP0 + TEMP1 + TEMP2 + TEMP3

TEMP0 ← ABS( DEST_BYTE5 - SRC_BYTE0)
TEMP1 ← ABS( DEST_BYTE6 - SRC_BYTE1)
TEMP2 ← ABS( DEST_BYTE7 - SRC_BYTE2)
TEMP3 ← ABS( DEST_BYTE8 - SRC_BYTE3)
DEST[95:80] ← TEMP0 + TEMP1 + TEMP2 + TEMP3

TEMP0 ← ABS( DEST_BYTE6 - SRC_BYTE0)
TEMP1 ← ABS( DEST_BYTE7 - SRC_BYTE1)
TEMP2 ← ABS( DEST_BYTE8 - SRC_BYTE2)
TEMP3 ← ABS( DEST_BYTE9 - SRC_BYTE3)
DEST[111:96] ← TEMP0 + TEMP1 + TEMP2 + TEMP3

TEMP0 ← ABS( DEST_BYTE7 - SRC_BYTE0)
TEMP1 ← ABS( DEST_BYTE8 - SRC_BYTE1)
TEMP2 ← ABS( DEST_BYTE9 - SRC_BYTE2)
TEMP3 ← ABS( DEST_BYTE10 - SRC_BYTE3)
DEST[127:112] ← TEMP0 + TEMP1 + TEMP2 + TEMP3
DEST[VLMAX-1:128] (Unmodified)
```

VMPSADBW (VEX.128 encoded version)

```
SRC2_OFFSET ← imm8[1:0]*32
SRC1_OFFSET ← imm8[2]*32
SRC1_BYTE0 ← SRC1[SRC1_OFFSET+7:SRC1_OFFSET]
SRC1_BYTE1 ← SRC1[SRC1_OFFSET+15:SRC1_OFFSET+8]
```

```
SRC1_BYTE2 ← SRC1[SRC1_OFFSET+23:SRC1_OFFSET+16]
SRC1_BYTE3 ← SRC1[SRC1_OFFSET+31:SRC1_OFFSET+24]
SRC1_BYTE4 ← SRC1[SRC1_OFFSET+39:SRC1_OFFSET+32]
SRC1_BYTE5 ← SRC1[SRC1_OFFSET+47:SRC1_OFFSET+40]
SRC1_BYTE6 ← SRC1[SRC1_OFFSET+55:SRC1_OFFSET+48]
SRC1_BYTE7 ← SRC1[SRC1_OFFSET+63:SRC1_OFFSET+56]
SRC1_BYTE8 ← SRC1[SRC1_OFFSET+71:SRC1_OFFSET+64]
SRC1_BYTE9 ← SRC1[SRC1_OFFSET+79:SRC1_OFFSET+72]
SRC1_BYTE10 ← SRC1[SRC1_OFFSET+87:SRC1_OFFSET+80]

SRC2_BYTE0 ←SRC2[SRC2_OFFSET+7:SRC2_OFFSET]
SRC2_BYTE1 ← SRC2[SRC2_OFFSET+15:SRC2_OFFSET+8]
SRC2_BYTE2 ← SRC2[SRC2_OFFSET+23:SRC2_OFFSET+16]
SRC2_BYTE3 ← SRC2[SRC2_OFFSET+31:SRC2_OFFSET+24]

TEMP0 ← ABS(SRC1_BYTE0 - SRC2_BYTE0)
TEMP1 ← ABS(SRC1_BYTE1 - SRC2_BYTE1)
TEMP2 ← ABS(SRC1_BYTE2 - SRC2_BYTE2)
TEMP3 ← ABS(SRC1_BYTE3 - SRC2_BYTE3)
DEST[15:0] ← TEMP0 + TEMP1 + TEMP2 + TEMP3
TEMP0 ← ABS(SRC1_BYTE1 - SRC2_BYTE0)
TEMP1 ← ABS(SRC1_BYTE2 - SRC2_BYTE1)
TEMP2 ← ABS(SRC1_BYTE3 - SRC2_BYTE2)
TEMP3 ← ABS(SRC1_BYTE4 - SRC2_BYTE3)
DEST[31:16] ← TEMP0 + TEMP1 + TEMP2 + TEMP3
TEMP0 ← ABS(SRC1_BYTE2 - SRC2_BYTE0)
TEMP1 ← ABS(SRC1_BYTE3 - SRC2_BYTE1)
TEMP2 ← ABS(SRC1_BYTE4 - SRC2_BYTE2)
TEMP3 ← ABS(SRC1_BYTE5 - SRC2_BYTE3)
DEST[47:32] ← TEMP0 + TEMP1 + TEMP2 + TEMP3
TEMP0 ← ABS(SRC1_BYTE3 - SRC2_BYTE0)
TEMP1 ← ABS(SRC1_BYTE4 - SRC2_BYTE1)
TEMP2 ← ABS(SRC1_BYTE5 - SRC2_BYTE2)
TEMP3 ← ABS(SRC1_BYTE6 - SRC2_BYTE3)
DEST[63:48] ← TEMP0 + TEMP1 + TEMP2 + TEMP3
TEMP0 ← ABS(SRC1_BYTE4 - SRC2_BYTE0)
TEMP1 ← ABS(SRC1_BYTE5 - SRC2_BYTE1)
TEMP2 ← ABS(SRC1_BYTE6 - SRC2_BYTE2)
TEMP3 ← ABS(SRC1_BYTE7 - SRC2_BYTE3)
DEST[79:64] ← TEMP0 + TEMP1 + TEMP2 + TEMP3
TEMP0 ← ABS(SRC1_BYTE5 - SRC2_BYTE0)
TEMP1 ← ABS(SRC1_BYTE6 - SRC2_BYTE1)
TEMP2 ← ABS(SRC1_BYTE7 - SRC2_BYTE2)
```

```
TEMP3 ← ABS(SRC1_BYTE8 - SRC2_BYTE3)
DEST[95:80] ← TEMP0 + TEMP1 + TEMP2 + TEMP3
TEMP0 ← ABS(SRC1_BYTE6 - SRC2_BYTE0)
TEMP1 ← ABS(SRC1_BYTE7 - SRC2_BYTE1)
TEMP2 ← ABS(SRC1_BYTE8 - SRC2_BYTE2)
TEMP3 ← ABS(SRC1_BYTE9 - SRC2_BYTE3)
DEST[111:96] ← TEMP0 + TEMP1 + TEMP2 + TEMP3

TEMP0 ← ABS(SRC1_BYTE7 - SRC2_BYTE0)
TEMP1 ← ABS(SRC1_BYTE8 - SRC2_BYTE1)
TEMP2 ← ABS(SRC1_BYTE9 - SRC2_BYTE2)
TEMP3 ← ABS(SRC1_BYTE10 - SRC2_BYTE3)
DEST[127:112] ← TEMP0 + TEMP1 + TEMP2 + TEMP3
DEST[VLMAX-1:128] ← 0
```

Intel C/C++ Compiler Intrinsic Equivalent

MPSADBW __m128i _mm_mpsadbw_epu8 (__m128i s1, __m128i s2, const int mask);

Flags Affected

None

Other Exceptions

See Exceptions Type 4; additionally

#UD	If VEX.L = 1.

MUL—Unsigned Multiply

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
F6 /4	MUL *r/m8*	A	Valid	Valid	Unsigned multiply (AX ← AL ∗ *r/m8*).
REX + F6 /4	MUL *r/m8**	A	Valid	N.E.	Unsigned multiply (AX ← AL ∗ *r/m8*).
F7 /4	MUL *r/m16*	A	Valid	Valid	Unsigned multiply (DX:AX ← AX ∗ *r/m16*).
F7 /4	MUL *r/m32*	A	Valid	Valid	Unsigned multiply (EDX:EAX ← EAX ∗ *r/m32*).
REX.W + F7 /4	MUL *r/m64*	A	Valid	N.E.	Unsigned multiply (RDX:RAX ← RAX ∗ *r/m64.*

NOTES:

* In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:r/m (r)	NA	NA	NA

Description

Performs an unsigned multiplication of the first operand (destination operand) and the second operand (source operand) and stores the result in the destination operand. The destination operand is an implied operand located in register AL, AX or EAX (depending on the size of the operand); the source operand is located in a general-purpose register or a memory location. The action of this instruction and the location of the result depends on the opcode and the operand size as shown in Table 3-66.

The result is stored in register AX, register pair DX:AX, or register pair EDX:EAX (depending on the operand size), with the high-order bits of the product contained in register AH, DX, or EDX, respectively. If the high-order bits of the product are 0, the CF and OF flags are cleared; otherwise, the flags are set.

In 64-bit mode, the instruction’s default operation size is 32 bits. Use of the REX.R prefix permits access to additional registers (R8-R15). Use of the REX.W prefix promotes operation to 64 bits.

See the summary chart at the beginning of this section for encoding data and limits.

Table 3-66. MUL Results

Operand Size	Source 1	Source 2	Destination
Byte	AL	r/m8	AX
Word	AX	r/m16	DX:AX
Doubleword	EAX	r/m32	EDX:EAX
Quadword	RAX	r/m64	RDX:RAX

Operation

```
IF (Byte operation)
    THEN
        AX ← AL ∗ SRC;
    ELSE (* Word or doubleword operation *)
        IF OperandSize = 16
            THEN
                DX:AX ← AX ∗ SRC;
            ELSE IF OperandSize = 32
                THEN EDX:EAX ← EAX ∗ SRC; FI;
            ELSE (* OperandSize = 64 *)
                RDX:RAX ← RAX ∗ SRC;
        FI;
FI;
```

Flags Affected

The OF and CF flags are set to 0 if the upper half of the result is 0; otherwise, they are set to 1. The SF, ZF, AF, and PF flags are undefined.

Protected Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
	If the DS, ES, FS, or GS register contains a NULL segment selector.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.
#UD	If the LOCK prefix is used.

Real-Address Mode Exceptions

#GP	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS	If a memory operand effective address is outside the SS segment limit.
#UD	If the LOCK prefix is used.

Virtual-8086 Mode Exceptions

#GP(0)	If a memory operand effective address is outside the CS, DS, ES, FS, or GS segment limit.
#SS(0)	If a memory operand effective address is outside the SS segment limit.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made.
#UD	If the LOCK prefix is used.

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#SS(0)	If a memory address referencing the SS segment is in a non-canonical form.
#GP(0)	If the memory address is in a non-canonical form.
#PF(fault-code)	If a page fault occurs.
#AC(0)	If alignment checking is enabled and an unaligned memory reference is made while the current privilege level is 3.

MULPD—Multiply Packed Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
66 0F 59 /*r* MULPD *xmm1*, *xmm2/m128*	A	V/V	SSE2	Multiply packed double-precision floating-point values in *xmm2/m128* by *xmm1*.
VEX.NDS.128.66.0F.WIG 59 /r VMULPD xmm1,xmm2, xmm3/m128	B	V/V	AVX	Multiply packed double-precision floating-point values from xmm3/mem to xmm2 and stores result in xmm1.
VEX.NDS.256.66.0F.WIG 59 /r VMULPD ymm1, ymm2, ymm3/m256	B	V/V	AVX	Multiply packed double-precision floating-point values from ymm3/mem to ymm2 and stores result in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs a SIMD multiply of the two or four packed double-precision floating-point values from the source operand (second operand) and the destination operand (first operand), and stores the packed double-precision floating-point results in the destination operand. The source operand can be an XMM register or a 128-bit memory location. The destination operand is an XMM register. See Figure 11-3 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for an illustration of a SIMD double-precision floating-point operation.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the destination YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

MULPD (128-bit Legacy SSE version)
DEST[63:0] ← DEST[63:0] * SRC[63:0]
DEST[127:64] ← DEST[127:64] * SRC[127:64]
DEST[VLMAX-1:128] (Unmodified)

VMULPD (VEX.128 encoded version)
DEST[63:0] ← SRC1[63:0] * SRC2[63:0]
DEST[127:64] ← SRC1[127:64] * SRC2[127:64]
DEST[VLMAX-1:128] ← 0

VMULPD (VEX.256 encoded version)
DEST[63:0] ← SRC1[63:0] * SRC2[63:0]
DEST[127:64] ← SRC1[127:64] * SRC2[127:64]
DEST[191:128] ← SRC1[191:128] * SRC2[191:128]
DEST[255:192] ← SRC1[255:192] * SRC2[255:192]

Intel C/C++ Compiler Intrinsic Equivalent

MULPD __m128d _mm_mul_pd (m128d a, m128d b)

VMULPD __m256d _mm256_mul_pd (__m256d a, __m256d b);

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 2

MULPS—Multiply Packed Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
0F 59 /*r* MULPS *xmm1*, *xmm2/m128*	A	V/V	SSE	Multiply packed single-precision floating-point values in *xmm2/mem* by *xmm1*.
VEX.NDS.128.0F.WIG 59 /r VMULPS xmm1,xmm2, xmm3/m128	B	V/V	AVX	Multiply packed single-precision floating-point values from xmm3/mem to xmm2 and stores result in xmm1.
VEX.NDS.256.0F.WIG 59 /r VMULPS ymm1, ymm2, ymm3/m256	B	V/V	AVX	Multiply packed single-precision floating-point values from ymm3/mem to ymm2 and stores result in ymm1.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Performs a SIMD multiply of the four packed single-precision floating-point values from the source operand (second operand) and the destination operand (first operand), and stores the packed single-precision floating-point results in the destination operand. See Figure 10-5 in the *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 1*, for an illustration of a SIMD single-precision floating-point operation.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The second source can be an XMM register or an 128-bit memory location. The destination is not distinct from the first source XMM register and the upper bits (VLMAX-1:128) of the corresponding YMM register destination are unmodified.

VEX.128 encoded version: the first source operand is an XMM register or 128-bit memory location. The destination operand is an XMM register. The upper bits (VLMAX-1:128) of the destination YMM register destination are zeroed.

VEX.256 encoded version: The first source operand is a YMM register. The second source operand can be a YMM register or a 256-bit memory location. The destination operand is a YMM register.

Operation

MULPS (128-bit Legacy SSE version)

```
DEST[31:0] ← SRC1[31:0] * SRC2[31:0]
DEST[63:32] ← SRC1[63:32] * SRC2[63:32]
DEST[95:64] ← SRC1[95:64] * SRC2[95:64]
DEST[127:96] ← SRC1[127:96] * SRC2[127:96]
DEST[VLMAX-1:128] (Unmodified)
```

VMULPS (VEX.128 encoded version)

```
DEST[31:0] ← SRC1[31:0] * SRC2[31:0]
DEST[63:32] ← SRC1[63:32] * SRC2[63:32]
DEST[95:64] ← SRC1[95:64] * SRC2[95:64]
DEST[127:96] ← SRC1[127:96] * SRC2[127:96]
DEST[VLMAX-1:128] ← 0
```

VMULPS (VEX.256 encoded version)

```
DEST[31:0] ← SRC1[31:0] * SRC2[31:0]
DEST[63:32] ← SRC1[63:32] * SRC2[63:32]
DEST[95:64] ← SRC1[95:64] * SRC2[95:64]
DEST[127:96] ← SRC1[127:96] * SRC2[127:96]
DEST[159:128] ← SRC1[159:128] * SRC2[159:128]
DEST[191:160]← SRC1[191:160] * SRC2[191:160]
DEST[223:192] ← SRC1[223:192] * SRC2[223:192]
DEST[255:224] ← SRC1[255:224] * SRC2[255:224].
```

Intel C/C++ Compiler Intrinsic Equivalent

```
MULPS     __m128 _mm_mul_ps(__m128 a, __m128 b)
VMULPS __m256 _mm256_mul_ps (__m256 a, __m256 b);
```

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 2

MULSD—Multiply Scalar Double-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F2 0F 59 /*r* MULSD *xmm1*, *xmm2/m64*	A	V/V	SSE2	Multiply the low double-precision floating-point value in *xmm2/mem64* by low double-precision floating-point value in *xmm1*.
VEX.NDS.LIG.F2.0F.WIG 59/r VMULSD xmm1,xmm2, xmm3/m64	B	V/V	AVX	Multiply the low double-precision floating-point value in xmm3/mem64 by low double precision floating-point value in xmm2.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Multiplies the low double-precision floating-point value in the source operand (second operand) by the low double-precision floating-point value in the destination operand (first operand), and stores the double-precision floating-point result in the destination operand. The source operand can be an XMM register or a 64-bit memory location. The destination operand is an XMM register. The high quadword of the destination operand remains unchanged. See Figure 11-4 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for an illustration of a scalar double-precision floating-point operation.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The first source operand and the destination operand are the same. Bits (VLMAX-1:64) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

MULSD (128-bit Legacy SSE version)
DEST[63:0] ← DEST[63:0] * SRC[63:0]
DEST[VLMAX-1:64] (Unmodified)

VMULSD (VEX.128 encoded version)
DEST[63:0] ← SRC1[63:0] * SRC2[63:0]
DEST[127:64] ← SRC1[127:64]
DEST[VLMAX-1:128] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

MULSD __m128d _mm_mul_sd (m128d a, m128d b)

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 3

MULSS—Multiply Scalar Single-Precision Floating-Point Values

Opcode/ Instruction	Op/ En	64/32-bit Mode	CPUID Feature Flag	Description
F3 0F 59 /*r* MULSS *xmm1*, *xmm2/m32*	A	V/V	SSE	Multiply the low single-precision floating-point value in *xmm2/mem* by the low single-precision floating-point value in *xmm1*.
VEX.NDS.LIG.F3.0F.WIG 59 /r VMULSS xmm1,xmm2, xmm3/m32	B	V/V	AVX	Multiply the low single-precision floating-point value in xmm3/mem by the low single-precision floating-point value in xmm2.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	ModRM:reg (r, w)	ModRM:r/m (r)	NA	NA
B	ModRM:reg (w)	VEX.vvvv (r)	ModRM:r/m (r)	NA

Description

Multiplies the low single-precision floating-point value from the source operand (second operand) by the low single-precision floating-point value in the destination operand (first operand), and stores the single-precision floating-point result in the destination operand. The source operand can be an XMM register or a 32-bit memory location. The destination operand is an XMM register. The three high-order doublewords of the destination operand remain unchanged. See Figure 10-6 in the *Intel® 64 and IA-32 Architectures Software Developer's Manual, Volume 1*, for an illustration of a scalar single-precision floating-point operation.

In 64-bit mode, use of the REX.R prefix permits this instruction to access additional registers (XMM8-XMM15).

128-bit Legacy SSE version: The first source operand and the destination operand are the same. Bits (VLMAX-1:32) of the corresponding YMM destination register remain unchanged.

VEX.128 encoded version: Bits (VLMAX-1:128) of the destination YMM register are zeroed.

Operation

MULSS (128-bit Legacy SSE version)

DEST[31:0] ← DEST[31:0] * SRC[31:0]
DEST[VLMAX-1:32] (Unmodified)

VMULSS (VEX.128 encoded version)
DEST[31:0] ← SRC1[31:0] * SRC2[31:0]
DEST[127:32] ← SRC1[127:32]
DEST[VLMAX-1:128] ← 0

Intel C/C++ Compiler Intrinsic Equivalent

MULSS __m128 _mm_mul_ss(__m128 a, __m128 b)

SIMD Floating-Point Exceptions

Overflow, Underflow, Invalid, Precision, Denormal.

Other Exceptions

See Exceptions Type 3

MWAIT—Monitor Wait

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
0F 01 *C9*	MWAIT	A	Valid	Valid	A hint that allow the processor to stop instruction execution and enter an implementation-dependent optimized state until occurrence of a class of events.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
A	NA	NA	NA	NA

Description

MWAIT instruction provides hints to allow the processor to enter an implementation-dependent optimized state. There are two principal targeted usages: address-range monitor and advanced power management. Both usages of MWAIT require the use of the MONITOR instruction.

A CPUID feature flag (ECX bit 3; CPUID executed EAX = 1) indicates the availability of MONITOR and MWAIT in the processor. When set, MWAIT may be executed only at privilege level 0 (use at any other privilege level results in an invalid-opcode exception). The operating system or system BIOS may disable this instruction by using the IA32_MISC_ENABLE MSR; disabling MWAIT clears the CPUID feature flag and causes execution to generate an illegal opcode exception.

This instruction's operation is the same in non-64-bit modes and 64-bit mode.

MWAIT for Address Range Monitoring

For address-range monitoring, the MWAIT instruction operates with the MONITOR instruction. The two instructions allow the definition of an address at which to wait (MONITOR) and a implementation-dependent-optimized operation to commence at the wait address (MWAIT). The execution of MWAIT is a hint to the processor that it can enter an implementation-dependent-optimized state while waiting for an event or a store operation to the address range armed by MONITOR.

ECX specifies optional extensions for the MWAIT instruction. EAX may contain hints such as the preferred optimized state the processor should enter.

For Pentium 4 processors (CPUID signature family 15 and model 3), non-zero values for EAX and ECX are reserved. Later processors defined ECX=1 as a valid extension (see below).

The following cause the processor to exit the implementation-dependent-optimized state: a store to the address range armed by the MONITOR instruction, an NMI or SMI, a debug exception, a machine check exception, the BINIT# signal, the INIT# signal, and the RESET# signal. Other implementation-dependent events may also cause the processor to exit the implementation-dependent-optimized state.

In addition, an external interrupt causes the processor to exit the implementation-dependent-optimized state if either (1) the interrupt would be delivered to software (e.g., if HLT had been executed instead of MWAIT); or (2) ECX[0] = 1. Implementation-specific conditions may result in an interrupt causing the processor to exit the implementation-dependent-optimized state even if interrupts are masked and ECX[0] = 0.

Following exit from the implementation-dependent-optimized state, control passes to the instruction following the MWAIT instruction. A pending interrupt that is not masked (including an NMI or an SMI) may be delivered before execution of that instruction. Unlike the HLT instruction, the MWAIT instruction does not support a restart at the MWAIT instruction following the handling of an SMI.

If the preceding MONITOR instruction did not successfully arm an address range or if the MONITOR instruction has not been executed prior to executing MWAIT, then the processor will not enter the implementation-dependent-optimized state. Execution will resume at the instruction following the MWAIT.

MWAIT for Power Management

MWAIT accepts a hint and optional extension to the processor that it can enter a specified target C state while waiting for an event or a store operation to the address range armed by MONITOR. Support for MWAIT extensions for power management is indicated by CPUID.05H.ECX[0] reporting 1.

EAX and ECX will be used to communicate the additional information to the MWAIT instruction, such as the kind of optimized state the processor should enter. ECX specifies optional extensions for the MWAIT instruction. EAX may contain hints such as the preferred optimized state the processor should enter. Implementation-specific conditions may cause a processor to ignore the hint and enter a different optimized state. Future processor implementations may implement several optimized “waiting” states and will select among those states based on the hint argument.

Table 3-67 describes the meaning of ECX and EAX registers for MWAIT extensions.

Table 3-67. MWAIT Extension Register (ECX)

Bits	Description
0	Treat masked interrupts as break events (e.g., if EFLAGS.IF=0). May be set only if CPUID.01H:ECX.MONITOR[bit 3] = 1.
31: 1	Reserved

Table 3-68. MWAIT Hints Register (EAX)

Bits	Description
3 : 0	Sub C-state within a C-state, indicated by bits [7:4]
7 : 4	Target C-state* Value of 0 means C1; 1 means C2 and so on Value of 01111B means C0 Note: Target C states for MWAIT extensions are processor-specific C-states, not ACPI C-states
31: 8	Reserved

Note that if MWAIT is used to enter any of the C-states that are numerically higher than C1, a store to the address range armed by the MONITOR instruction will cause the processor to exit MWAIT only if the store was originated by other processor agents. A store from non-processor agent might not cause the processor to exit MWAIT in such cases.

For additional details of MWAIT extensions, see Chapter 14, “Power and Thermal Management,” of *Intel® 64 and IA-32 Architectures Software Developer’s Manual, Volume 3A*.

Operation

```
(* MWAIT takes the argument in EAX as a hint extension and is architected to take the argument in
ECX as an instruction extension MWAIT EAX, ECX *)
{
WHILE ( ("Monitor Hardware is in armed state")) {
    implementation_dependent_optimized_state(EAX, ECX); }
Set the state of Monitor Hardware as triggered;
}
```

Intel C/C++ Compiler Intrinsic Equivalent

MWAIT void _mm_mwait(unsigned extensions, unsigned hints)

Example

MONITOR/MWAIT instruction pair must be coded in the same loop because execution of the MWAIT instruction will trigger the monitor hardware. It is not a proper usage to execute MONITOR once and then execute MWAIT in a loop. Setting up MONITOR without executing MWAIT has no adverse effects.

Typically the MONITOR/MWAIT pair is used in a sequence, such as:

```
EAX = Logical Address(Trigger)
ECX = 0 (*Hints *)
```

```
EDX = 0 (* Hints *)

IF ( !trigger_store_happened) {
    MONITOR EAX, ECX, EDX
    IF ( !trigger_store_happened ) {
        MWAIT EAX, ECX
    }
}
```

The above code sequence makes sure that a triggering store does not happen between the first check of the trigger and the execution of the monitor instruction. Without the second check that triggering store would go un-noticed. Typical usage of MONITOR and MWAIT would have the above code sequence within a loop.

Numeric Exceptions

None

Protected Mode Exceptions

#GP(0)	If ECX[31:1] ≠ 0.
	If ECX[0] = 1 and CPUID.05H:ECX[bit 3] = 0.
#UD	If CPUID.01H:ECX.MONITOR[bit 3] = 0.
	If current privilege level is not 0.

Real Address Mode Exceptions

#GP	If ECX[31:1] ≠ 0.
	If ECX[0] = 1 and CPUID.05H:ECX[bit 3] = 0.
#UD	If CPUID.01H:ECX.MONITOR[bit 3] = 0.

Virtual 8086 Mode Exceptions

#UD	The MWAIT instruction is not recognized in virtual-8086 mode (even if CPUID.01H:ECX.MONITOR[bit 3] = 1).

Compatibility Mode Exceptions

Same exceptions as in protected mode.

64-Bit Mode Exceptions

#GP(0)	If RCX[63:1] ≠ 0.
	If RCX[0] = 1 and CPUID.05H:ECX[bit 3] = 0.
#UD	If the current privilege level is not 0.
	If CPUID.01H:ECX.MONITOR[bit 3] = 0.