A very simple form of conditional execution can be represented by this syntax:

if condition then dosomething

condition would of course be specified in polish notation, and *dosomething* would be some kind of simple statement. The whole thing would be reduced to assembly language in these steps:

Think of a new label that's never been used before, call it L Perform jump to L if the condition is false Do the something Put the label L here.

For example:

if $\langle x 7$ then set x = + x 1

produces (on the assumption that x is our first local variable):

So it all comes down to needing a function that can read a conditional expression in polish notation and generate a jump to a given label if that condition turns out to be false.

To make the reading easier, the example expressions will be in normal infix non-polish notation, considering the various cases one by one.

How to jump to L if false is false: produce JUMP L

How to jump to L if true is false:

How to jump to L if NOT A is false: jump to L if A is true

How to jump to L if A AND B is false: jump to L if A is false jump to L if B is false

```
How to jump to L if A OR B is false:
              make up a totally new label, call it M
              jump to M if A is true
              jump to L if B is false
              produce M:
How to jump to L if A < B is false:
              use the normal polish function to get the value of A into a register
              use the normal polish function to get the value of B into the next register
              produce COMP first register, second register
              produce JCOND GEQ, L
How to jump to L if A = B is false:
              use the normal polish function to get the value of A into a register
              use the normal polish function to get the value of B into the next register
              produce COMP first register, second register
              produce JCOND NEQ, L
How to jump to L if the value of the variable X is false:
              load X into a register
              produce COMP register, 0
              produce JCOND EQL, L
etc.
```

So a function to read a polish conditional and jump to a label if it is false would look something like this:

```
void jumpifpolishfalse(istringstream & sin, int label, int reg)
{ get first symbol;
```

```
if first symbol is the reserved word "false"
    output "JUMP _L", label
else if first symbol is the reserved word "true"
    do nothing
else if first symbol is the reserved word "not"
    jumpifpolishtrue(sin, label, reg);
else if first symbol is the reserved word "and"
{ jumpifpolishfalse(sin, label, reg);
    jumpifpolishfalse(sin, label, reg); }
else if first symbol is the reserved word "or"
{ int mylabel = nextfreelabel;
    nextfreelabel+=1;
    jumpifpolishtrue(sin, mylabel, reg);
    jumpifpolishfalse(sin, label, reg);
    jumpifpolishtrue(sin, mylabel, reg);
    jumpifpolishfalse(sin, label, reg);
    output "_L", mylabel, ":" }
```

```
else if first symbol is the operator "<"
{ polish(sin, reg);
   polish(sin, reg+1);
   output "COMP R", reg, ", R", reg+1
   output "JCOND GEQ, _L", label }
else
   etc</pre>
```

L:

BUT... what would happen if a boolean expression appeared outside of its natural environment of ifs and whiles, just in a normal expression, maybe like

set x = and > a 0 <= a 10

meaning that the variable x is to be set to 1 if a is between 0 and 10, and to zero otherwise.

The expression in an assignment is always processed by the original **polish** function, so the question is really what should **polish** do if it sees an operator like **and**, **or**, **>**, **=**?

The answer is quite simple. It must produce code that would get the value 0 or 1 into the appropriate register depending on the condition, and it has a friend jumpifpolishfalse that can deal with conditional execution, so:

```
LOAD reg, 0
reserve next free label, L.
jumpifpolishfalse(sin, L, reg+1);
LOAD reg, 1
```